An introduction to

enumerative algebraic bijective

combinatorics

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Chapter 1 Ordínary generating functions (4)

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From the previous lecture

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Euler's pentagonal theorem

$$\begin{aligned} \prod_{i \geqslant 1} (1 - q^{i}) &= \lambda - q - q^{2} + q^{3} + q^{2} - q^{2} - q^{3} + \cdots \\ &= \sum_{n \geqslant 1} (-1)^{n} \binom{n(3n-1)/2}{q} + q^{n(3n+1)/2} \end{aligned}$$

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5

12

22

pentagonal n(3n-1) numbers 2

A

The coefficient of q in TT (1-q')

= (number of partitions of n into an even number of distinct parts) (number of partitions of n into) (an odd number of distinct parts)

Euler pentagonal identity



construction of an involution changing the parity of the number of rows of the Ferrers diagram F

r(F) = number of cells of the top row of F

d(F) = longest sequence of cells in diagonal position on the NE border of F

(i) $r(F) \leq d(F)$ involution q

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case where q is not defined



case where q is not defined









$$\prod_{i \ge 1} (1 - q^{i}) = \lambda - q - q^{2} + q^{5} + q^{2} - q^{12} - q^{15} + \dots \\
= \sum_{n \ge 1} (-1)^{n} \binom{n(3n-1)}{2} + q^{n(3n+1)}/2 \\
= \sum_{n \ge 1} (-1)^{n} \binom{n(3n-1)}{2} + q^{n(3n+1)}/2$$

Euler pentagonal identity

Franklin (1881) bjection

sign-reversing involution

another exercise with a sign-reversing involution

With the construction of a sign-reversing involution prove that the generating function for heaps of dimens on the segment [0, k] exercise ID (t)(Fibonacci polynomial)



(enumerated by the number of dimers)

more about rational series

$$\frac{\operatorname{Proposition}}{\operatorname{k=C}} \quad \begin{array}{l} \mathcal{K}=\mathcal{C}, \quad \mathcal{K}>1, \quad \mathcal{F}_{\mathcal{K}}\neq 0, \\ \begin{array}{l} \begin{array}{l} \mathcal{E}a_{n} \ \mathcal{F}_{n20} \end{array} & \text{with} \quad a_{n} \in \mathbb{C}. \end{array} \end{array}$$

$$The 3 \quad \mathcal{F}elowing \quad \text{conditions} \quad \text{are} \quad equivelent: \end{array}$$

$$\begin{array}{l} (i) \quad \sum_{n20} a_{n} t^{n} = \frac{\mathcal{N}(t)}{\mathcal{D}(t)} \quad \text{with} \quad \mathcal{D}=1+\mathcal{F}_{\mathcal{A}}t+\ldots+\mathcal{F}_{\mathcal{K}}t^{k} \\ and \quad \mathcal{N} \quad pslynomial, \quad degree < k \end{array}$$

$$\begin{array}{l} (ii) \quad \mathcal{F}or \quad eveny \quad n20, \quad a_{n+k} + \mathcal{F}_{\mathcal{A}} \quad a_{n+k-1} + \ldots + \mathcal{F}_{\mathcal{A}} \quad a_{n} = 0 \end{array}$$

$$\begin{array}{l} (iii) \quad \ellet \quad \mathcal{D} = \prod_{i=1}^{r} \left(1-\lambda_{i}t\right)^{k_{i}} \quad \text{For} \quad every \quad n \ge 0, \\ a_{n} = \sum_{i=1}^{r} \quad \mathcal{F}_{i}(n) \quad \lambda_{i}^{n}, \quad \text{with} \quad \mathcal{F}_{i} \quad \mathcal{F}elynomial \\ \quad degree < k_{i} \end{array}$$

$$F_{n} = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^{n+1} - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^{n+1} \quad (n_{2}^{n})$$

Transition A $\sum_{n \ge 0} a_n t^n = \frac{N_{ij}}{det(I_n - At)}$

an expression with the inverse of the zeros of the polynomial det (In-At) = the eigenvalues of A

Zeros of
$$F_n(z)$$
 Fibonacci
 $sin((n+1)\Theta) = sinU_n(aS\Theta)$
 $U_n(z) = M_n(2z)$ $M_n^{*}(z) = z^n M_n(1/z) = F_n(z^2)$

zeros of $U_n(x)$: $\sum_{n=1}^{\infty} \cos\left(\frac{k}{(n+1)}T\right) = \frac{1}{n-1}, n f$ inverse of zeros of $F_n(x)$: $\sum_{n=1}^{\infty} 4\cos\left(\frac{k}{n+1}T\right), k=1, n f$

Fr (t) bounded Fr (t) Dyck paths F. (t)

 $T_{n}(x) = \frac{1}{2}C_{n}(2x)$ $C_{n}^{*} = L_{n}(x^{2}) \cos(n\theta) = T_{n}(\cos\theta)$

zeros of $T_n(x)$: $\int cos(\frac{(k-1)}{2n}\pi)$, k=1,..,n

inverse of zeros of $L_n(x)$ Lucas polynomial $\begin{cases} 4 \cos^2\left(\frac{(2k-1)}{2n}\pi\right), k=1,...,n \end{cases}$

Lagrange inversion formula

Lagrange inversion formula

IK field characteristic O f(t) has a neuprocal power series $y = f^{(-1)}(t)$ if f(0) = 0 and $f'(0) \neq 0$

(et $f(t) = \frac{t}{\varphi(t)}$ with $\varphi(0) \neq 0$

y is the unique solution of $y = t \varphi(y)$ (with y(0)=0)

Proposition (Lagrange inversion formula) Let $g(t) \in [K[[t]]$, the coefficient of t^n in g(y) is

 $[t^n]g(y) = \frac{1}{n} [t^{n-1}]g'(t)q(t)^n$

analytic proof with residu of meromorph
 combinatorial proof in Ch3

example Catalan generating function $y = \sum_{n \ge 1} C_n t^n$, $y = 1 + ty^2$, z = 1 + y $n \ge 1$ then $z = t \cdot q(z)$ with $q(t) = (1+t)^2$. We have $C_n = \frac{1}{n} [t^{n-1}](1+t)^{2n} = \frac{4}{n} {2n \choose n-1}$

 $C_n = \frac{1}{(n+1)} \binom{2n}{n}$

algebraícíty with hidden decomposable structures

example: planar maps





Tutte (1960) an number of rooted planar maps with n edges

 $y = \frac{1-4z}{(1-3z)^2}$ $= \sum_{n \geq 0} a_n t^n$ $z = \frac{E}{(1-3z)}$



 $y = \frac{1-4z}{(1-3z)^2}$ $\mathbf{Z} = \frac{E}{(1-3\mathbf{Z})}$

 $z = t + 3z^2$

z = t f(t) $h = 1 + 3th^2$



 $= h - th^3$

 $\begin{cases} h = 1 + 3th^2 \\ y = h - th^3 \end{cases}$

Tutte (1960) an number of rooted planar maps with n edges



Catalan $\int h = 1 + 3th^2 \quad 3^n C_n$ $y = h - th^3$

Cori, Vauquelin (1970)

> methodobgy (lebw)

 $\begin{cases} h = 1 + 3th^2 \\ y = h - th^3 \end{cases}$

coding planar maps with words in the difference of two algebraic languages

Cori, Vauquelin (1970)

Lagrange inversion formula

8 = 2 ant"

 $y = \frac{1-4z}{(1-3z)^2}$ $\begin{cases} z = t \varphi(z) \\ y = g(z) \end{cases}$ with $\int \frac{f(t)}{t} = \frac{1}{1-3t}$ $\int \frac{f(t)}{(1-3t)^2} = \frac{1-4t}{(1-3t)^2}$ $\mathbf{Z} = \frac{E}{(1-3\mathbf{Z})}$

 $a_{n} = \frac{\lambda}{n} [t^{n-1}] (t^{n-1}) (t^{n-1})$ $=\frac{2}{n}\left[t^{n-1}\right]\left((1-6t)(1-3t)^{-n-3}\right)$ $= \frac{2}{n} \left[t^{n-1} \right] \left(1 - 3t \right)^{n-3} - 6 \left[t^{n-2} \right] \left(1 - 3t \right)^{n-3}$ $=\frac{2}{n}\left(3_{\times}^{n-1}\frac{(n+3)-(2n+1)}{(n-1)!}-6_{\times}3_{\times}^{n-2}\frac{(n+3)-(2n)}{(n-2)!}\right)$ $= \frac{2}{n} \frac{3^{n-1}}{(n+3)-(2n)} \left(\frac{2n+1}{n-1}-2\right)$ $= 2 \times 3^{n} (n+3) \dots (2n)$ = $2 \times 3^n \frac{(2n)!}{(n+2)! n!}$ $= 2 \times \frac{3^n}{(n+2)} C_n$





Schaeffer (1997)



Schaeffer (1997)



Schaeffer (1997)








Cori, Vauquelin (1970)

Schaeffer (1997)

Bouttier, Di Francesco, Guitter (2002) · · · many other



complements

algebraicity with hidden decomposable structures

example: directed animals





Y generating function for the number of directed animal with n points satisfies the system of algebraic equations:

y= z+ yz $z = t + tz + tz^2$

exercise algebraic equations for Motekin paths and prefix of Motekin paths



Motzkin path



 $z = t + tz + tz^2$





an anecdote





J. Vannimenus (1982, 1983) J.P. Nadal

B. Derrida



directed animals on a circular strip

 $\sum_{n=1}^{k} = \frac{1}{k} \sum_{p=0}^{k-1} (-1)^{p} \sin \alpha_{p} \frac{k-1}{\prod_{i=1}^{k-1} \left(\frac{\sin(i+\frac{1}{2}) \alpha_{p}}{2} \right)^{k}}{\sin \frac{\alpha_{p}}{2}} (1+2\cos \alpha_{p})^{k}$ animals $d_{p} = \frac{2p+1}{2k} \pi$ circular strip

width fr







J. Vannimenus

J.P. Nadal

(1982, 1983)

B. Derrida

 $T_n(x) = \frac{1}{2}C_n(2x)$ $C_n^* = L_n(x^2) \cos(n\theta) = T_n(\cos\theta)$

zeros of $T_n(x)$: $\int cos(\frac{(k-1)}{2n}\pi)$, t=1,..,n

complements

algebraicity with hidden decomposable structures

example: convex polyomínoes



Polyominoes enumeration





















convex polyominoes

Pen = number of convex polyominoes with perimeter 2n



 $\sum_{n_{1/2}} \operatorname{Pen} t^{2n} \frac{t^4 (1 - 6t^2 + 11t^4 - 4t^6)}{(1 - 4t^2)^2} - 4t^8 (1 - 4t^2)^{-3/2}$ Penes = (2n+11)4" - 4(2n+1)(2n) (n20)

Delest, X.V. (1984)

convex polyominoes

Pen = number of convex polyominoes with perimeter 2n

> methodology (lebw)

convex Polyominoes with words algébraic languages



random convex polyominoes

with fixed perimeter





rational and algebraic generating functions

A flavor of theoretical computer science with

Schütengerger methodology coding with words of algebraic languages

Schützenberger methodology oding combinatorial objects with words of algebraic language

context-free language (algebraic) language

algebraic grammar D→TD or D→E (empty word) $T \rightarrow \chi D \bar{\chi}$

context-free language (algebraic)

algebraic grammar (empty word) or $\mathcal{D} \rightarrow \varepsilon$ $\mathcal{D} \rightarrow \mathcal{T}\mathcal{D}$ $T \rightarrow \chi D \bar{\chi}$

D TD ZDZD ZTDZD x z D x D x D XXZ DZD XXX TD XD xxx xJx Jx XXX X X X D ママズ マズズ エコ XXX X XXX DX x x z x x x x x x z

(restricted) Dyck language DS fx, 73* (i) $|w|_{\chi} = |w|_{\overline{\chi}}$ (nb of occurrences of χ in the word w) (ii) for every factorization w = uv $|u|_{\chi} \ge |u|_{\bar{\chi}}$ Coding of Dyck paths



only (i): bilateral Dyck language coding bilateral Dyck paths (2n) enumerated by
algebraic grammar D->TD or D->E (empty word) $T \rightarrow \chi D \bar{\chi}$



zD z D TPZD 2 xxDxDxD XXI DID XXX TD XD XXXXXXX XXX x X x 222 222 XXX X XXX DX x x z x x x x x x x



derivation tree XDX x え

XXXXXXXXX

ambiguous grammar if there exist a word wEL having 2 distinct derivation trees

D->DD or E D-2DZ ambiguous grammer

Prop Chomoky - Schützenberger L'algebraic language, $L \subseteq X^*$ having a non-ambiguous grammar. Let $a_n = L \cap X^n$ (nb of words of L of length n) Then the power series $L = \sum a_n t^n$ is algebraic (on (x))

non-commutative power series $\mathbf{P} = 1 + \mathbf{x} \mathbf{P} \mathbf{\overline{z}} \mathbf{P}$

exercise algebraic equations for Motekin paths and prefix of Motekin paths



Rational langage words accepted by a finite automaton witial state $s \in S$ $\mathcal{A} = (S, X, \Theta, S_{0}, F)$ states alphabet transition final states FCSfunction $\Theta: (S, X) \rightarrow t \in S$

analog of matrix metho do bgy

Rational Algebraic languages



bijection

matchings - paths a of E1, nJ length n going from S, to SA



such that $v(\omega) = (-\infty)^k t^n$ k = number of dimers of the matching



initial state s final states F= tyf



$$X = \{a, b\}^{*}$$

$$L = (a + bb)^{*}$$
(product of words

a and bb)

automaton



hint for the exercise:

think in terms L'automaton

an number of such polyominoes with n cells $a_n = F_{2n-2}$ Filonacci numbers

generating functions

algebraic D-finite or not D-finite?



 $P_{k}(n)a_{n+k} + P_{k}(n)a_{n+k-1} + \cdots + P_{k}(n)a_{n} = 0$

rational (recurrence with B. ..., Re constants

example paths in a quadrant the elementary steps



23 are D-finite 56 not D-finite reduced to 79 cases

4 are algebraic

Cessel Kneweras



4 are algebraic

Introduction of catalytic variables kernel methodology

example: - vertically convex polypoinnes enumerated by perimenter, area and number of cells in the last column catalytic variable

- planer maps (Tutte)



one catalytic -> algebraic





23 are D-finite 56 not D-finite reduced to 79 cases

Group associated to a path D-finite (=> the group is finite





C(t)not D-finite connected g. f. heap (Bousquet-Neibu, Rechnitzer) 2002 $C(t) = \frac{Q}{(1-Q)} \left[1 - \sum_{k>1} \frac{Q^{k+1}}{1-Q^{k}(1+Q)} \right]$ $Q = \sum_{n > 1} C_n t^n$ Catalan number

