

Growth diagrams, local rules and beyond

(part I)

21st Ramanujan Symposium:
National Conference on algebra and its applications

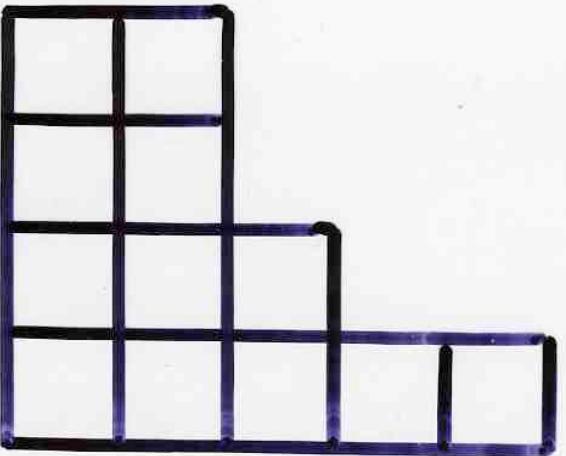
Ramanujan Institute,
University of Madras
28 February 2018

Xavier Viennot
CNRS, LaBRI, Bordeaux
and IMSc, Chennai
www.viennot.org

mirror website
www.imsc.res.in/~viennot

RS

The Robinson-Schensted correspondence



7	12			
6	10			
3	5	9		
1	2	4	8	11

Young
tableau

shape λ

$f_\lambda =$ number of
Young tableaux
with
shape λ

$$\sigma = (1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10)$$

6	10			
3	5	8		
1	2	4	7	9

P

8	10			
2	5	6		
1	3	4	7	9

Q



The Robinson-Schensted correspondence
between permutations and pairs of
(standard) Young tableaux with the same shape

Representation theory of groups

Case of the group G_n permutations

irreducible
representations



partition
of n

dimension
of the irreducible
representation
 $(=$ order of the
matrices $)$

=

\sum_{λ}
number of Young
tableaux
with shape λ

finite group G

$$|G| = \sum_R (\deg R)^2$$

irreducible
representation

for the symmetric
group G_n
(permutations)

$$n! = \sum_\lambda (\ell_\lambda)^2$$

partition
of n

“local” algorithm on a grid or “growth diagrams”

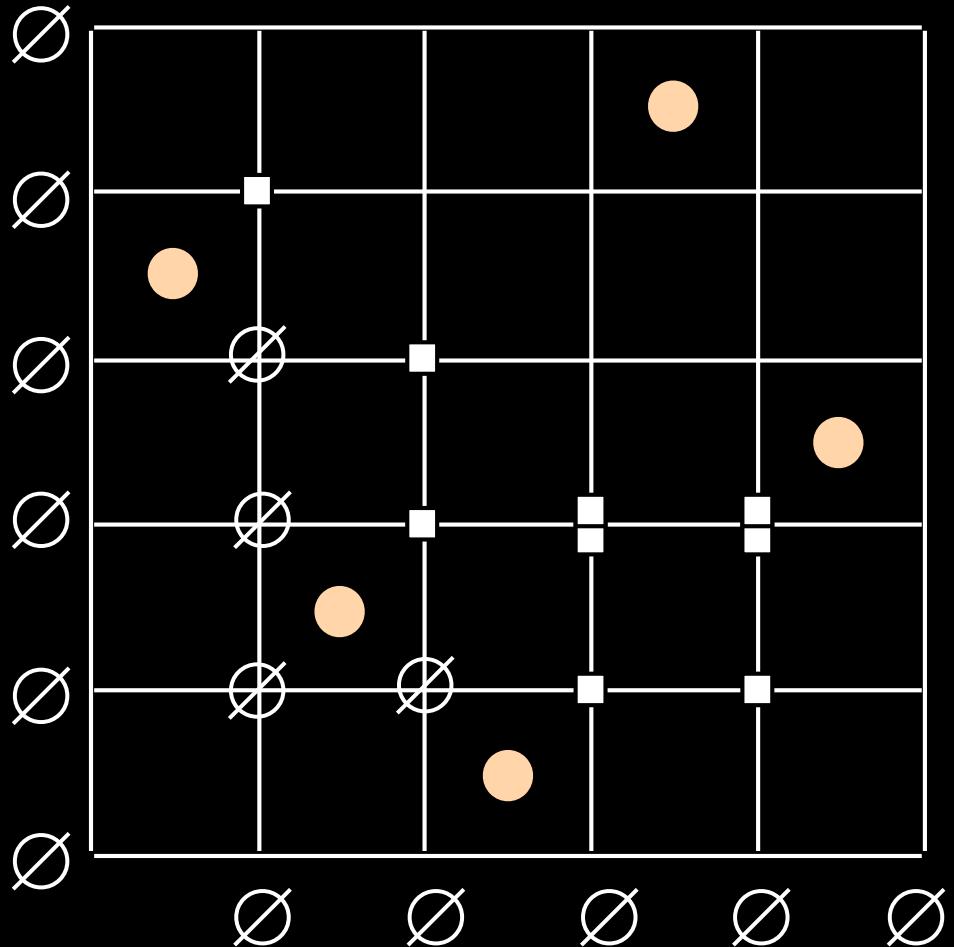
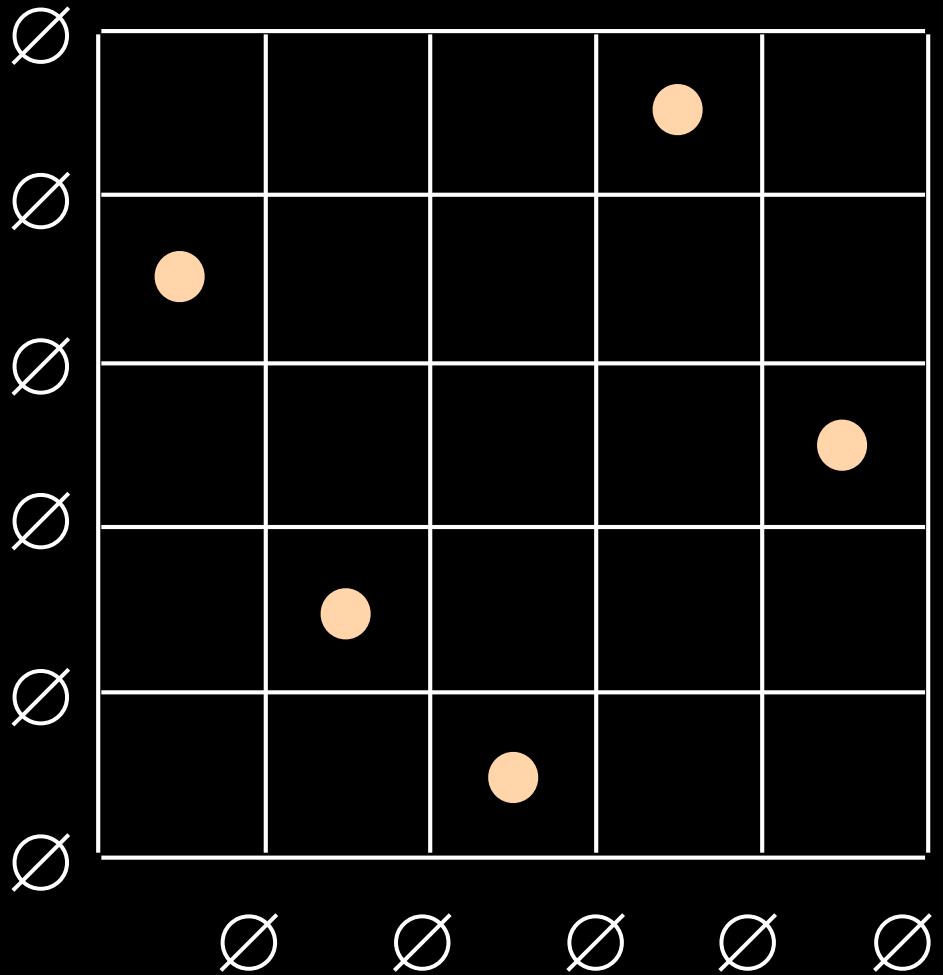
S. Fomin, 1986, 1994

C.Krattenthaler



initial
state

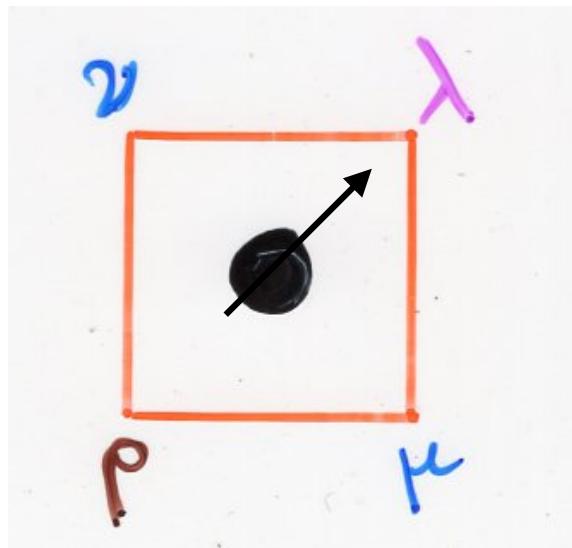
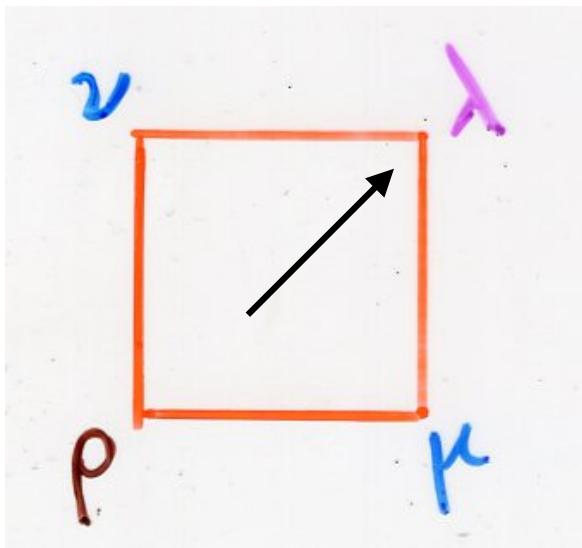
during the
labeling process



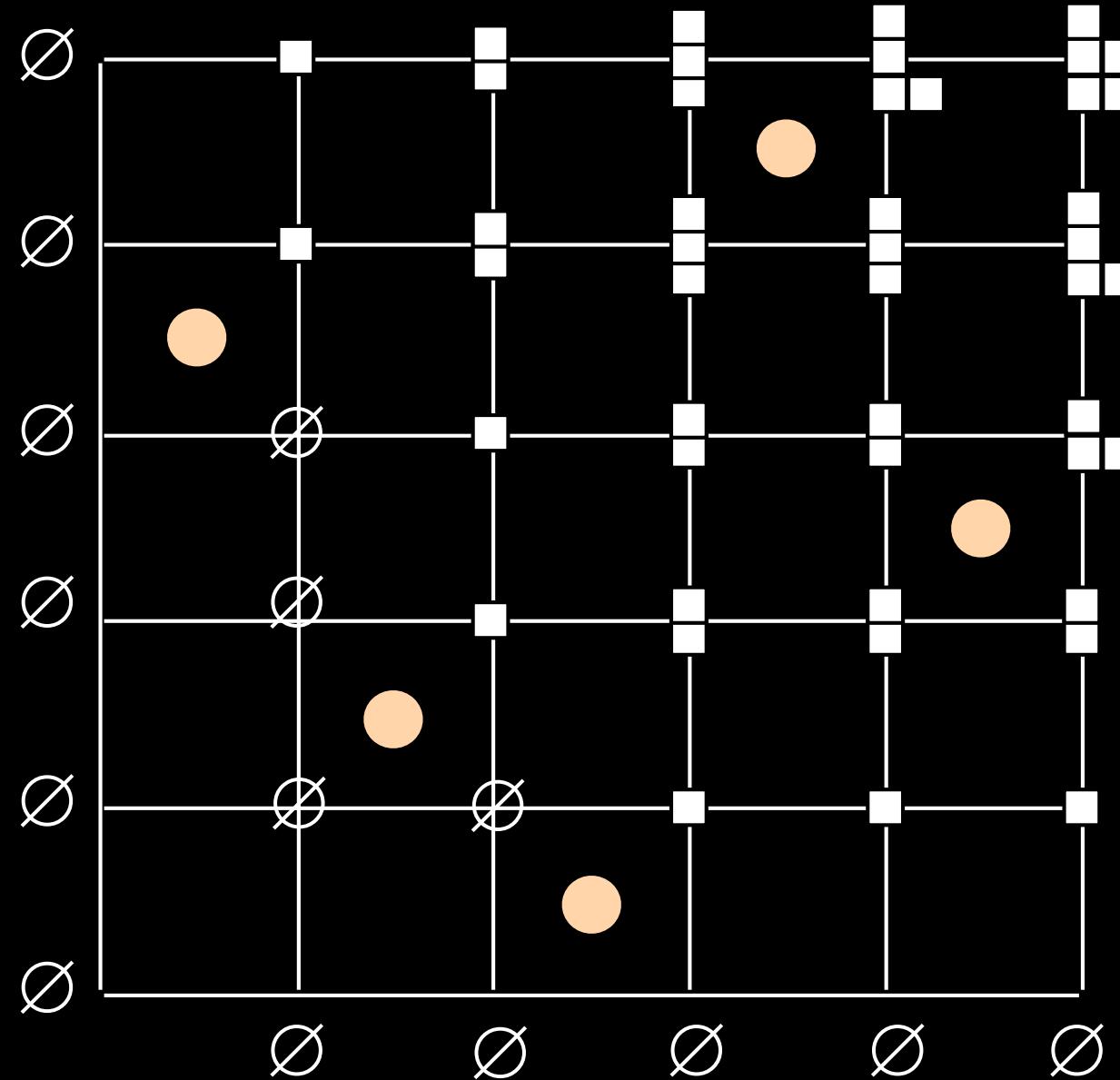
$$\sigma = 4, 2, 1, 5, 3$$

"growth diagrams"

"local rules"

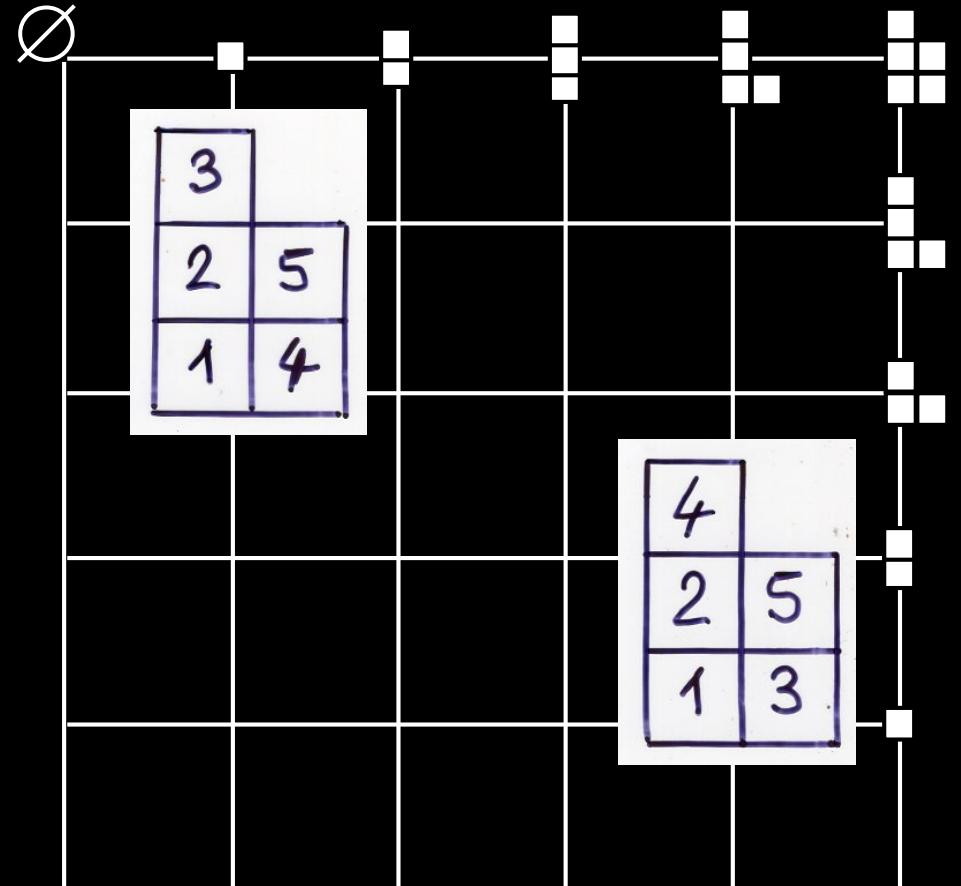
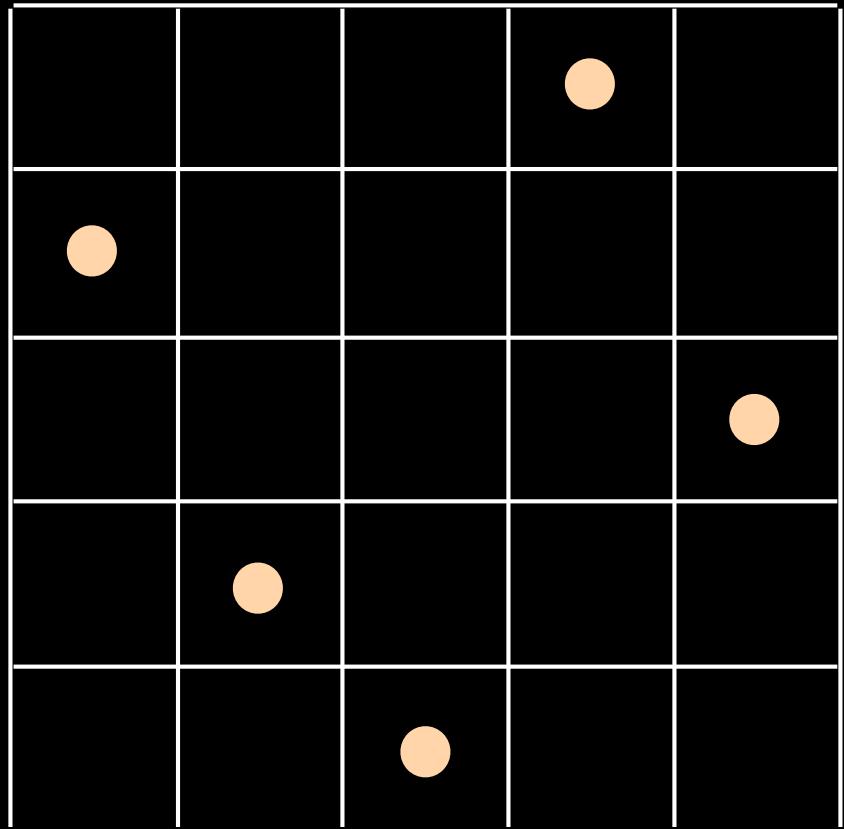


*final
state*



"growth diagrams"

"local rules"

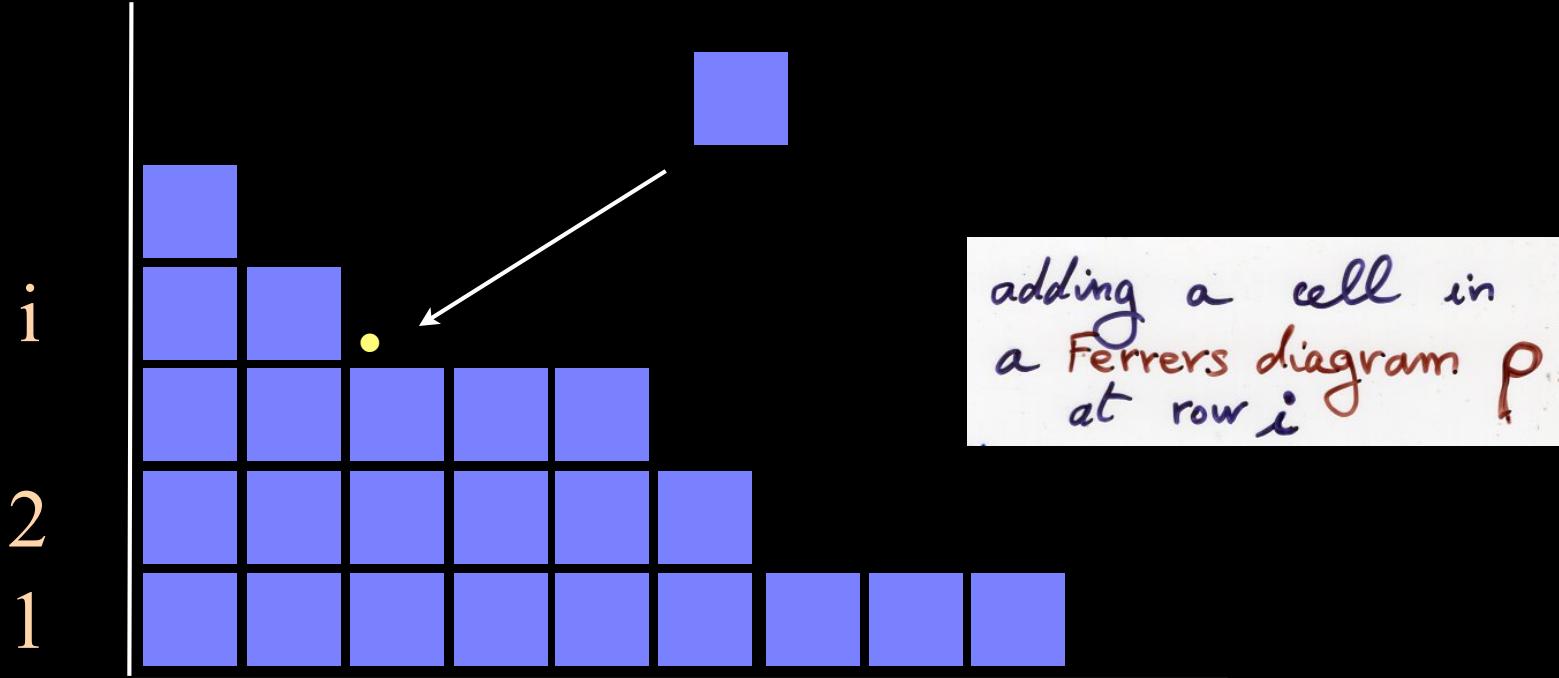


$$\sigma = 4, 2, 1, 5, 3$$

- this bijection is the same as the Robinson-Schensted correspondence

notations

operator U_i



$$U_i(\rho) = \rho + (i)$$

"local rules"

(i) $\rho = \mu = \nu$ and $\begin{array}{c} \square \\ \rho \quad \mu \end{array}$ then $\lambda = \rho$

(ii) $\rho = \mu \neq \nu$, then $\lambda = \nu$

(iii) $\rho = \nu \neq \mu$, then $\lambda = \mu$

(iv) ρ, μ, ν pairwise \neq , then $\lambda = \mu \cup \nu$

(v) $\rho \neq \mu = \nu$, then $\lambda = \mu + (i+1)$

given that $\mu = \nu$ and ρ differ in the i -th row
[in fact $\mu = \nu = \rho + (i)$]

(vi) $\rho = \mu = \nu$ and $\begin{array}{c} \square \\ \rho \quad \mu \end{array}$, then $\lambda = \mu + (1)$

"local rules"

(ii) $\rho = \mu \neq \nu$, then $\lambda = \nu$

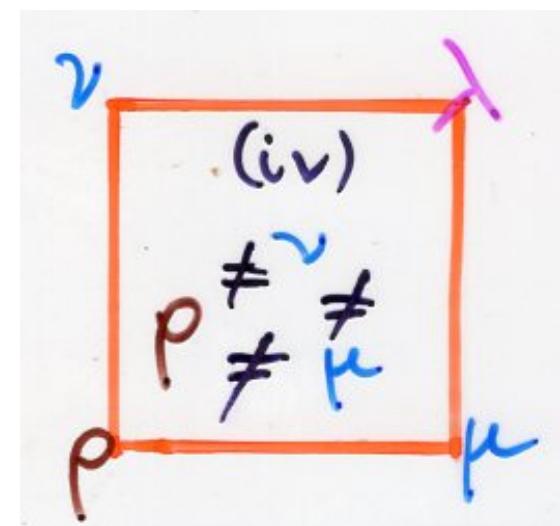
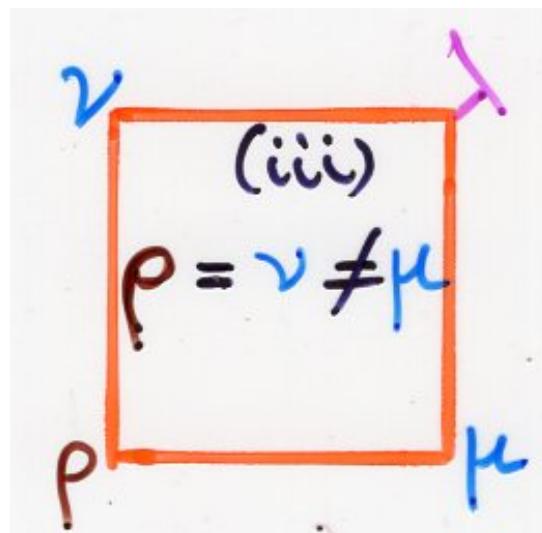
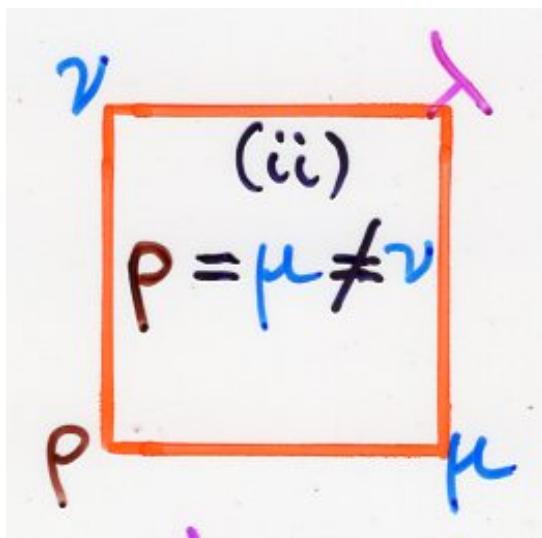
(iii) $\rho = \nu \neq \mu$, then $\lambda = \mu$

(iv) ρ, μ, ν pairwise \neq , then $\lambda = \mu \cup \nu$

$\mu \neq \nu$

"local rules"

$$\mu \neq \nu$$



$$\lambda = \nu$$

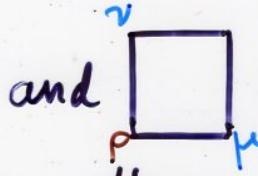
$$\lambda = \mu$$

$$\lambda = \mu \cup \nu$$

"local rules"

(i)

$$\rho = \mu = \nu$$



$$\lambda = \rho$$

$$\mu = \nu$$

(v)

$$\rho \neq \mu = \nu, \text{ then } \lambda = \mu + (i+1)$$

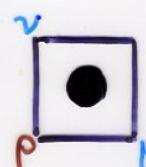
given that $\mu = \nu$ and ρ differ in the i -th row

[in fact $\mu = \nu = \rho + (i)$]

$$\mu = \nu$$

(vi)

$$\rho = \mu = \nu$$

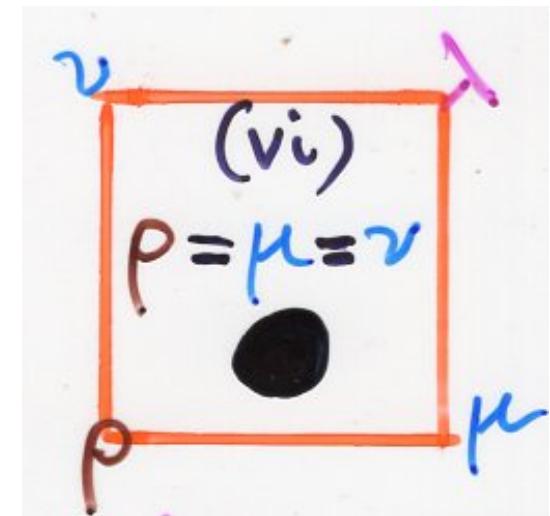
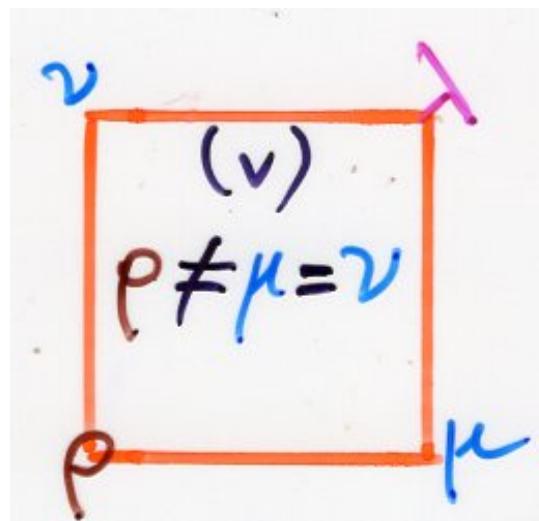
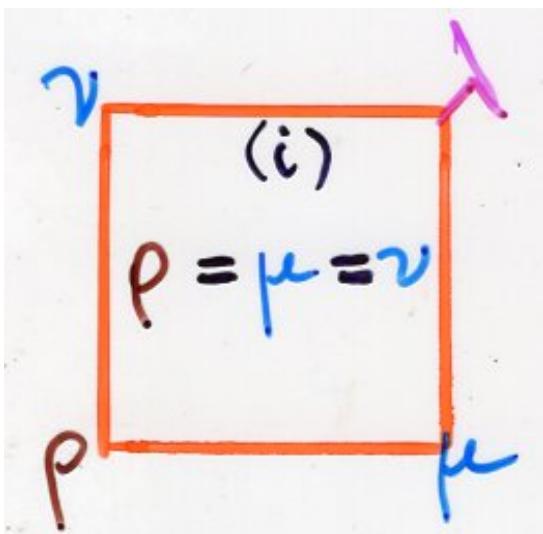


$$\lambda = \mu + (1)$$

$$\mu = \nu$$

"local rules"

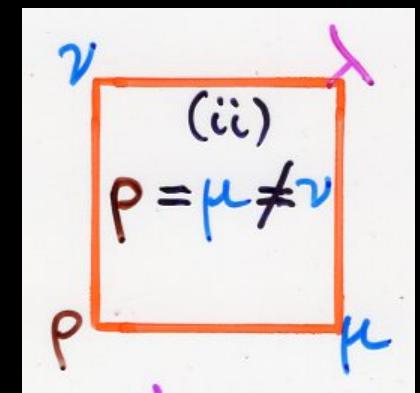
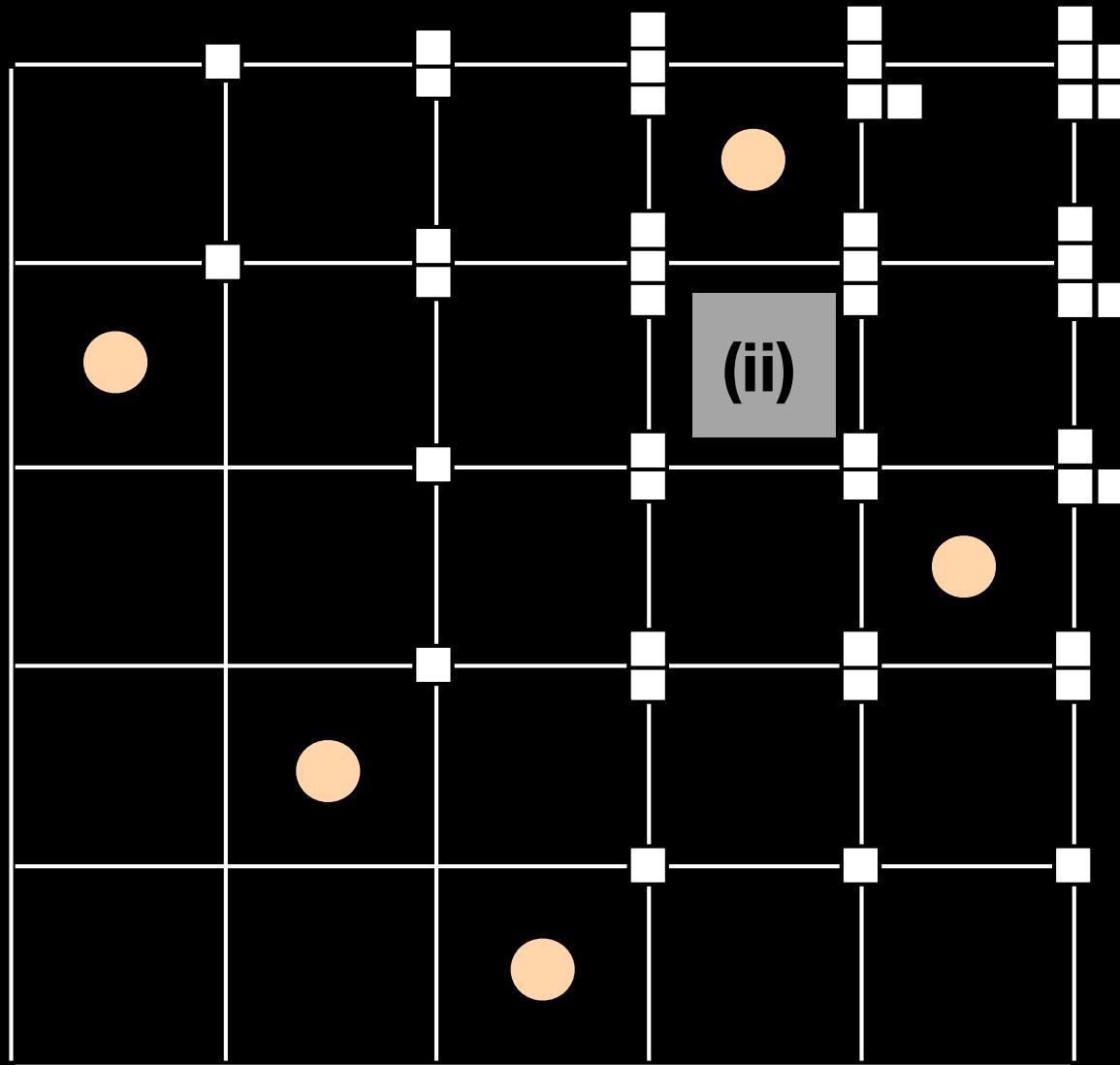
$$\mu = \nu$$



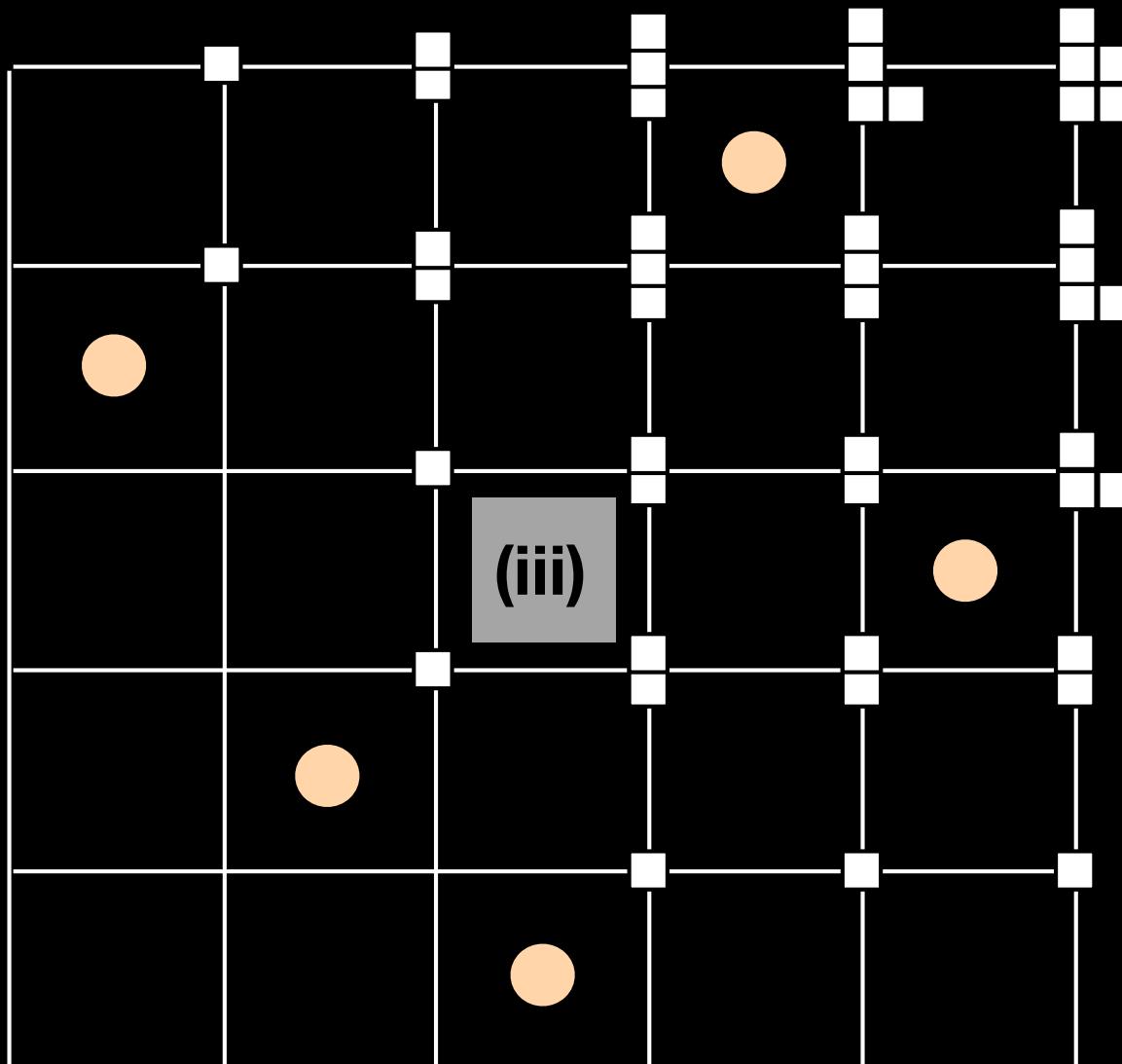
$$\lambda = \rho$$

$$\lambda = \begin{cases} \mu & + (i+1) \\ \nu & \end{cases}$$

$$\lambda = \begin{cases} \rho & + (1) \\ \mu & \\ \nu & \end{cases}$$

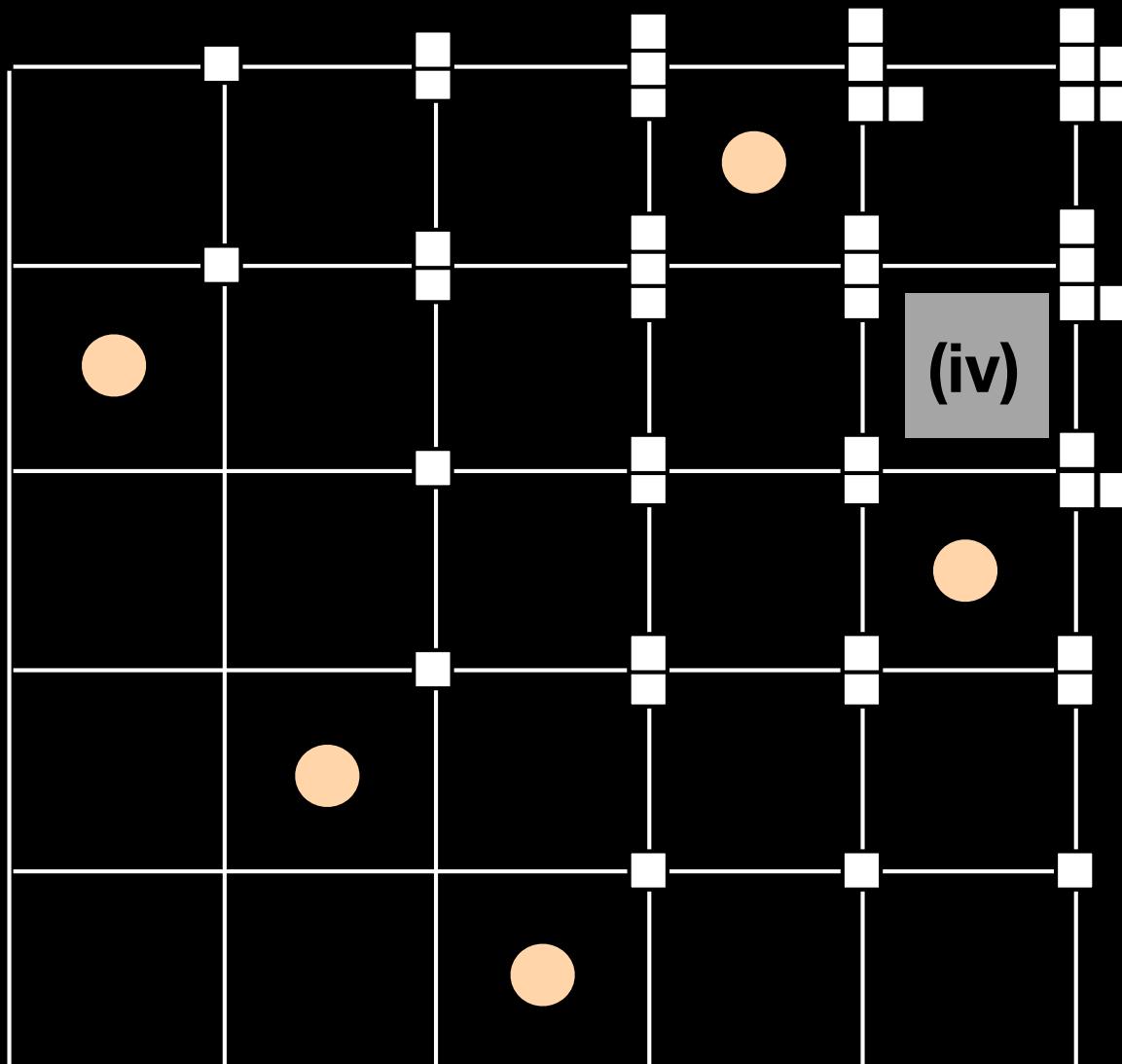


$$\lambda = \nu$$



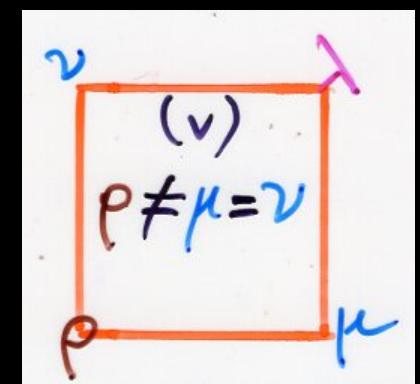
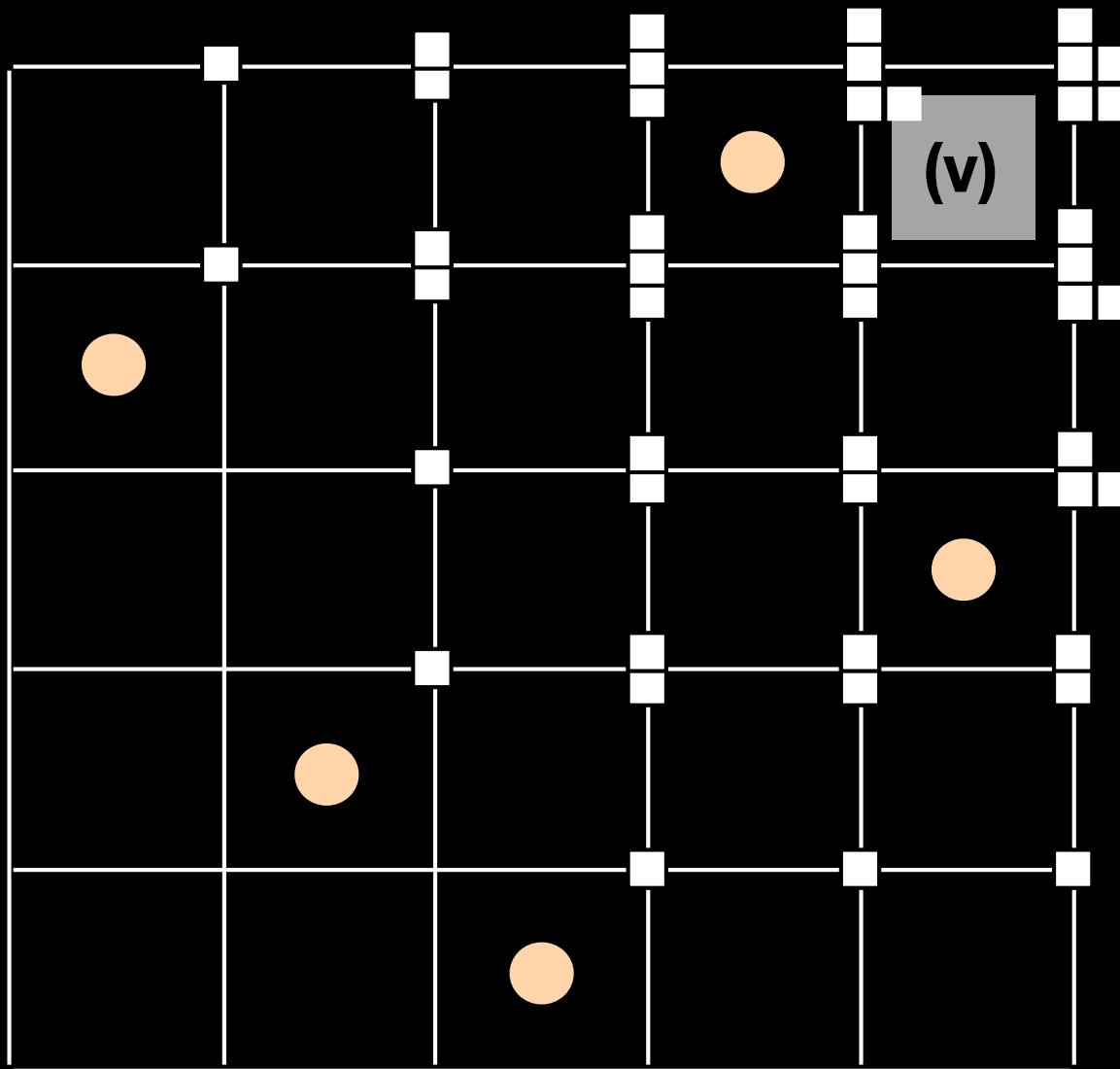
A hand-drawn diagram of a rectangle. The top edge is labeled with the Greek letter ν (nu) in blue. The right edge is labeled with the Greek letter λ (lambda) in pink. The bottom edge is labeled with the Greek letter μ (mu) in blue. The left edge is labeled with the Greek letter ρ (rho) in pink. Inside the rectangle, the text "(iii)" is written in blue, and the equation $\rho = \nu \neq \mu$ is written in blue.

A hand-drawn equation consisting of two pink arrows pointing towards each other, forming a loop, followed by the equals sign and the Greek letter μ in blue.

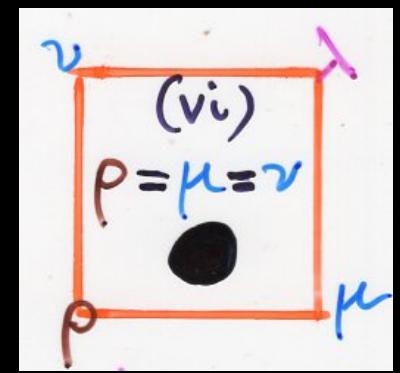
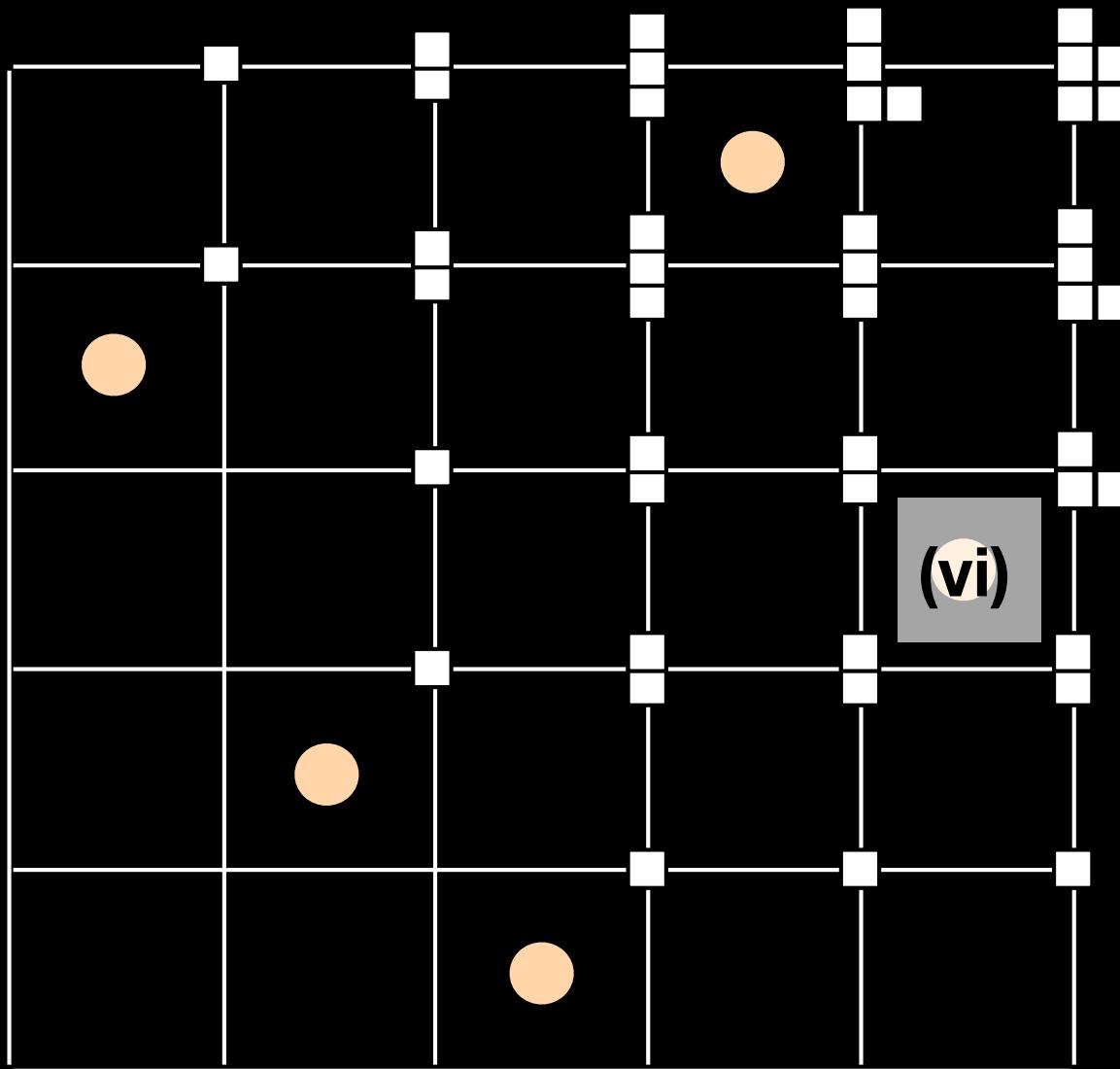


A hand-drawn diagram of a rectangle with vertices labeled p , μ , ν , and λ . The top edge is labeled ν , the bottom edge is labeled μ , the left edge is labeled p , and the right edge is labeled λ . Inside the rectangle, the text "(iv)" is written above the top edge, and below the bottom edge, there are two pairs of inequality signs: $\nu \neq p$ and $\mu \neq p$.

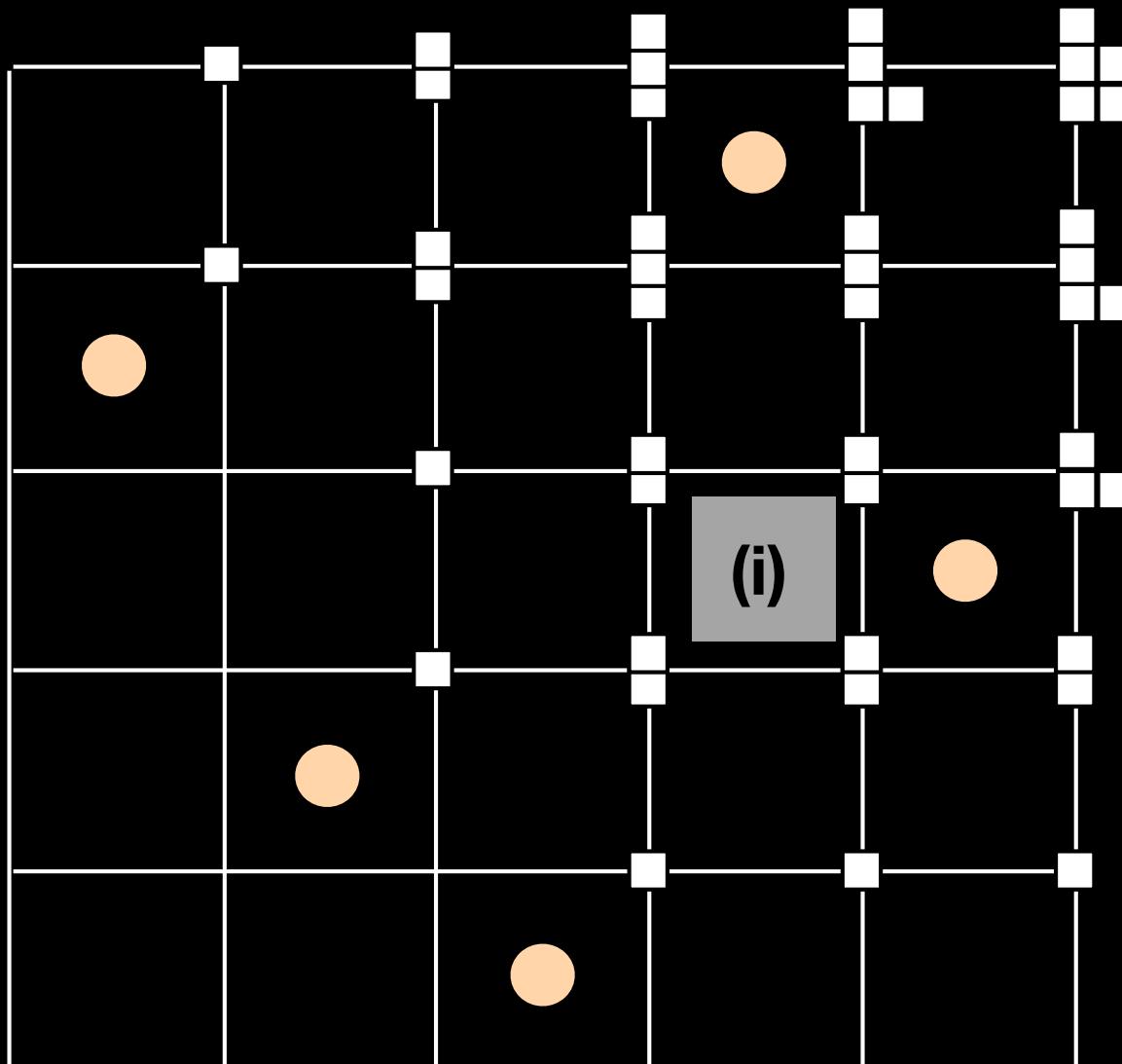
$$\lambda = \mu \cup \nu$$



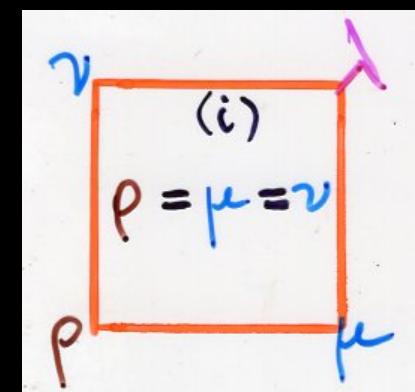
$$\lambda = \begin{cases} \mu & + (i+1) \\ v & \end{cases}$$



$$\lambda = \begin{cases} \rho \\ \mu + (1) \\ \nu \end{cases}$$



(i)



$$\lambda = \rho$$

« local rules on vertices »

Marc A. A. van Leeuwen (1996)

The Robinson-Schensted and Schützenberger algorithms, an elementary approach

C.Krattentheler, (2006).

GROWTH DIAGRAMS, AND INCREASING AND DECREASING CHAINS IN FILLINGS OF FERRERS SHAPES

M.Rubey. (2007)

Increasing and Decreasing Sequences in Fillings of Moon Polyominoes

C.Krattentheler, (2016).

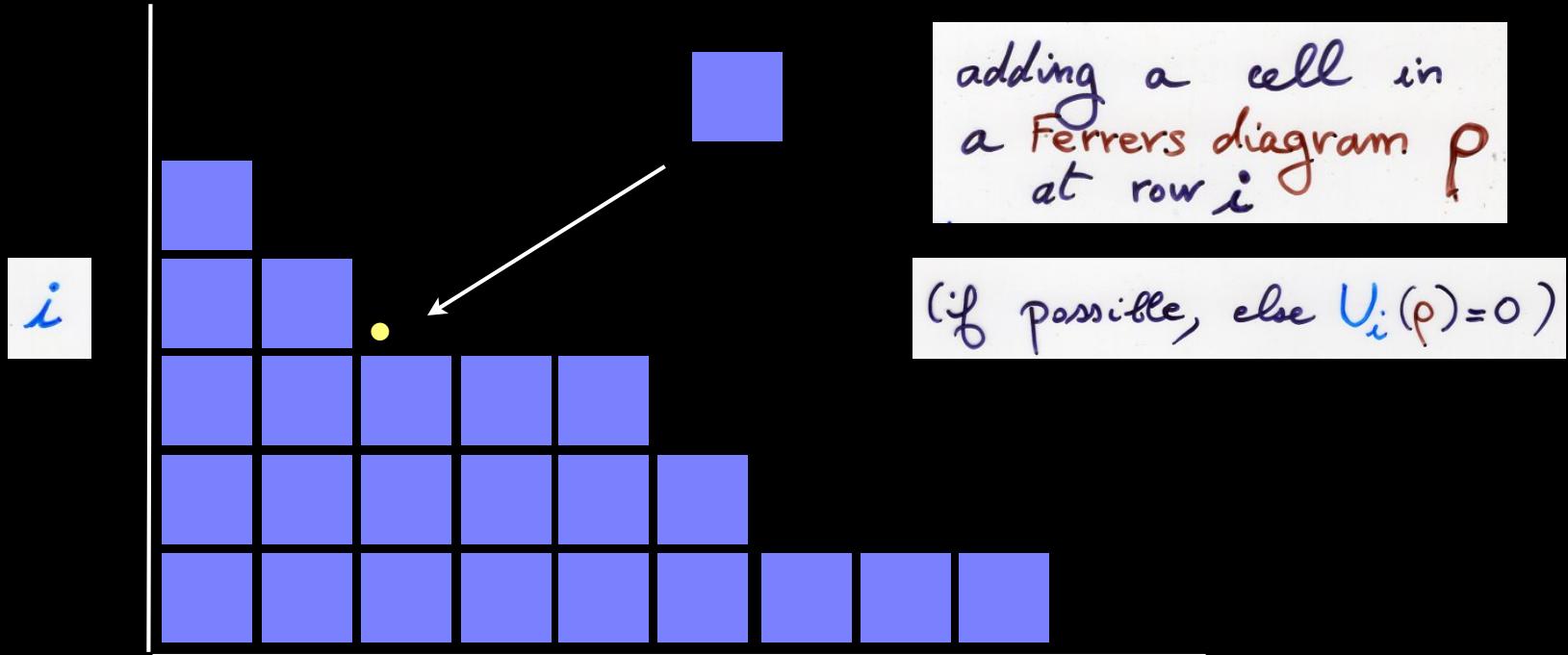
Proof of a conjecture of Burrill on oscillating tableaux

Combinatorial representation of
the algebra

$$UD = DU + \text{Id}$$

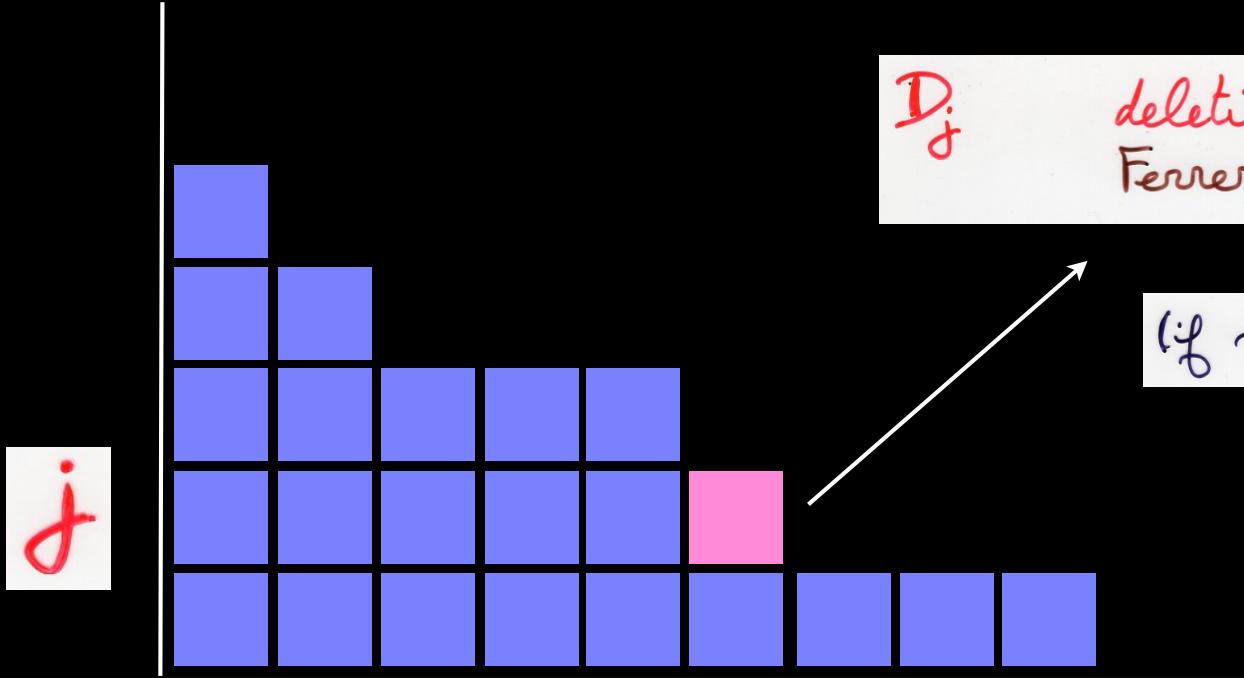
notations

operator U_i



$$U_i(\rho) = \rho + (i)$$

$$D_j(\rho) = \rho - (j)$$



D_j

deleting a cell in a
Ferrers diagram ρ at row j

(if possible, else $D_j(\rho)=0$)

$$U = \sum_{i \geq 1} U_i$$

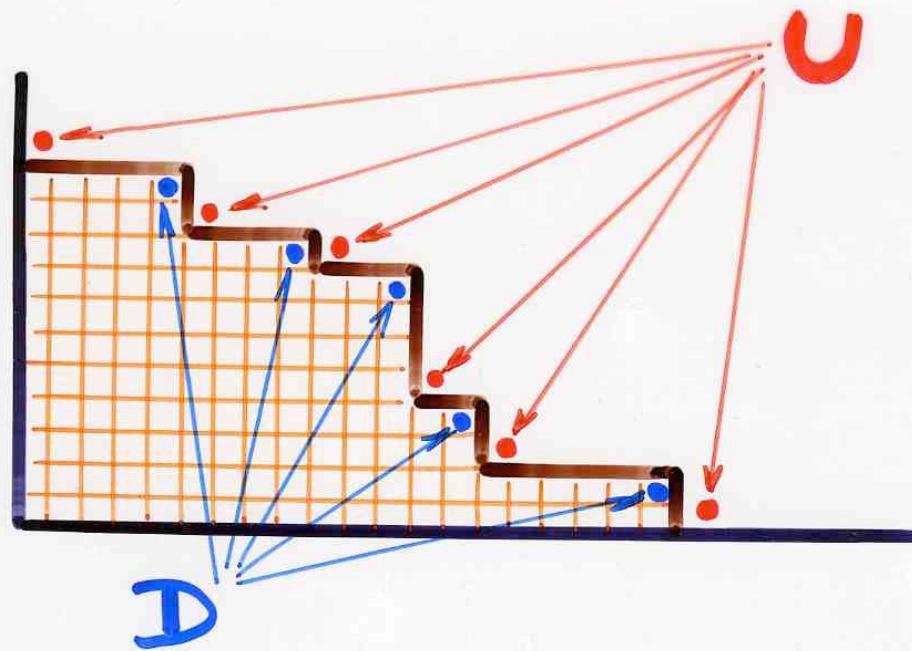
$$D = \sum_{i \geq 1} D_i$$

U and D are operators acting on
the vector space generated by Ferrers
diagrams.

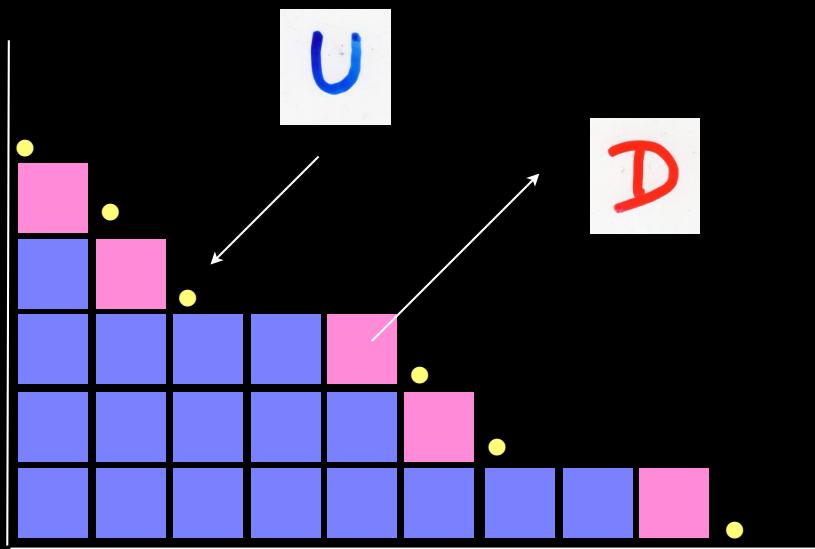
$$U \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|}\hline & & \\ \hline & & \\ \hline & & \\ \hline \end{array} \end{array} = \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|c|}\hline & & & \\ \hline & & & \\ \hline & & & \\ \hline & & & \color{blue} \blacksquare \\ \hline \end{array} \end{array} + \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|c|}\hline & & & \\ \hline & & & \\ \hline & & & \color{blue} \blacksquare \\ \hline & & & \color{blue} \blacksquare \\ \hline \end{array} \end{array} + \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|c|}\hline & & & \\ \hline & \color{blue} \blacksquare & & \\ \hline \end{array} \end{array}$$

$$D \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|}\hline & & \\ \hline & & \\ \hline & & \\ \hline \end{array} \end{array} = \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|}\hline & & \\ \hline & & \\ \hline & & \color{red} \bullet \\ \hline \end{array} \end{array} + \begin{array}{c} \\ \text{Ferrers diagram: } \begin{array}{|c|c|c|}\hline & & \\ \hline & & \\ \hline & & \\ \hline \end{array} \end{array} \color{red} \bullet$$

$$UD = DU + I$$

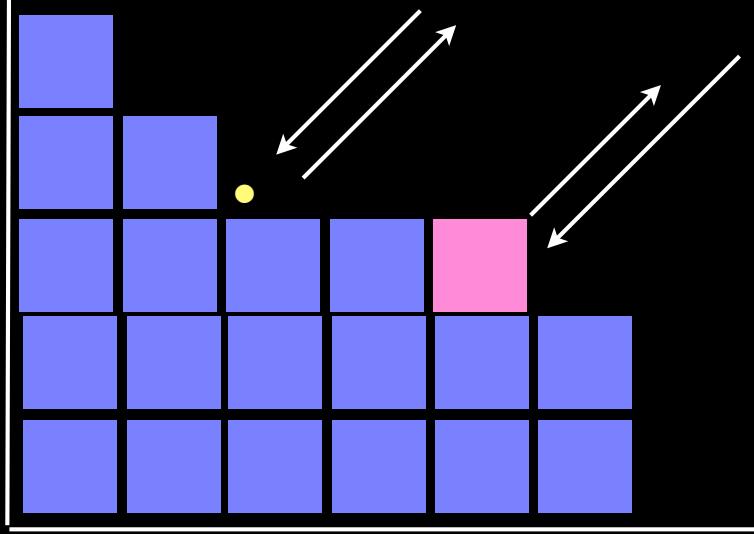
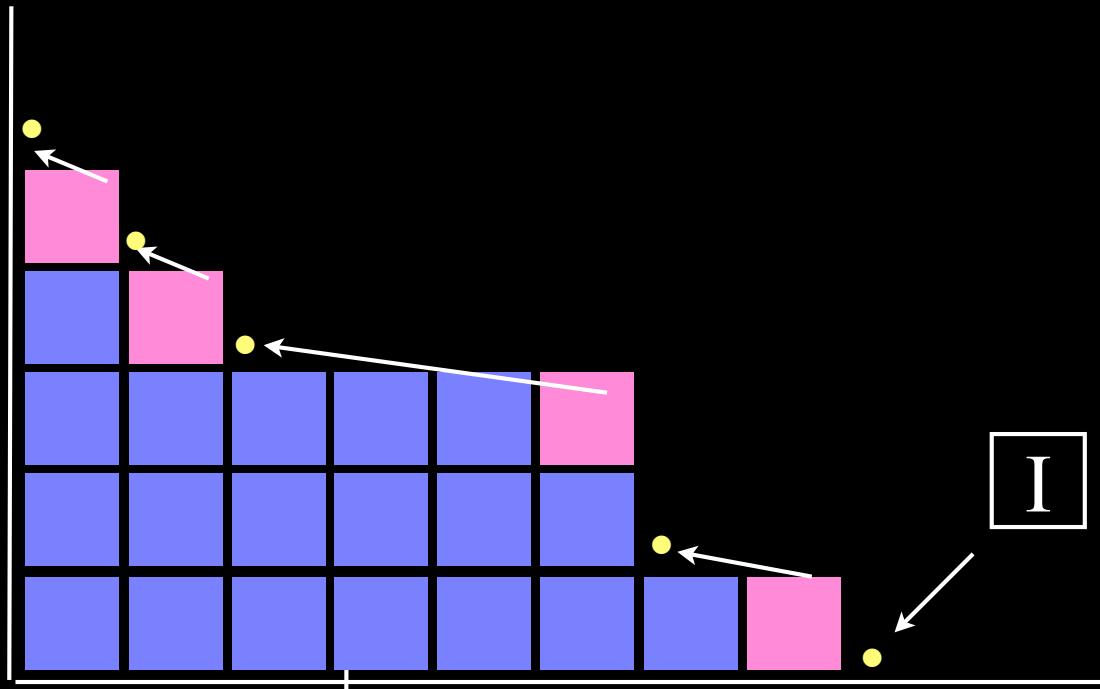
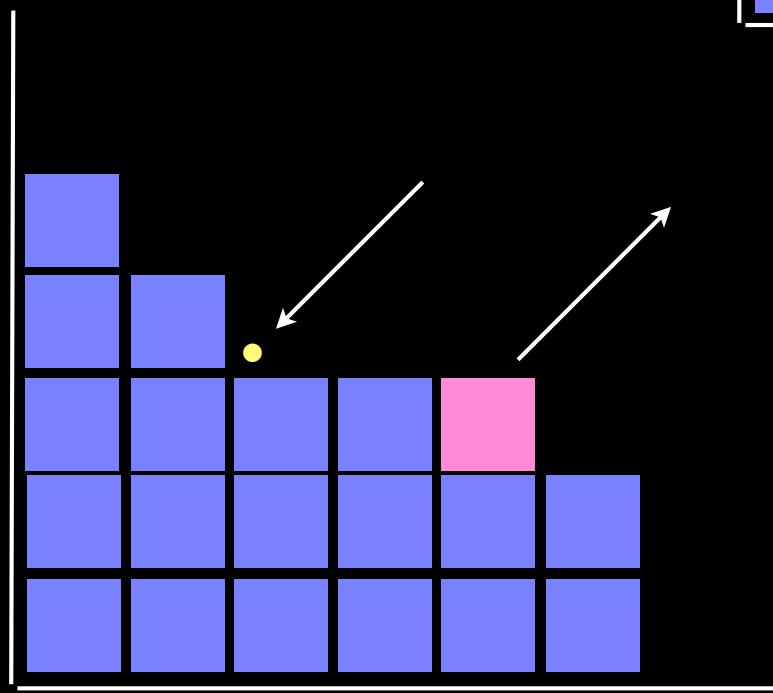


operators
 U and D



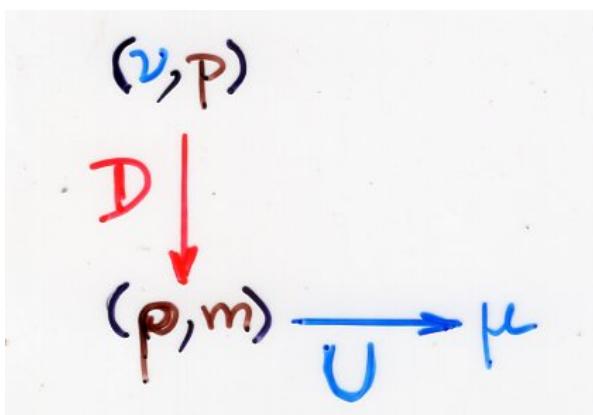
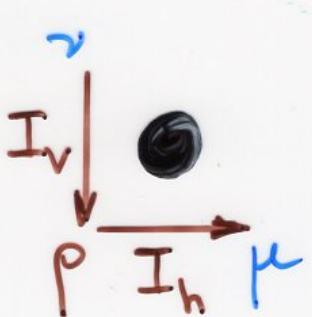
Young lattice

{ U adding
 D deleting a cell in a Ferrers
 diagram

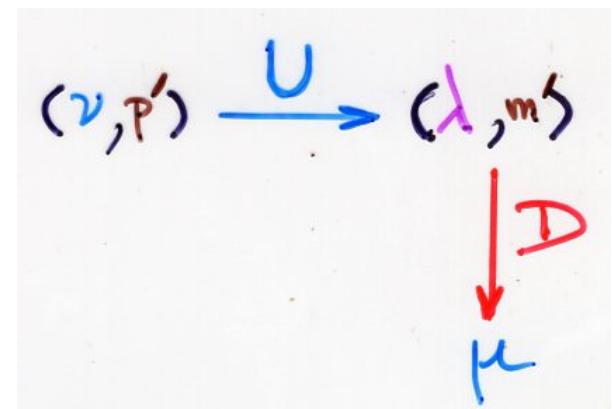


$$UD = DU + I_v I_h$$

"commutation diagrams"

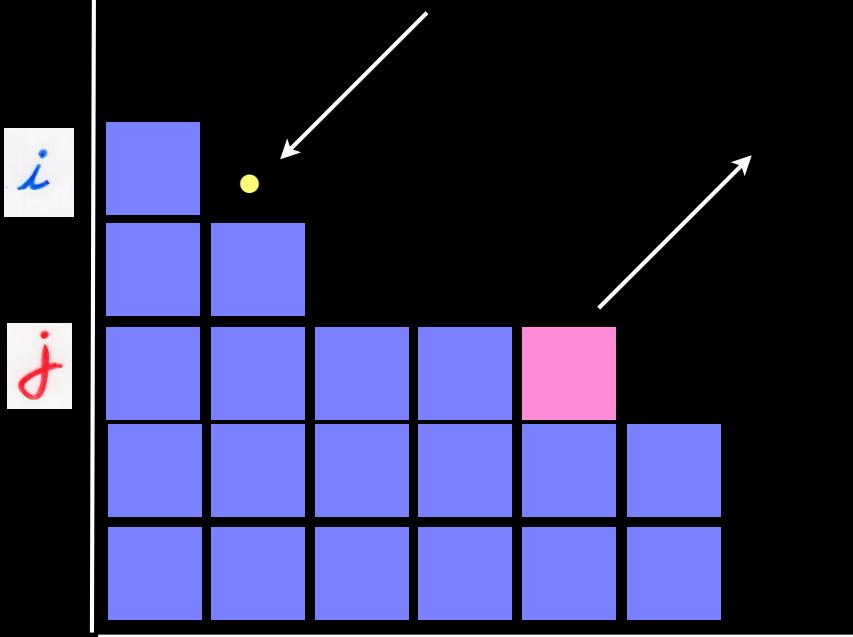
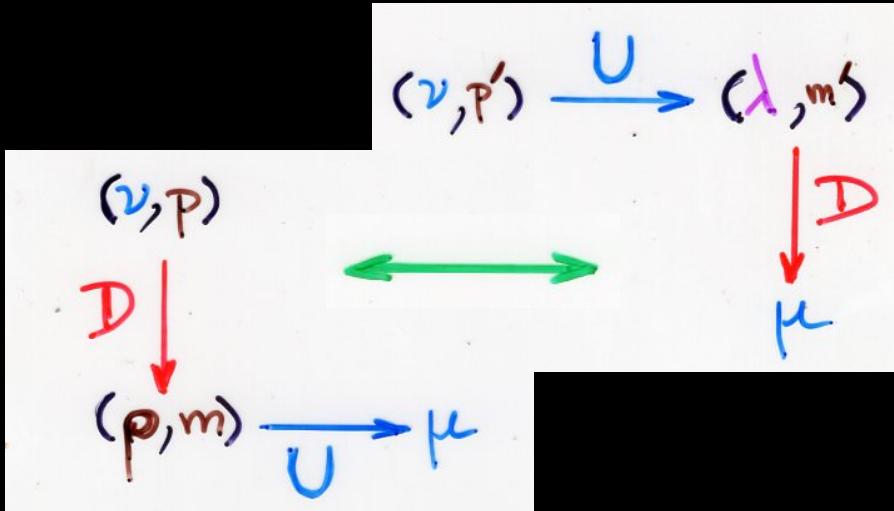


bijection



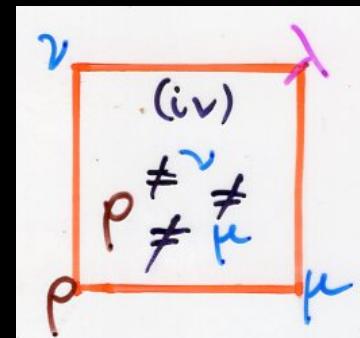
p, m, p', m' are "positions"

in v, p, v, λ respectively

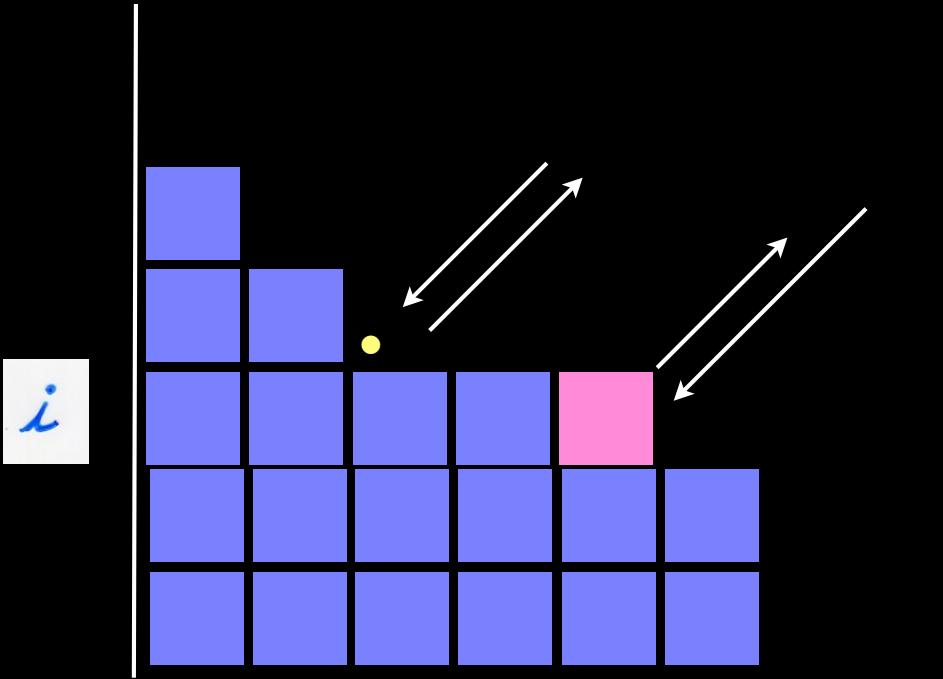
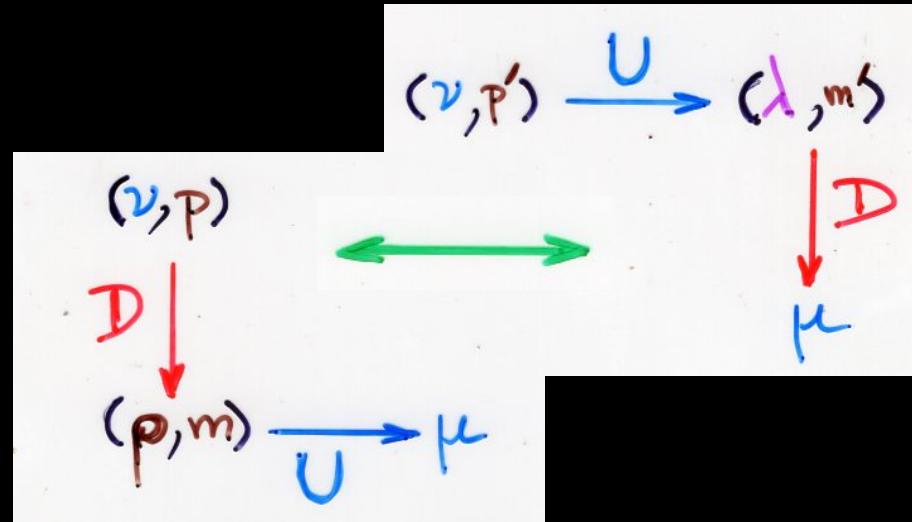


$$\begin{aligned}
 p &= j \\
 m &= i
 \end{aligned}$$

$$\begin{aligned}
 p' &= i \\
 m' &= j
 \end{aligned}$$

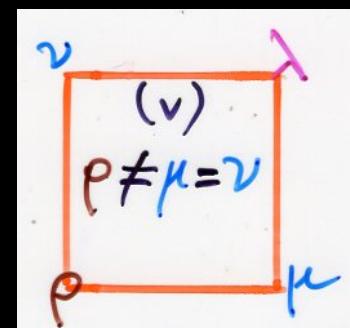


$$\begin{aligned}
 \nu &= p + (j) \\
 \mu &= p + (i) \\
 \lambda &= p + (i) + (j)
 \end{aligned}$$

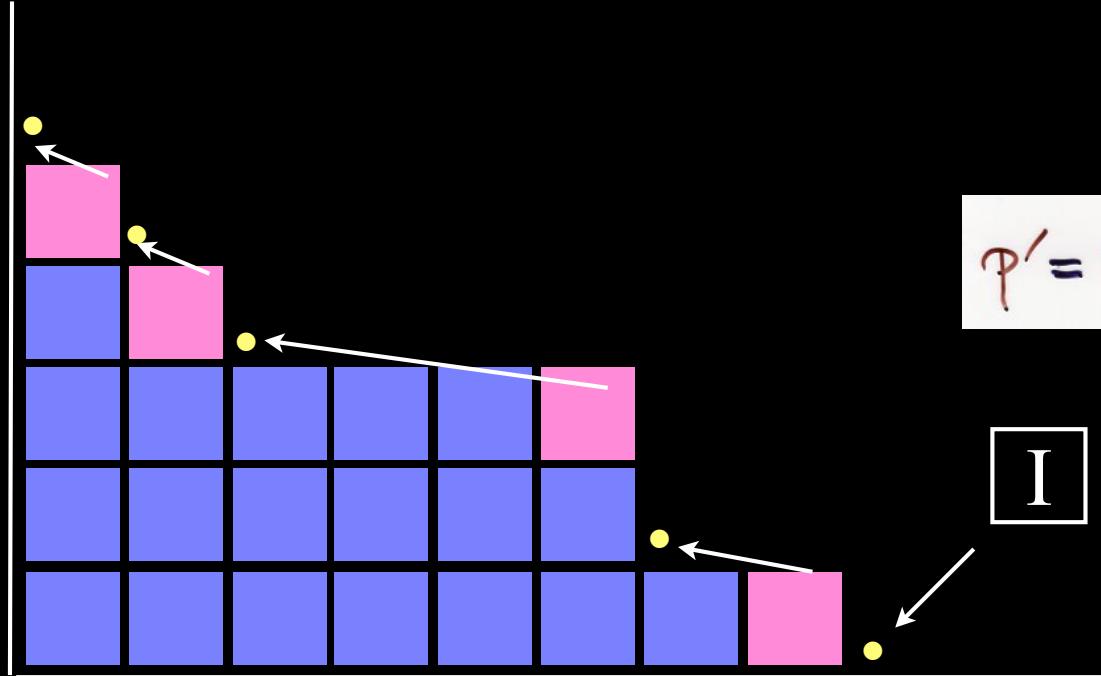
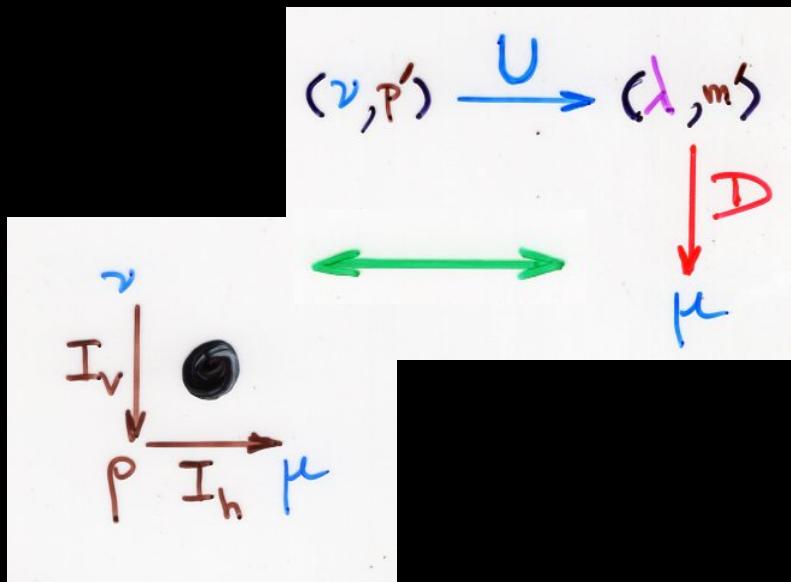


$$\begin{aligned}
 p &= i \\
 m &= i
 \end{aligned}$$

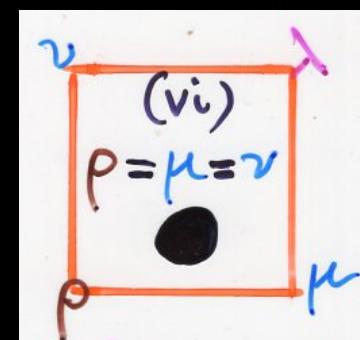
$$\begin{aligned}
 p' &= i+1 \\
 m' &= i+1
 \end{aligned}$$



$$\begin{aligned}
 \mu &= v = p + (i) \\
 \lambda &= \mu + (i+1)
 \end{aligned}$$



$$p' = m' = 1$$



$$\lambda = \begin{cases} p \\ \mu + (1) \\ v \end{cases}$$

3	
2	5
1	4

1

2

3

1

2

2

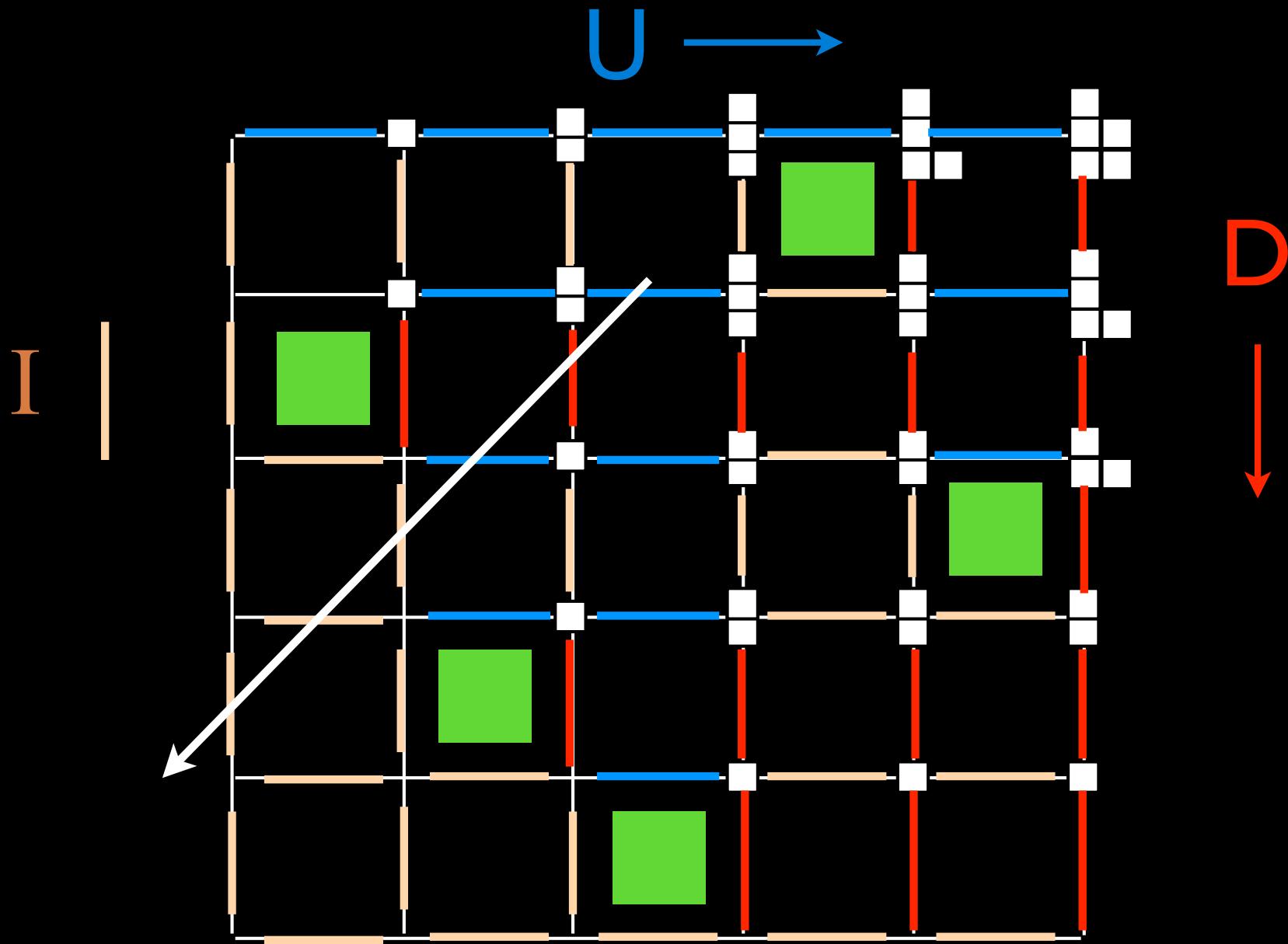
3

1

2

1

4	
2	5
1	3

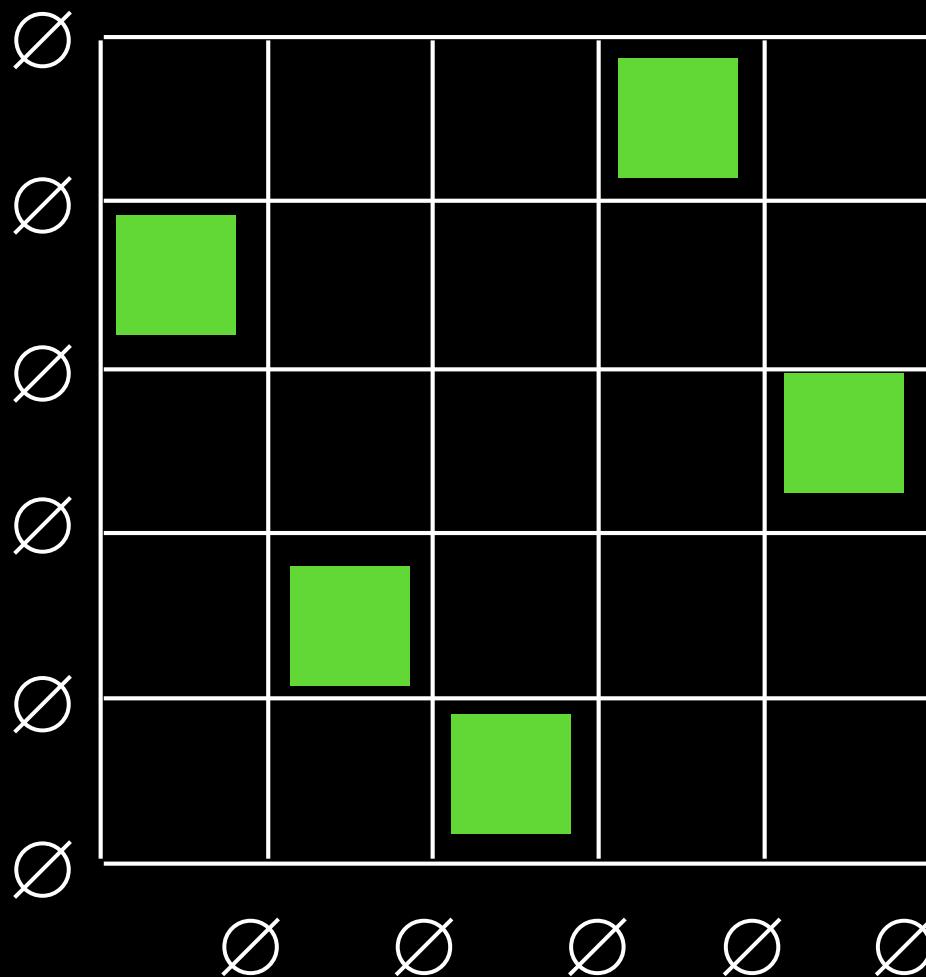


This "propagation" algorithm is
exactly the reverse of Fomin's "growth
diagrams"

I

3	
2	5
1	4

1 2 3 1 2



4	
2	5
1	3

"The **cellular** ansatz."

quadratic
algebra **Q**

$$UD = DU + \text{Id}$$

Q-tableaux

combinatorial objects
on a 2D lattice

permutations

towers placements

representation of **Q**
by combinatorial
operators

bijections

RSK

pairs of
Young tableaux

(i) first step

(ii) second step

quadratic
algebra **Q**

Q-tableaux

$$UD = DU + \text{Id}$$

permutations

Planarization of the rewriting rules

$$UD = DU + Id$$

commutations

$$UD \rightarrow DU \quad UD \rightarrow Id$$

rewriting rules

UUDD

$$UUDD = UDUD + UD$$

$$= D U U D + 2 U D$$

$$= (DUDU + DU) + 2(DU + Id)$$

$$= (DDUU + 2DU) + 2(DU + Id)$$

$$= DDUU + 4DU + 2 Id$$

$$UD = D U + Id$$

commutations

Lemma Every word w with letters U and D can be written in a unique way

$$w = \sum_{i,j \geq 0} c_{ij}(w) D^i U^j$$

normal ordering
in physics

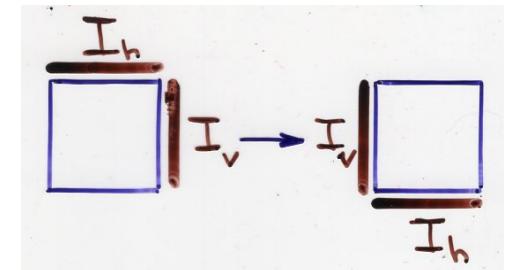
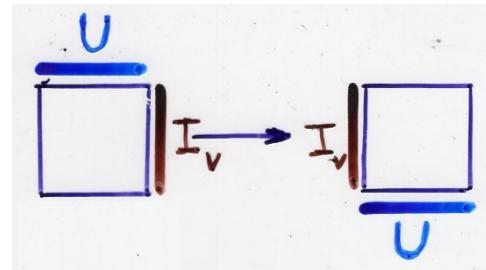
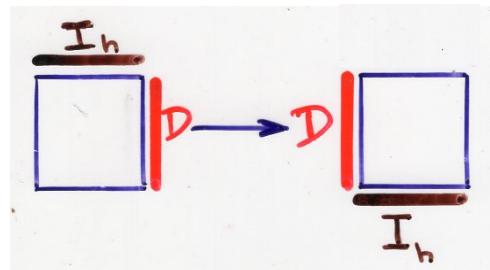
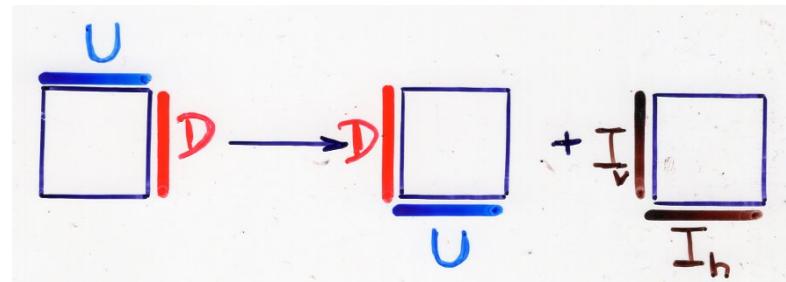
$$UD = DU + Id$$

commutations

$$UD \rightarrow DU \quad UD \rightarrow Id$$

rewriting rules

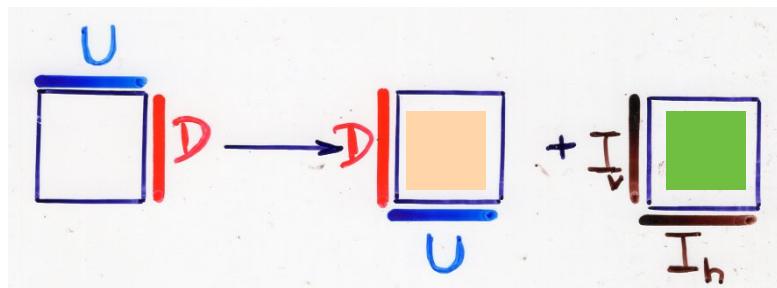
planarization of the rewriting rules



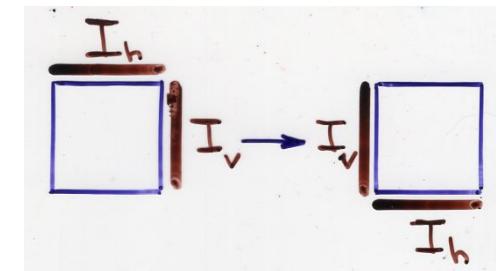
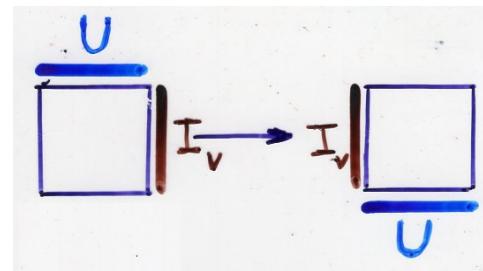
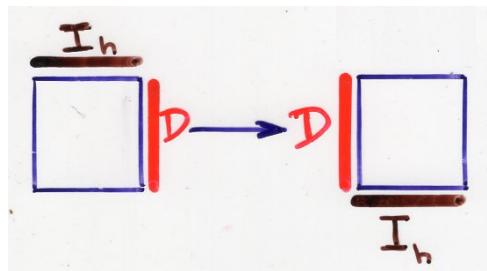
homogenization
of the system
of commutations
relations

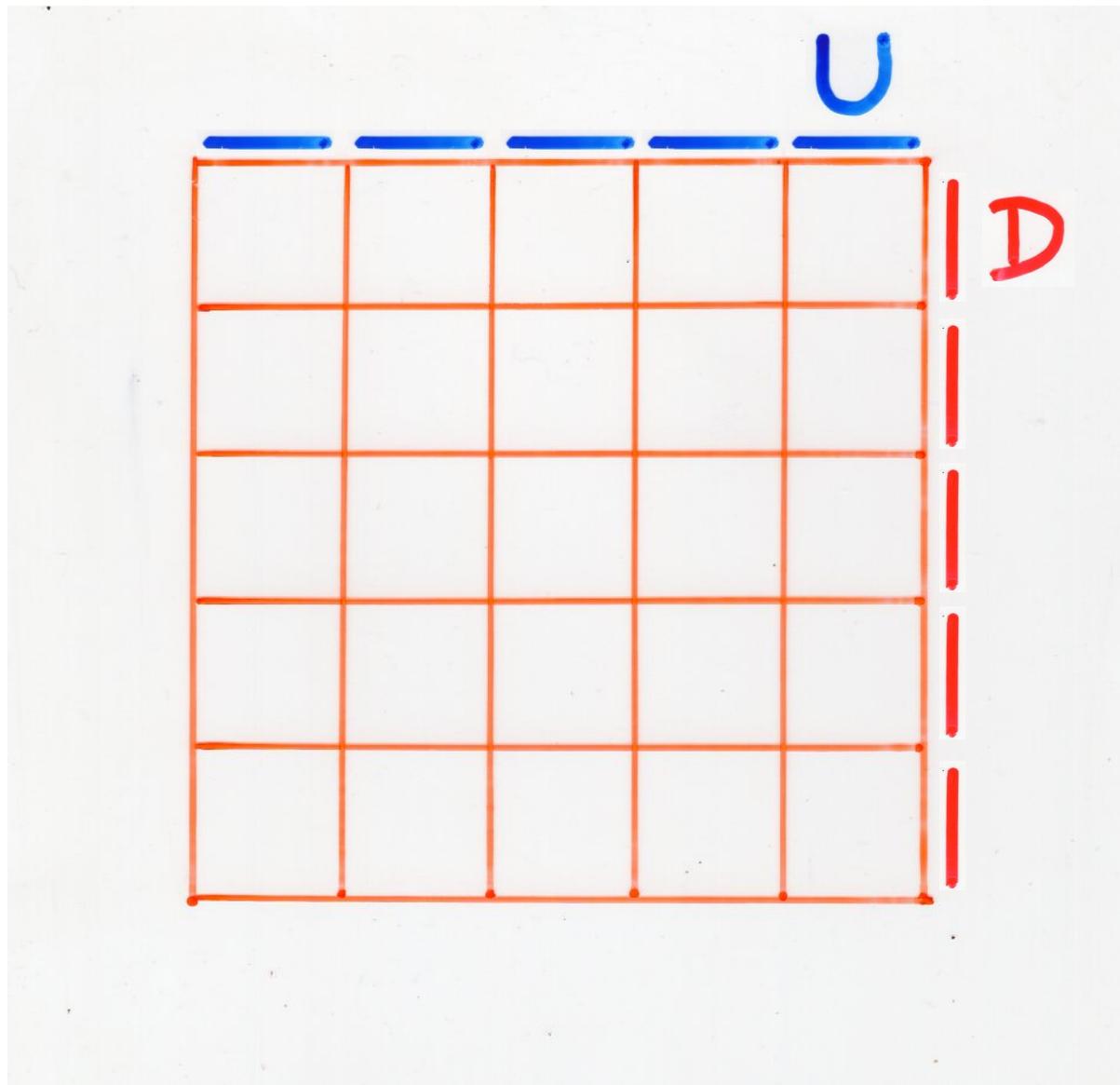
$$\left\{ \begin{array}{l} UD = DU + I_v I_h \\ UI_v = I_v U \\ I_h D = DI_h \\ I_h I_v = I_v I_h \end{array} \right.$$

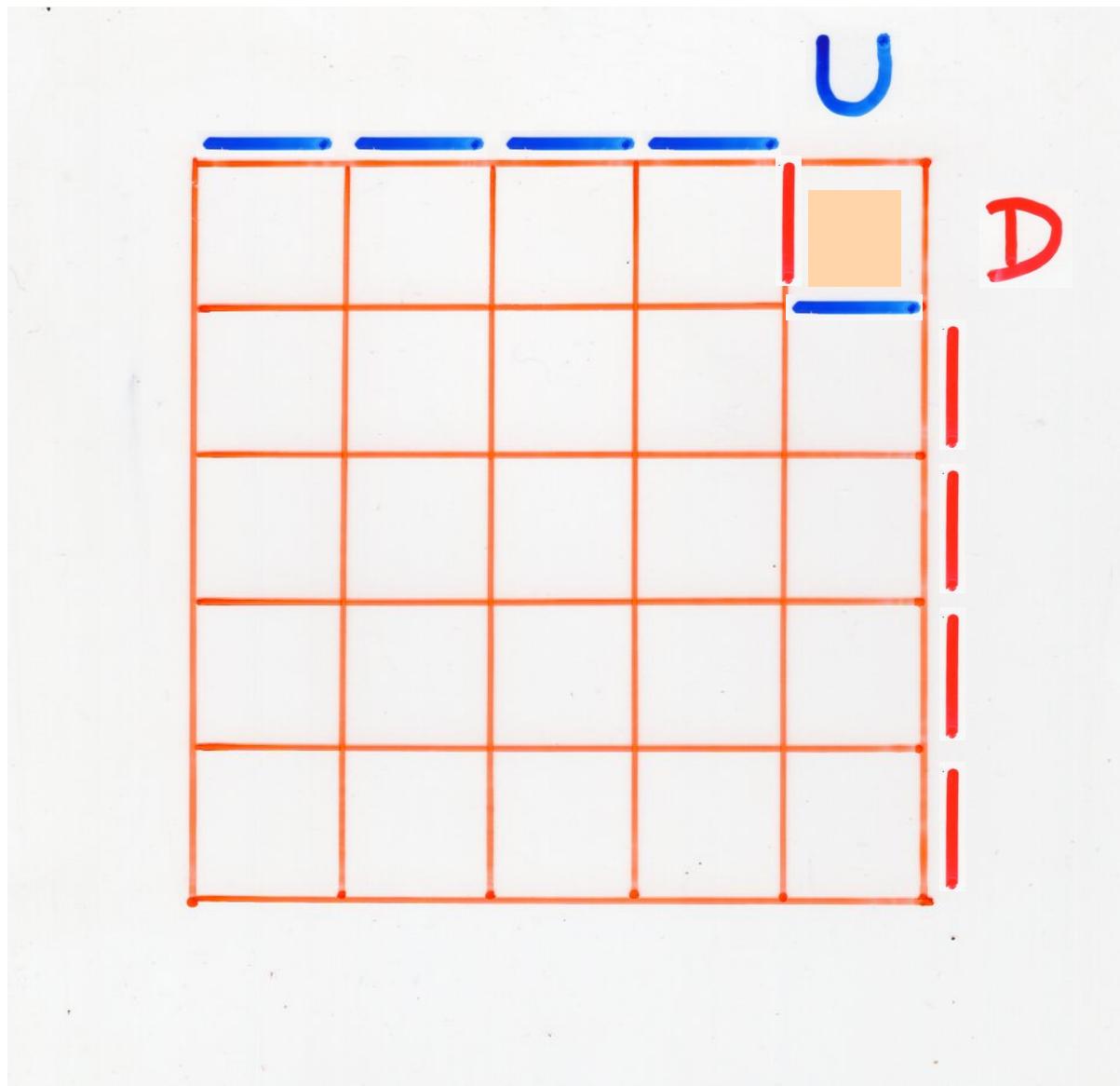
$$\left\{ \begin{array}{l} UD \rightarrow DU \\ UI_v \rightarrow I_v U \\ I_h D \rightarrow DI_h \\ I_h I_v \rightarrow I_v I_h \end{array} \right. \quad \begin{array}{l} UD \rightarrow I_v I_h \\ \text{rewriting rules} \end{array}$$

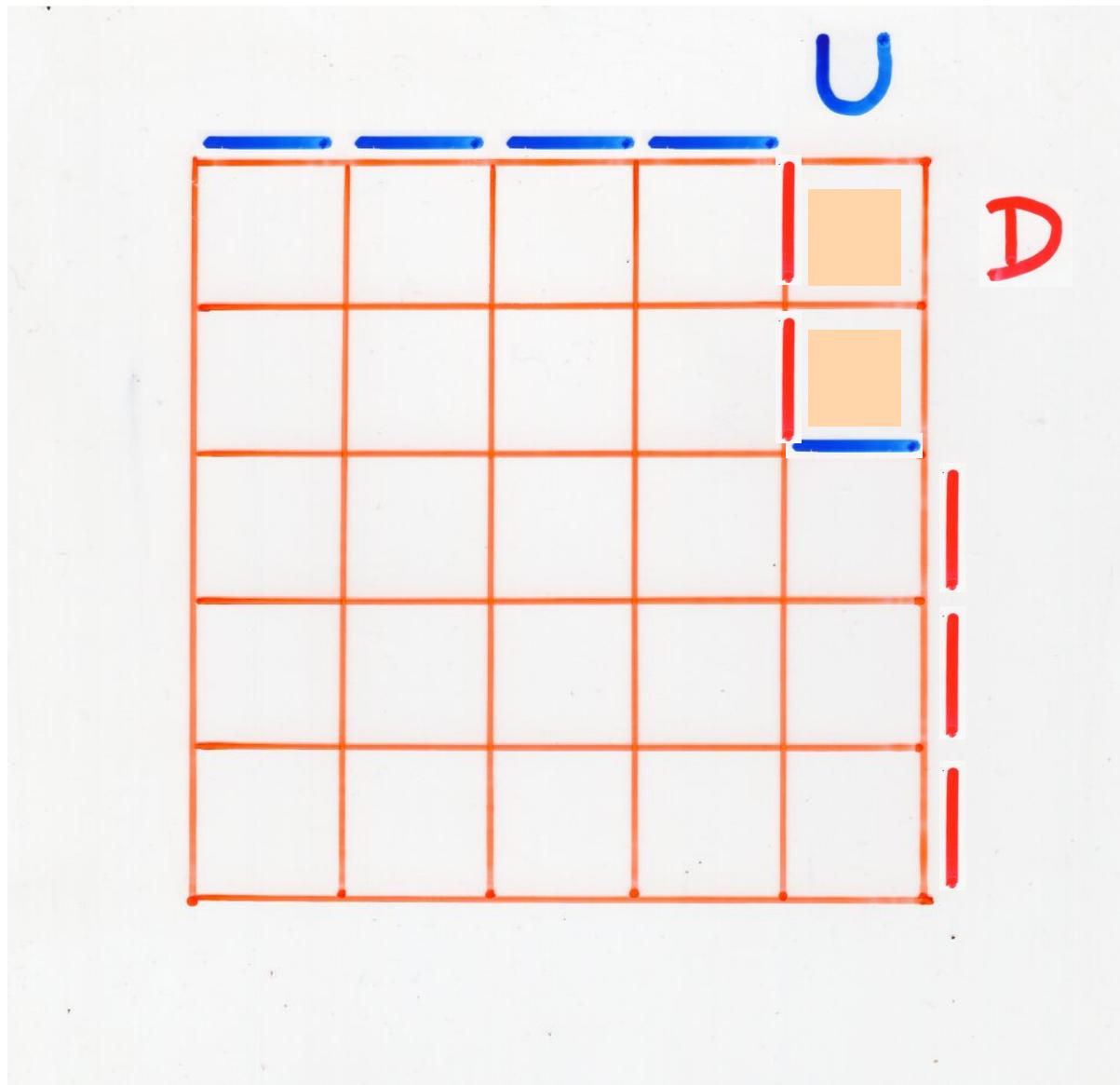


"planarization" of the "rewriting rules"

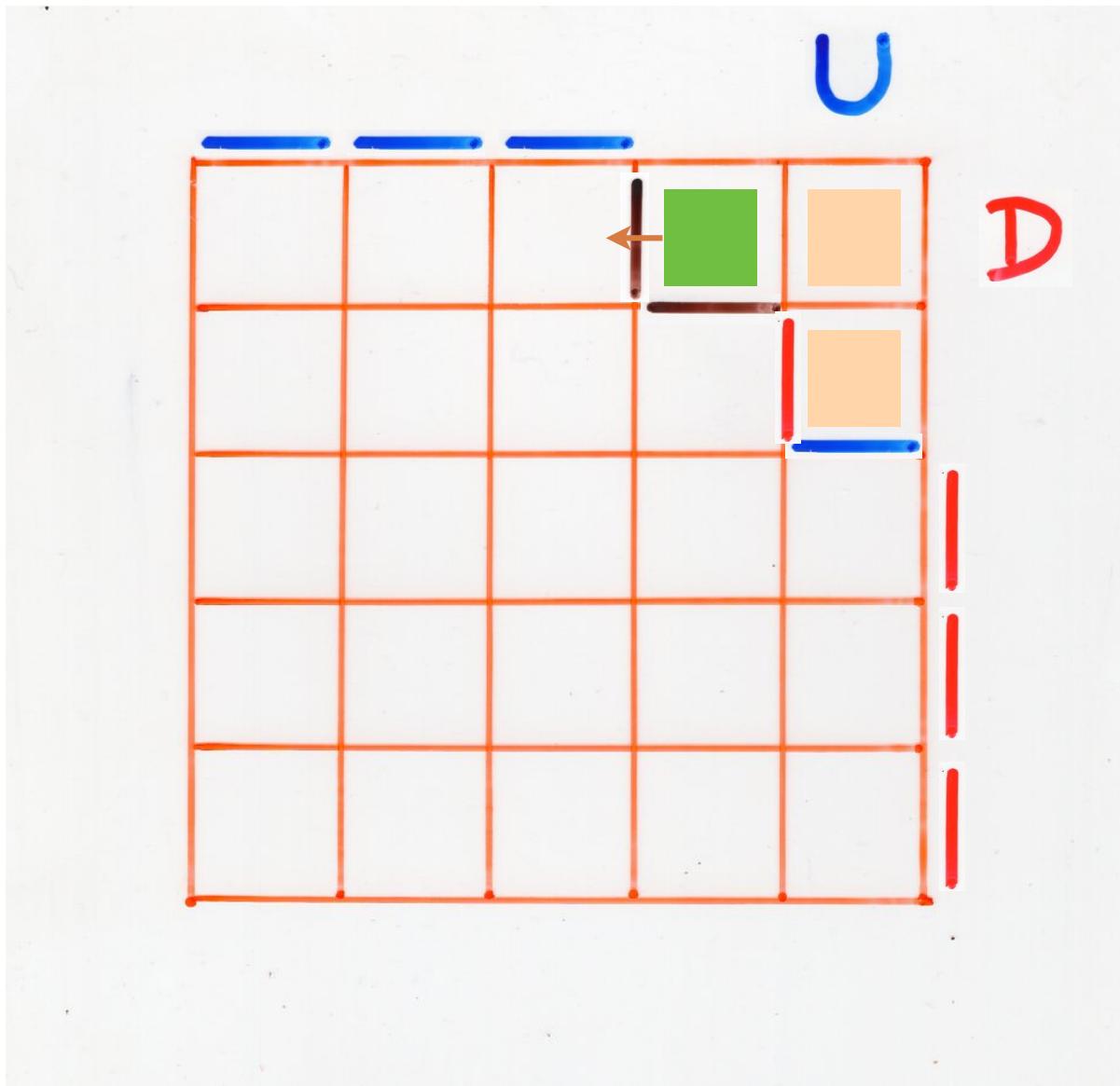


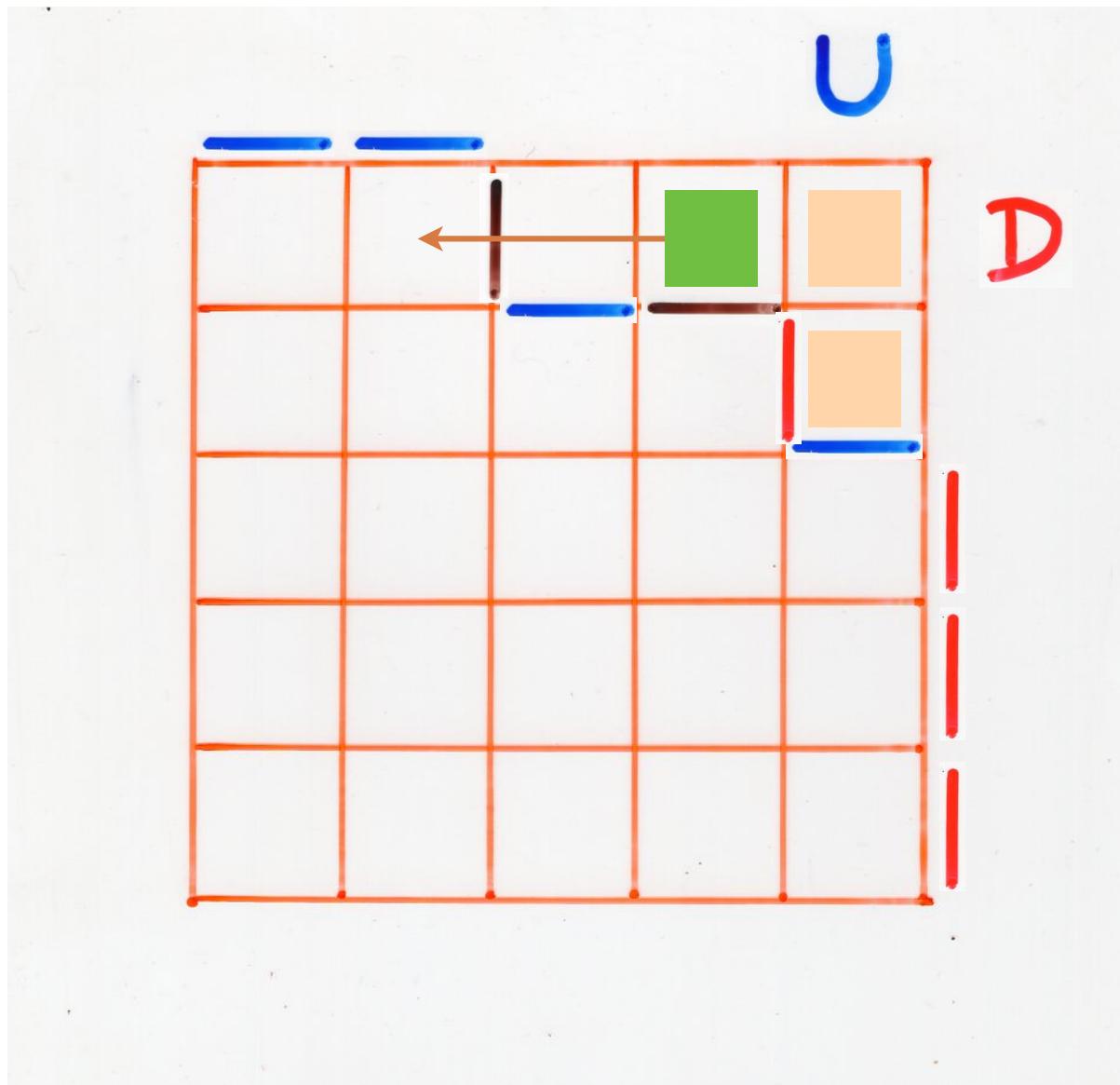


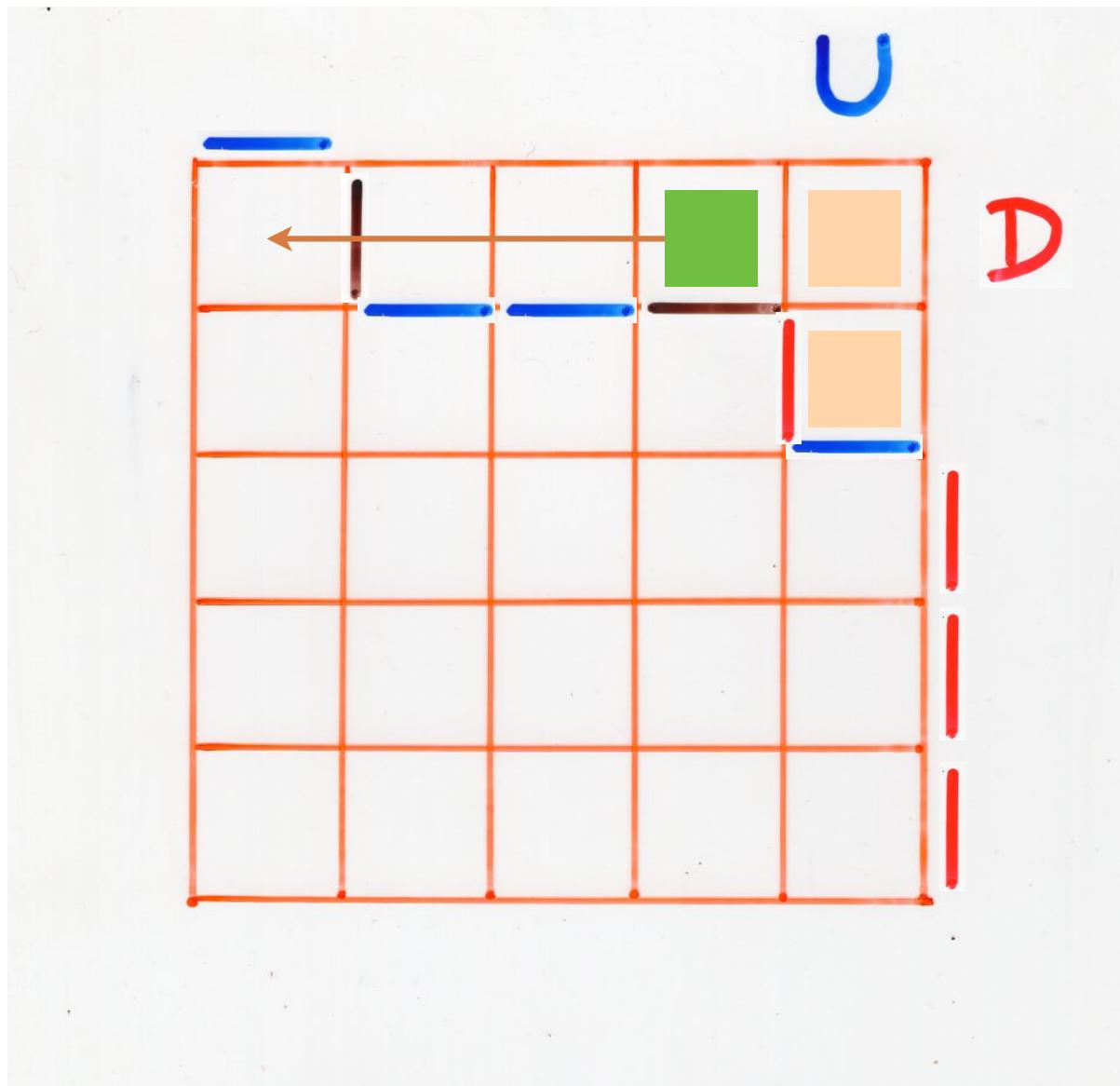


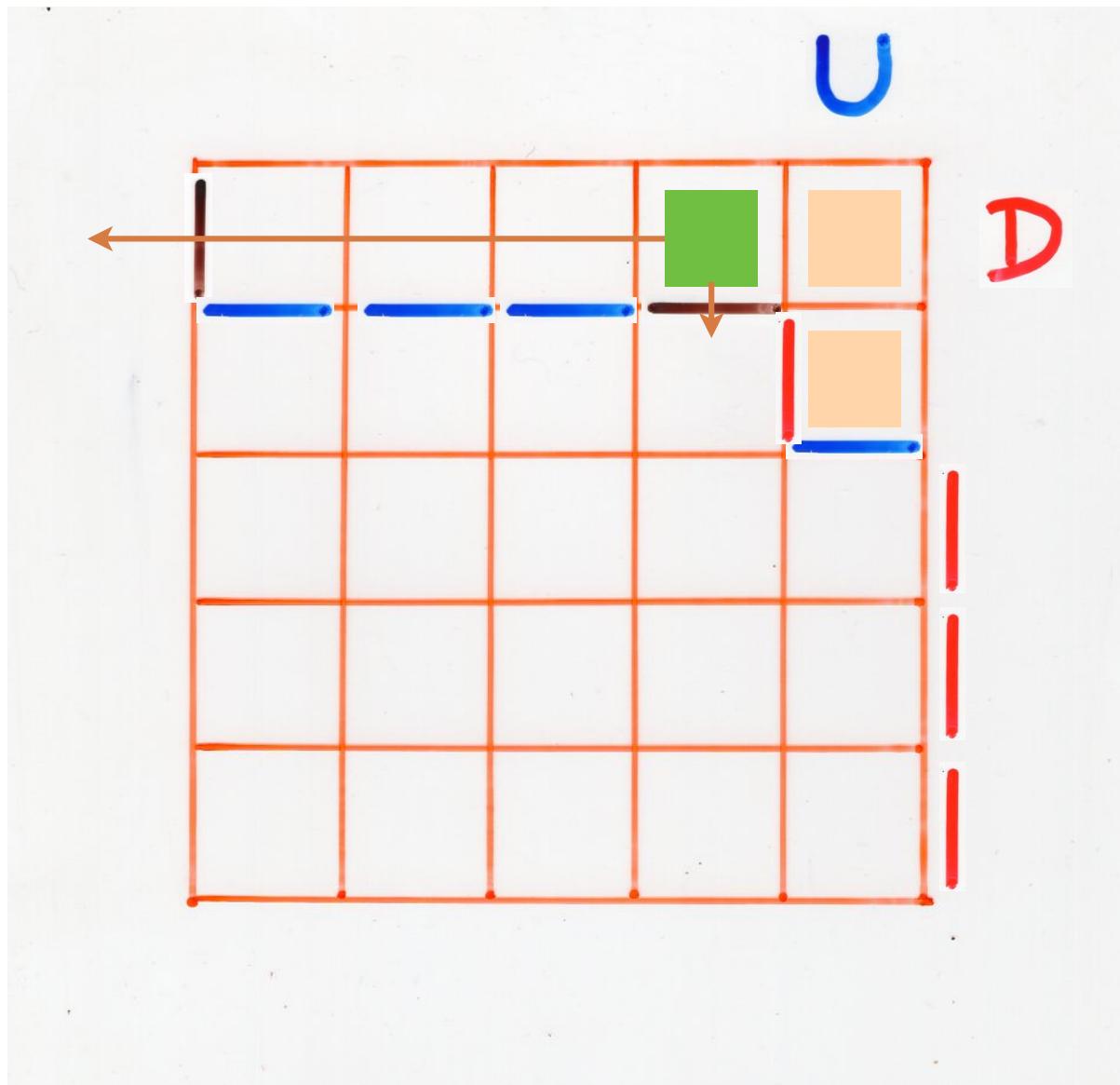


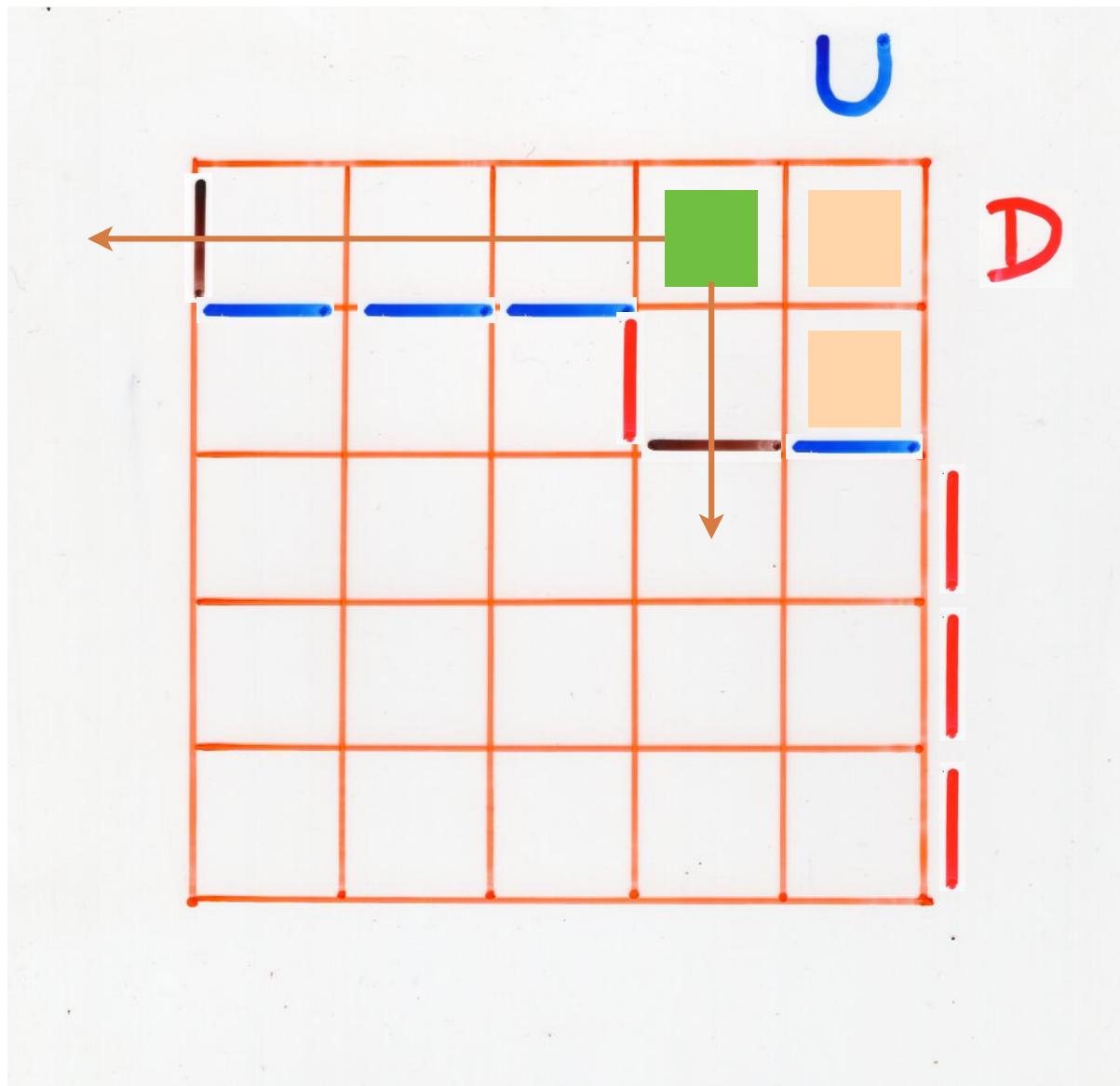
$\text{---} \quad I_h$

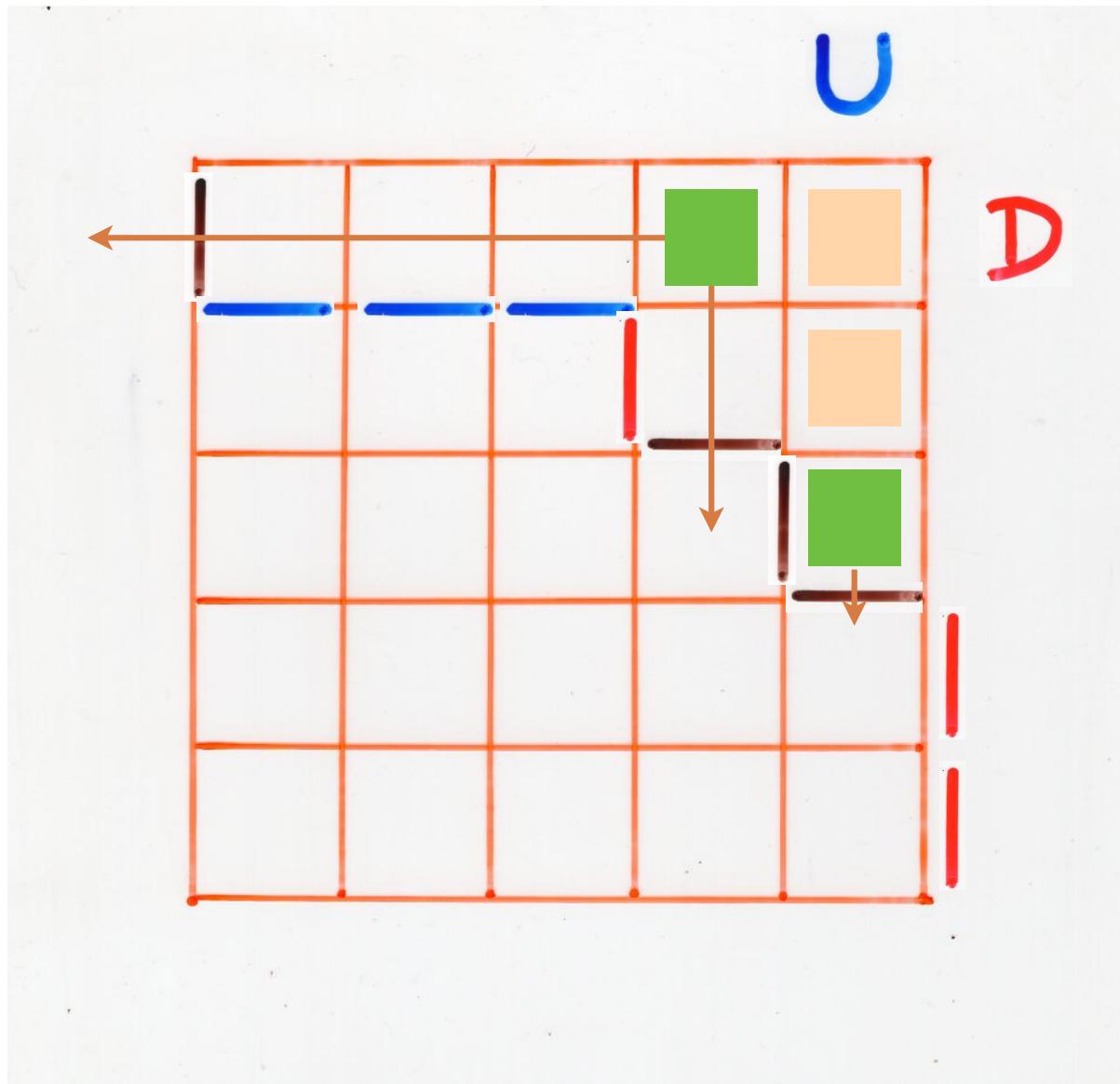


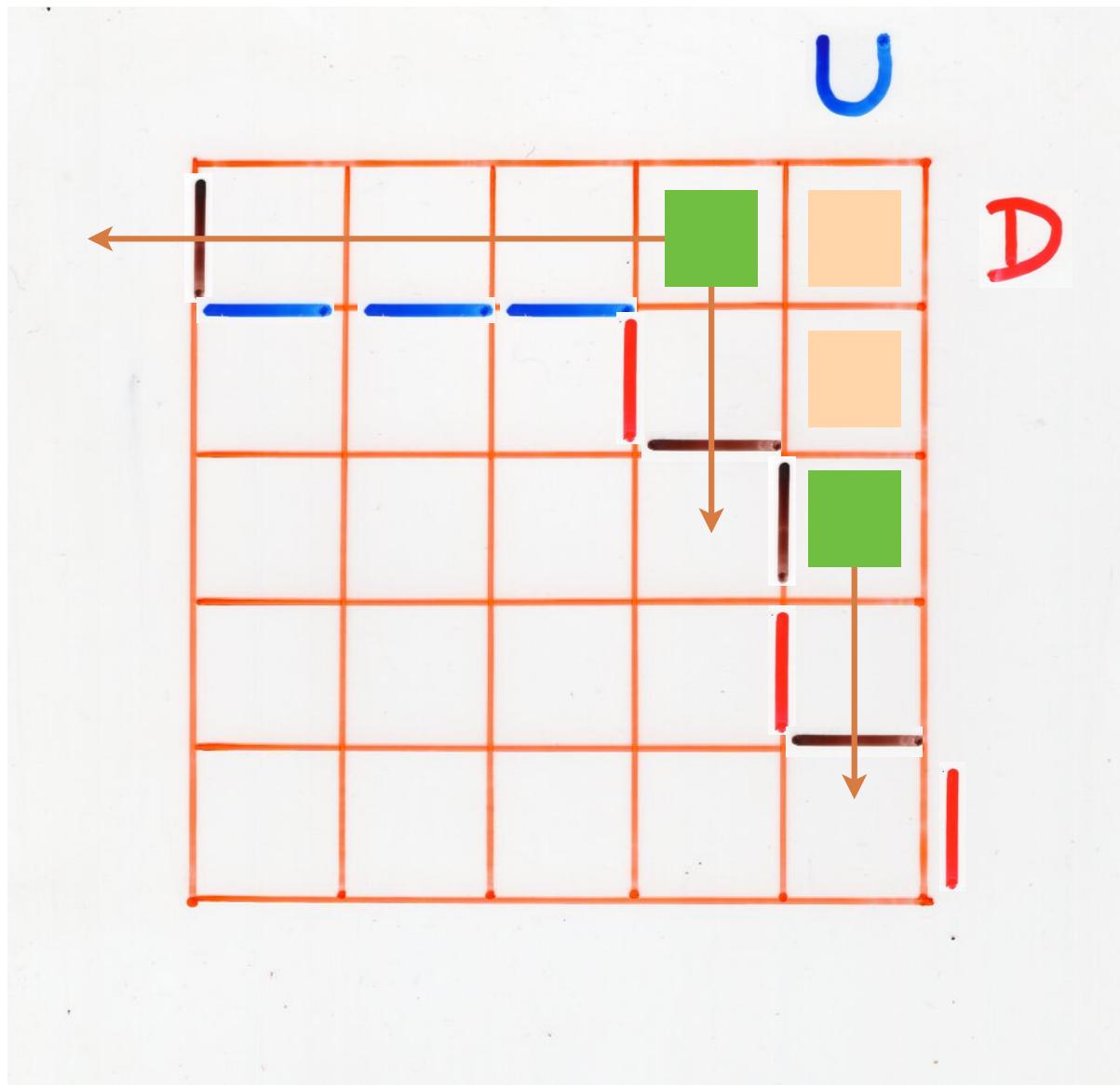


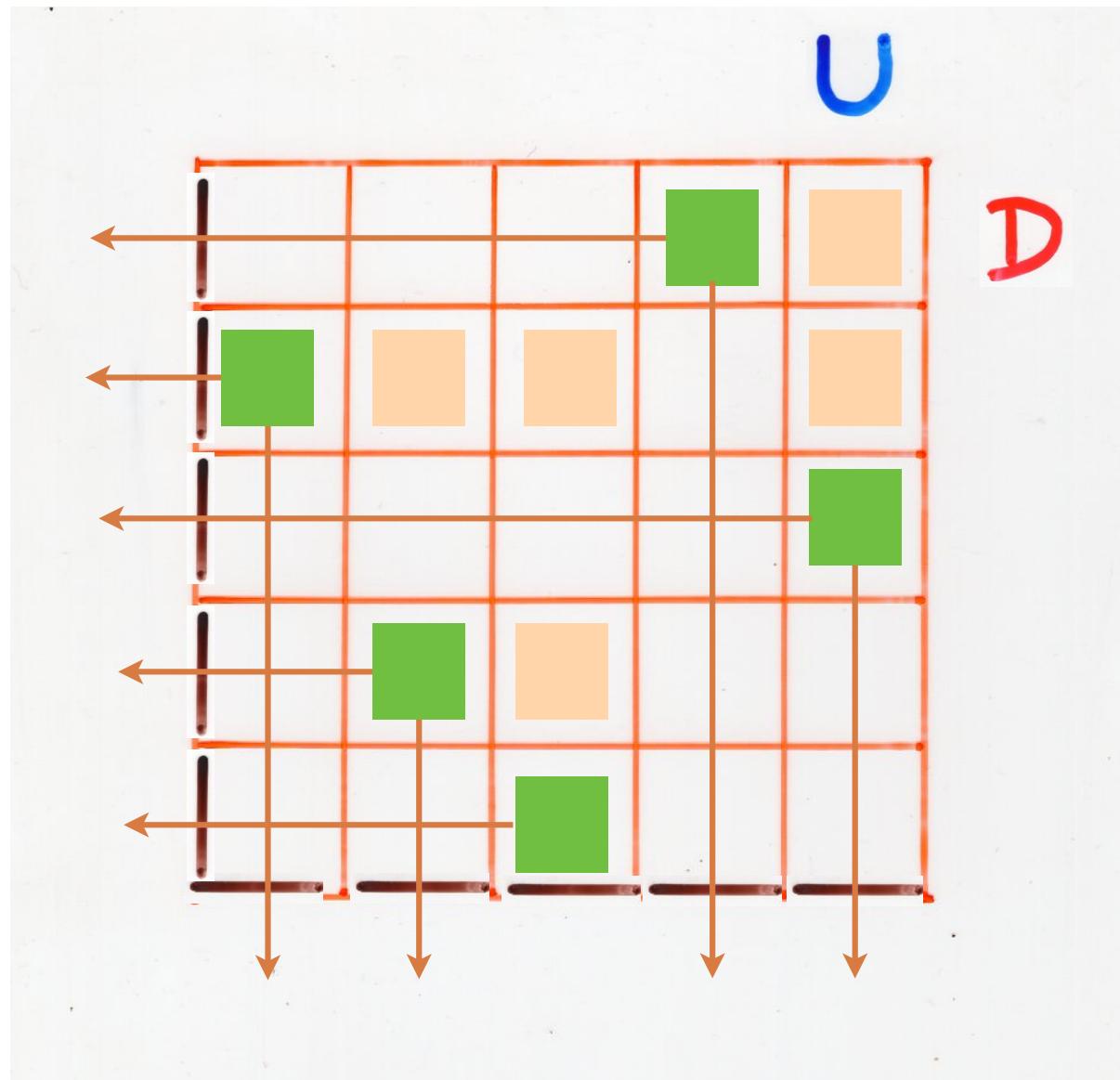


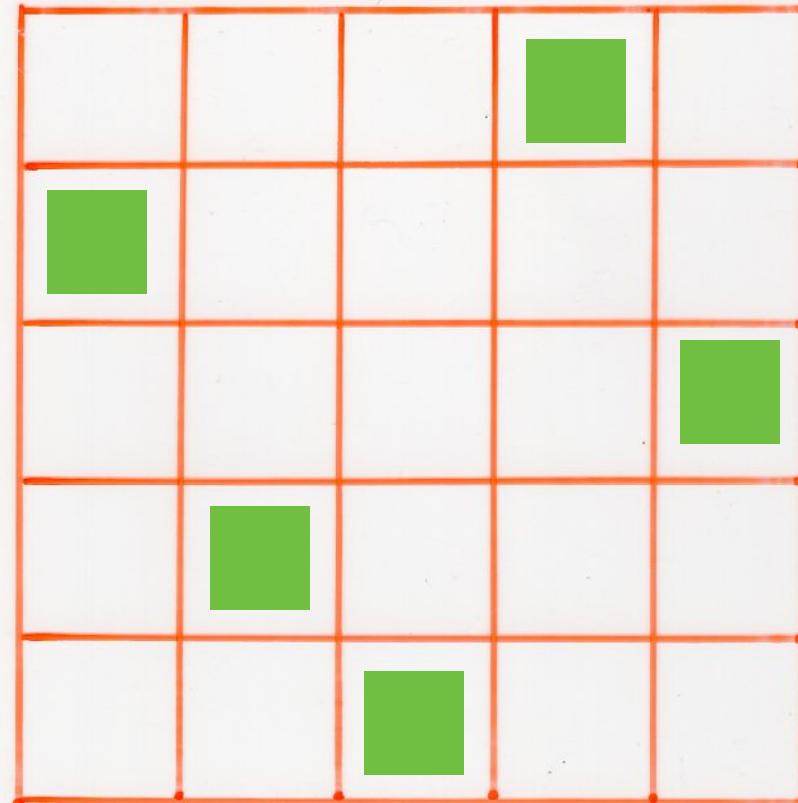












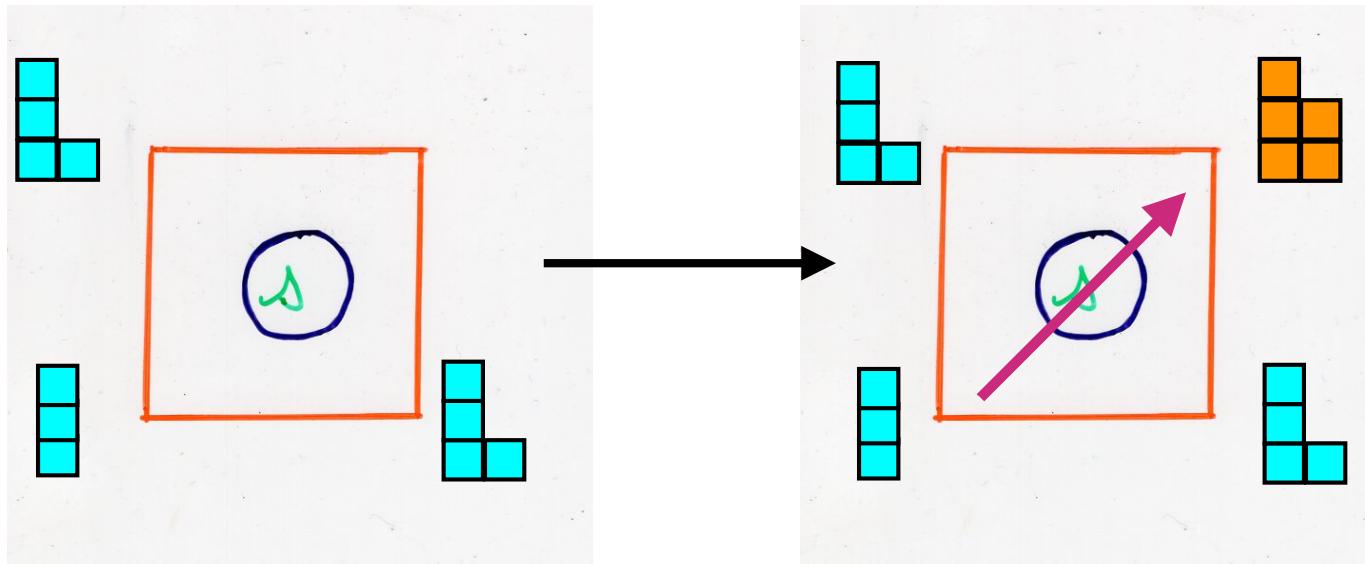
quadratic
algebra **Q**

$$UD = DU + \text{Id}$$

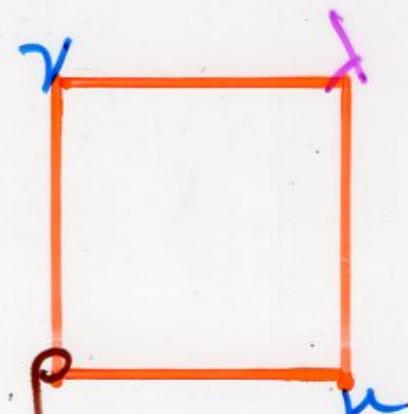
permutation
as a **Q**-tableau

From « local rules » on vertices
to
« local rules » on edges

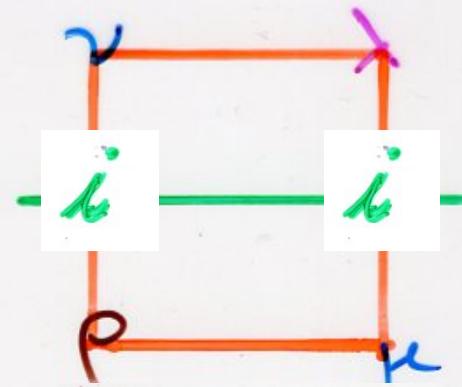
Fomin's
"local rules"
"growth diagrams"



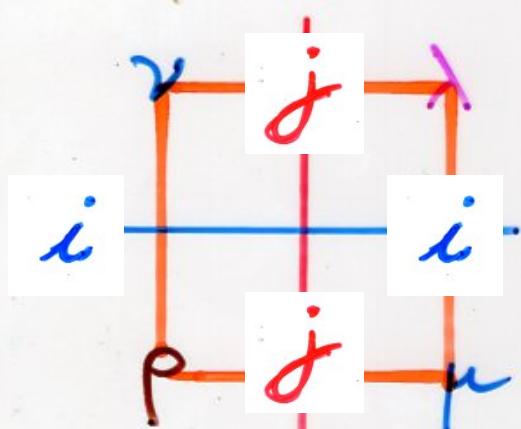
"local rules"
on the vertices



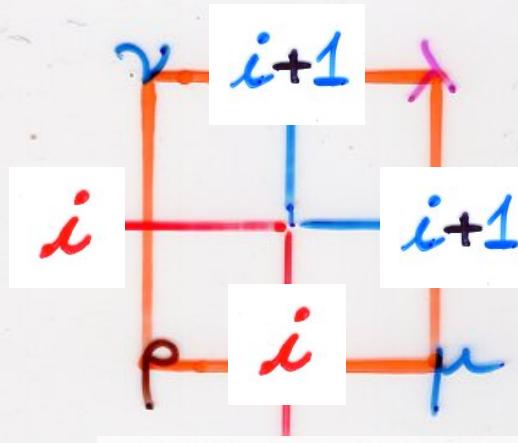
$$\lambda = \rho = \mu = \nu$$



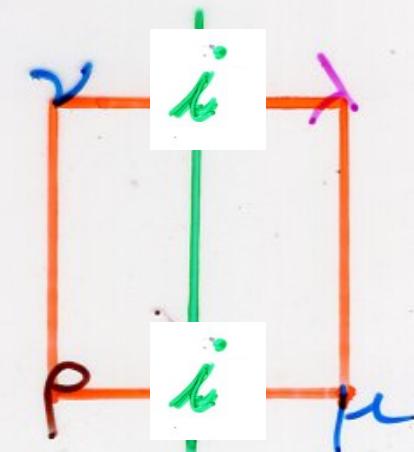
$$\begin{aligned} \rho &= \mu \\ \lambda &= \nu = \rho + (i) \end{aligned}$$



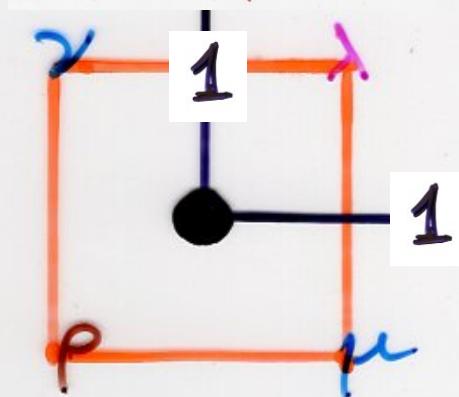
$$\begin{aligned} \nu &= \rho + (i) \\ \mu &= \rho + (j) \\ \lambda &= \rho + (i) + (j) \end{aligned}$$



$$\begin{aligned} \mu &= \nu = \rho + (i) \\ \lambda &= \mu + (i+1) \end{aligned}$$



$$\begin{aligned} \rho &= \nu \\ \lambda &= \mu = \rho + (j) \end{aligned}$$

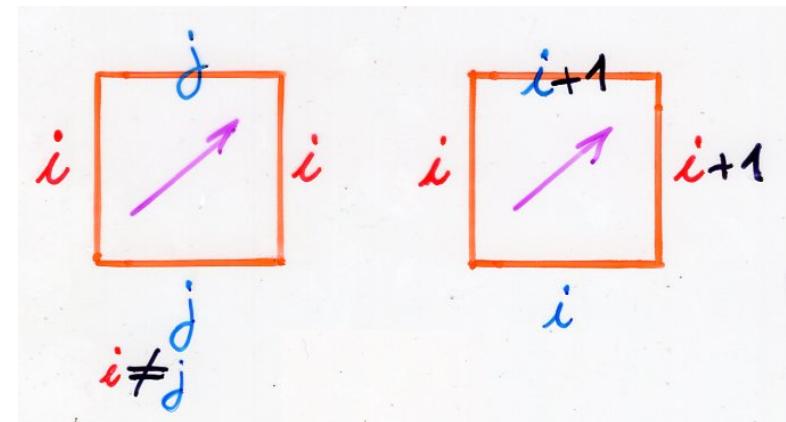
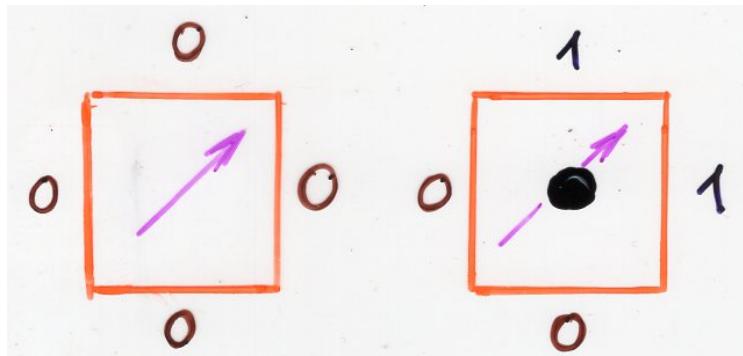


$$\lambda = \begin{cases} \rho \\ \mu + (1) \\ \nu \end{cases}$$

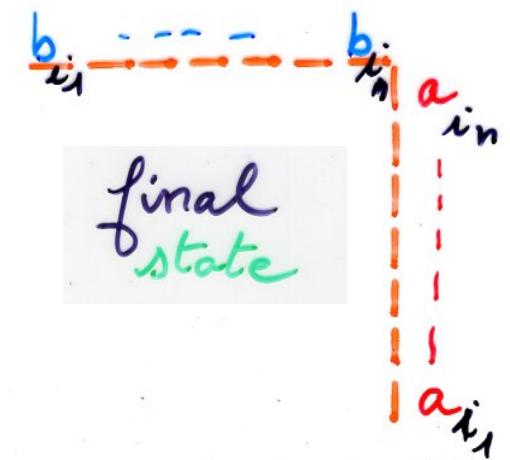
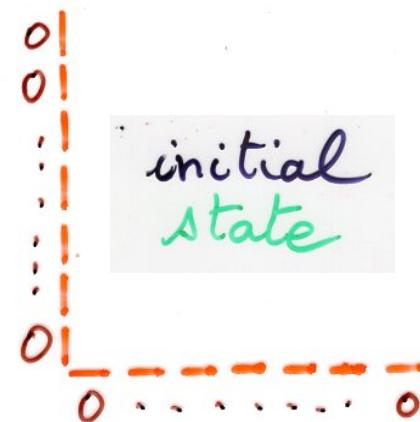
"local rules"
on the edges

<u>state</u>	$\{0, 1, 2, \dots\}$
state	$\{0, 1, 2, \dots\}$

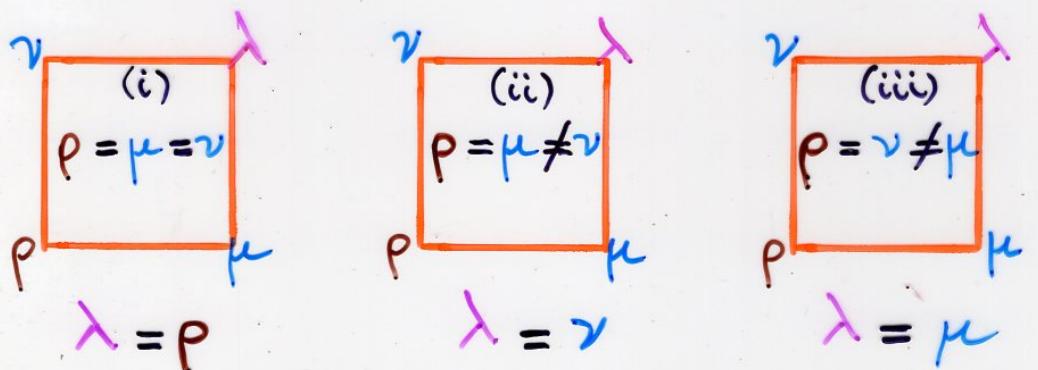
set of labels
 $L = \{\square, \bullet\}$



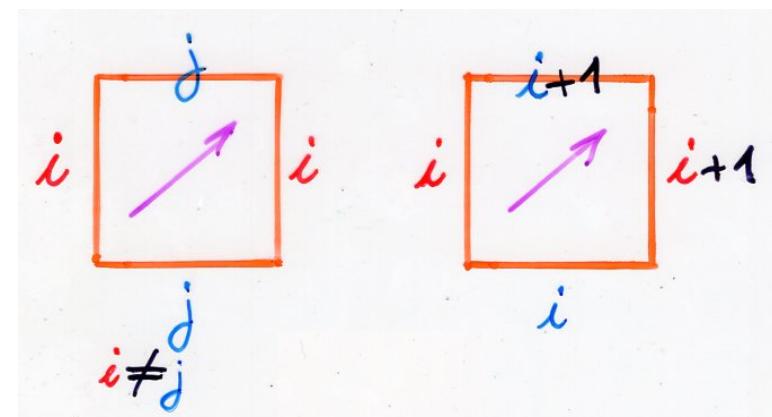
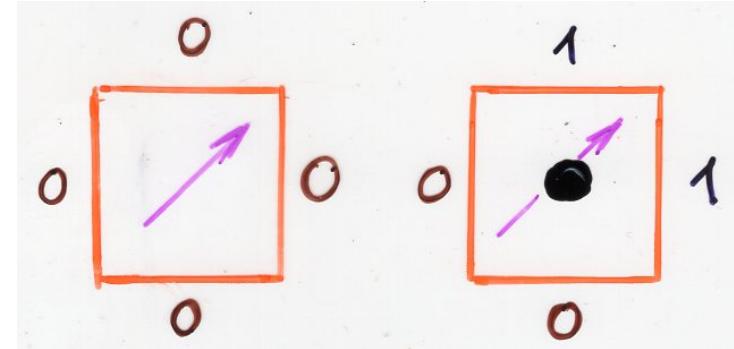
"planar
rewriting"

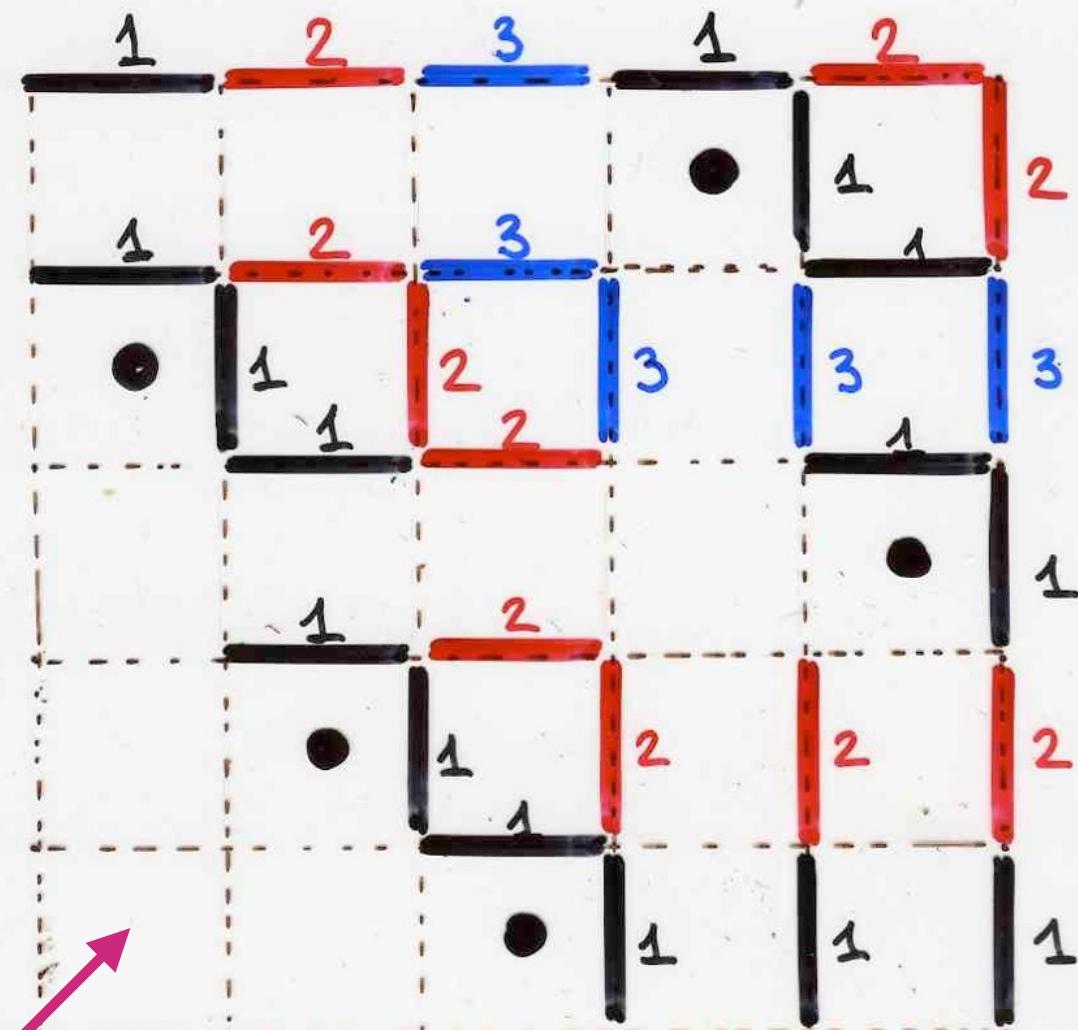


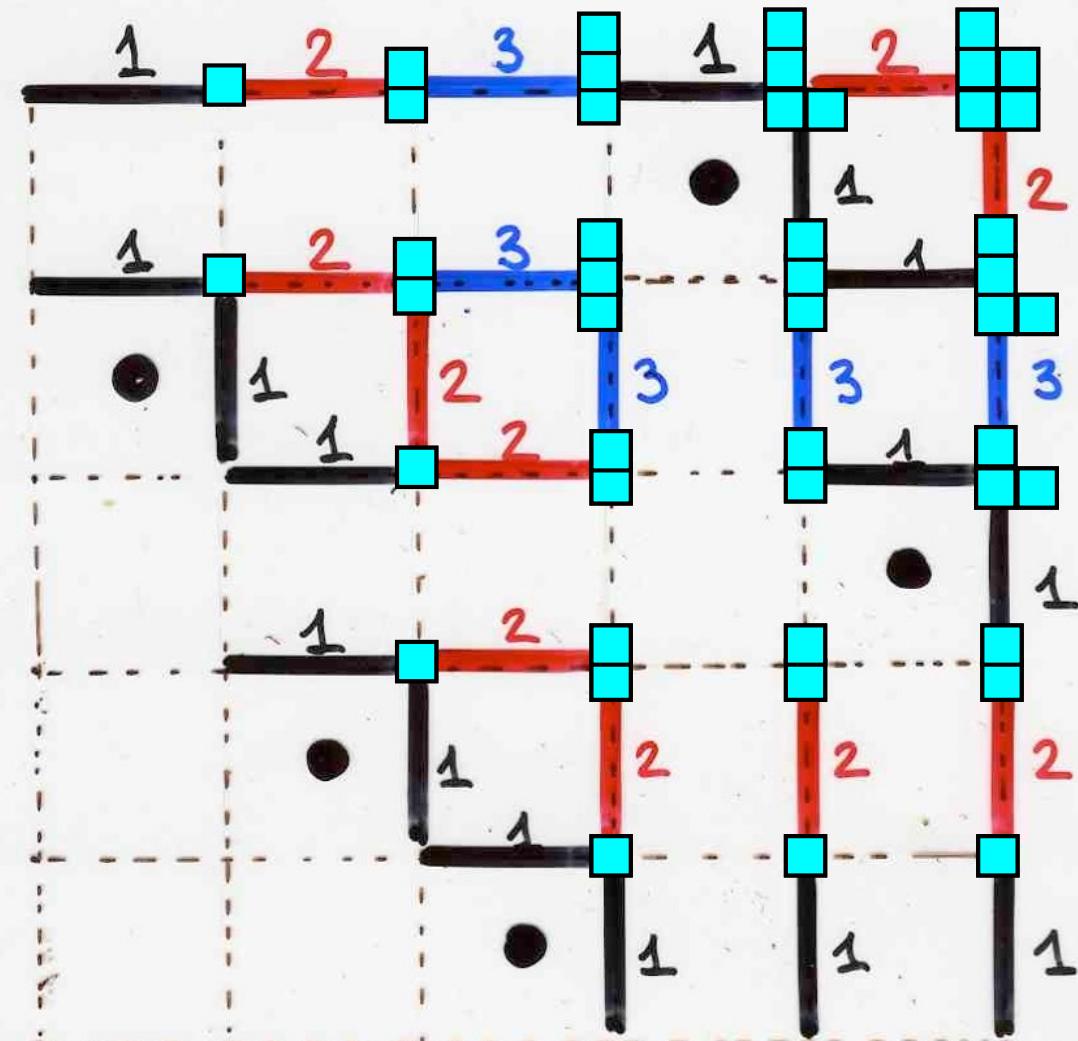
"local rules"
on the vertices

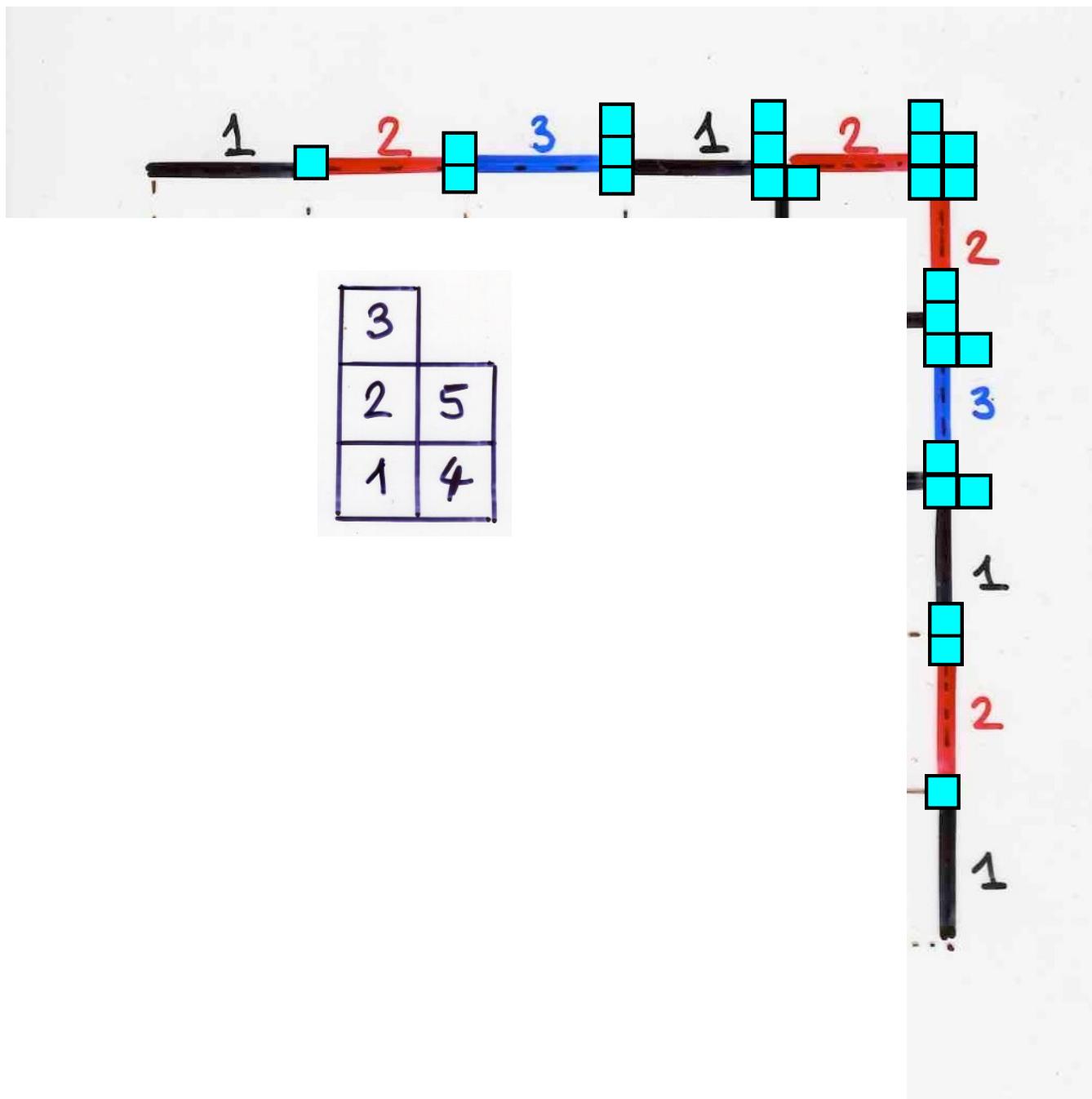


"local rules"
on the edges







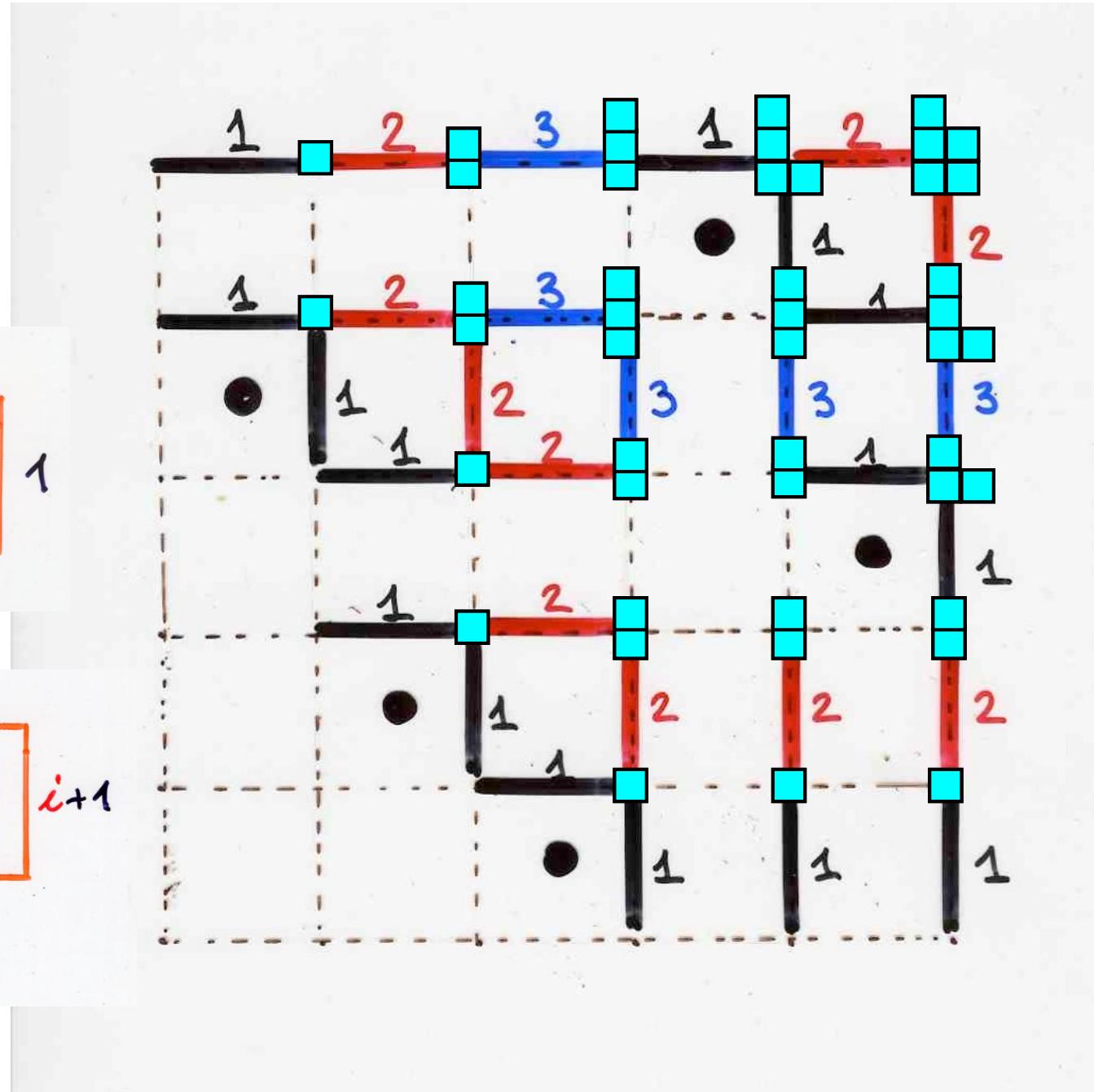
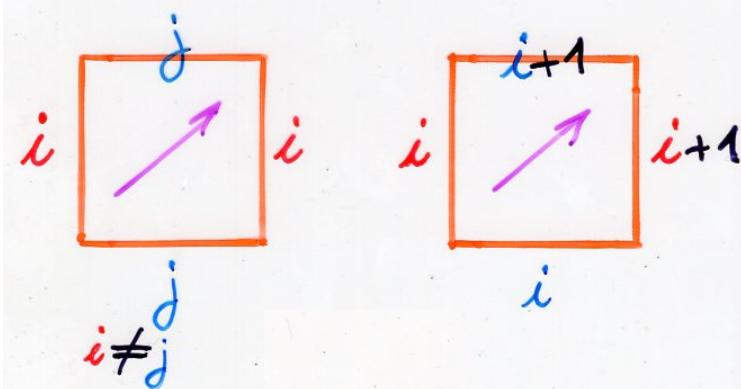
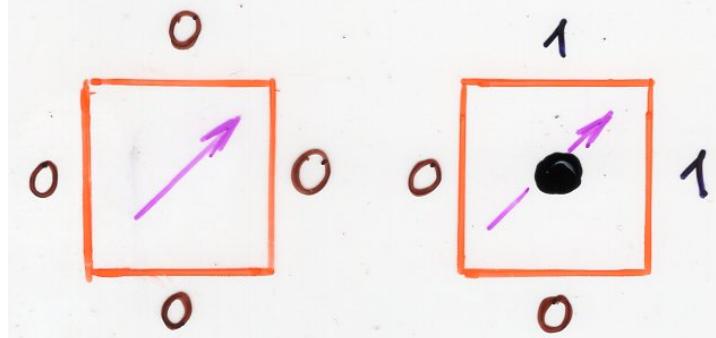


I claim that much attention should be given to the « local rules on edges » rather than « local rules on vertices ».

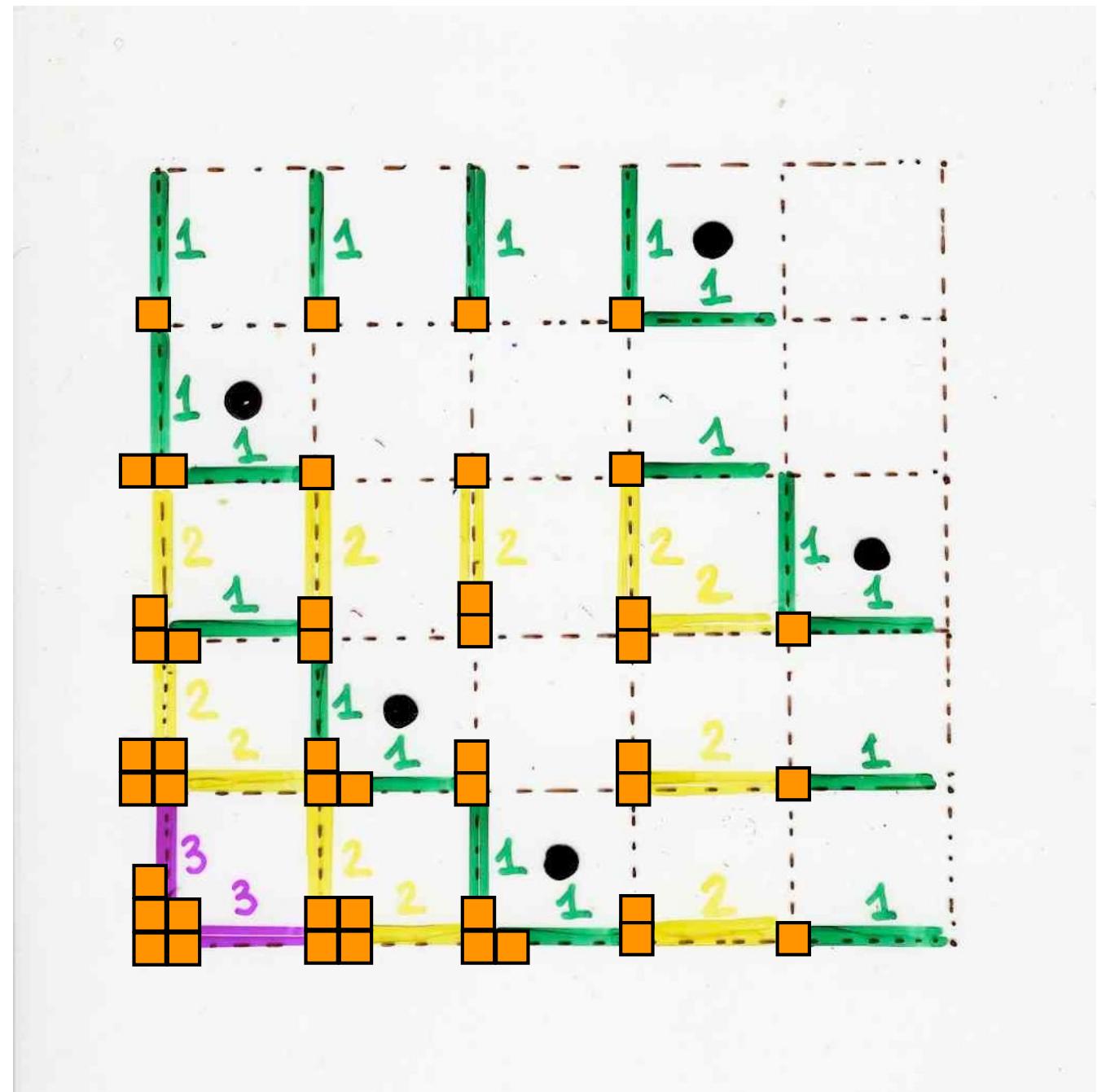
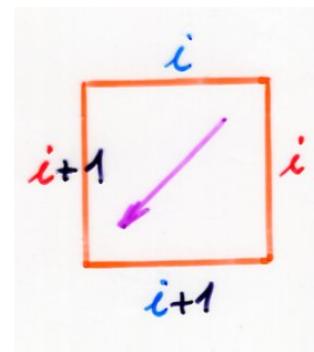
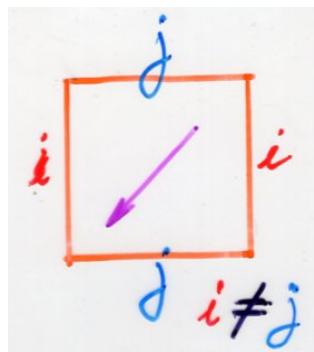
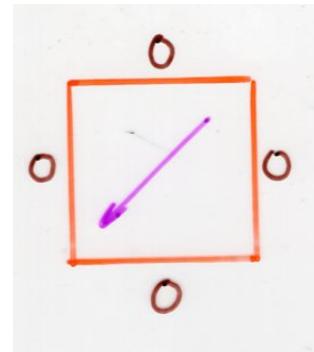
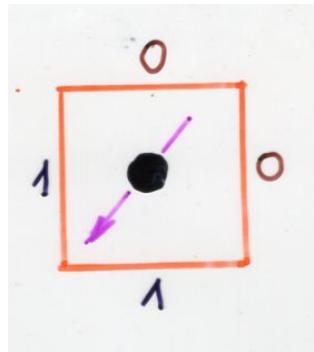
This is part of the philosophy of the « cellular ansatz »

The bilateral
RSK planar automaton

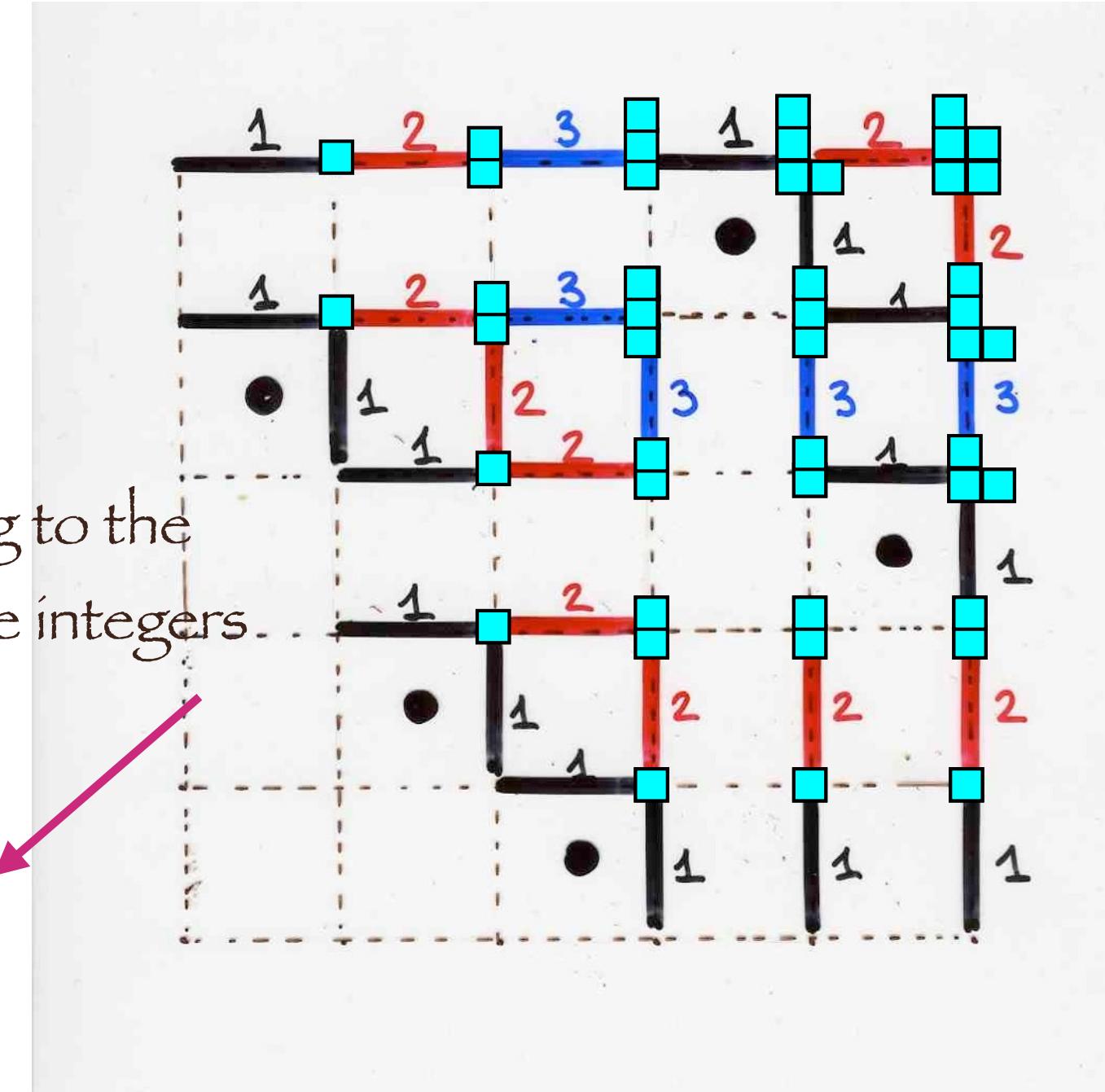
"local rules"
on the edges



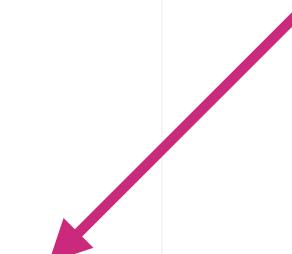
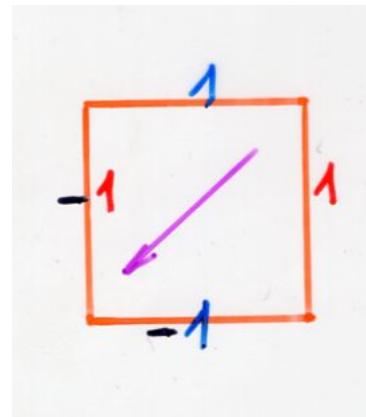
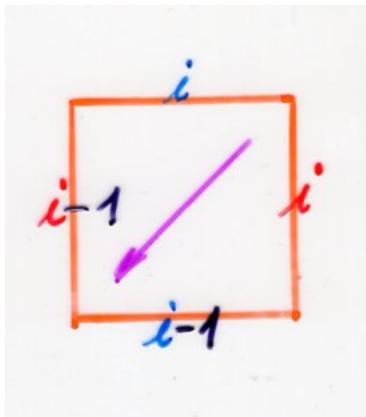
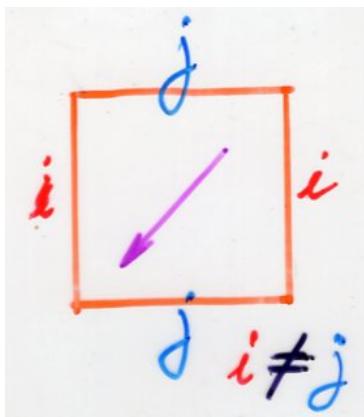
"local rules"
on the edges



Going to the
negative integers...

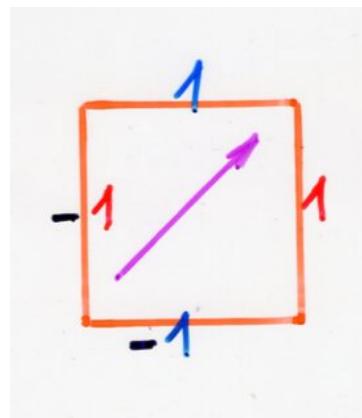
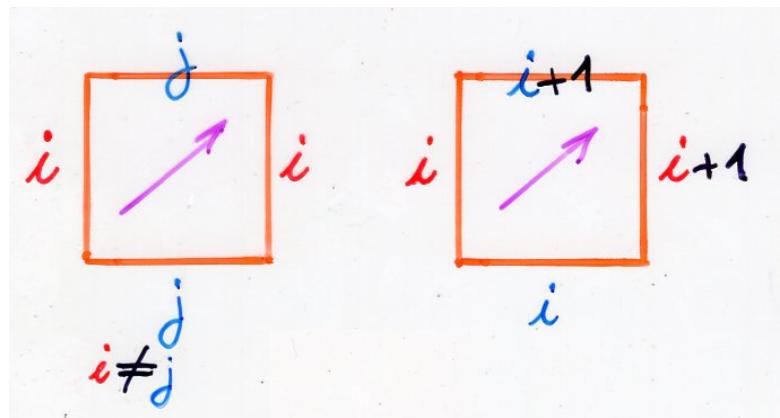


bilateral
planar automaton RSK



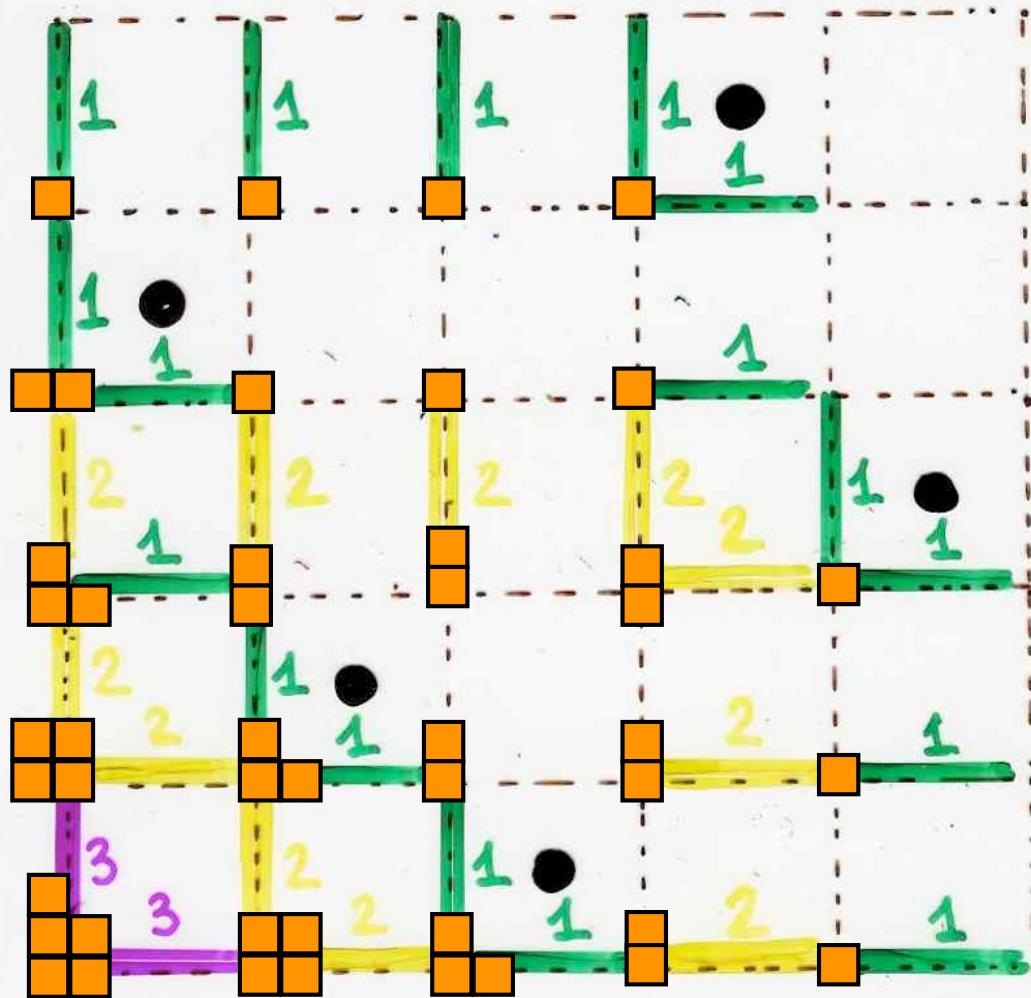
$$i \in \mathbb{Z} - \{0\}$$

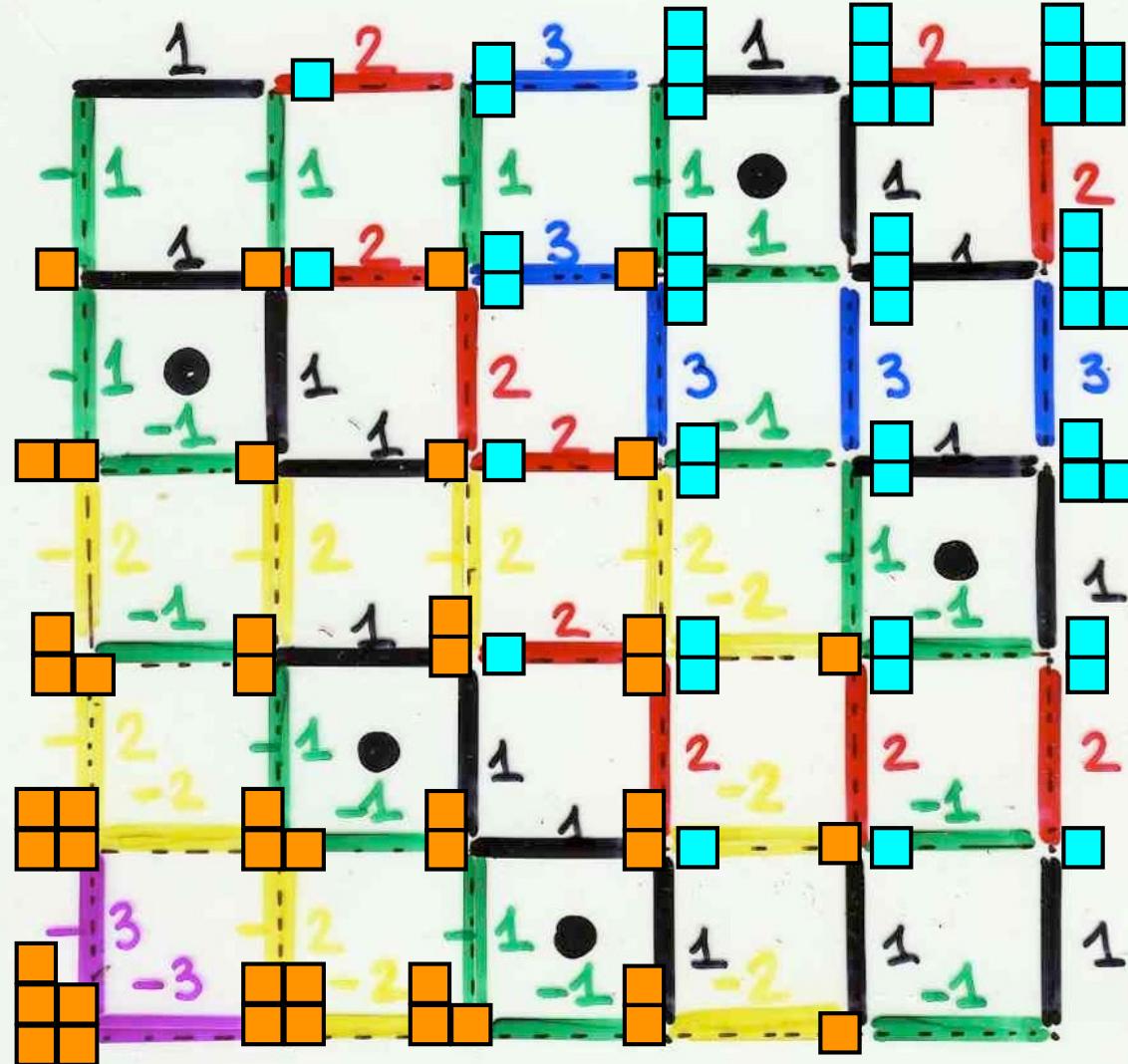
$$j \in \mathbb{Z} - \{0\}$$



1	2	3	1	2	
-1 1	-1 2	-1 3	-1 1	1 1	2
-1 -1	1 1	2 2	3 -1	3 1	3
-2 -1	2 1	2 2	2 -2	1 -1	1
2 -2	-2 1	1 1	2 -2	2 -1	2
3 -3	2 -2	1 -1	1 -2	1 -1	1







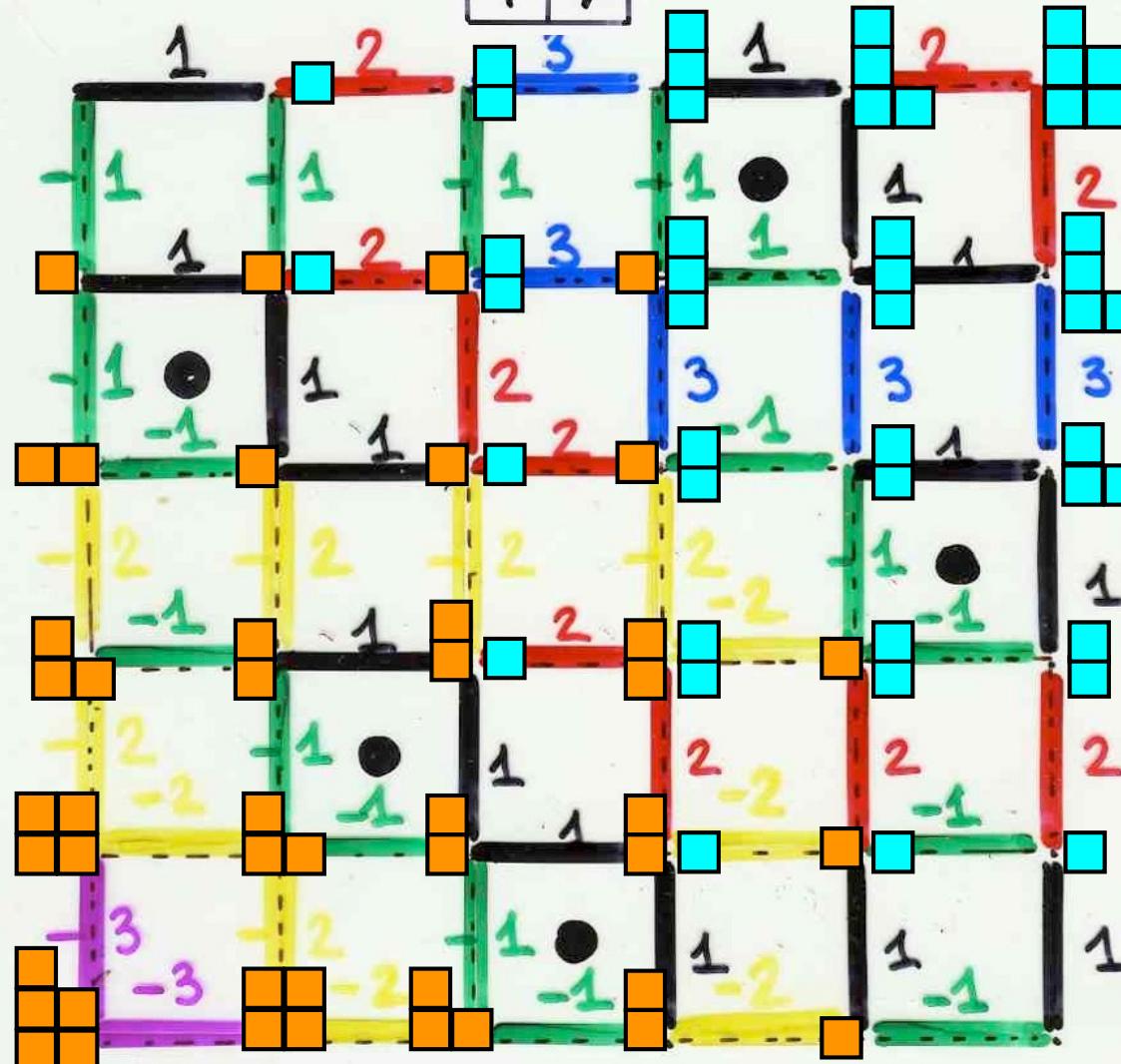
Schützenberger

Duality!

5	
3	4
1	2



3	
2	5
1	4



4	
2	5
1	3



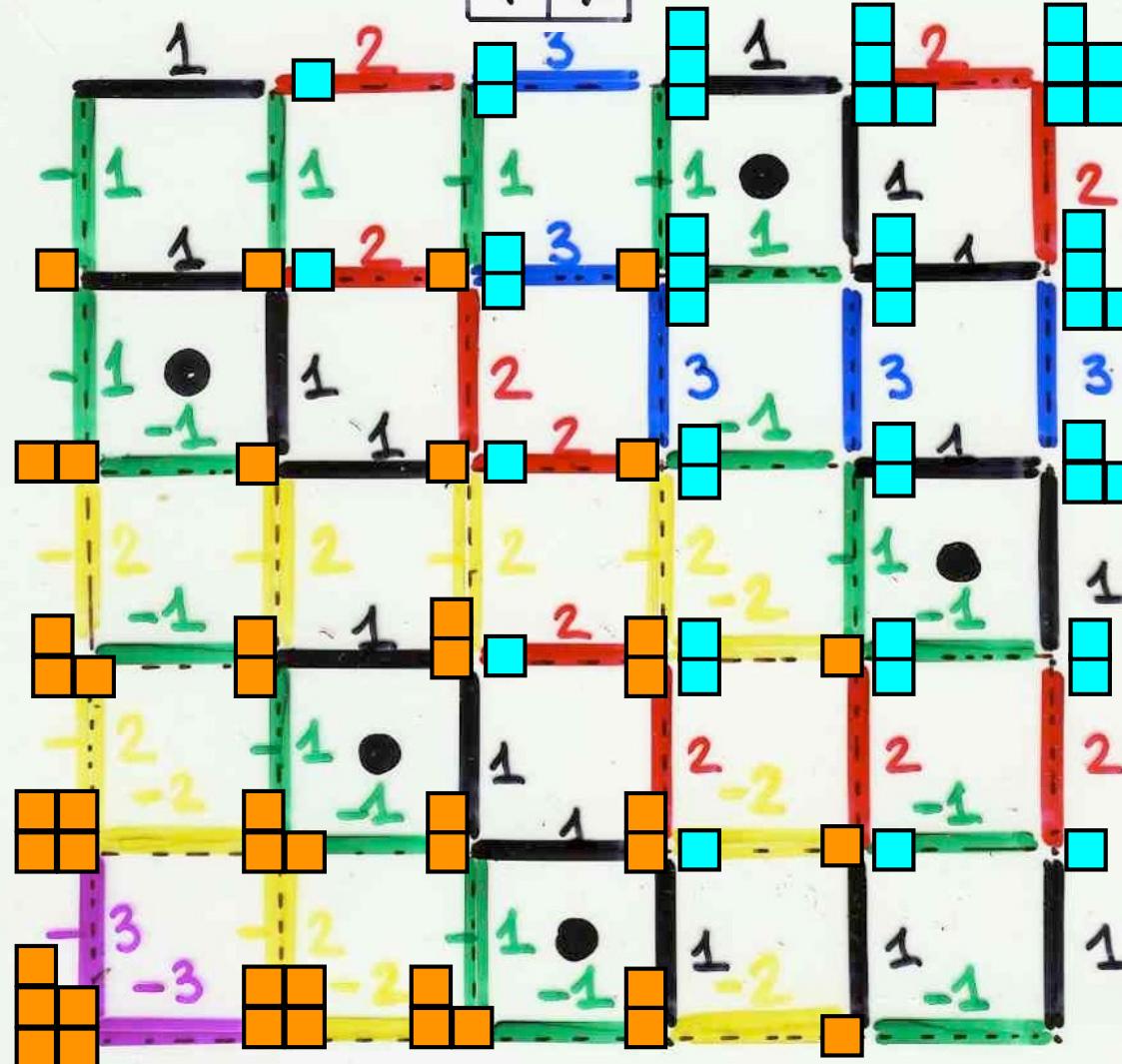
5	
2	4
1	3

Schützenberger

Duality!

P^* =
dual

5	
3	4
1	2

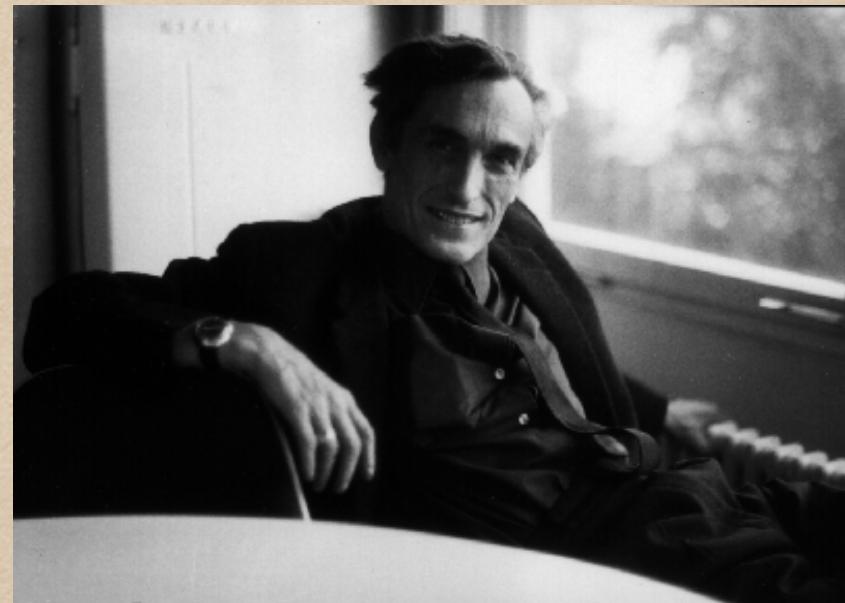


5	
2	4
1	3

$P =$

4	
2	5
1	3

dual of a Young tableau



M.P. Schützenberger

6	10			
3	5	8		
1	2	4	7	9

6	10			
3	5	8		
	2	4	7	9

6	10			
3	5	8		
2		4	7	9

6	10			
3	5	8		
2	4		7	9

6	10			
3	5	8		
2	4	7		9

6	10			
3	5	8		
2	4	7	9	

6	10			
3	5	8		
2	4	7	9	1

6	10			
3	5	8		
	4	7	9	1

6	10			
	5	8		
3	4	7	9	1

6	10			
5		8		
3	4	7	9	1

6	10			
5	8	2		
3	4	7	9	1

6	10			
5	8	2		
	4	7	9	1

6	10			
5	8	2		
4		7	9	1

6	10			
5	8	2		
4	7		9	1

6	10			
5	8	2		
4	7	9	3	1

6	10			
5	8	2		
	7	9	3	1

6	10			
	8	2		
5	7	9	3	1

	10			
6	8	2		
5	7	9	3	1

10	4			
6	8	2		
5	7	9	3	1

10	4			
6	8	2		
	7	9	3	1

10	4			
	8	2		
6	7	9	3	1

10	4			
8	5	2		
6	7	9	3	1

10	4			
8	5	2		
	7	9	3	1

10	4			
8	5	2		
7		9	3	1

10	4			
8	5	2		
7	9	6	3	1

10	4			
8	5	2		
	9	6	3	1

10	4			
	5	2		
8	9	6	3	1

7	4			
10	5	2		
8	9	6	3	1

7	4			
10	5	2		
	9	6	3	1

7	4			
10	5	2		
9	8	6	3	1

7	4			
10	5	2		
8	6	3	1	

7	4			
9	5	2		
10	8	6	3	1

7	4			
9	5	2		
8	6	3	1	

7	4			
9	5	2		
10	8	6	3	1

7	4			
9	5	2		
10	8	6	3	1

complement

$$(i)^c = n+1-i$$

P^* =

dual

4	7	
2	6	9
1	3	5

8 10

P =

6	10	
3	5	8
1	2	4

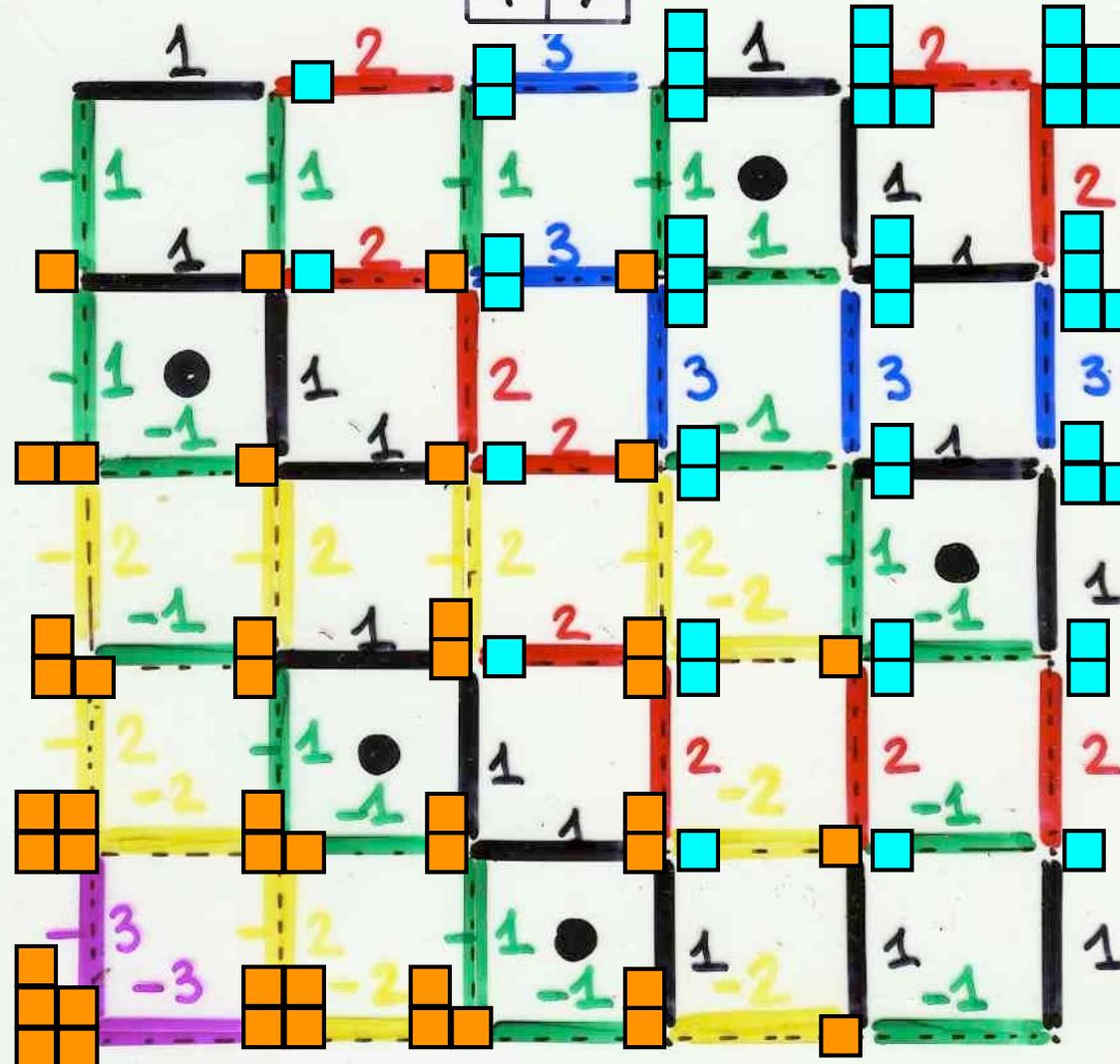
7 9

Schützenberger

Duality!

P^* =
dual

5	
3	4
1	2



5	
2	4
1	3

$P =$

4	
2	5
1	3

