## Zeta function on graphs revisited with the theory of heaps of pieces

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zeta function of a graph

I hava-Selberg zeta function of a graph

Riemann zeta function

$$\leq (\Delta) = \sum_{n \geq 1} \frac{1}{n^{\Delta}}$$

for 
$$N = P_1 \cdots P_k$$

prime numbers
decomposition

Euler identity

$$\frac{2}{2}(\Delta) = \frac{1}{1-p^{-\Delta}}$$
Prime

number

I hava-Selberg zeta function of a graph

some "prime" over the graph G

some "prime" over the graph G

equivalence class of a circuit C

prime

path on X

#### permutations

$$= \begin{pmatrix} 123456789 \\ 439758162 \end{pmatrix}$$

#### path on X

$$\omega = (\Delta_0, \dots, \Delta_i, \Delta_{i+1}, \dots, \Delta_n)$$

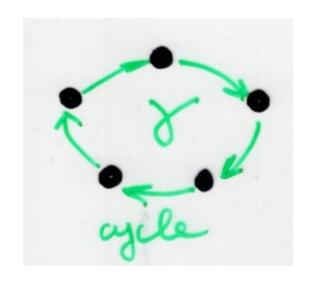
$$\Delta_i \in X \quad i=0,\dots, n$$

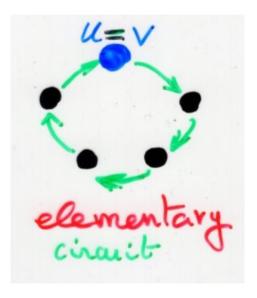
notation som

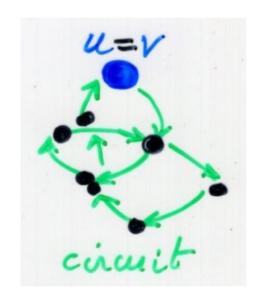
so starking vertex son ending vertex (si, six) elementary

length |w| = n (number of elementary steps) n+1 vertices

groduct two paths



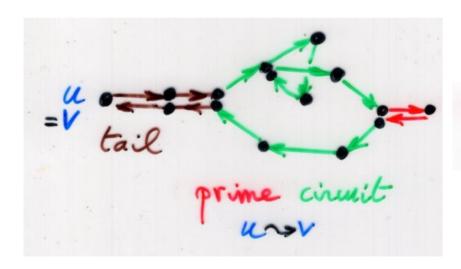




(i) 
$$\leq_6(t) = \prod_{[c]} \frac{1}{(1-t^{[c]})}$$

equivalence class prime circuit

no backtracking



· back tracking

(-no back tracking

I hava-Selberg zeta function of a graph

(i) 
$$\leq_6(t) = \frac{11}{(1-t^{|c|})}$$

(ii) 
$$\zeta_{G}(t) = \frac{1}{\det(4-Ht)}$$

(iii) 
$$\zeta_{6}^{(t)} = \frac{1}{(1-t^{2})^{m-n}} \frac{1}{\det(I-tA+t^{2}(D-I))}$$

Bass formula

Bass (1992) Hashimoto (1989) Venkou, Nikitin (1994) Sunada (1986,88)

Stark, Terras (1996, 2000) book Northshield (1999) Foata, Zeiberger (1999) lyective proof

Bartholdi (1999)

Mizumo, Sato (2000, ..., 2009)

many others

- quantum walks

extending number theory to paths on Graphs

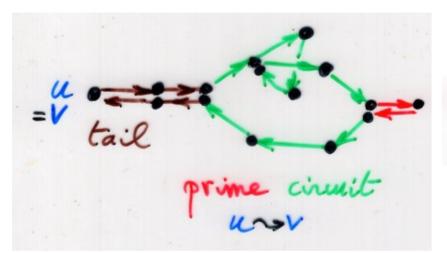
(i) 
$$\frac{1}{(1-t^{|c|})} \log \zeta_{G}(t) = \sum_{[c]} \sum_{PM} \frac{1}{P} t^{|c|}$$

$$t d \log \zeta_{6}(t) = \sum_{[C]} \sum_{p>1} |c| t^{p|C|} equivalence class circuit no back tracking$$

no backtracking

$$= \sum_{[c]} |c| t^{|\omega|}$$

td log < 5 (t) = \( \tau \) (-no tail (-no back tracking



(ii) 
$$\zeta_{G}(t) = \frac{1}{\det(4-Ht)}$$

(iii) 
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#### Ramanujan graphs

spectrum of a graph G

of A digacency matrix

spectrum of A related to the poles of the zeta function  $Z_G(t)$ 

G is Ramanyan  $\iff$  the poles of  $\underset{\leftarrow}{\swarrow}(q^{-s})$  occur only on  $\underset{\leftarrow}{\lessgtr} R(s) = 1/2 \underset{\leftarrow}{\lessgtr}$  and at  $\underset{\leftarrow}{\i}$ 

Ramanujan graph

$$G$$
  $(q+1)$ -regular  $\Rightarrow$   $(q+1)$  is an eigenvalue

G connected, these are the unique largest eigenvalue

(Alon-Boppana) as IVI inveases, vertices tend to be constrained in [-2\sq , 2\sq ].

G (q+1)-regular graph is called Ramanyan if all its eigenvalues lie in the interval  $\left[-2\sqrt{q}, 2\sqrt{q}\right]$ , except possibly  $\pm (q+1)$ 



first infinite family (Lubotsty, Phillips, Sarnak) was constructed using Ramanyain conjecture on coefficients of the Delekind y-functions

$$\eta(\tau) = q^{1/24} \prod_{n=1}^{\infty} (1-q^n)$$
 $q = e^{2\pi i \tau}$ 

Dedekind eta function

$$\eta(\tau) = q^{1/24} \prod_{n=1}^{\infty} (1-q^n)$$
 $q = e^{2\pi i \tau}$ 

Dedekind eta function

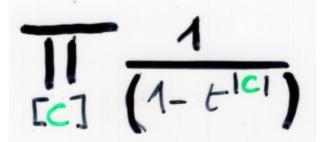
$$\Delta(\tau) = (2\pi)^{12} (\eta(\tau))^{24}$$
discriminant

$$\frac{1}{(2\pi)^{42}} \Delta(\tau) = 9 \pi (1-9^{11})^{24}$$

$$=\sum_{n=1}^{\infty} T(n)q^n$$

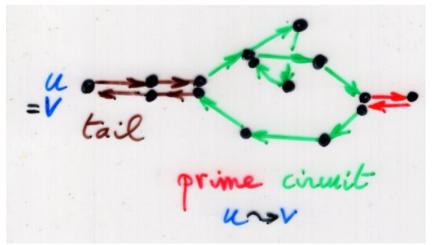
> Ramanyan a conjecture 11/2





equivalence class prime circuit

no backtracking



Paths in graphs and linear algebra

Lemma 
$$X = \{1, 2, ..., k\}$$

$$A = (a_{ij})_{1 \le i,j \le k} \quad \text{matrix}$$

$$(I - A)^{-1} = \sum_{\alpha} v(\alpha)$$

$$\text{path on } S$$

$$\text{with } v(i,j) = a_{ij}$$

$$\sum_{i \in \mathcal{N}_j} V(\omega) =$$

(i) 
$$\leq_{6}(t) = \frac{11}{(1-t^{-101})}$$

equivalence class prime circuit

no backtracking

(ii) 
$$\zeta_{G}(t) = \frac{1}{\det(1-Ht)}$$

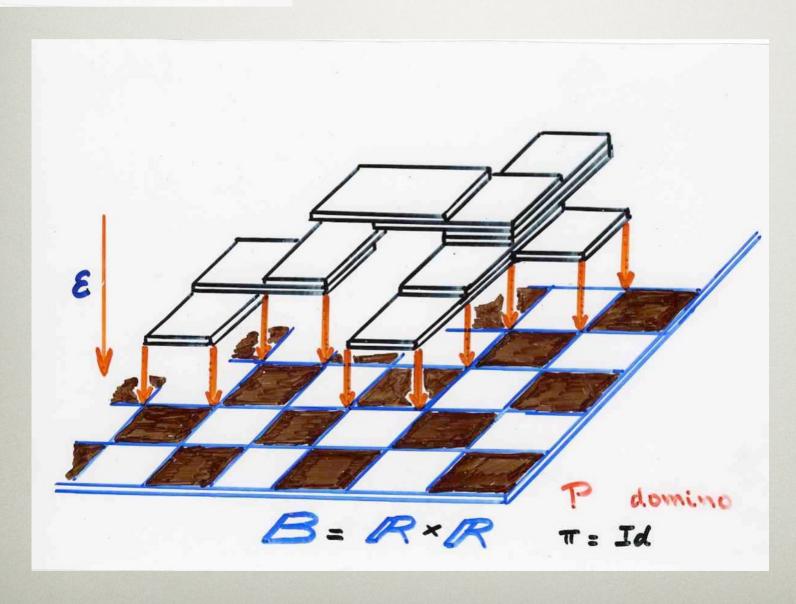
(iii) 
$$\zeta_{6}^{(t)} = \frac{1}{(1-t^{2})^{m-n}} \frac{1}{\det(I-tA+t^{2}(D-I))}$$

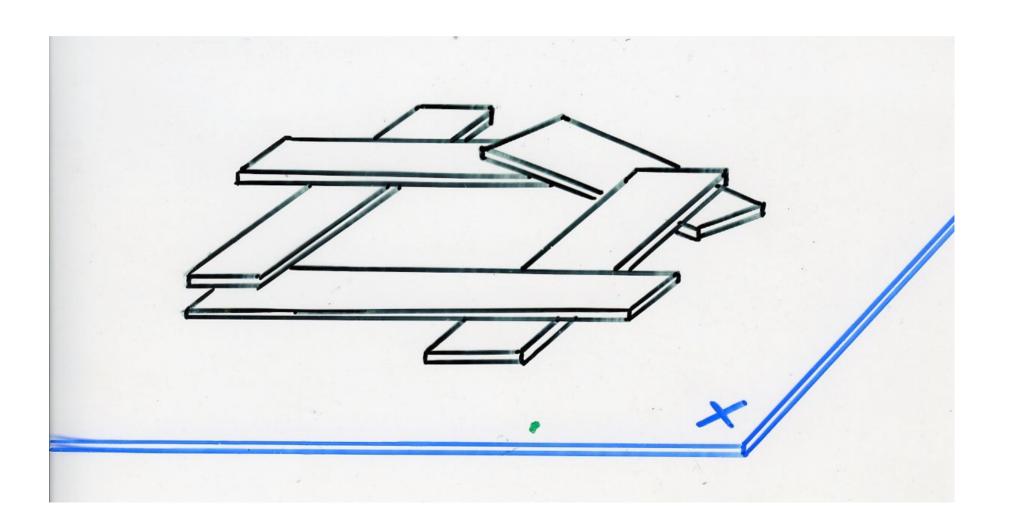
# det (I-A)?

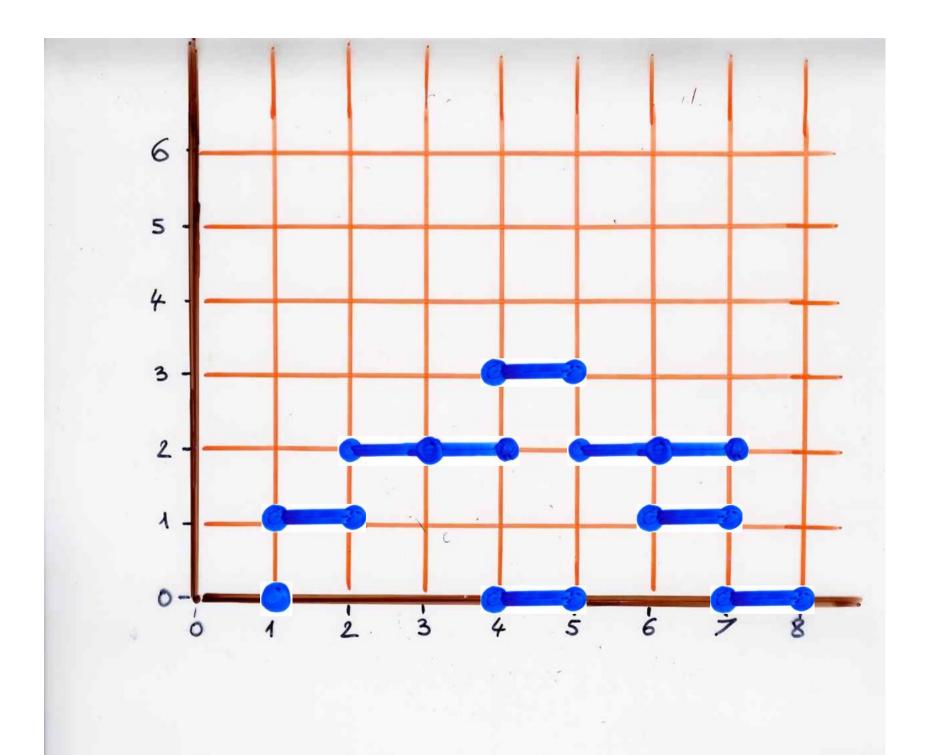


Heaps of pieces: definition with levels

### Introduction







heap definition • P set (of basic pieces) binary relation on P symmetric (dependency relation) heap E, finite set of pairs

(d, i) & EP, i & N (called pieces)

projection level (i) (ii)

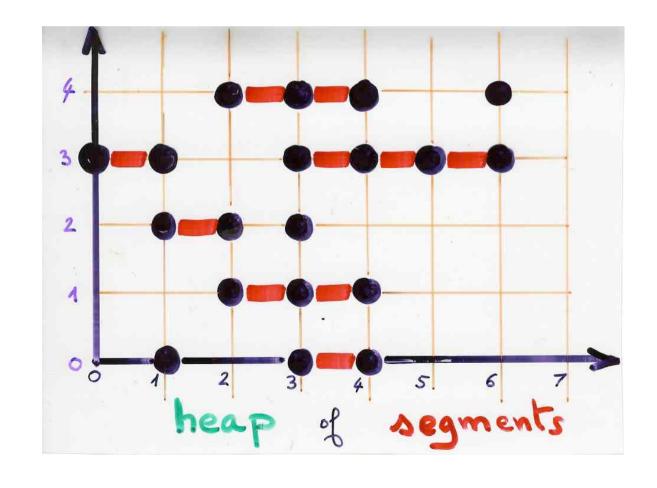
heap definition • P set (of basic pieces) binary relation on P symmetrice (dependency relation) heap E, finite set of pairs

(d, i) & EP, i & N (called pieces)

projection level (i) (a,i), (B,j)∈ E, ~ ℃ B ⇒ i ≠ j (ii) (d, i) ∈ E, i>0 => ∃r∈P, abp, (B, i-1) E E

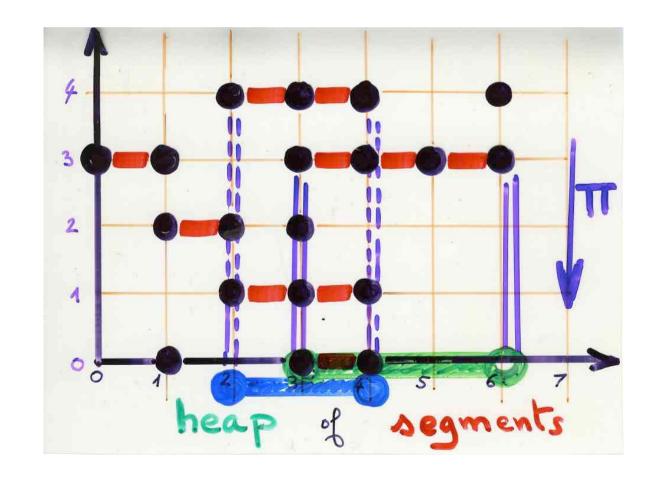
ex: heap of segments over IN

$$P = \{ [a,b] = \{a,ad,...,b\}, 0 \le a \le b \}$$
 $\{a,b\} \in [a,d] \iff [a,b] \cap [a,d] \neq \emptyset$ 

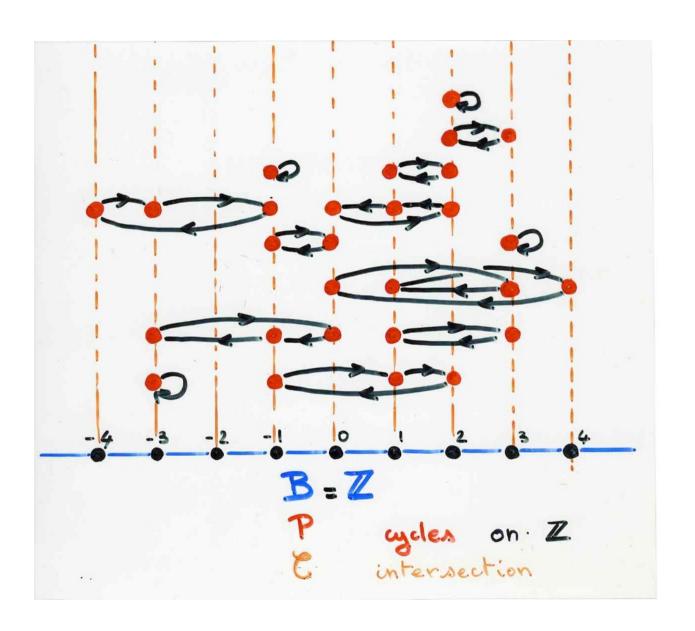


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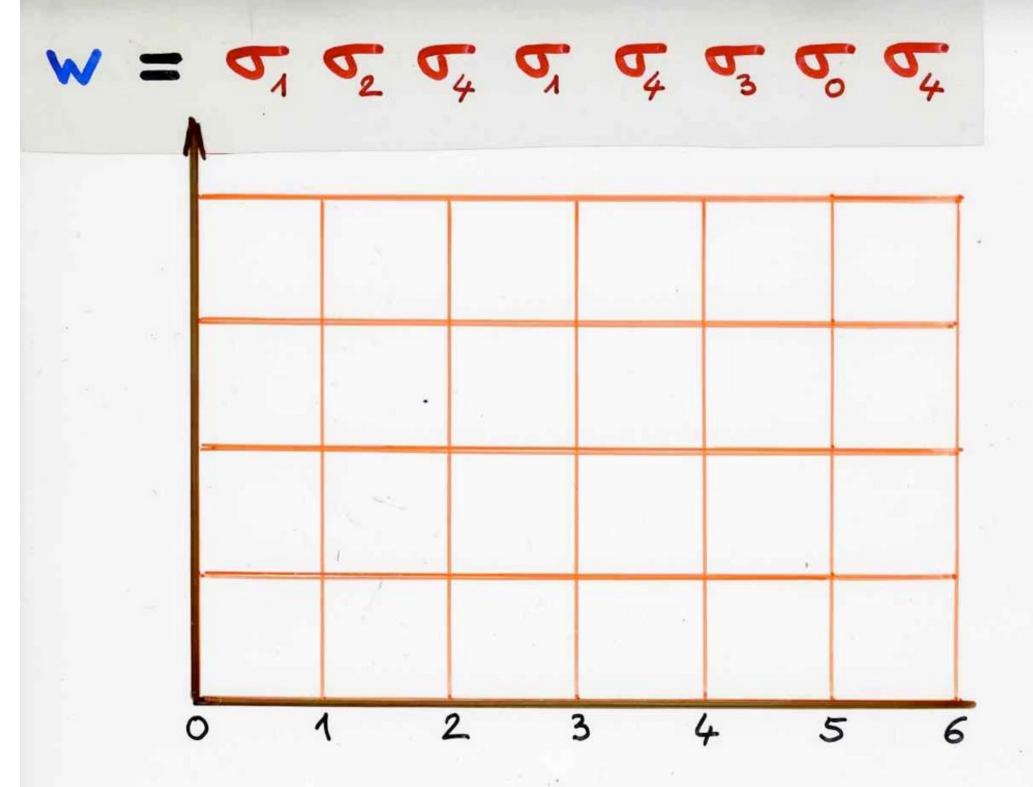


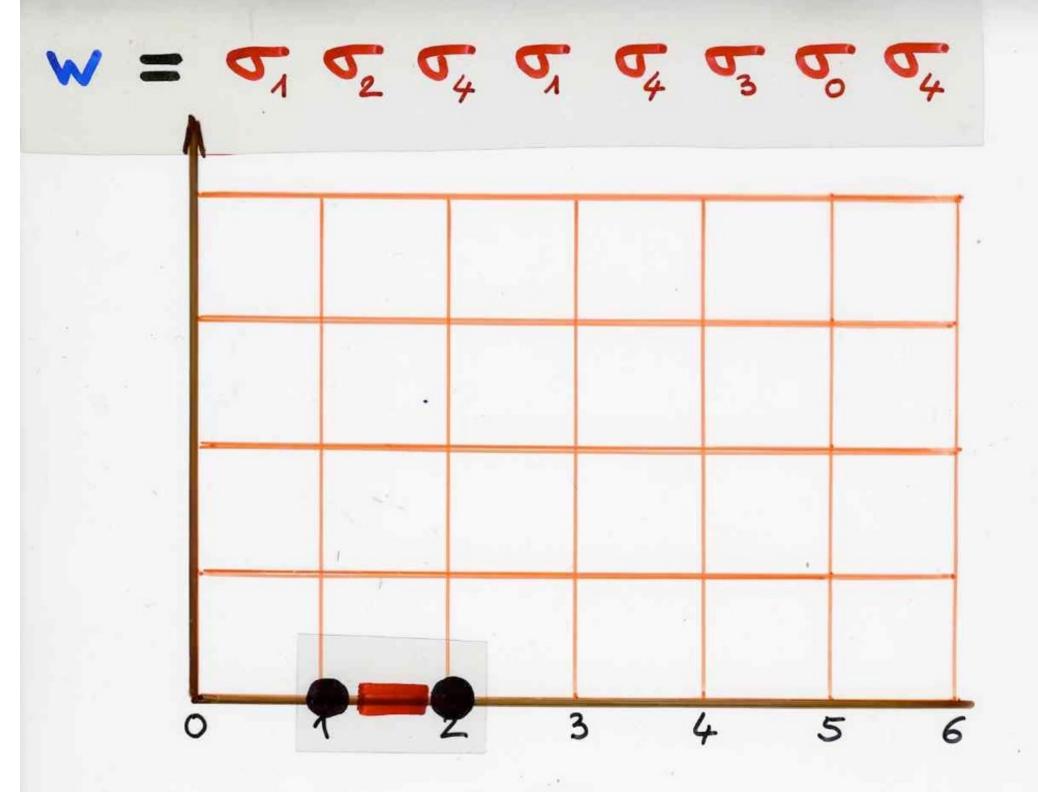
## basic pieces P = { cycles on Z} Supp(8) = {-7,-3, 2,5,9} 8 6 5 supp (8) A Supp (8) # \$

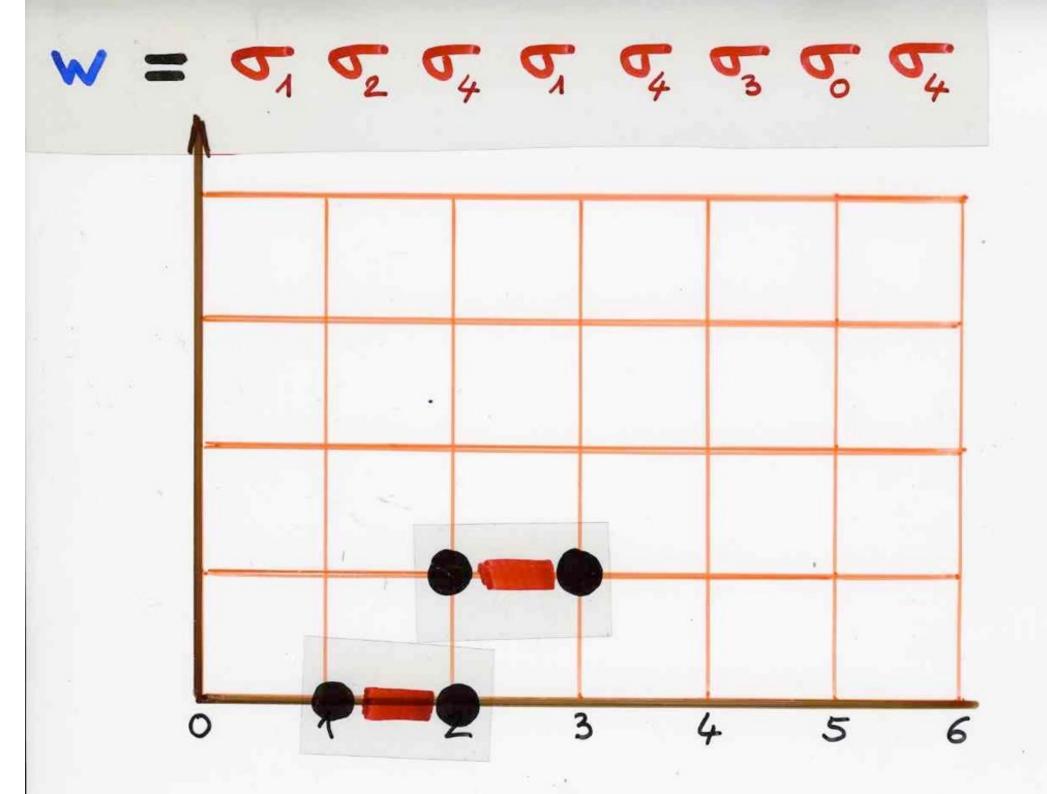


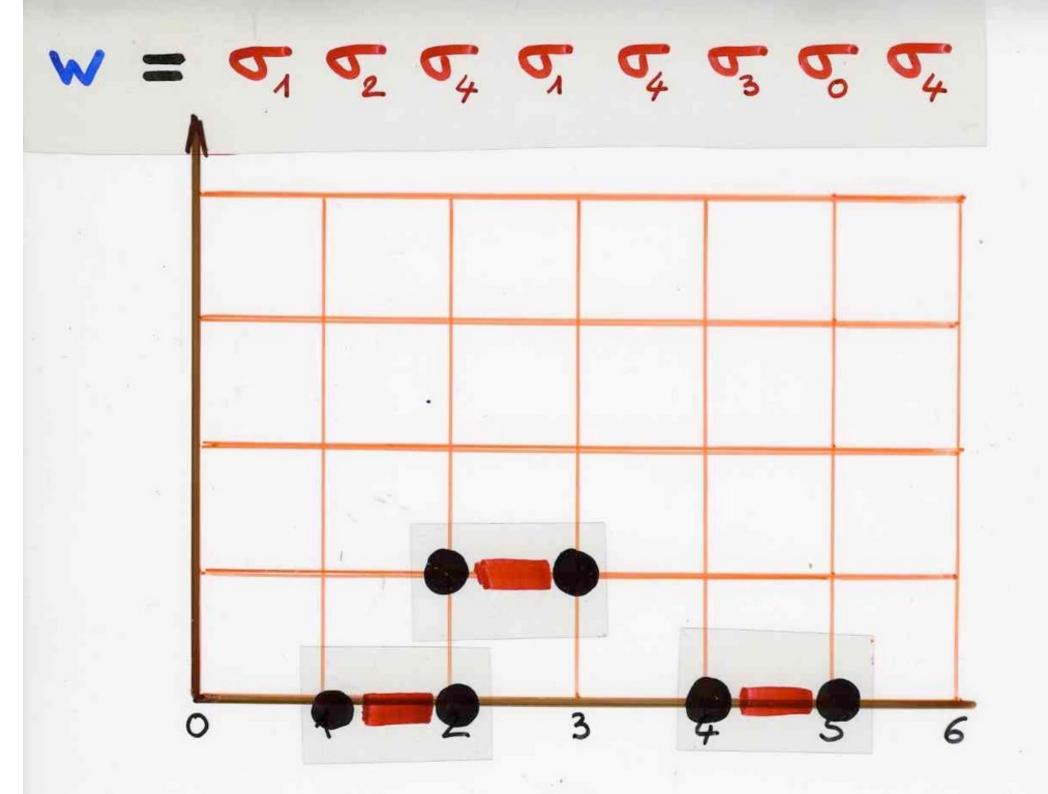
Heaps of pieces and commutation monoids

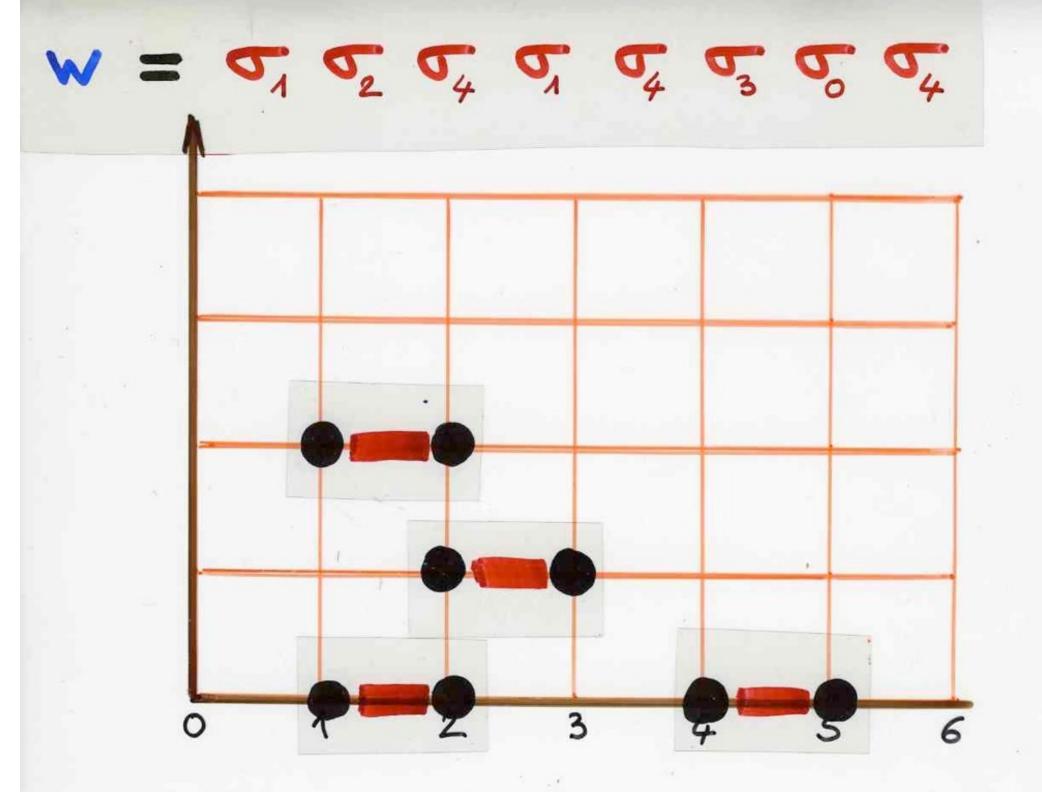
ex: heaps of dimers on  $\mathbb{N}$   $P = \{ [i, ai] = \sqrt{i}, i \neq 0 \}$ 

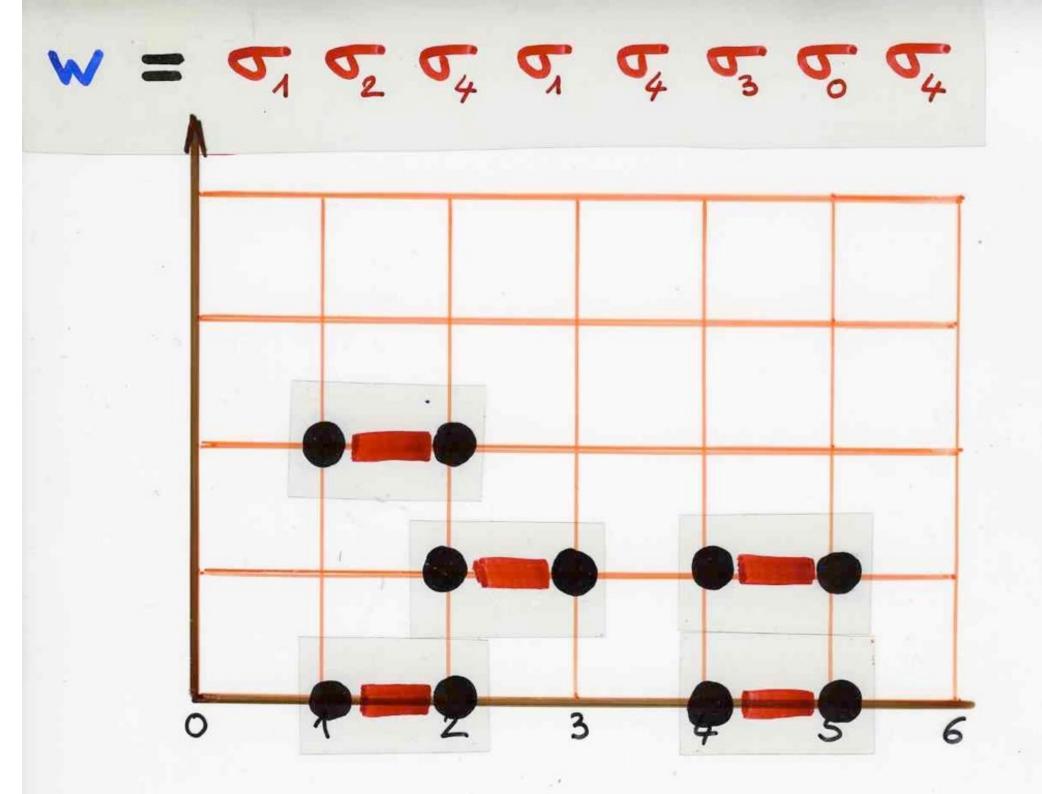


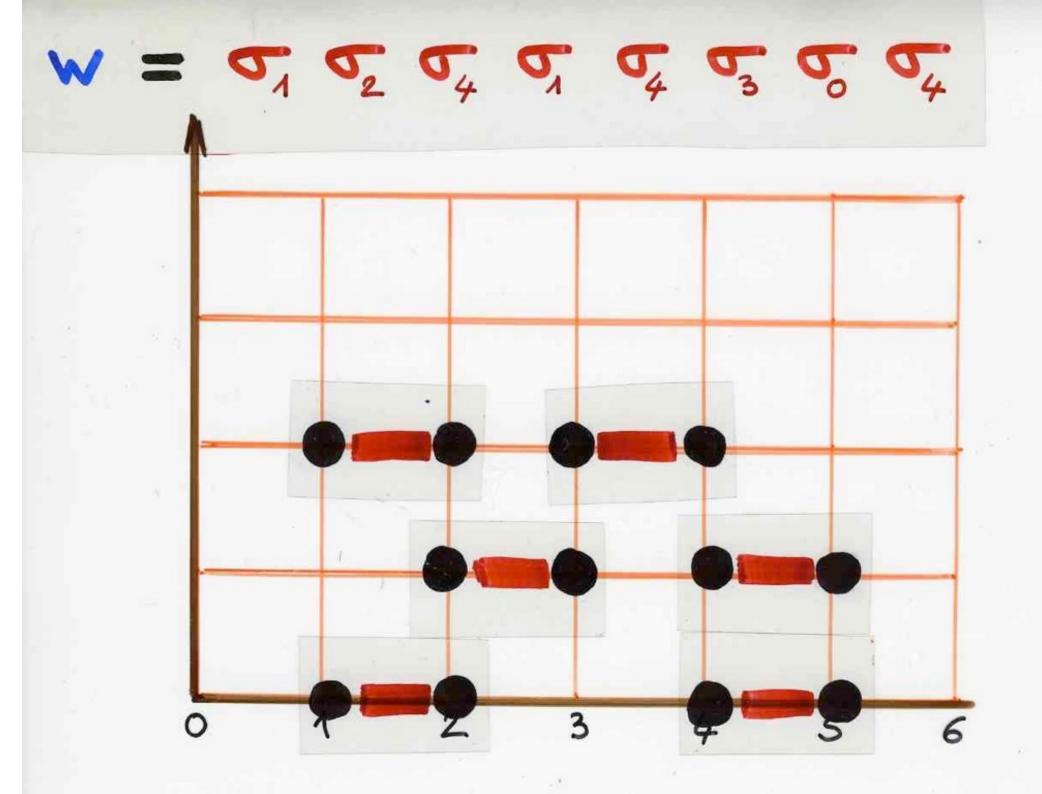


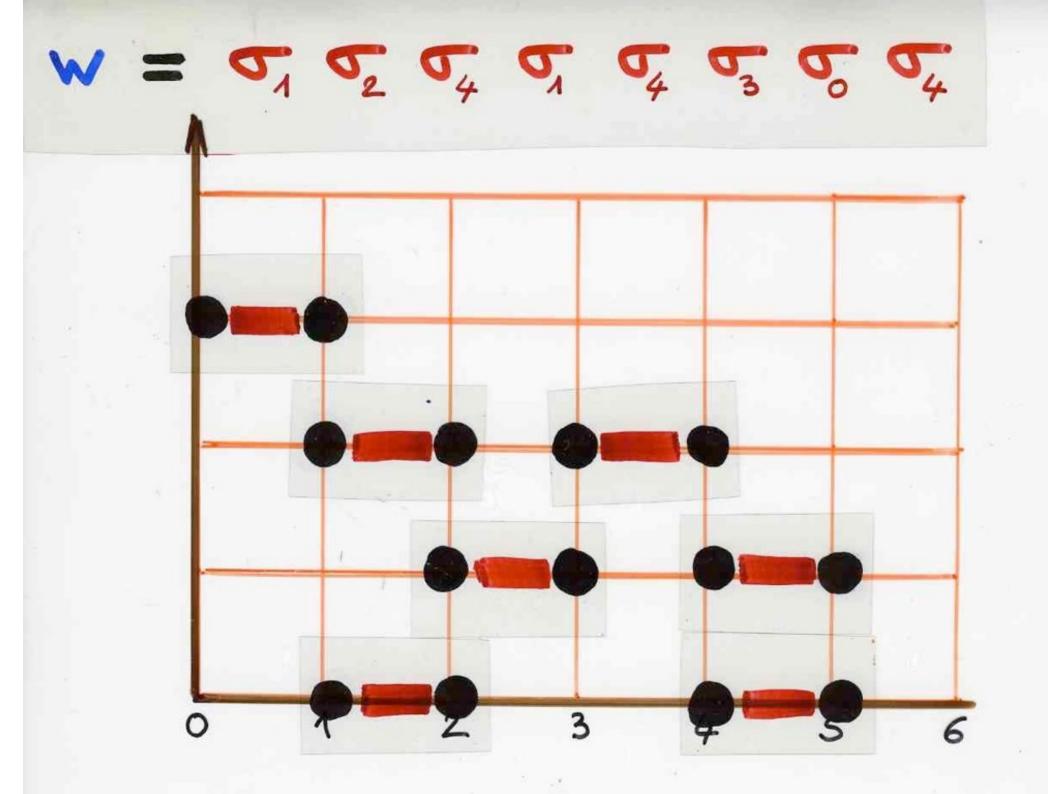


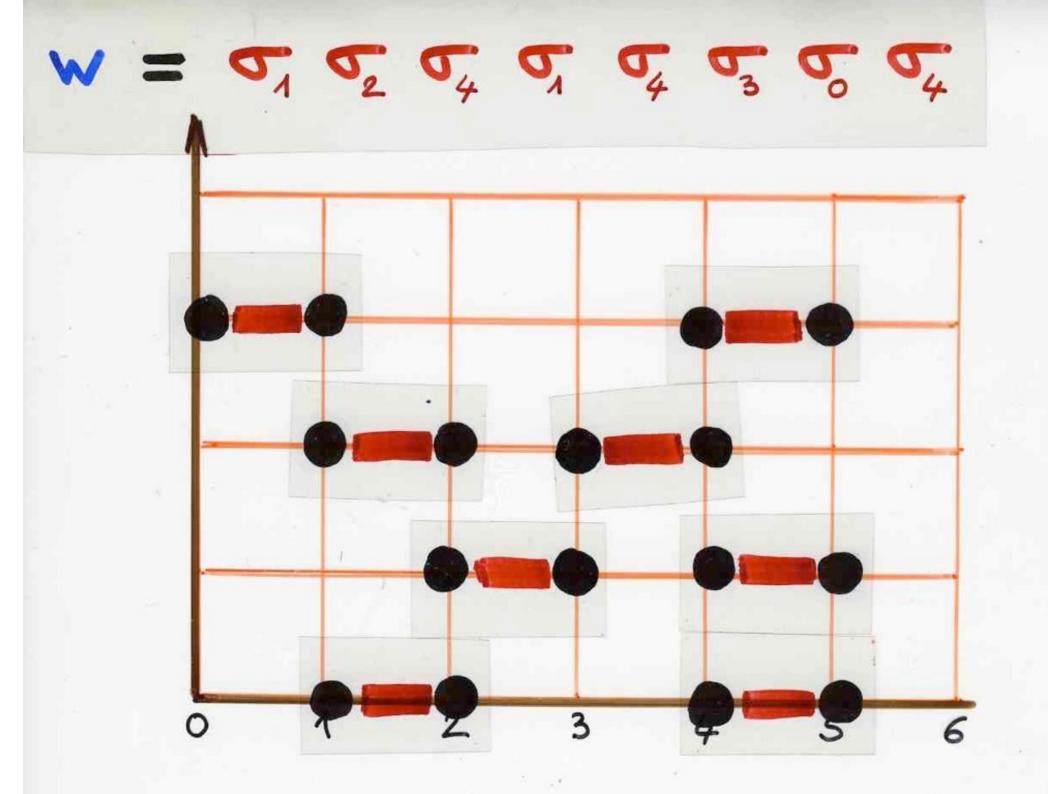


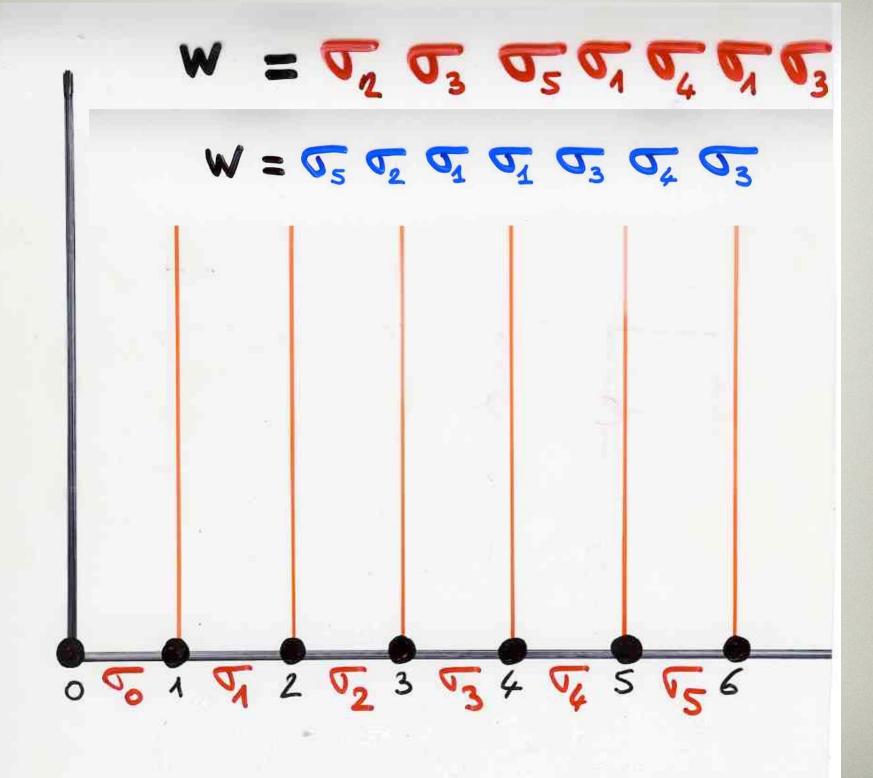


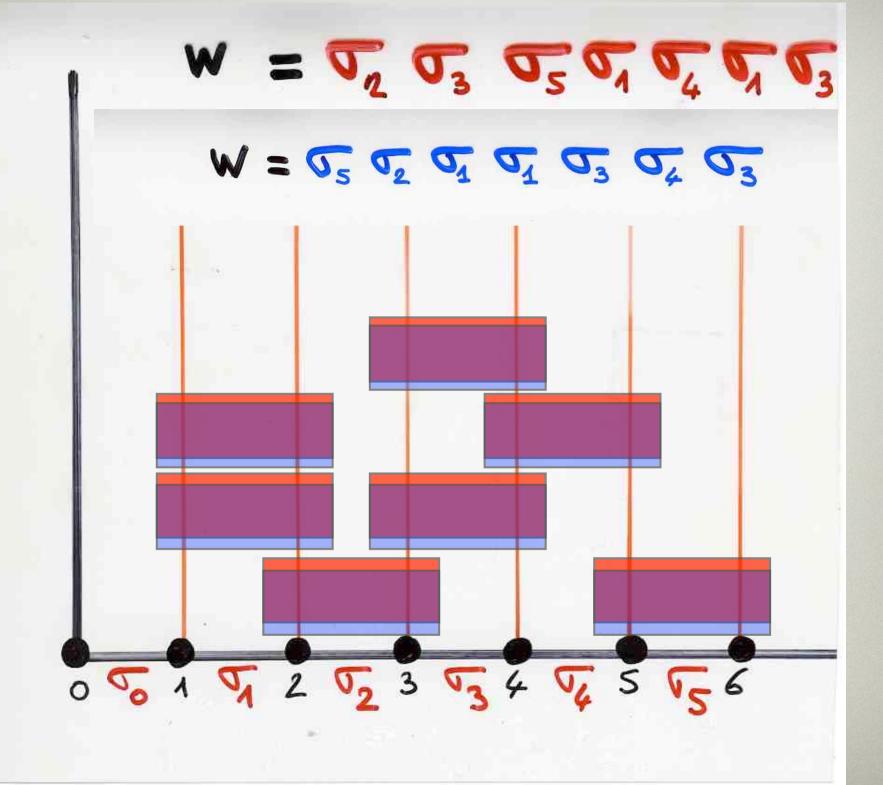








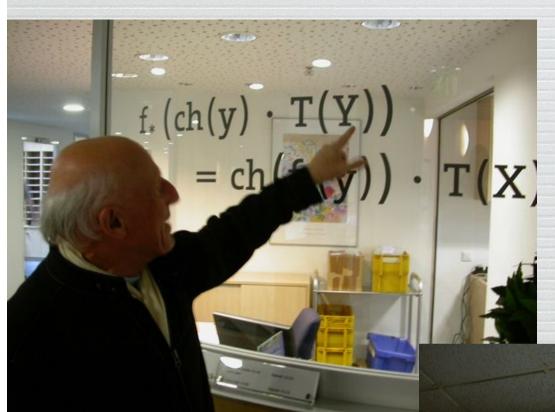




ex: heaps of dimers on  $\mathbb{N}$ P = { [i,41] =  $\sqrt{i}$ , iso}

C commutations

Ti  $\sqrt{i}$  =  $\sqrt{i}$  iff |i-j|>2



Cartíer-Foata commutation Monoids (1969)

Heaps of pieces (X.V. 1985)



Trace monoids

Computer Science

model for parallelism

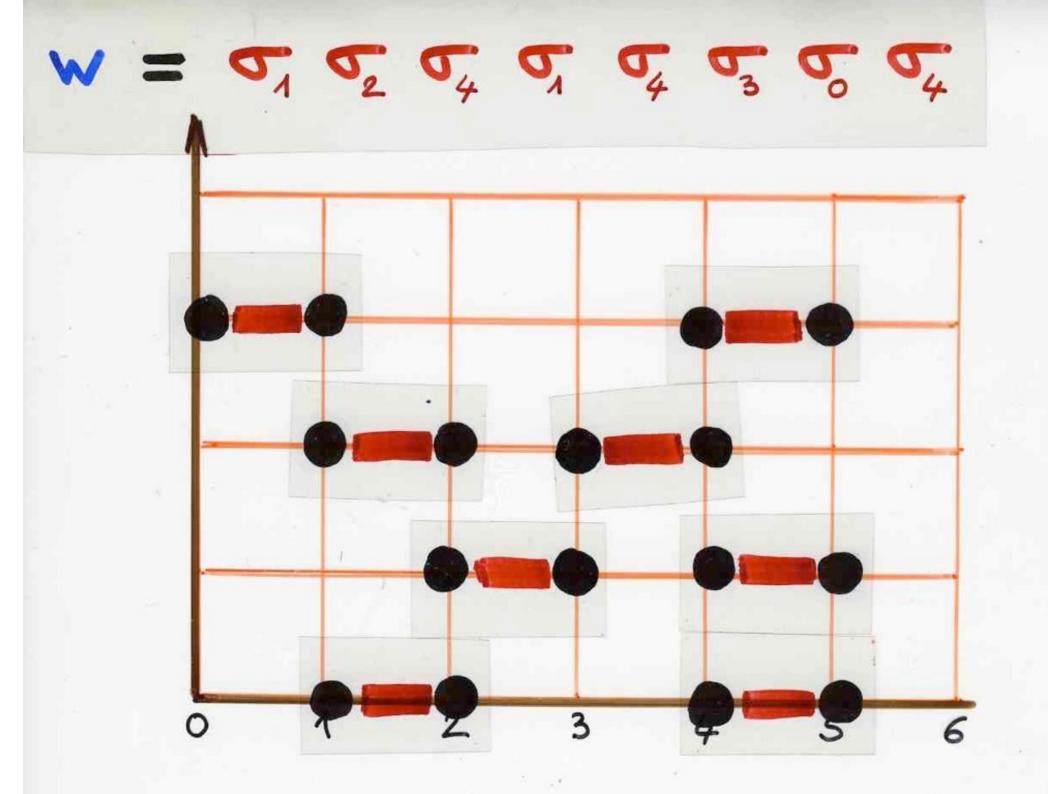
concurrency access to

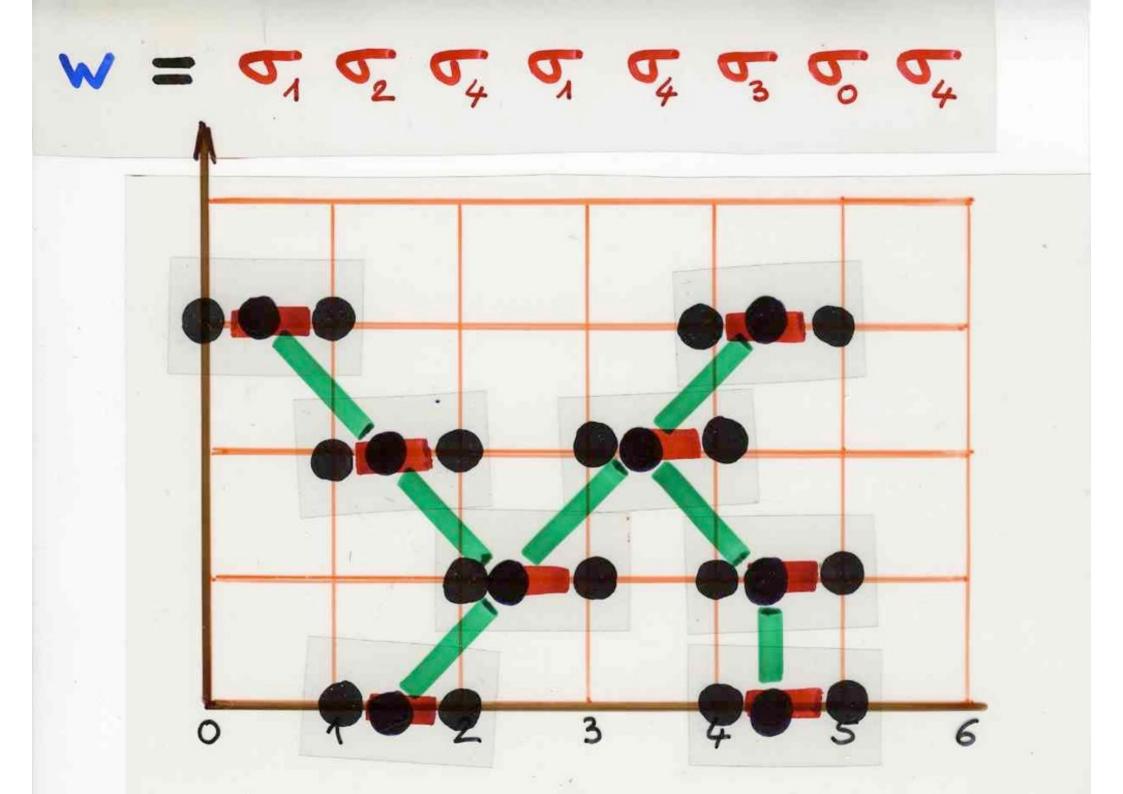
data structures

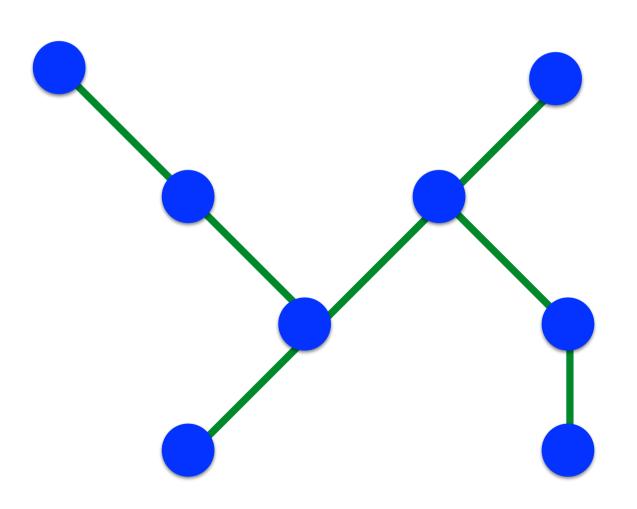
Trace

Mazurkiewicz (1977)
model of the legical behavior
of safe Petri nets

Diekert, Rosenberg ed. (1995) The book of traces Heaps of pieces as a poset



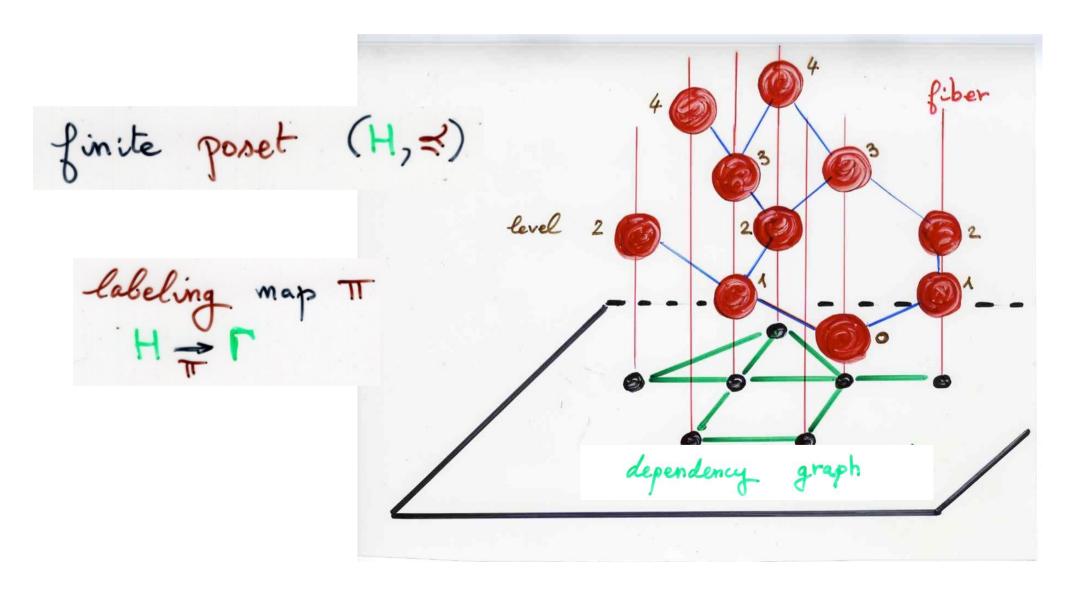




Second definition: heaps as a poset on a graph

$$\Gamma=(S,E)$$

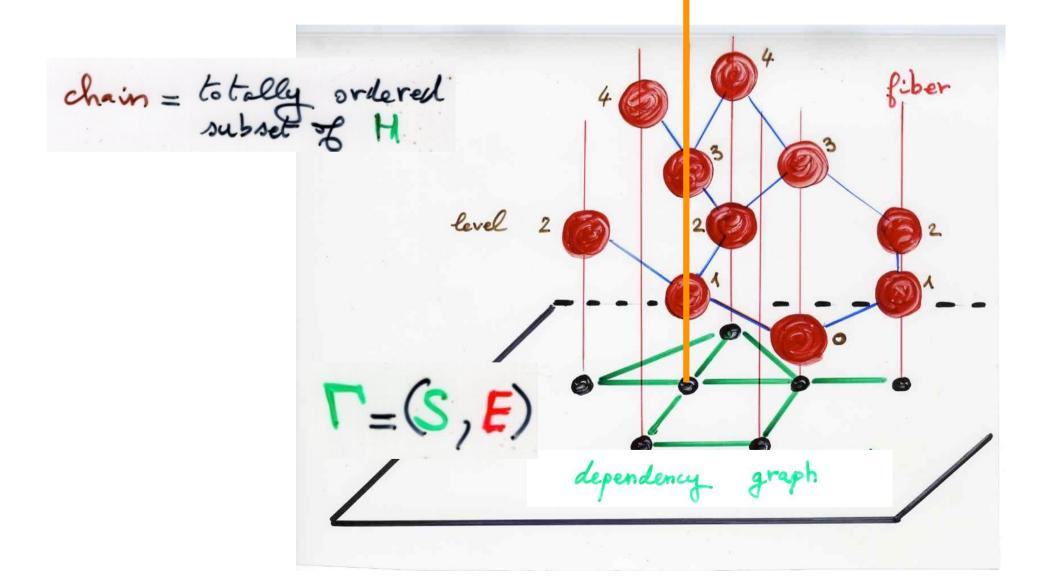
is the dependency graph



for every vertex SES

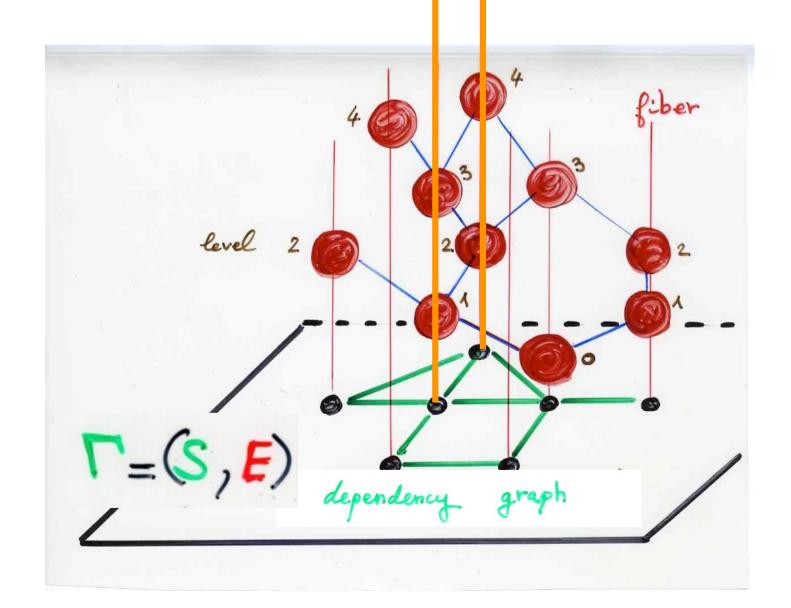
Hs = TF-1 (753) is a chain

fiber over DES



for any edges is, if of [
H<sub>s,t</sub> = TF-1(is,ti) is a chain

fiber over 7s, to



The order relation  $\leq$  is the transitive closure of the relations given by all chains of (i)!

Hs Hs,t

(i.e. the smallest partial ordering containing these chains)

## Algebraic graph theory revisited with heaps of pieces

an example:

chromatic polynomial and and acyclic orientations of a graph

chromatic polynomial

number of (proper) coloring of the graph 6 with \ colors

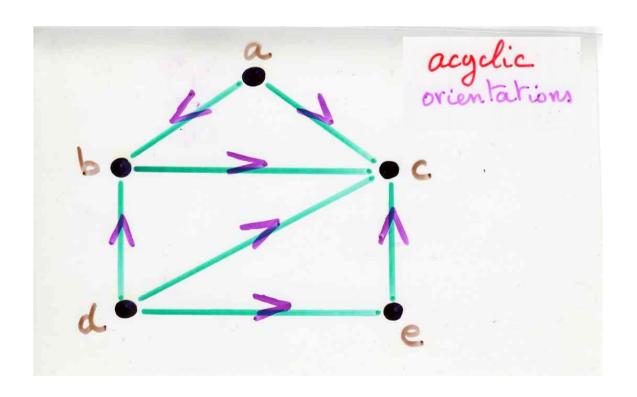


$$\alpha(G)$$

acyclic orientations of 6

Proposition (Stanley, 1973)
$$\alpha(G) = (-1)^{n(G)} \chi_{G}(-1)$$

Troposition (Stanley, 1973)
$$\alpha(G) = (-1)^{n(G)} \chi_{G}(-1)$$

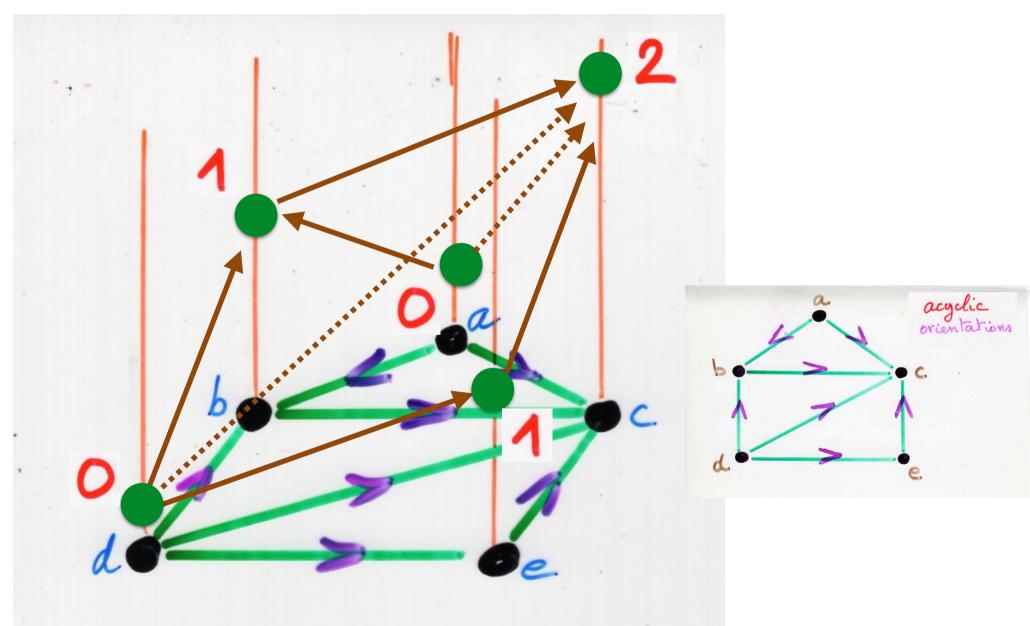


commutation (Carrier-Foata) monoid

from Gessel (1985)?

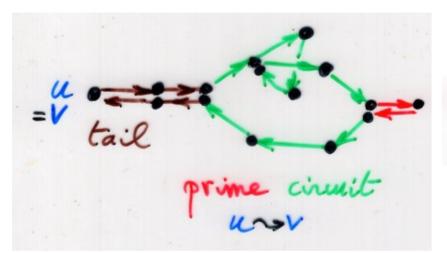
Bijection





3 basic lemma on heaps

td log < 5 (t) = \( \tau \) (-no tail (-no back tracking



(ii) 
$$\zeta_{G}(t) = \frac{1}{\det(4-Ht)}$$

(iii) 
$$\zeta_{6}^{(t)} = \frac{1}{(1-t^{2})^{m-n}} \frac{1}{\det(I-tA+t^{2}(D-I))}$$

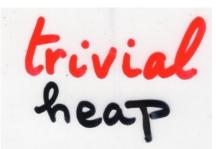
- 3 basic lemma:
- Inversion lemma
- Logarithmic lemma
- circuit = heap of cycles

## First basic lemma on heaps: the inversion lemma

1/D

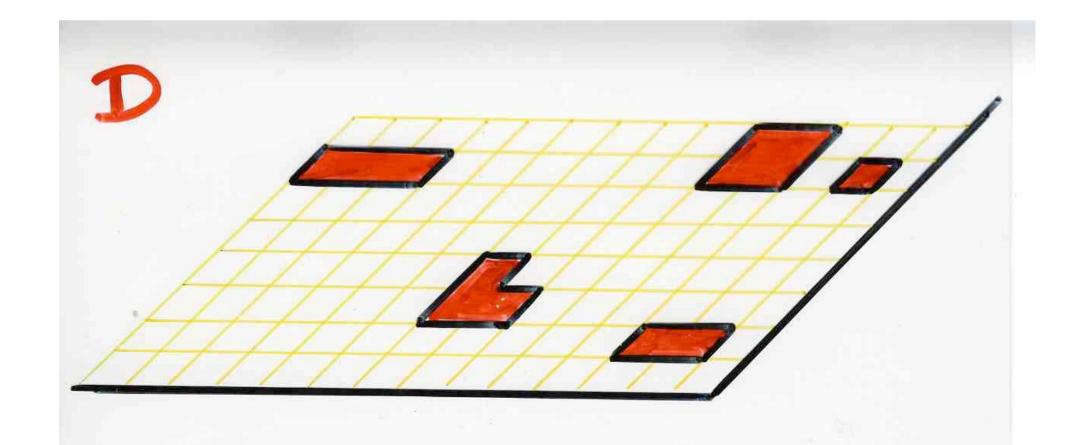
## the inversion lemma

all pieces (d, i) at level 0





all pieces (4,i) at level 0



valuation

$$V: P \longrightarrow K[x,y,...]$$

lasic

piece

$$V(E) = \prod V(\alpha i)$$

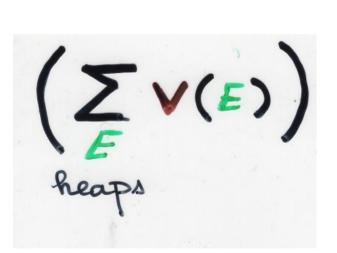
heap  $(\alpha,i) \in E$ 

# the inversion lemma

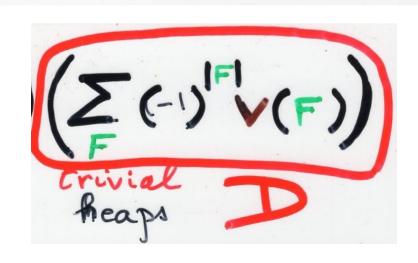
1

(Z(-1)FV(F)) trivial fleaps

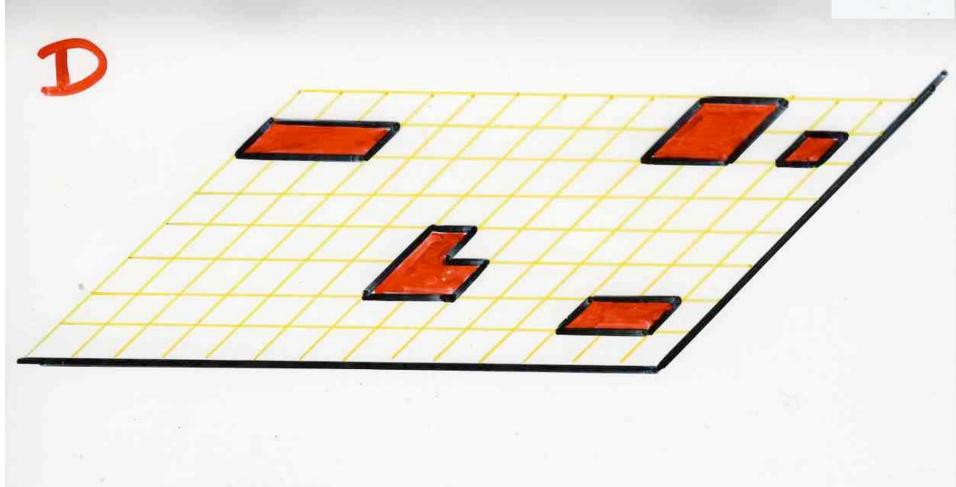
## the inversion lemma

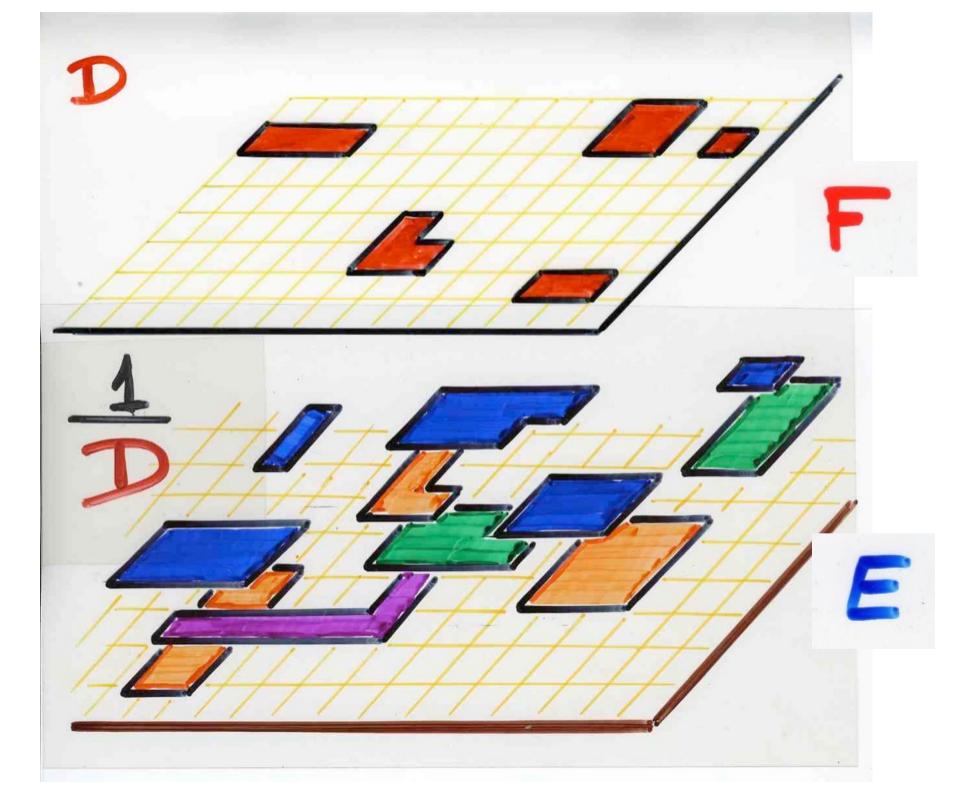


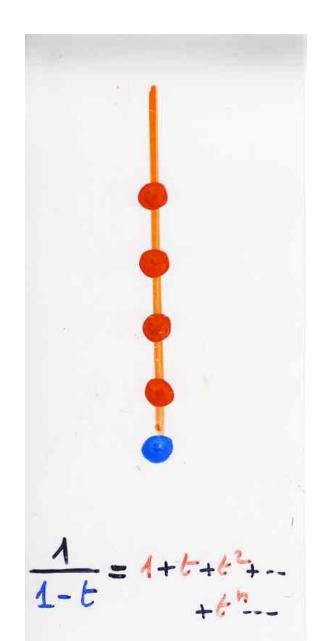
1











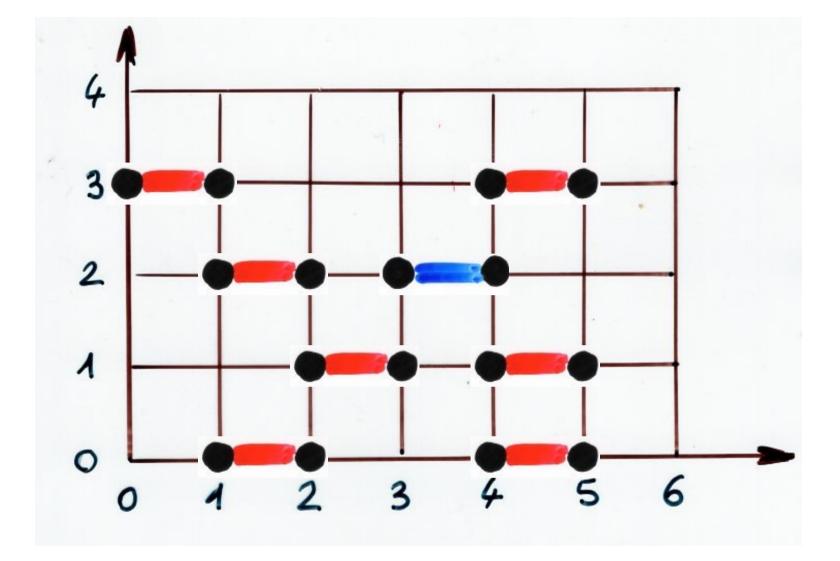
$$\frac{1}{(1-\frac{t_{1}}{4})(1-\frac{t_{2}}{5})(1-\frac{t_{3}}{5})(1-\frac{t_{4}}{5})} = \sum_{\substack{d_{1},d_{2},d_{3},\\d_{3},d_{4},d_{5},\\d_{5},d_{5},d_{5},\\d_{4},d_{5},d_{5},\\d_{4},d_{5},d_{5},d_{5},d_{6},d_{5},\\d_{5},d_{$$

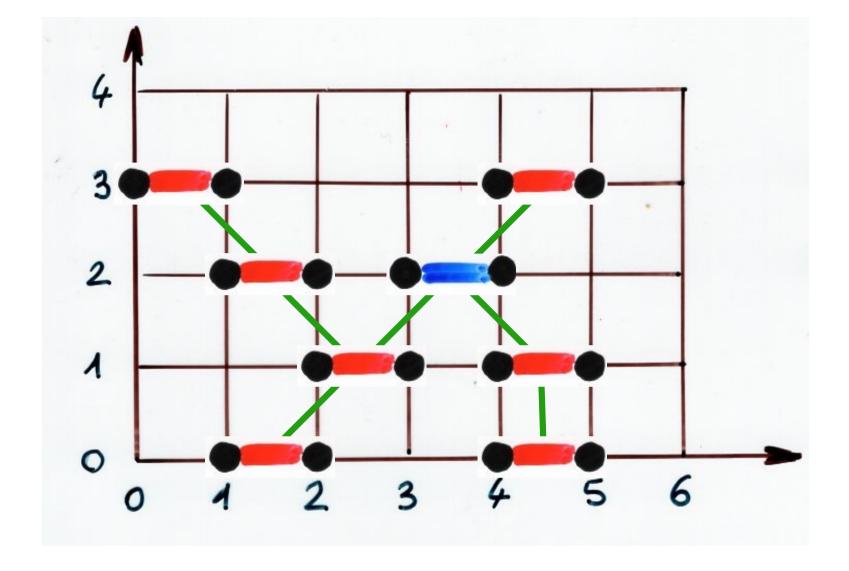
the logarithmic lemma

weight

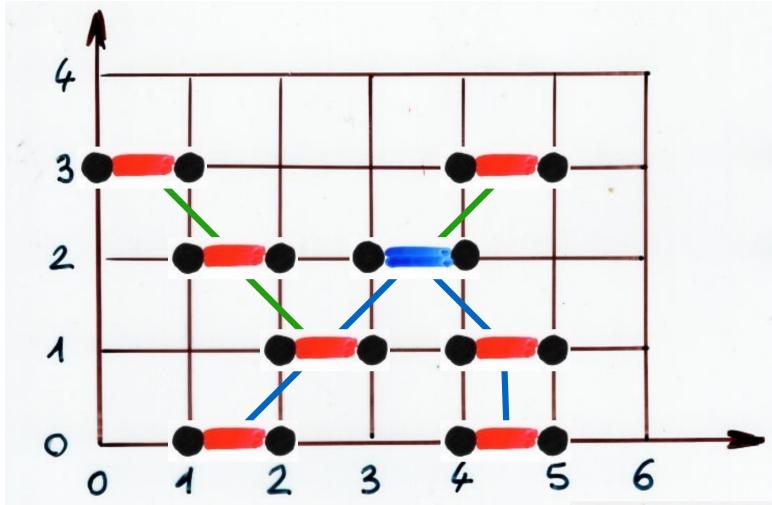
The logarithmic. Lemma

$$\frac{t}{dt} \frac{d}{dt} \left( \sum_{k \in A} V(E) t^{|E|} \right) = \sum_{k \in A} V(P) t^{|P|}$$
Pyramid



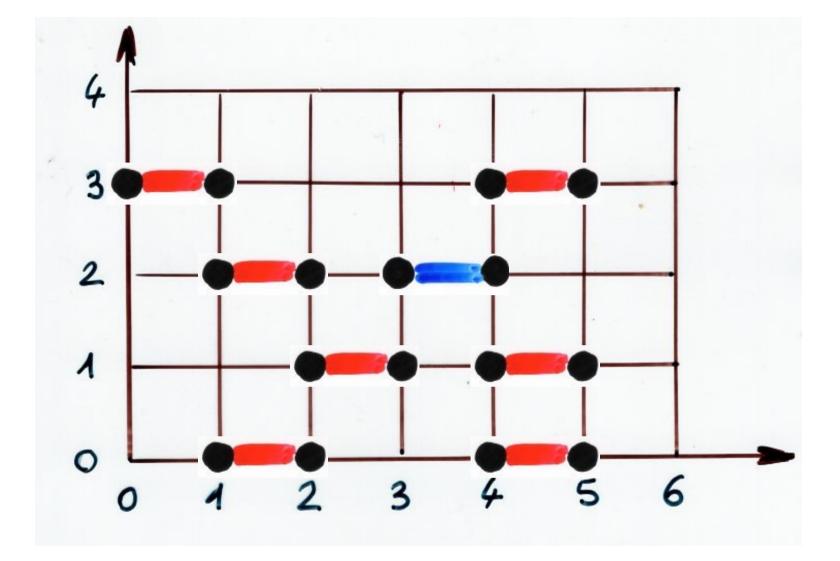


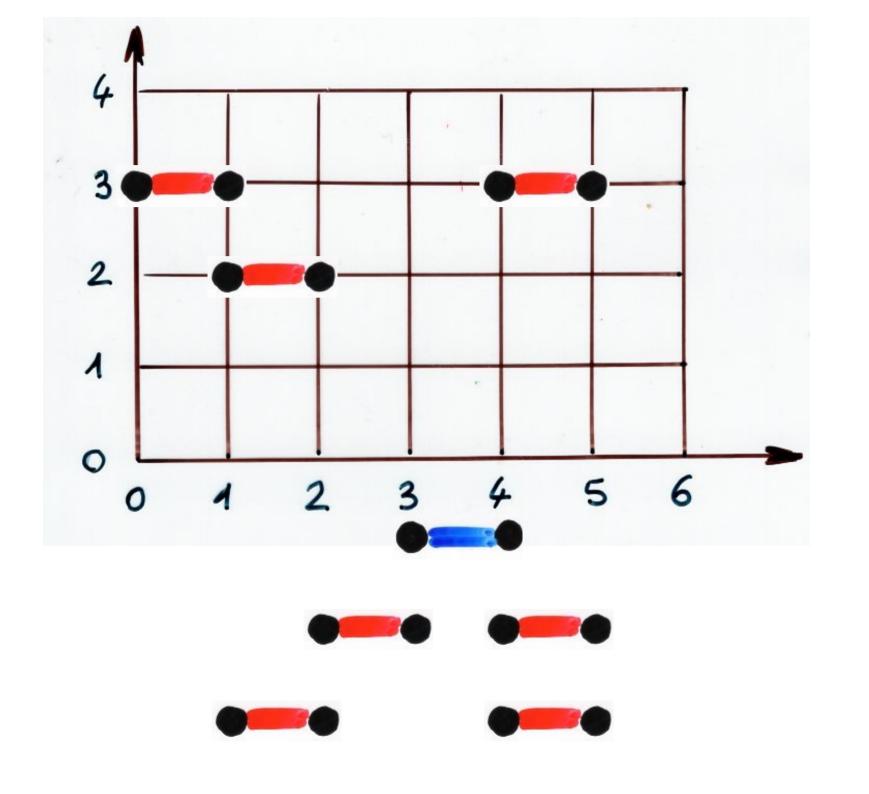
poset ing the heap E

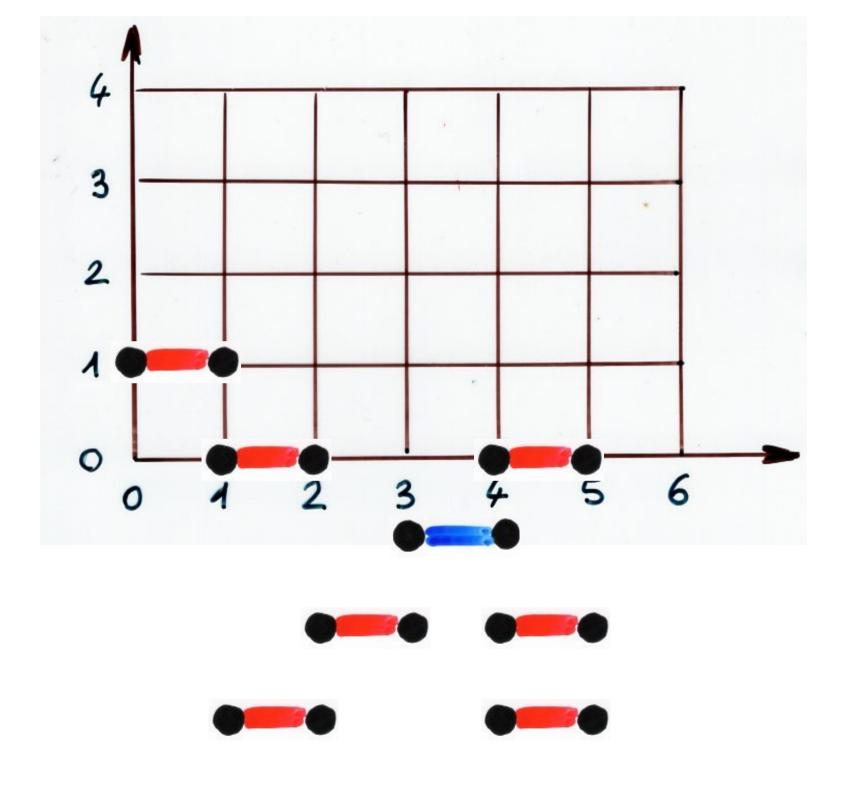


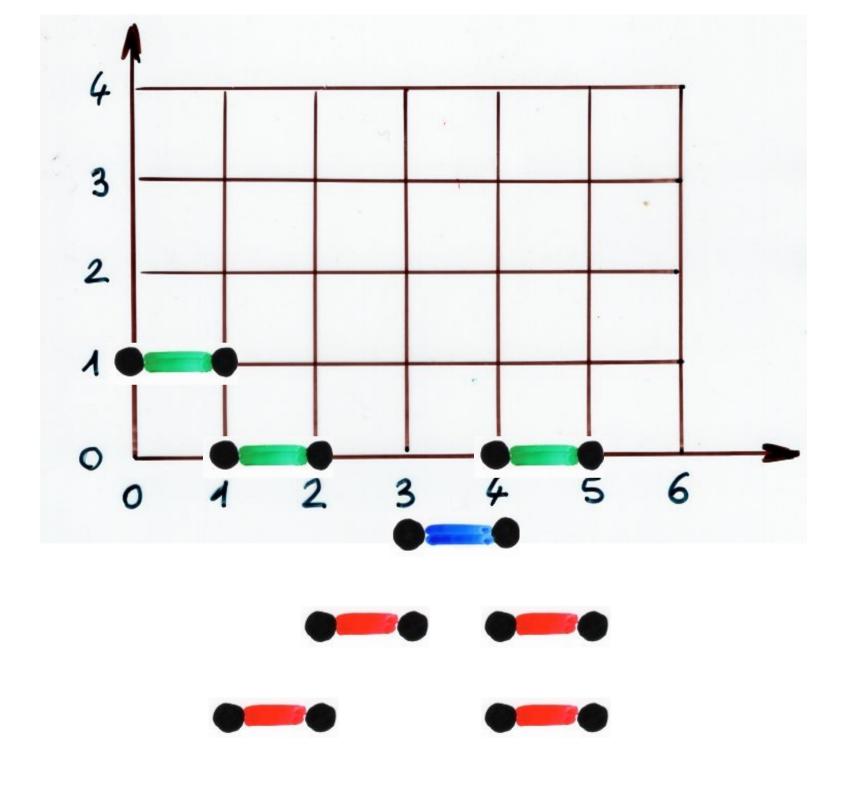
XEE ideal generated by x

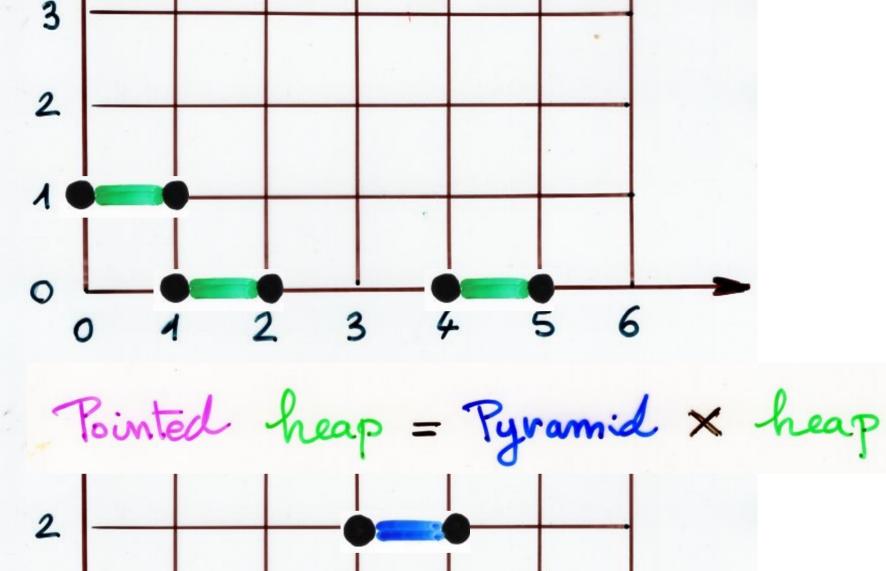
pyramid

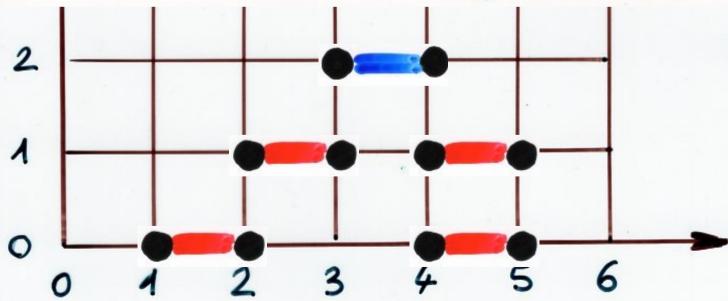












ty' = = 4

Pointed heap = Pyramid x heap

$$z = \sum_{v(P)} v(P)$$

Pyramid

$$\frac{ty'}{y} = z$$

1E1 of elements

= \(\text{V(P)} t^{\text{P}}\)

Pyramid

Pyramid

The logarithmic Lemma

#### The logarithmic Lemma

equivalent form

$$log(\sum_{k \in A} V(E) t^{|E|})$$
also:

$$\frac{\log \left(\sum V(E) t^{|E|}\right)}{\text{Reap}} = \sum V(P) \frac{t}{|P|}$$
Pyramid

Pyramid

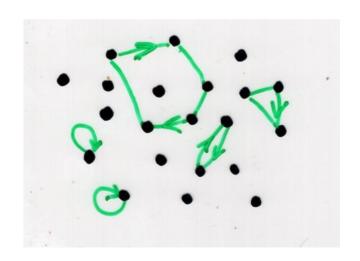
## Interpretation of

(ii) 
$$\zeta_{G}(t) = \frac{1}{\det(4-Ht)}$$

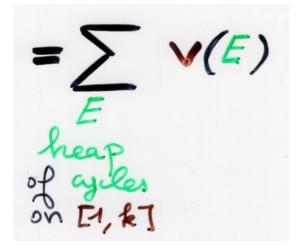
$$A = (a_{ij})_{1 \le i,j \le k}$$

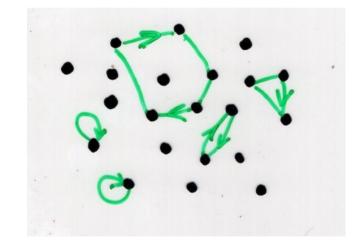
$$\sum_{(-1)}^{(-1)} a_{1}\sigma(a) \cdots a_{k}\sigma(k)$$

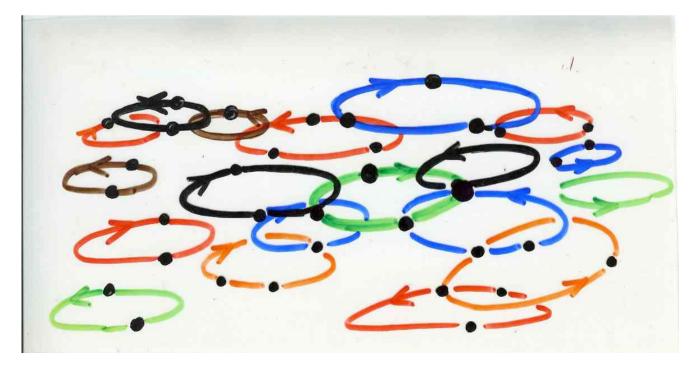
$$\sigma \in G_{k}$$
permutation

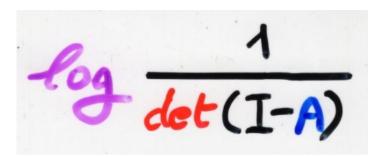


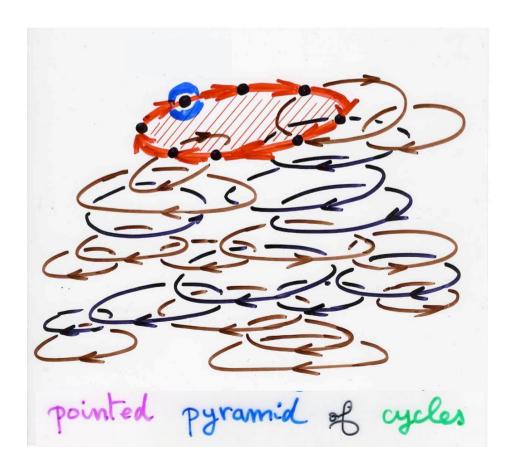
#### inversion lemma

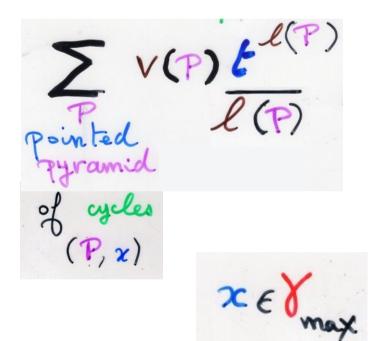






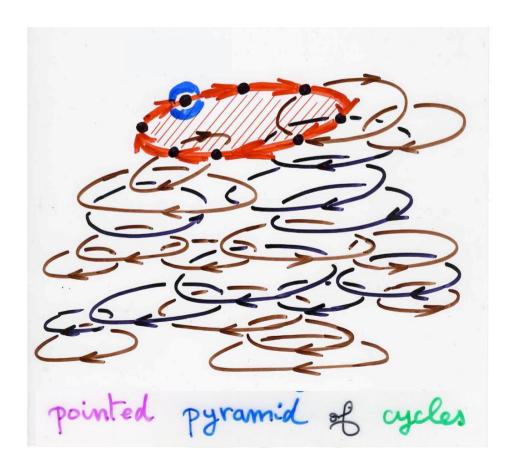


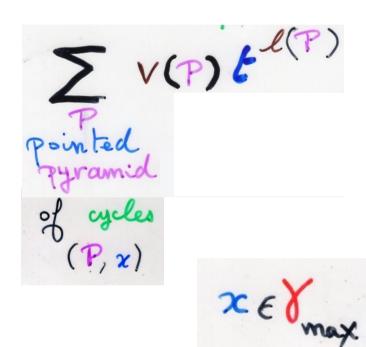




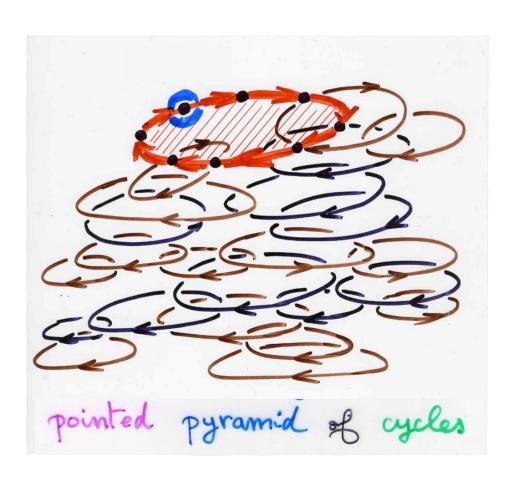
- the unique cycle maximal piece has a distinguished vertex (or edge)

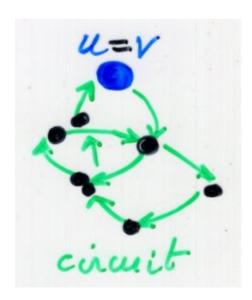
t d log det (I-A)





#### From the third lemma:





are in bijection with pointed pyramics of cycles

The third lemma:

Paths and heaps of cycles

Bijection

 $u, v \in X$ 

going from u to v

- · n self-avoiding path going from u to v
- E heap of cycles such that

  the projections  $\alpha = \pi(m)$  of the

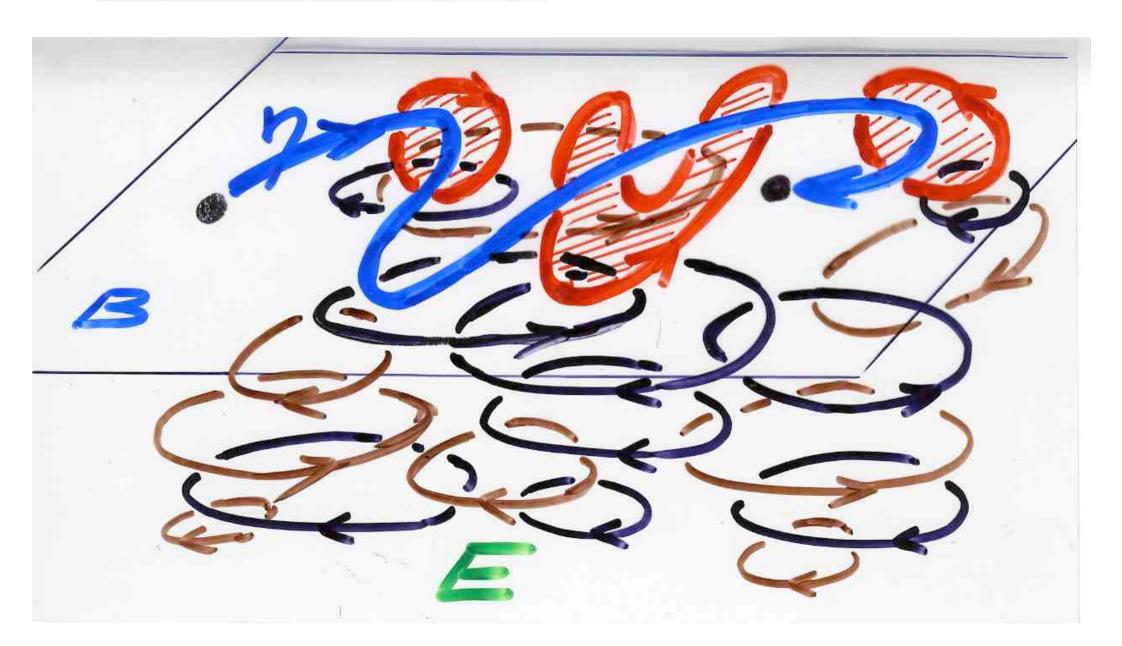
  maximal pieces intersect  $\eta$  (  $\alpha$  and  $\eta$  has a common vertex)

for any 1, t ∈ X

the numbers of occurrences of the edge (s, E) in cv and in  $(\eta, E)$  are the same.

$$\Rightarrow$$
  $\vee(\omega) = \vee(\gamma)\vee(E)$ 

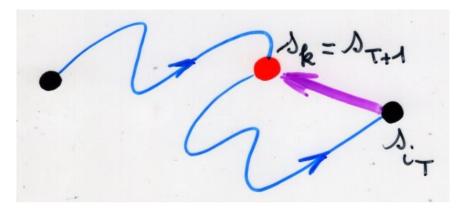
# The bijection X



• suppose 
$$\begin{cases} \text{Cut}_{T}(\omega) = (s_0 = u, ..., s_{i_T}) \end{cases}$$
  
 $E_{T}(\omega)$  heap of cycles

$$\begin{cases} Cut_{T+1}(\omega) = (s_0=u,...,s_{i_T},s_{T+1}) \\ E_{T+1}(\omega) = E_{T}(\omega) \end{cases}$$

(ii) if 
$$S_{T+1} \in Cut_{T}(\omega)$$
,  $S_{T+1} = S_{k}$ 



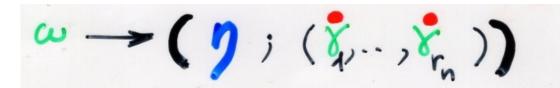
$$\begin{cases} Cut_{T+1}(\omega) = (s_0=u,...,s_k) \\ E_{T+1}(\omega) = E_T(\omega)o\delta \end{cases}$$

$$\omega \xrightarrow{\chi} (\eta, E)$$

$$E = E_n(\omega)$$

loop-erased process LERW

Lawler, 1987



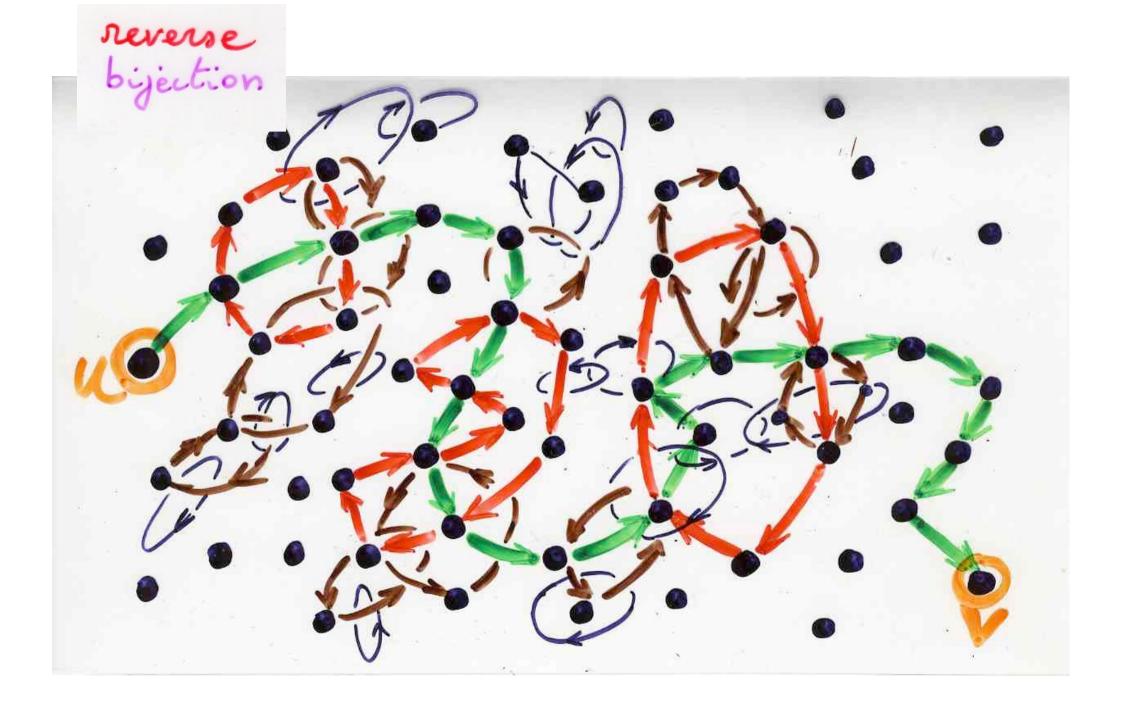
self-avoiding path wasv

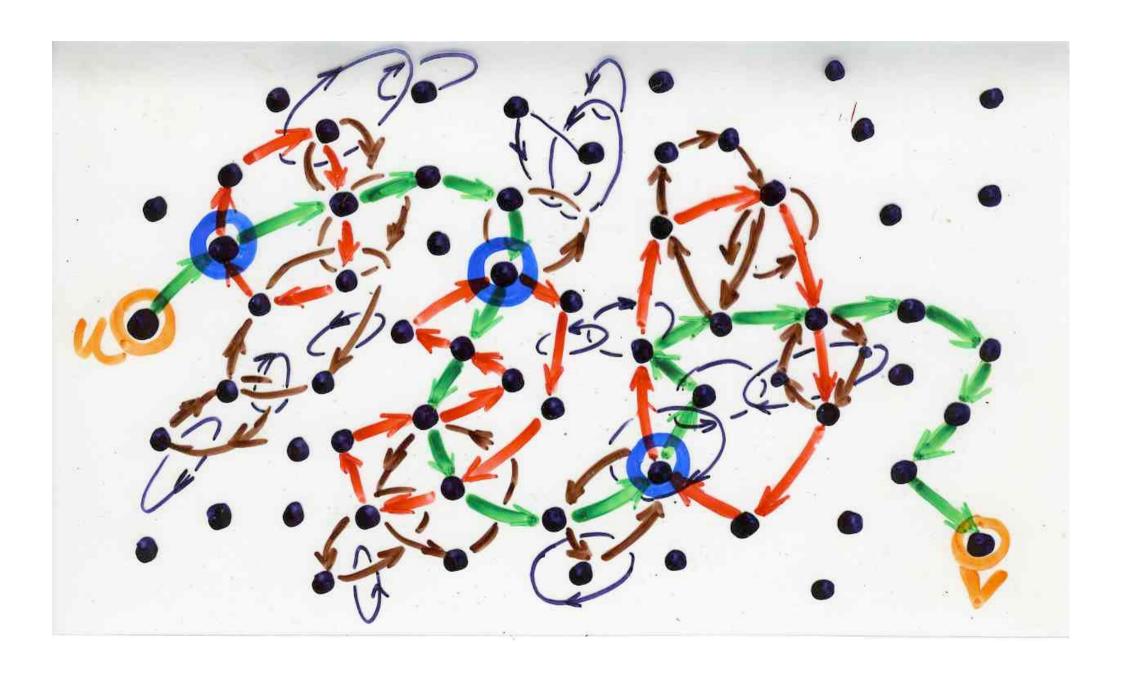
sequence of pointed cycles

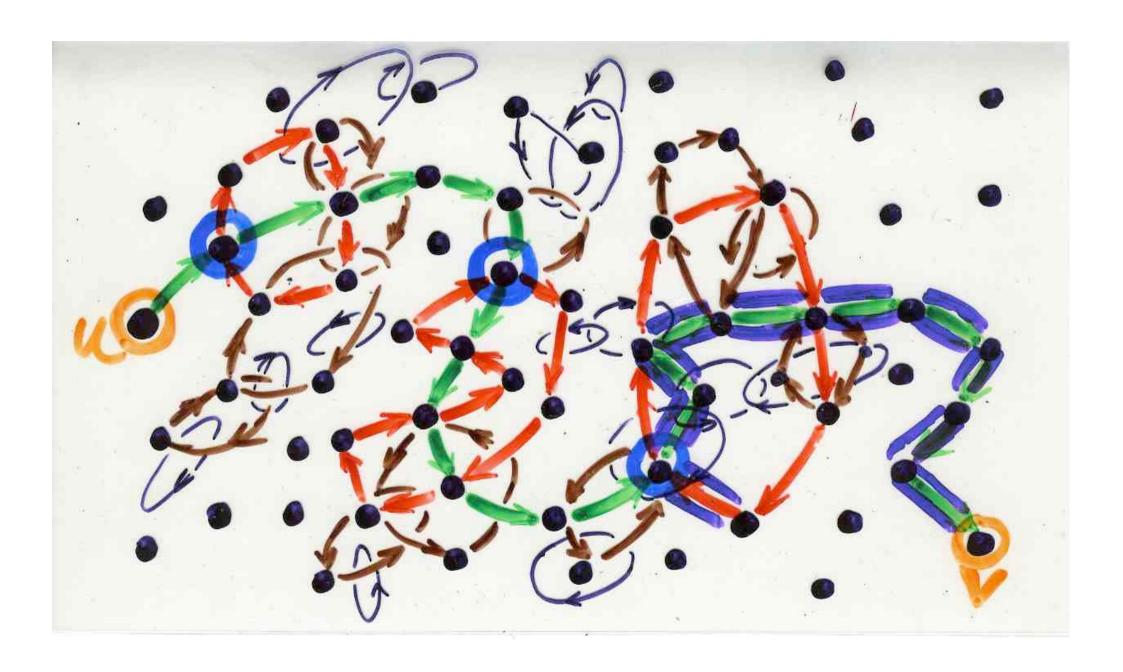
from the pair (n; (1,..., 1,...))
we can reconstruct the path w

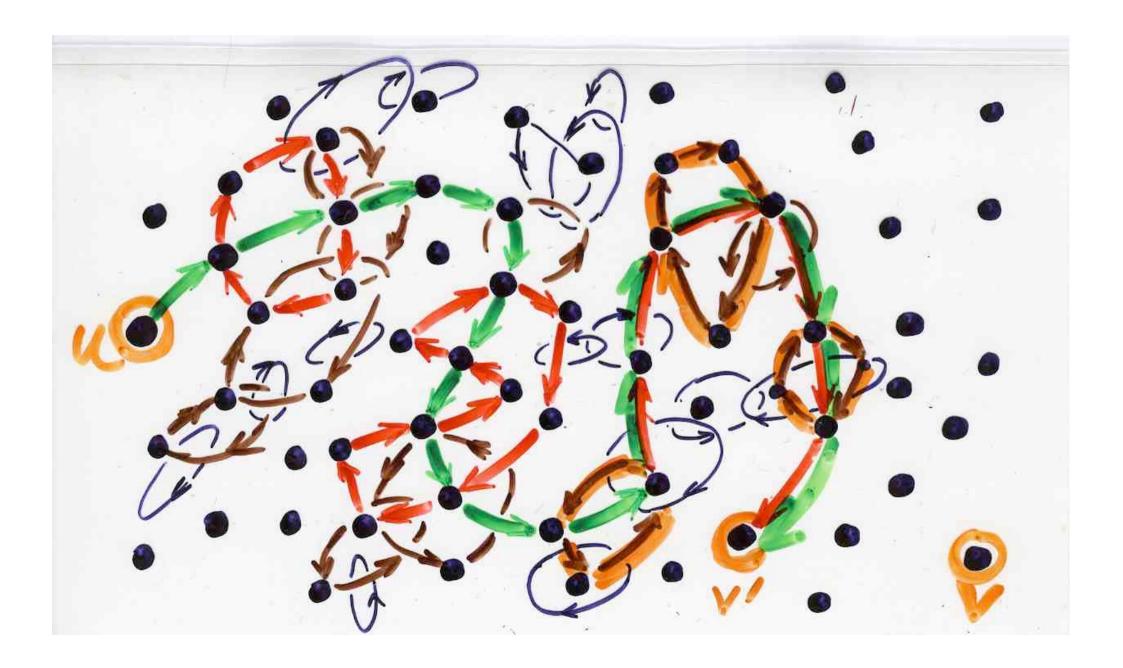
$$\omega \longrightarrow (\eta, E)$$

heaps of cycles on X

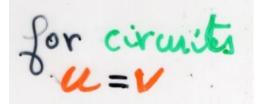




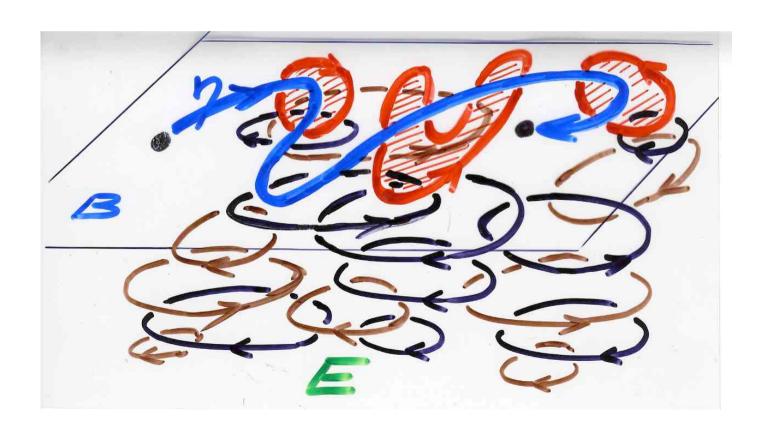




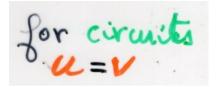
### The bijection X

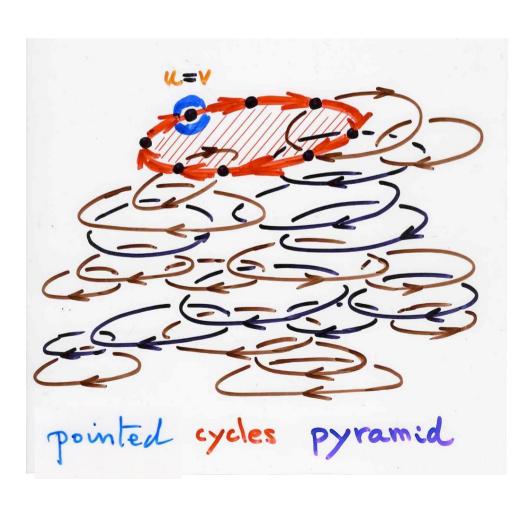


I is reduced to the vertex u=v

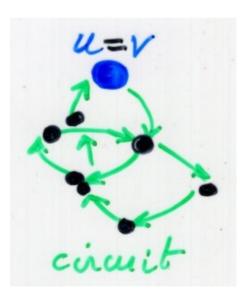


### The bijection X





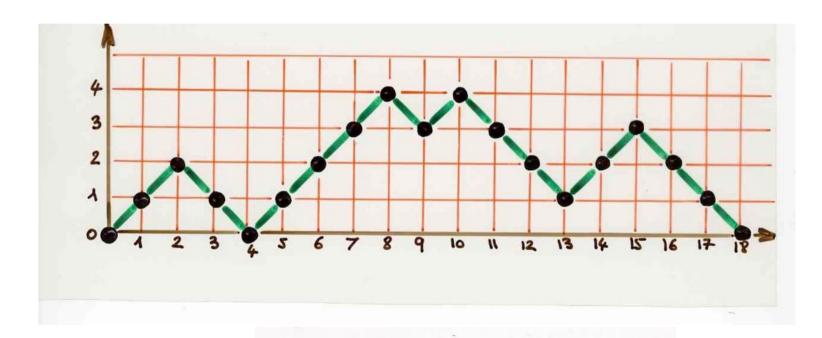
Corollary Circuits on X are in bijection with pointed pyramics of cycles



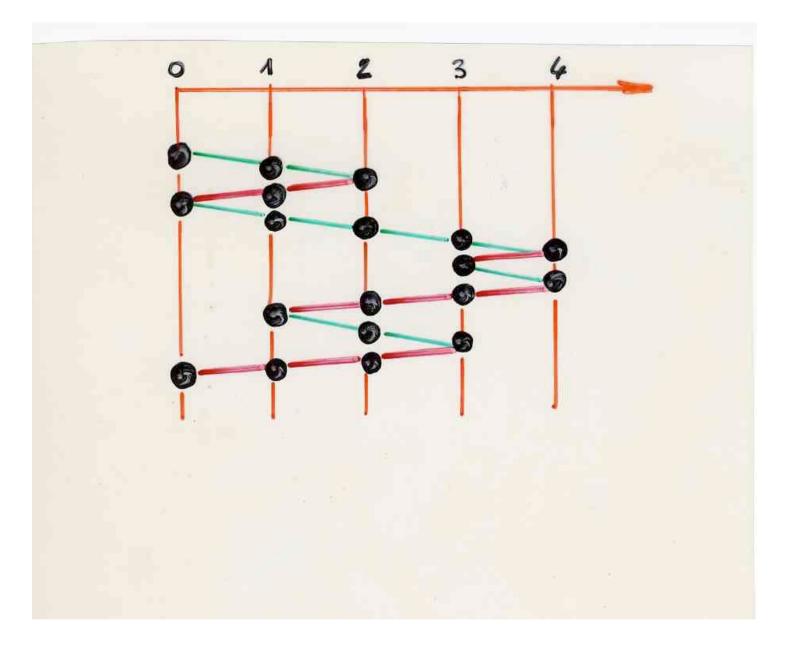
an example with Dyck paths

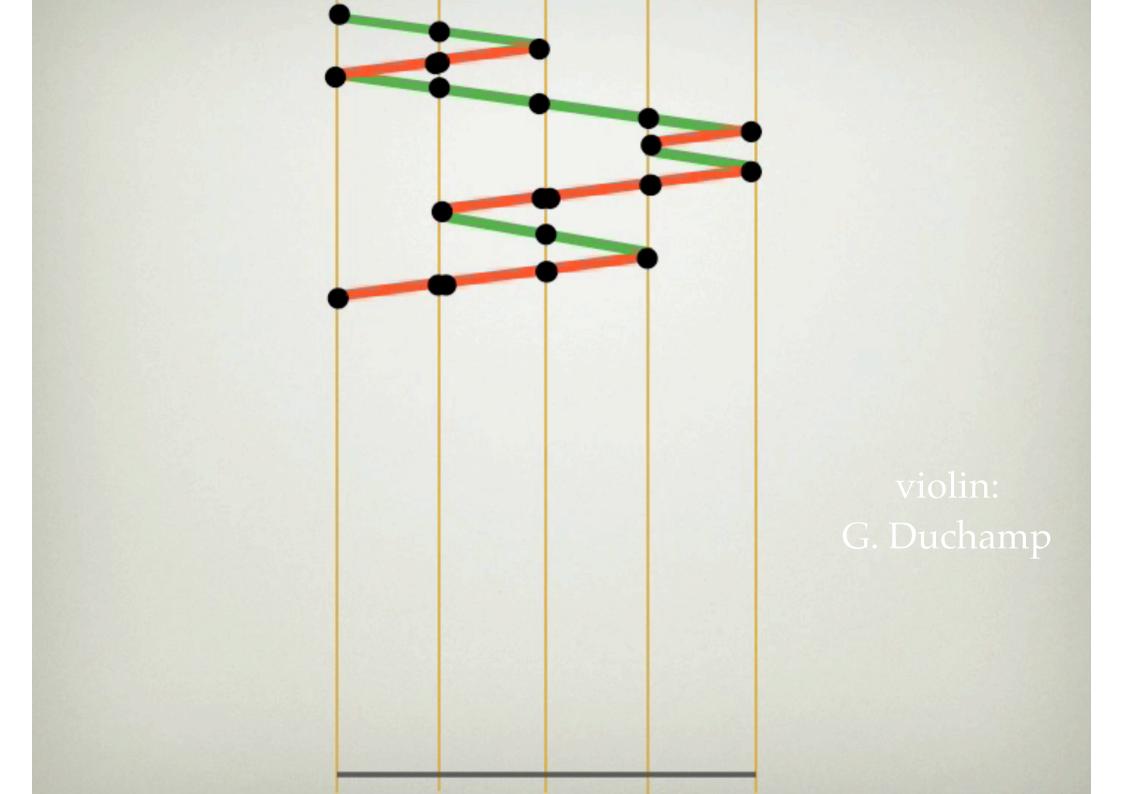
## 

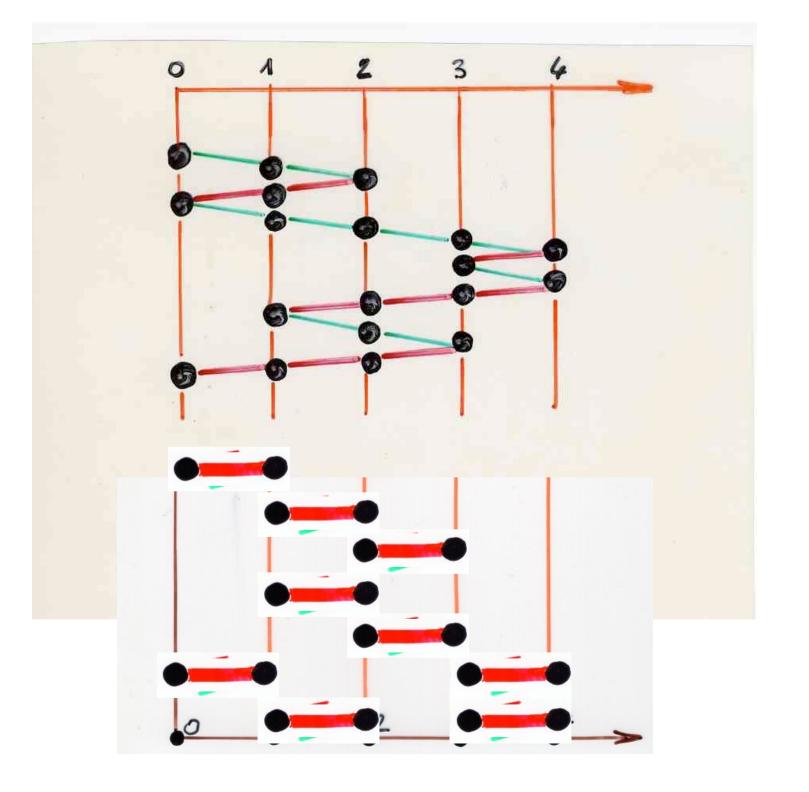
#### Dyck Path

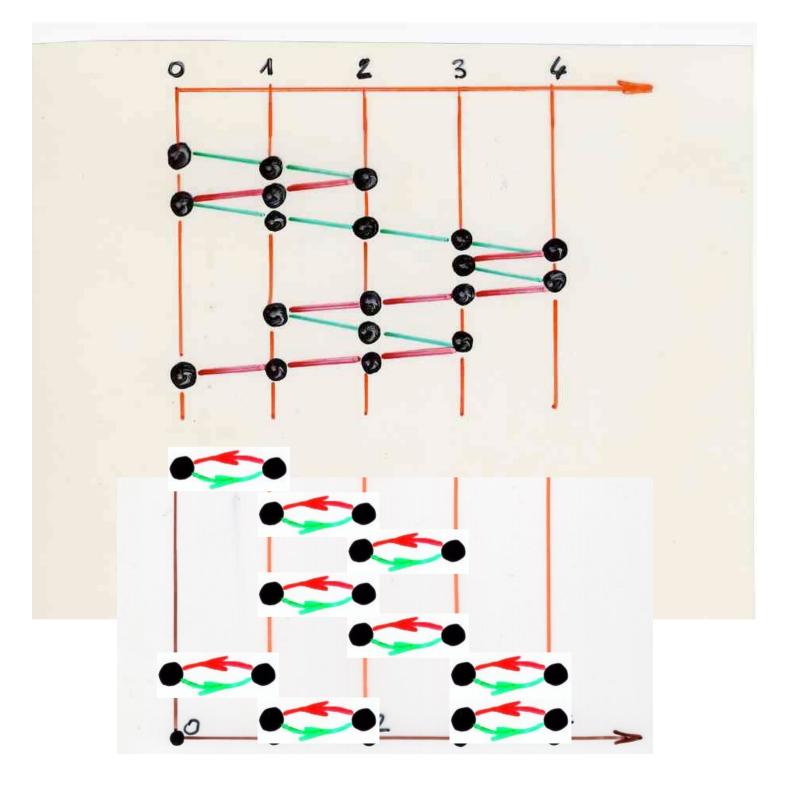


$$C_n = \frac{1}{(2n+1)} \binom{2n+1}{n}$$







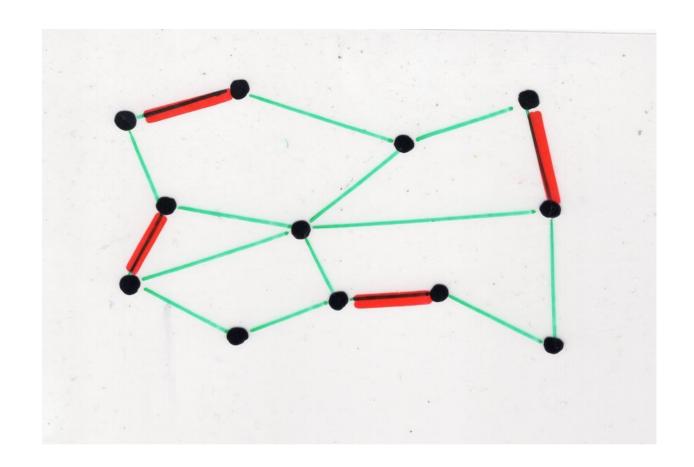


paths on a tree

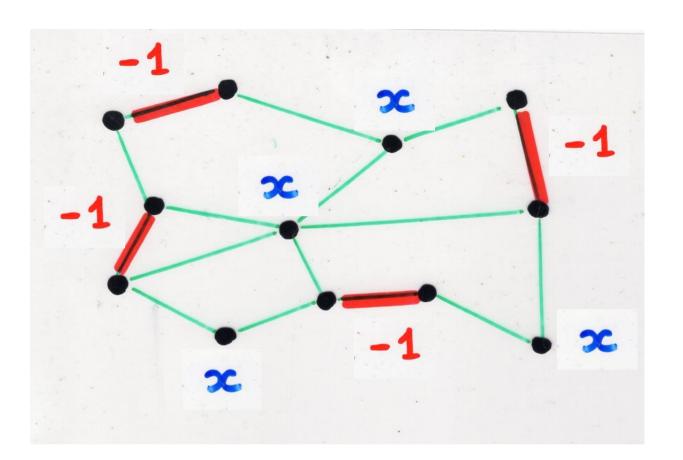
Godsil (1981)

tree-like paths

G graph, χ
w path on G with ω——(η, Ε).
w is tree-like iff the heap Ε contains only agles of length 2.



matching
of a graph G = set of 2 by 2
disjoint edges



matching polynomial of a graph G

$$\chi(x) = det(Ix - A)$$

# Linear algebra revisited with heaps of pieces

of lassical theorem
in linear algebra

with the 3 basic lemma:

- Inversion lemma
- Logarithmic lemma
- circuit = heap of cycles

```
MacMahon "master theorem"
Cartier-Foata (1969)

Matrix inversion
Foata (1979)

Jacobi identity
(log det)

Jackson (1977)
Foata (1980)
```

Cayley- Hamilton theorem Stranbing (1983)

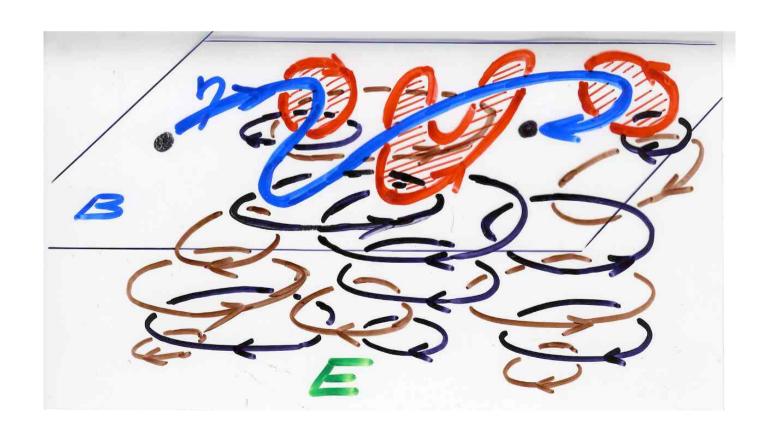
Zeilberger (1985)

Jacobi identity (duality)

Lalonde (1990, 1996)

Fomin (2001), Talaska (2012)

$$\sum_{i \in \mathcal{N}_j} V(\omega) =$$

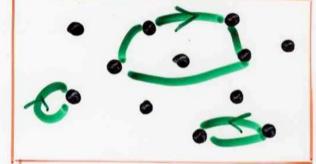


Prop. 
$$\sum_{i} V(\omega) = \frac{N_{ij}}{D}$$

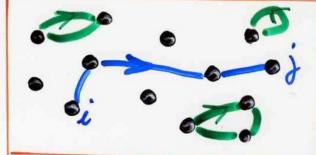
Ny =  $\sum_{i} V(y) N_y$ 

self-avoiding path inaj

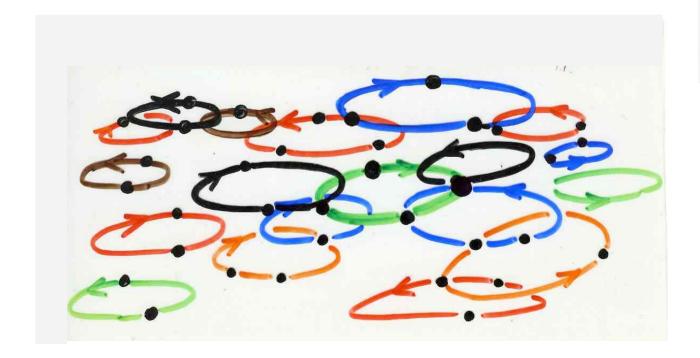
$$D = \sum_{\substack{\{X_1, \dots, X_T\}\\ \text{2-ly 2 disjoint wiles}}} (-1)^T v(X_1) \dots v(X_T)$$



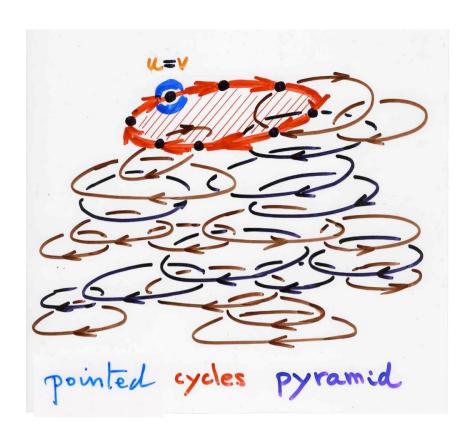
$$N_{ij} = \sum_{\{\gamma; \delta x_i, \delta x_i\}} (-1)^r v(\gamma) v(\delta) \cdots v(\delta)$$



Jacobi identity

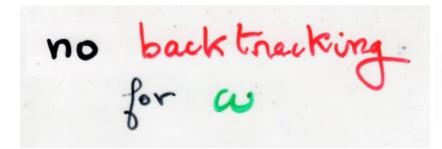


Paths with no backtracking



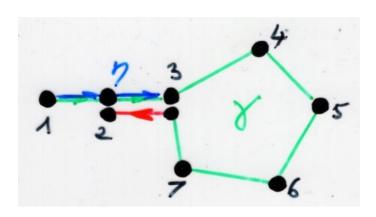
$$= \sum_{circuit} t^{|w|}$$

(-no back tracking





no cycle length 2 in E

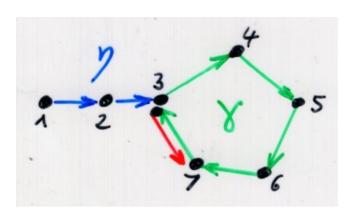


$$\omega \xrightarrow{\chi} (\eta, E)$$

$$\omega \rightarrow (\bullet, \bullet, do V)$$

no backtnecking

dength 2 in E



$$\omega \xrightarrow{\chi} (\eta, E)$$

for a

no cycle length 2 in E

#### second bijection

www (n,F)

 $w = (A_0, A_1, ..., A_i, A_n)$   $w = (A_0, A_1, ..., A_i, A_n)$   $w = (A_0, A_1, ..., A_i, A_n)$   $w = (A_0, A_1, ..., A_i, A_n)$ 

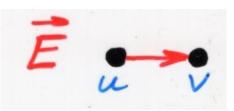
a path on V

- I (w)

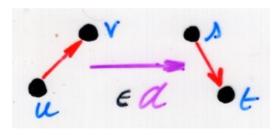
((so,sa),(sa,se)..., (si,si+1),.., (sn-1,sn))

path of IG



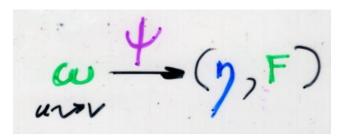


oriented line graph





second bijection 4



y trail

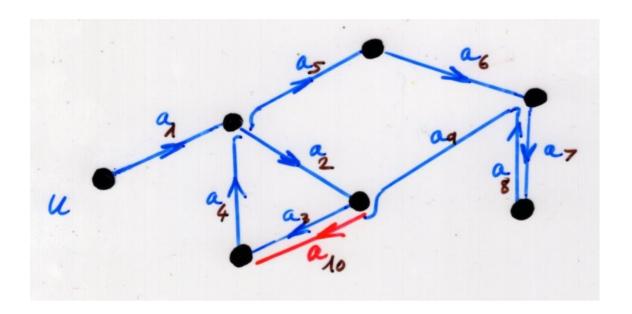
trail = path having all oriented bedges distinct

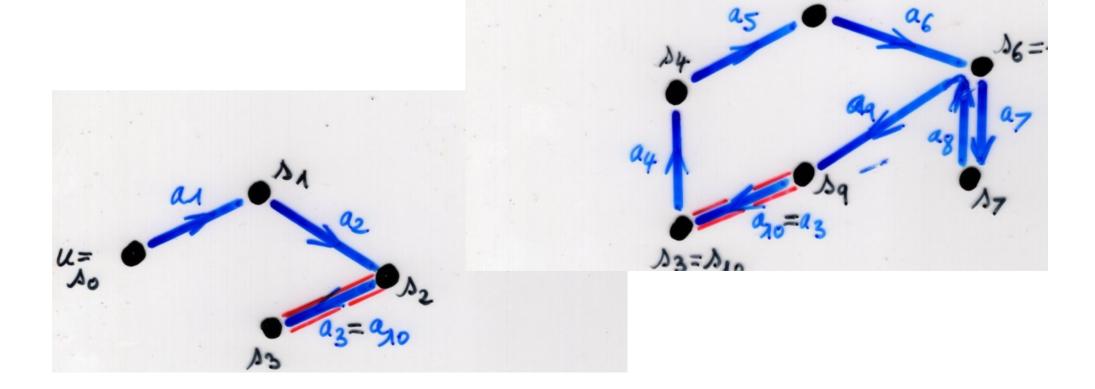
F heap of "oriented loops"

oriented

equivalence class of trail trail y up to a circular germutation of its edges

the byection 4





$$\omega$$
 path on  $V$ 

$$\omega = (\Delta_0, ..., \Delta_n) \quad \begin{array}{c} u = \Delta_0 \\ v = \Delta_n \end{array}$$

$$\overline{L}(\omega) = (e_1, ..., e_n)$$

$$e_i = (s_{i-1}, s_i) \text{ oriented edges}$$

- suppose 
$$w_{+} = (s_{0}, y_{+}) \rightarrow (y_{+}, f_{+})$$

$$y_{\tau} = (a_{1/7}, a_{i\tau})$$
 trail going  $a_{i} = (u, s_{i\tau})$  from  $a_{i} = (u, s_{i\tau})$  from  $s_{i\tau} = s_{i\tau}$ 

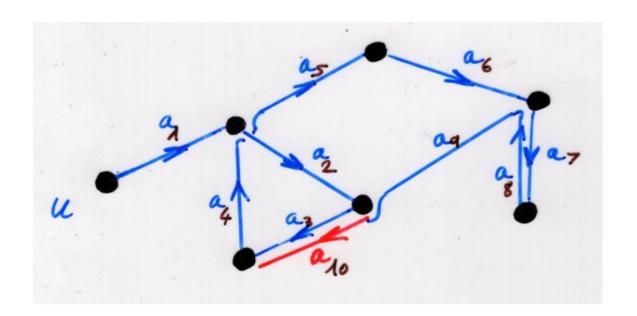
at time T+1, two cases

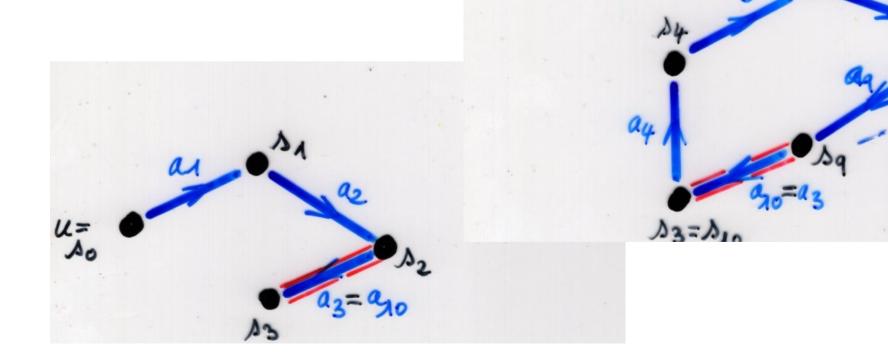
(6) (
$$\Delta T$$
,  $\Delta T+1$ ) does not appear in  $\gamma_T$   
then  $\gamma_{TH} = (\alpha_1, \gamma_1, \alpha_2, (\Delta_T, \Delta_{TH}))$   
 $F_{TH} = F_T$ 

at time T+1, two cases

(ii) 
$$\frac{1}{2}$$
  $(s_T, s_{T+1}) = a$ , edge of  $j_T$ 

$$\frac{\psi(\omega)}{\tau=n} = (\eta_n, F_n)$$





PS





```
each oriented loops of F
is non backtracking
```

Proof of the second formula for the zeta function of a graph

(ii) 
$$\zeta_{G}(t) = \frac{1}{\det(1-Ht)}$$

$$T = adjacency matrix
of the oriented line graph
$$\overline{LG} = (\overline{E}, \alpha)$$$$

$$T = (t_{(i,j),(k,\ell)})$$
 $t_{(i,j),(k,\ell)} = \{i,j\} \atop i\neq k$ 

B submatrix of T

$$\mathbf{B} = \left(b_{(i,j)(k,\ell)}\right)$$

pointed pyramids
of non book tracking
oriented loops

t d la 1 det (1-Ht)

td log Zo(t)

(-no tail

The third formula for zeta

$$G = (V, E)$$

$$D = (d_{ii})$$

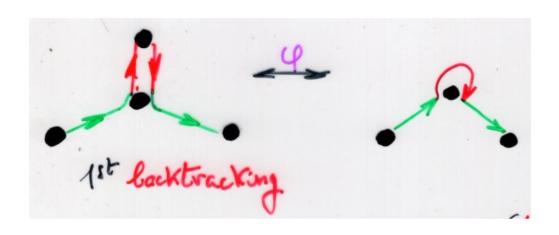
$$d_{ii} = deg v_i$$

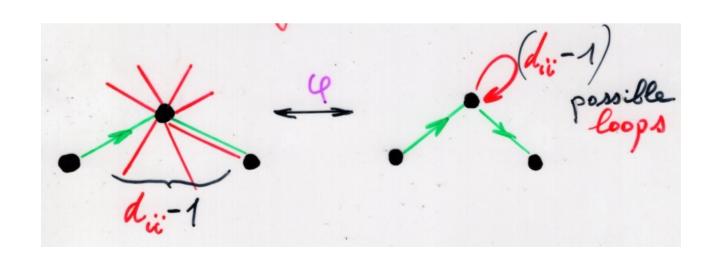
$$\frac{t \, d \, det}{dt} \frac{1}{dt} \frac{det(I-tA+t^2(D-I))}$$

= 
$$\sum_{\alpha} v(\alpha)$$

$$\begin{cases} V(i,j) = t \\ V(i,i) = -t^2 ((degi)-1) \end{cases}$$

#### The idea ...





Back to number theory

$$\frac{2}{2}(3) = \frac{1}{1-p^{-3}}$$
Prime
number

$$\leq_{6}(t) = \frac{11}{(1-t^{1c_{1}})}$$

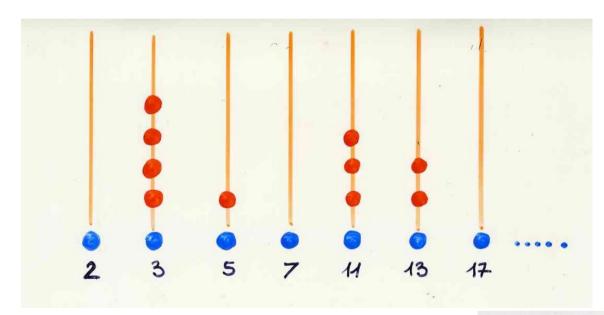
some "prime" over the graph G

$$\frac{2}{2}(3) = \frac{1}{1-p^{-3}}$$
Prime
number

equivalence class prime circuit

no backtracking

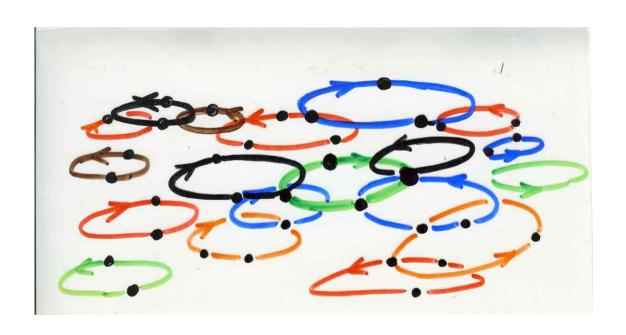
$$\sum_{n\geq 1} n^{-s} = \left(\sum_{n\geq 1} \mu(n) n^{-s}\right)^{-1}$$



$$\frac{2}{2}(3) = \frac{1}{1-p^{-3}}$$
Prime

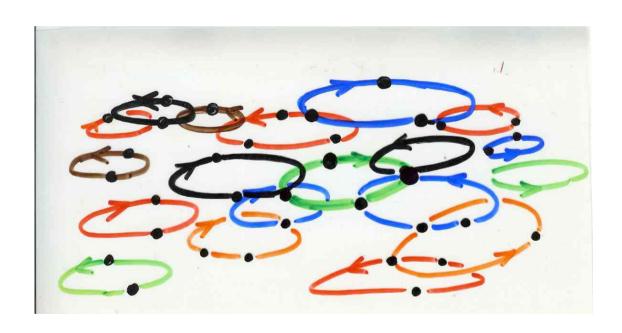
number

$$\sum_{n \ge 1} n^{-s} = \left( \sum_{n \ge 1} \mu(n) n^{-s} \right)^{-1}$$

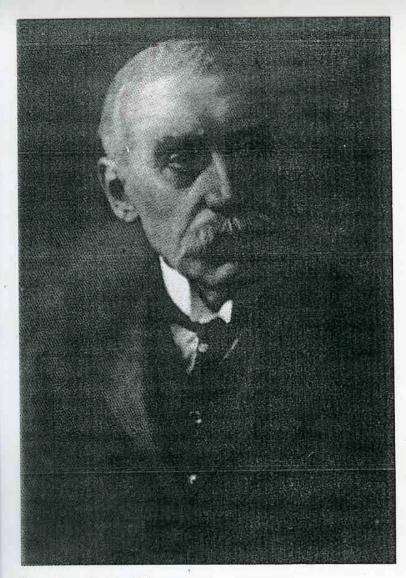


equivalence class prime circuit

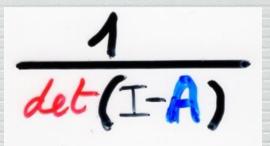
extending number theory to paths on Graphs



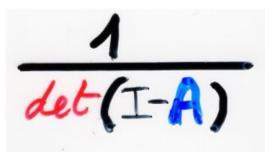
equivalence class prime circuit MacMahon Master theorem

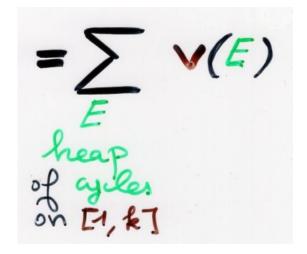


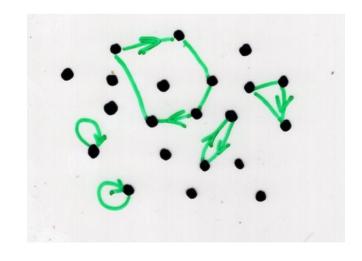
Percy Alexander MacMahon

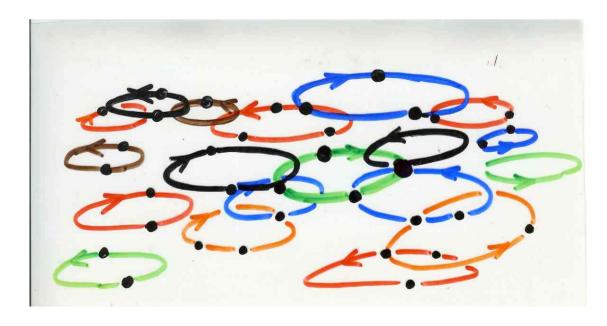


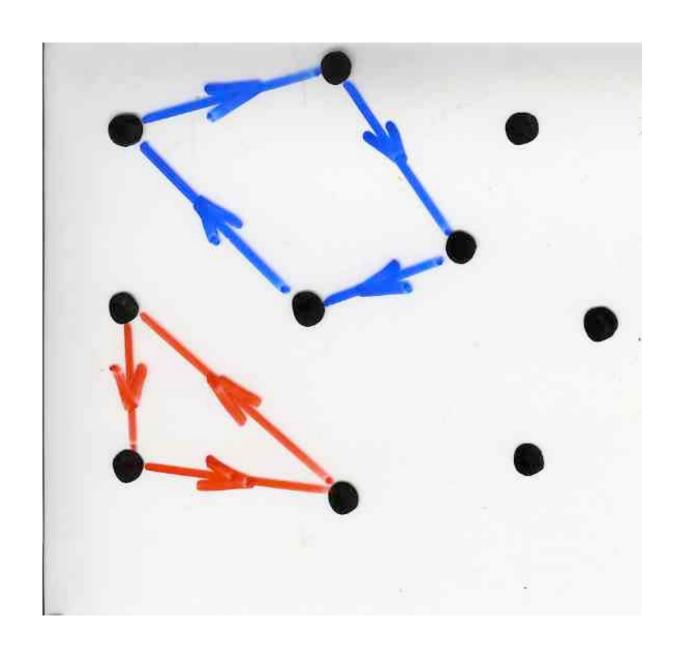
## inversion

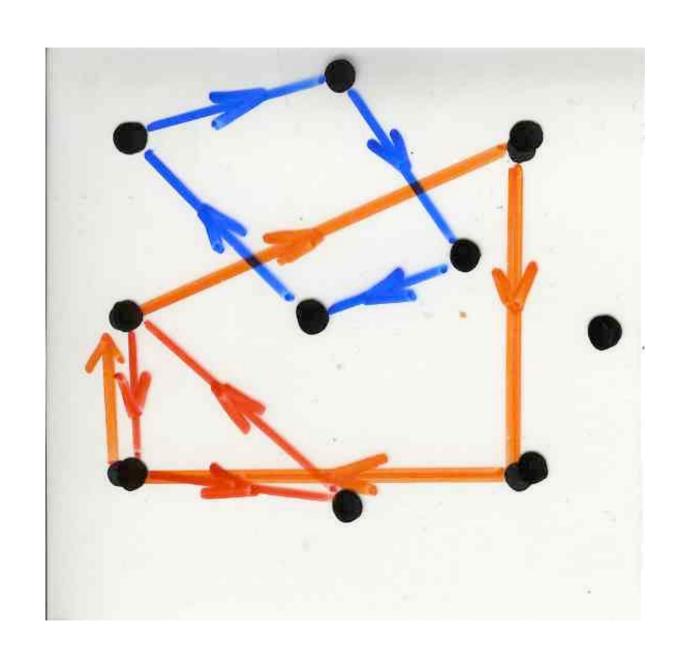


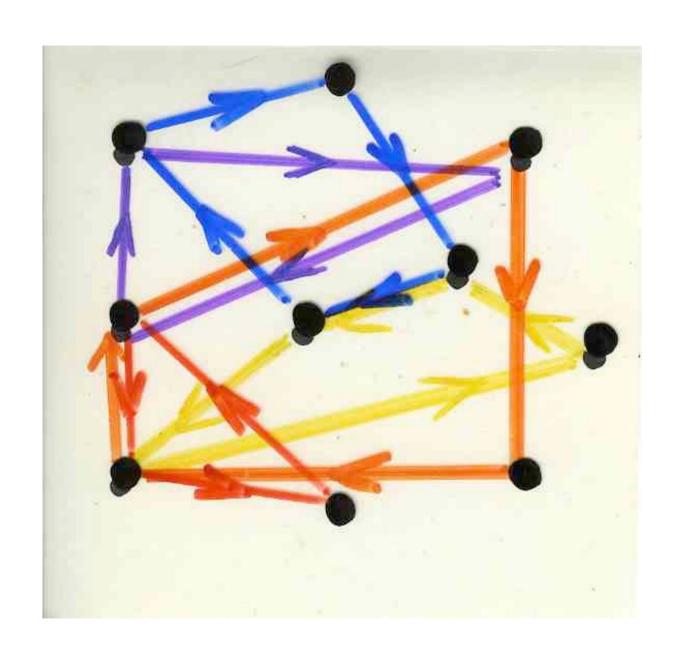


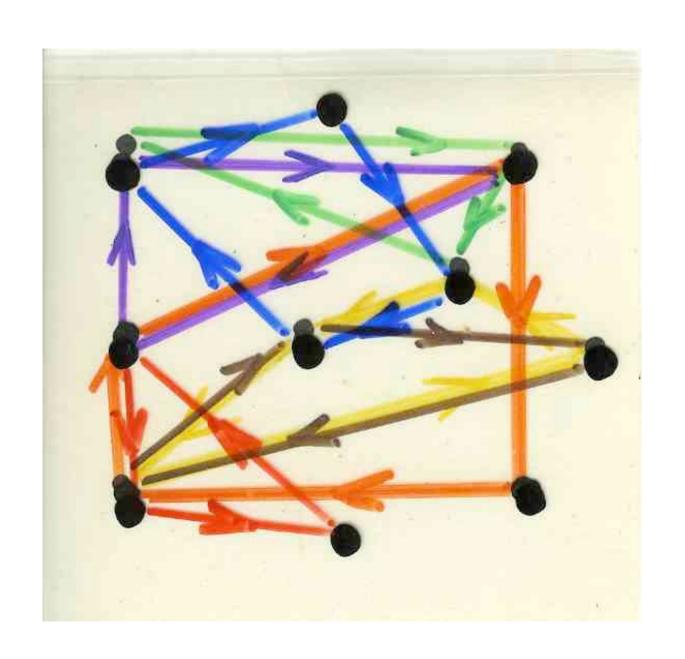


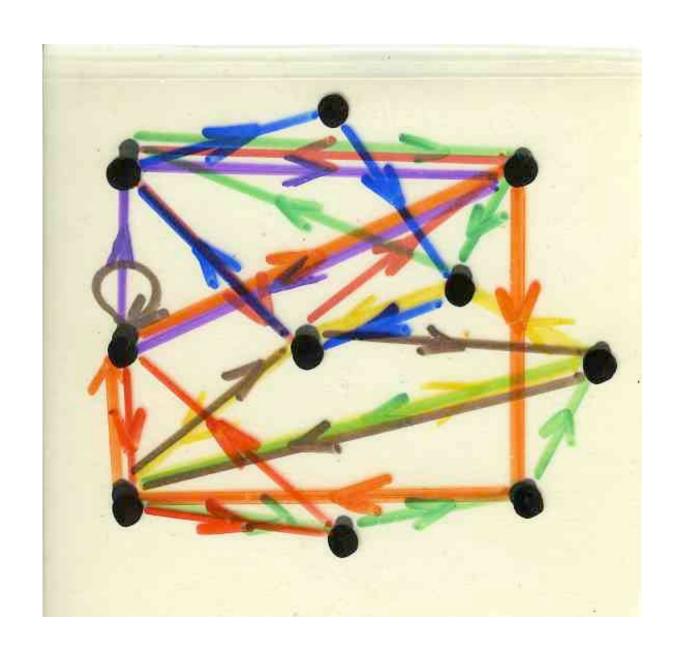


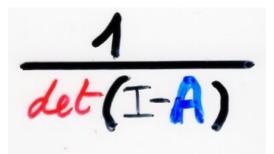


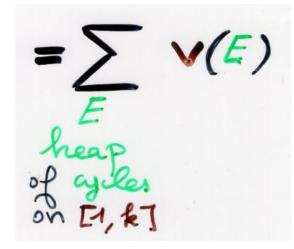


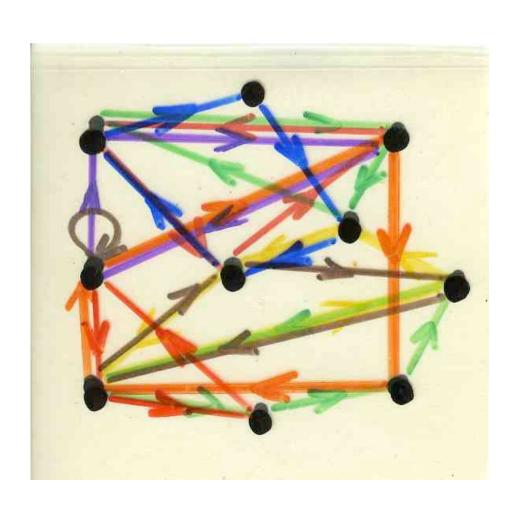


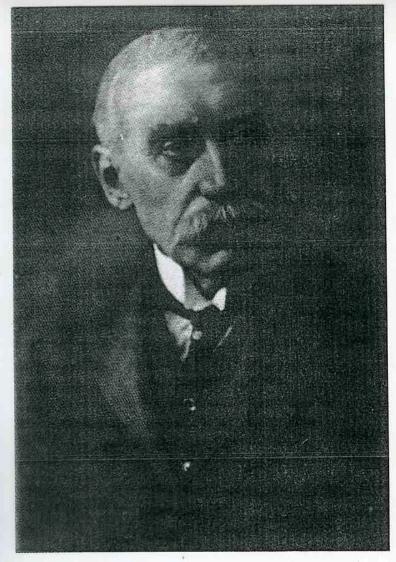








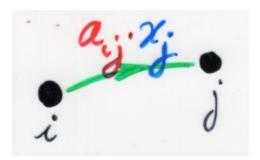




Percy Alexander MacMahon

Where is my MASTER THEOREM

#### Mac Mahon master theorem



$$A = (a_{ij})_{n \times n}$$

$$X = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

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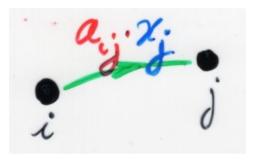
$$X = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$X = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

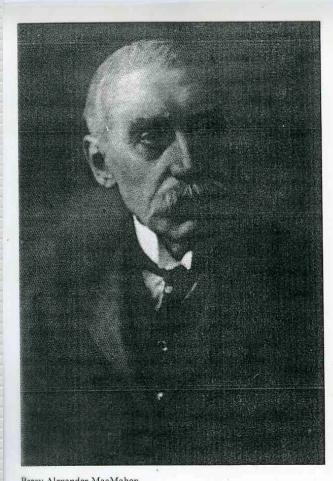
$$X = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$X = \begin{bmatrix}$$

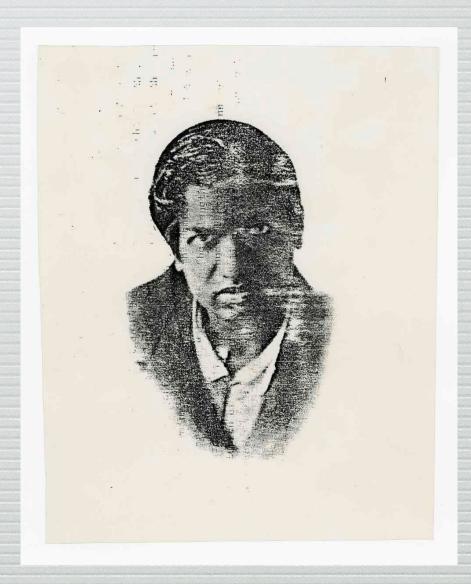
### Mac Mahon master theorem



The coefficient of  $x_1^{\alpha_1} \dots x_n^{\alpha_n}$  in  $\frac{1}{\det(\mathbf{I} - AX)}$  the coefficient of  $x_1^{\alpha_1} \dots x_n^{\alpha_n}$  in  $x_1^{\alpha_1} \dots x_n^{\alpha_n}$ 



Percy Alexander MacMahon





ॐ सरस्वत्ये नमः।



# The Art of Bijective Combinatorics Part II, commutations and heaps of pieces (video-book, course IMSc Chennai, 2017)

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Chapter 4 Heaps and linear algebra

Chapter 5 Heaps and algebraic graph theory

Chapter 5b, zeta function of a graph

Chapter 6 Heaps and Coxeter groups

