

Course IMSc, Chennai, India



January-March 2018

The cellular ansatz:
bijective combinatorics and quadratic algebra

Xavier Viennot

CNRS, LaBRI, Bordeaux

www.viennot.org

mirror website

www.imsc.res.in/~viennot

Chapter 1

RSK

The Robinson-Schensted-correspondence (Ch1c)

IMSc, Chennai
January 18, 2018

Xavier Viennot
CNRS, LaBRI, Bordeaux
www.viennot.org

mirror website
www.imsc.res.in/~viennot

From Ch 1a, 1b:

The Robinson-Schensted correspondence

- Schensted's insertions
- geometric version with "shadow lines »
- Fomin "local rules" or "growth diagrams »

$$\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 \\ 3 & 1 & 6 & 10 & 2 & 5 & 8 & 4 & 9 & 7 \end{pmatrix}$$

6	10			
3	5	8		
1	2	4	7	9

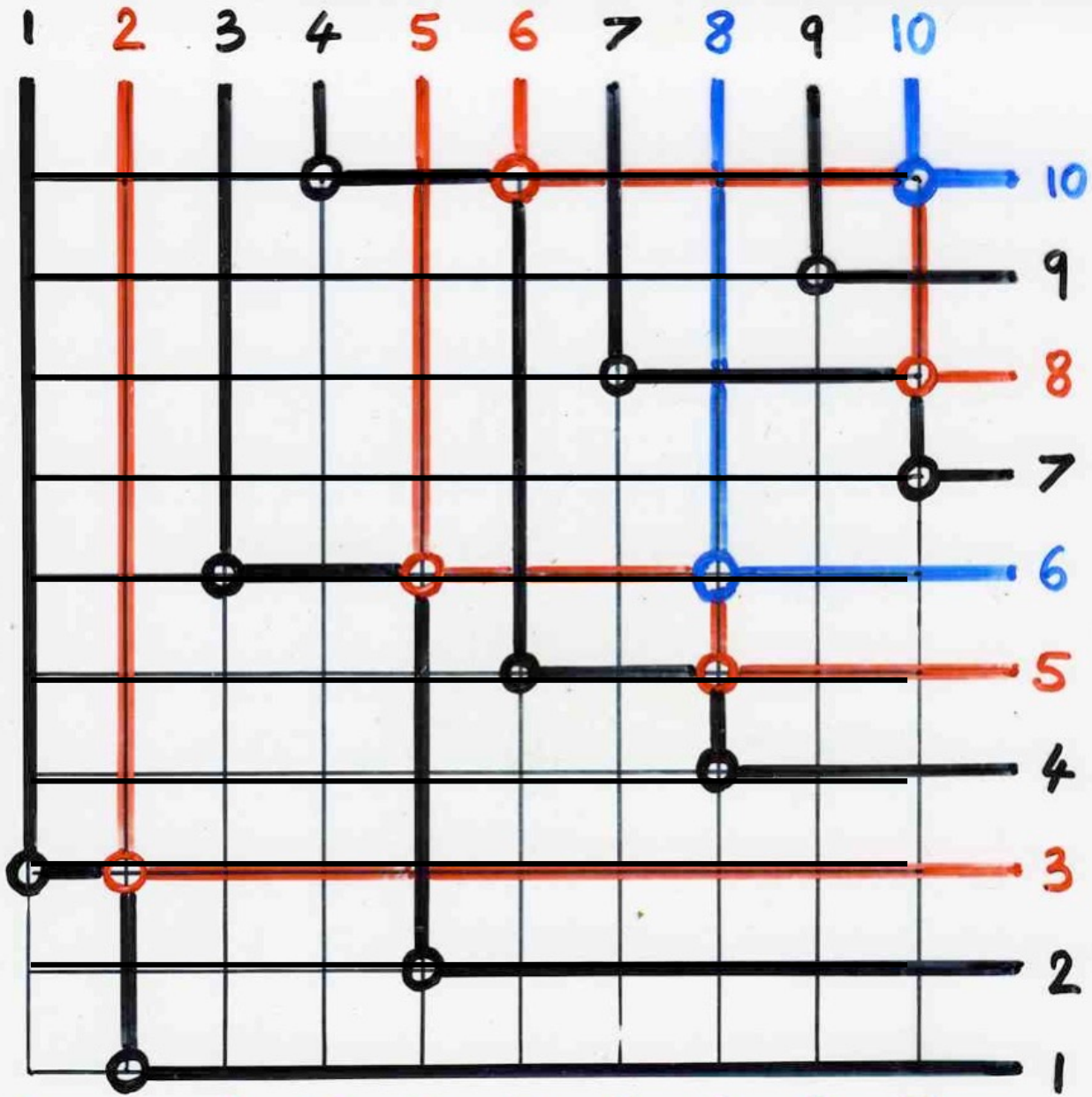
P



8	10			
2	5	6		
1	3	4	7	9

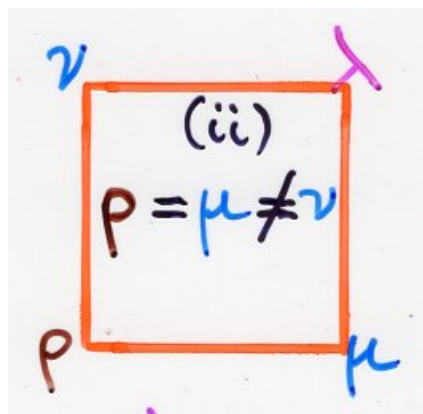
Q

The Robinson-Schensted correspondence between permutations and pairs of (standard) Young tableaux with the same shape

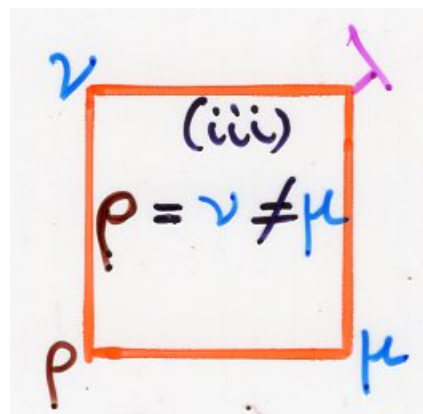


$\sigma = 3 \quad 1 \quad 6 \quad 10 \quad 2 \quad 5 \quad 8 \quad 4 \quad 9 \quad 7$

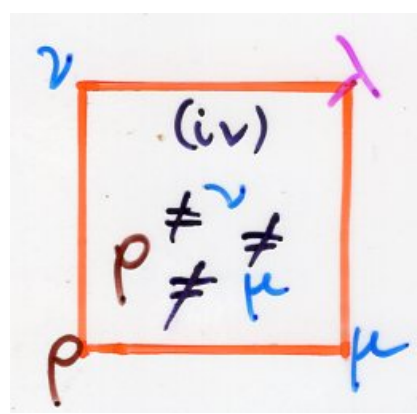
"local rules"



$$\lambda = v$$

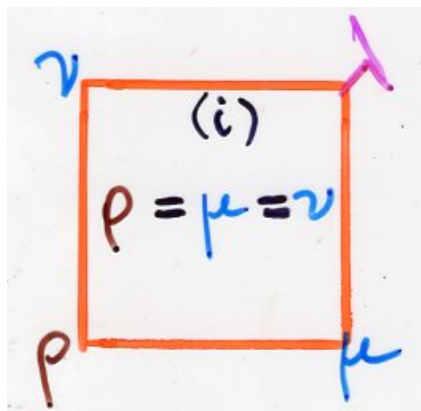


$$\lambda = \mu$$

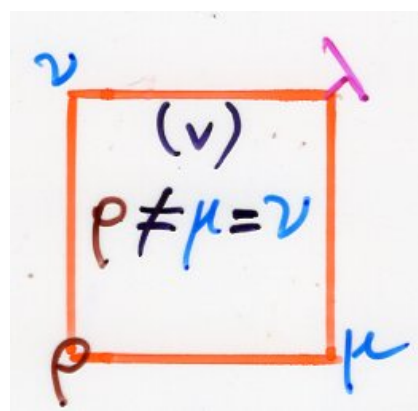


$$\lambda = \mu \cup v$$

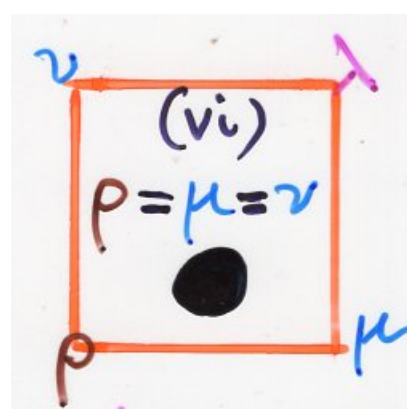
$$\mu \neq v$$



$$\lambda = \rho$$

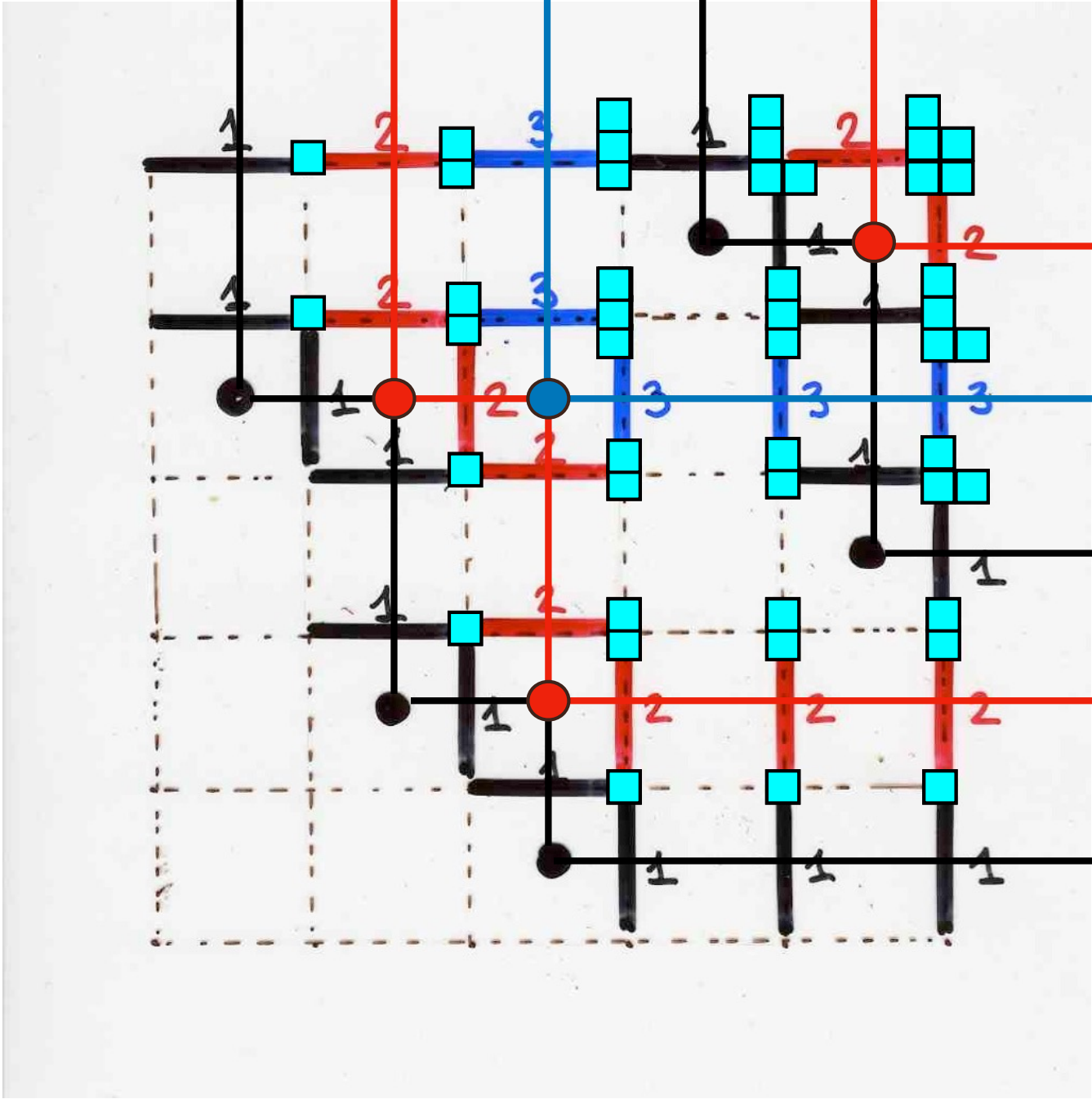


$$\lambda = \left\{ \begin{matrix} \mu \\ v \end{matrix} \right\} + (i+1)$$



$$\lambda = \left\{ \begin{matrix} \rho \\ \mu \\ v \end{matrix} \right\} + (1)$$

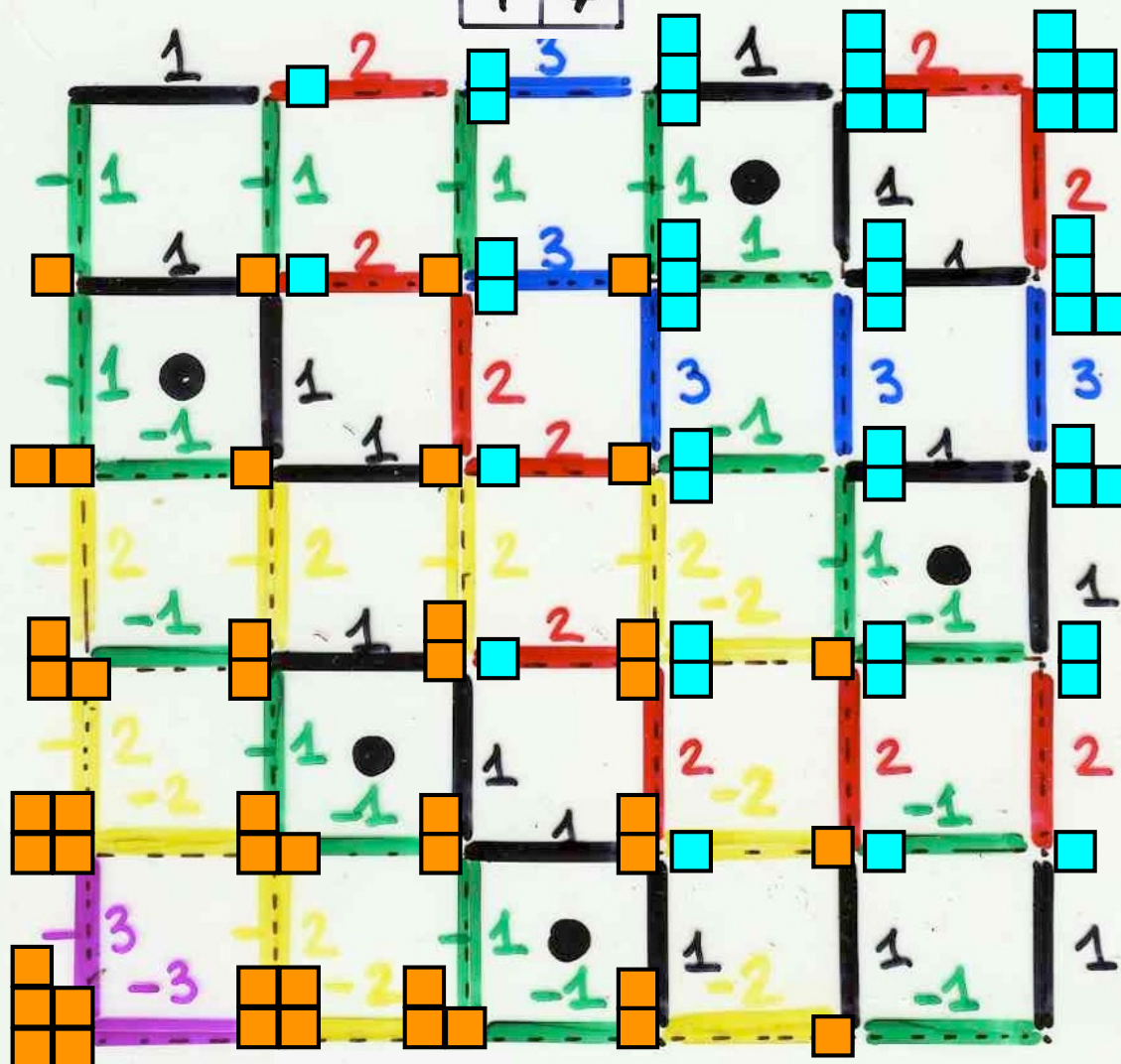
$$\mu = v$$



Schützenberger

Duality!

3	
2	5
1	4



4	
2	5
1	3



5	
3	4
1	2



5	
2	4
1	3

"The cellular ansatz"

quadratic algebra Q

Q -tableaux

representation of Q
by combinatorial operators

$$UD = DU + Id$$

combinatorial objects
on a 2D lattice

bijections

permutations

RSK

pairs of
Young tableaux

towers placements



(i) first step

(ii) second step

commutations

rewriting rules

planarization

an example

Heisenberg operators U, D

creation and annihilation operators

quantum mechanics

$$UD = DU + Id$$

commutations

Lemma Every word w with letters U and D can be written in a unique way

$$w = \sum_{i,j \geq 0} c_{ij}(w) D^i U^j$$

normal ordering
in physics

The monomials $\{D^i U^j\}_{i,j \geq 0}$
form a basis of the
Weyl-Heisenberg algebra

$$Q = \mathbb{C}\langle U, D \rangle / \mathcal{J}$$

non-commutative polynomials
in variable U and D
(free associative algebra)

\mathcal{J} ideal generated by
the relation $UD = DU + I$

$$UD = DU + Id$$

commutations

$$UD \rightarrow DU$$

$$UD \rightarrow Id$$

rewriting rules

UUDD

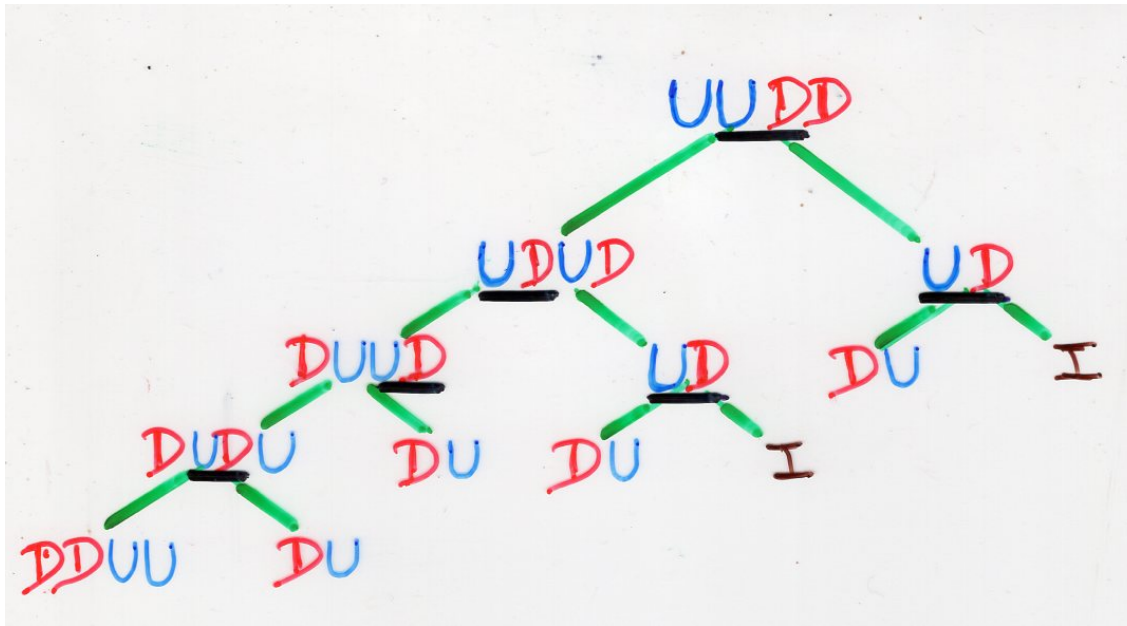
$$UUDD = UDUD + UD$$

$$= DUUD + 2UD$$

$$= (DU DU + DU) + 2(DU + Id)$$

$$= (DDUU + 2DU) + 2(DU + Id)$$

$$= DDUU + 4DU + 2Id$$



$$U^2 D^2 = D^2 U^2 + 4DU + 2I$$

this polynomial is independent
of the order of the substitutions

$$U^n D^n = \sum_{0 \leq i \leq n} c_{n,i} D^i U^i$$

$$c_{n,0} = n!$$

permutations

Planarization of the rewriting rules

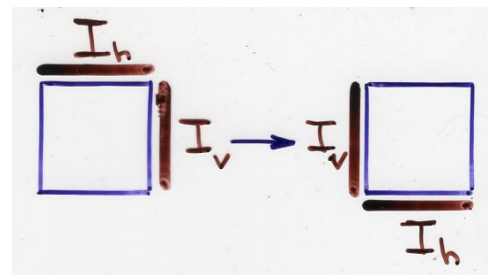
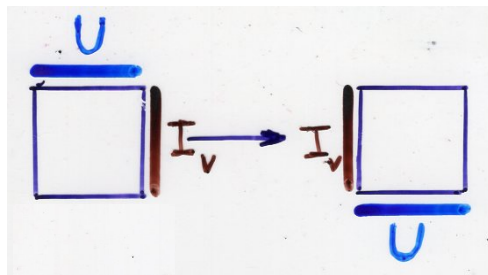
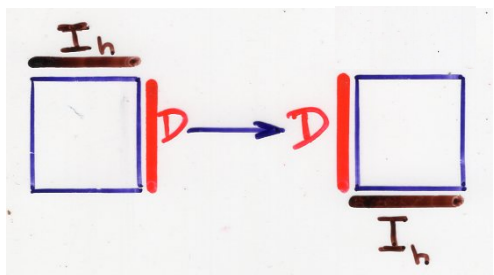
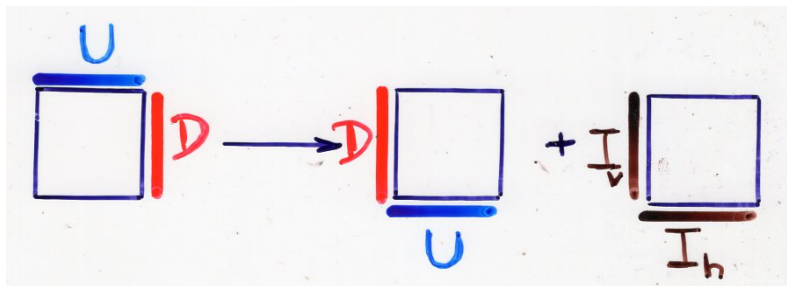
$$UD = DU + Id$$

commutations

$$UD \rightarrow DU \quad UD \rightarrow Id$$

rewriting rules

planarization of the rewriting rules

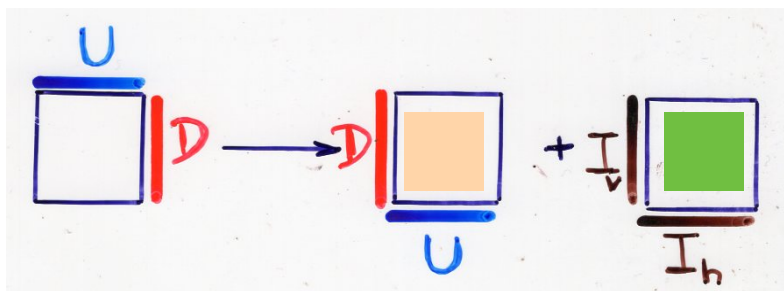


homogenization
of the system
of commutations
relations

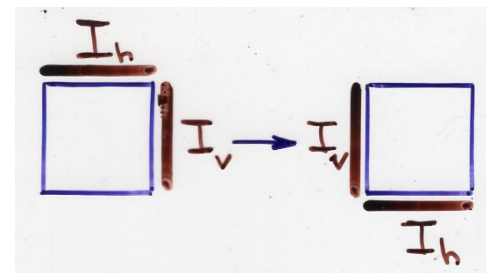
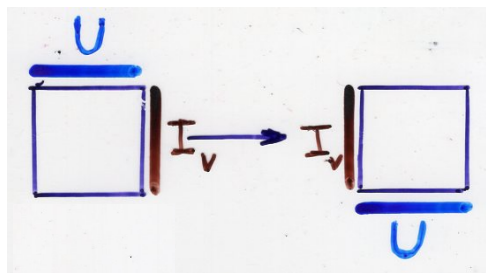
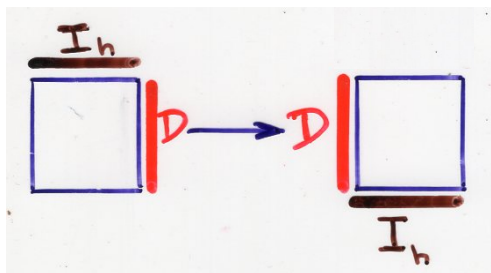
$$\left\{ \begin{array}{l} UD = DU + I_v I_h \\ UI_v = I_v U \\ I_h D = D I_h \\ I_h I_v = I_v I_h \end{array} \right.$$

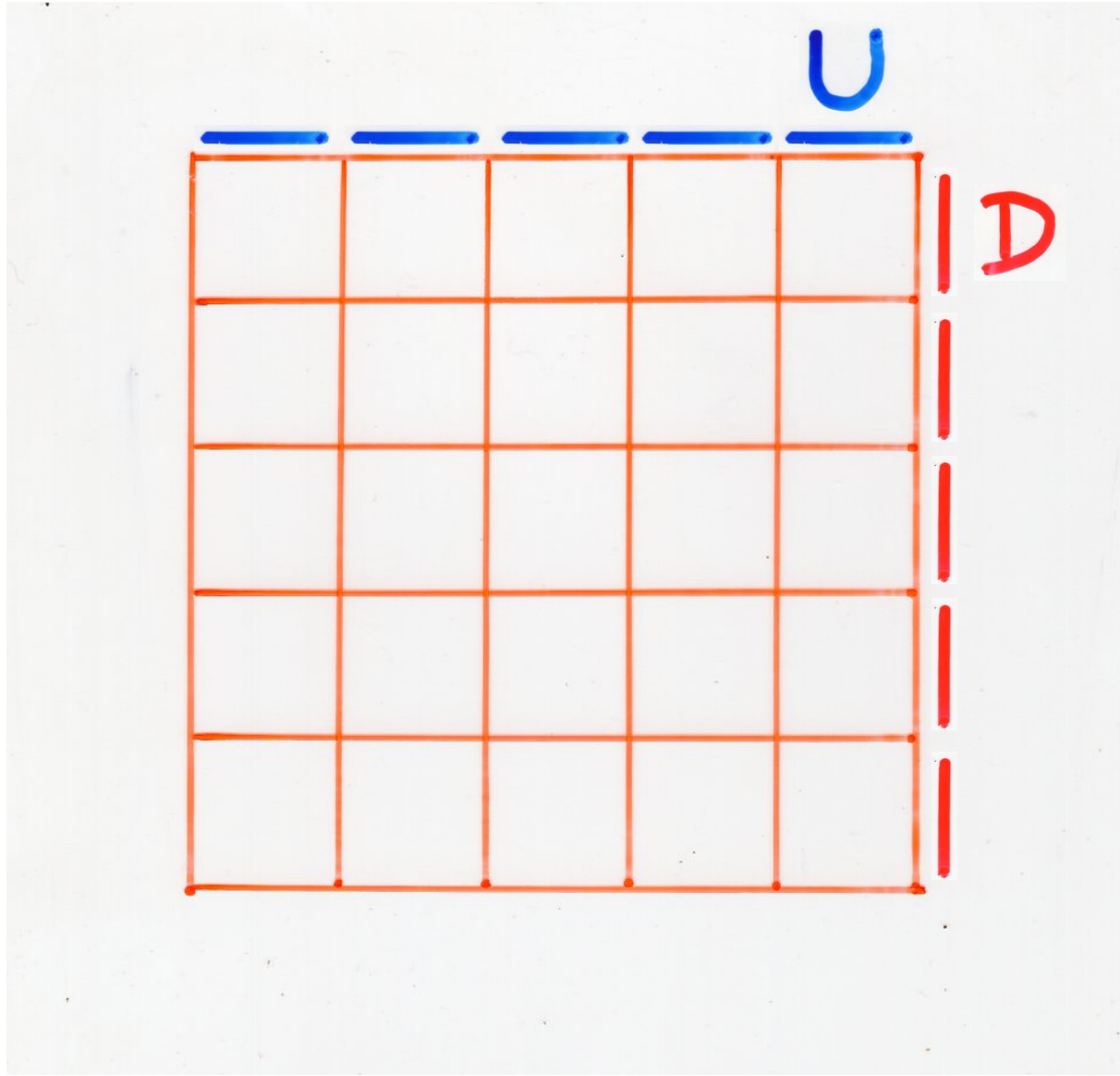
$$\left\{ \begin{array}{l} UD \rightarrow DU \\ UI_v \rightarrow I_v U \\ I_h D \rightarrow D I_h \\ I_h I_v \rightarrow I_v I_h \end{array} \right. \quad UD \rightarrow I_v I_h$$

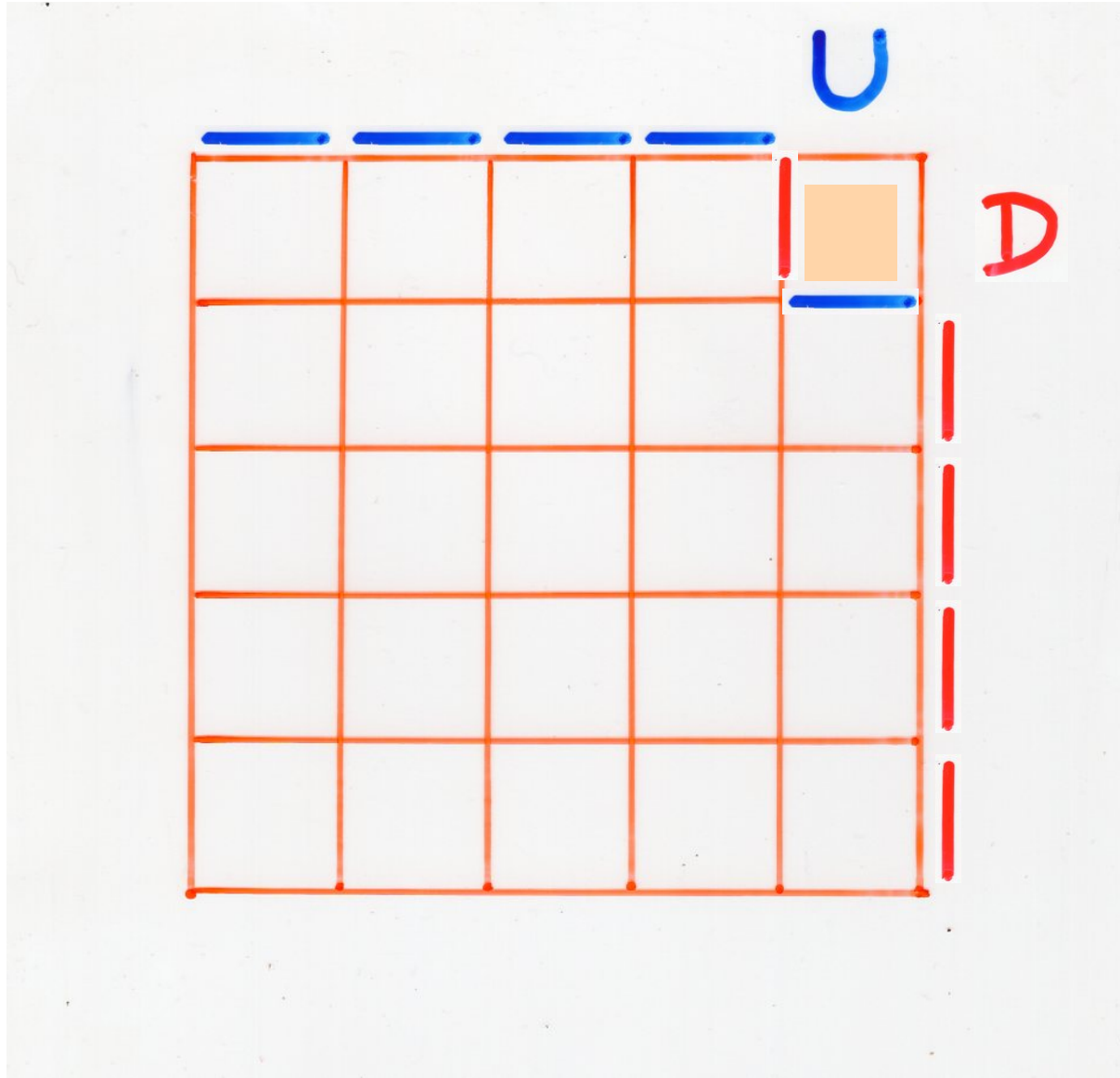
rewriting rules

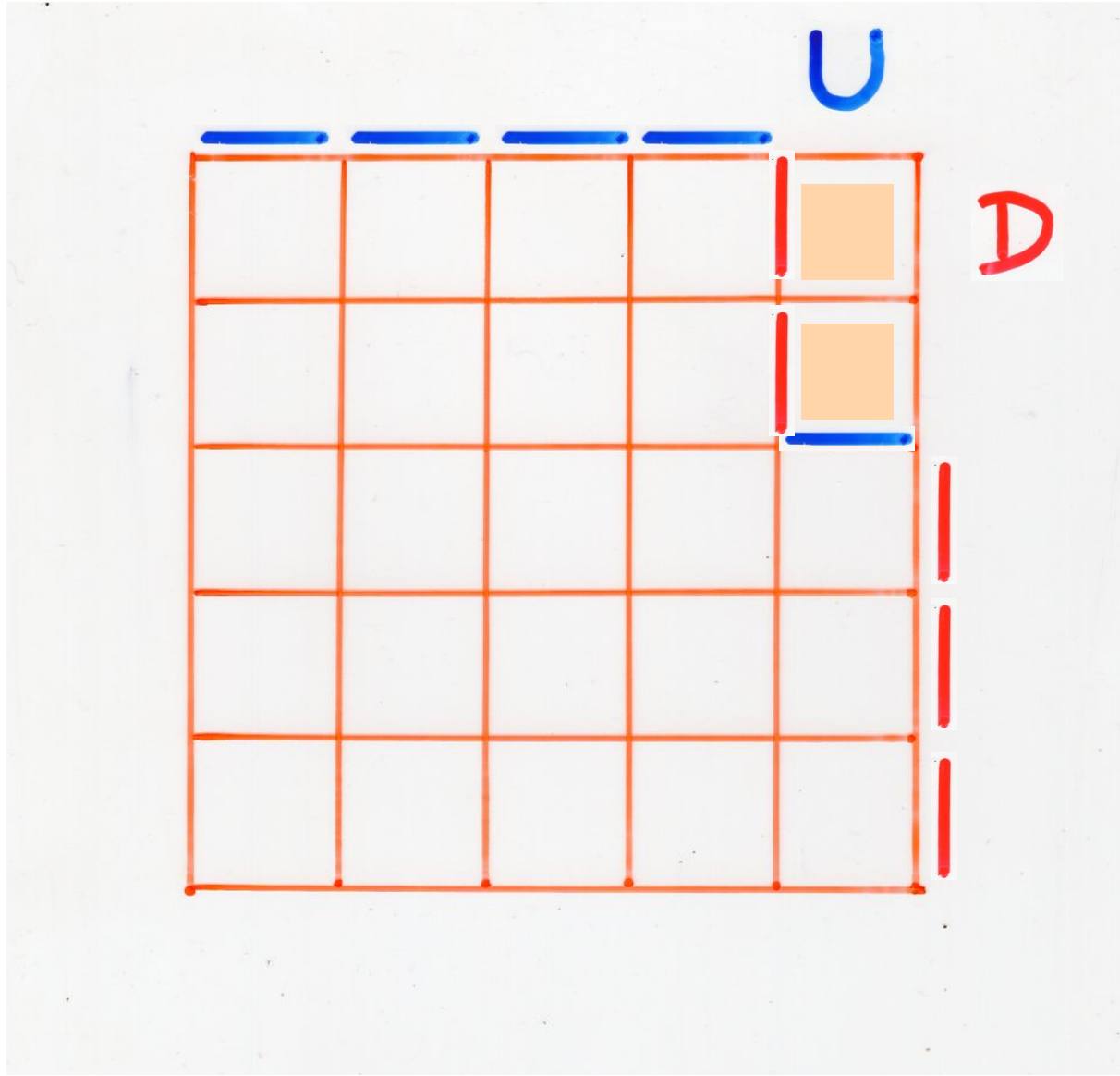


"planarization" of the "rewriting rules"

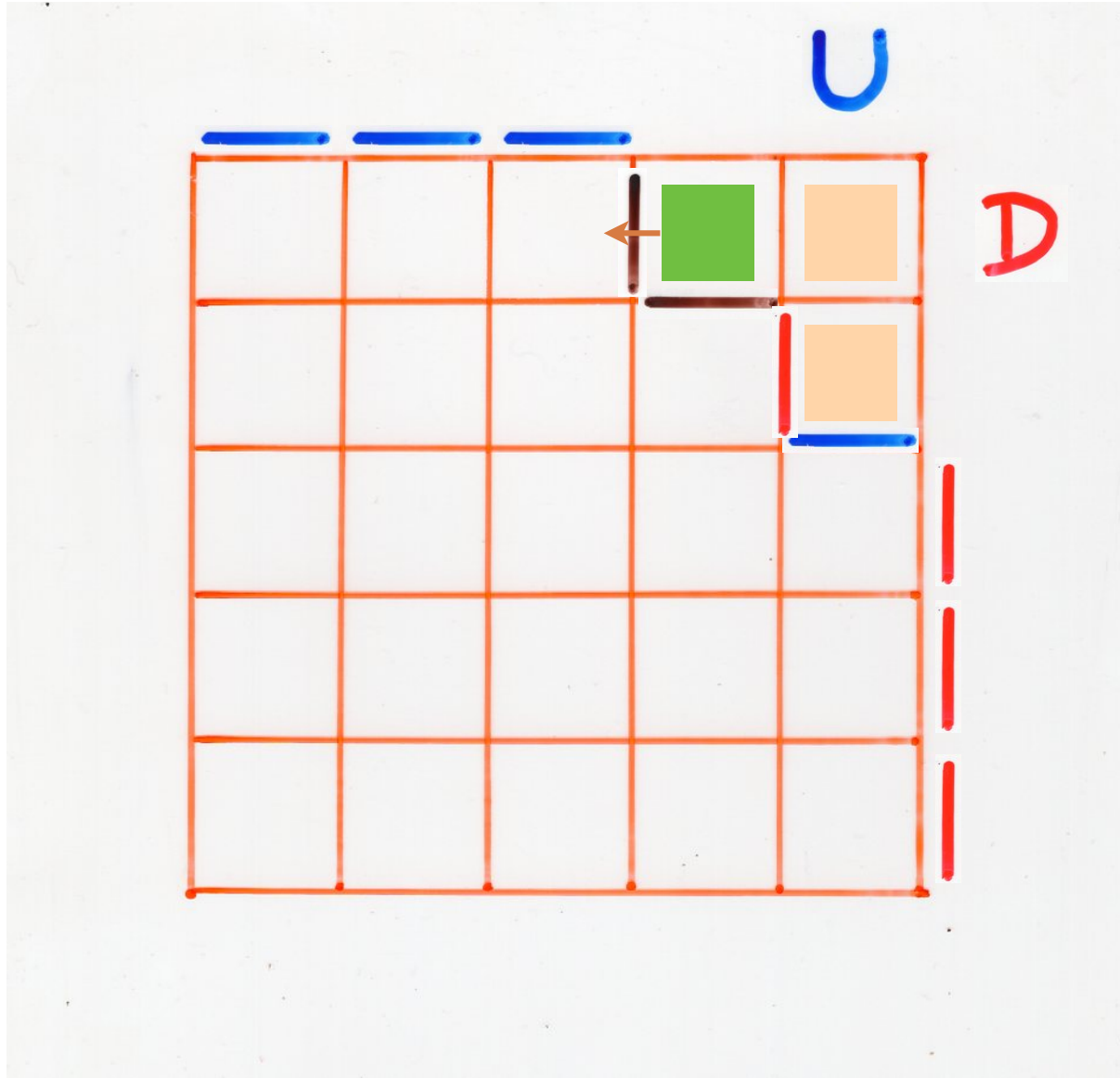




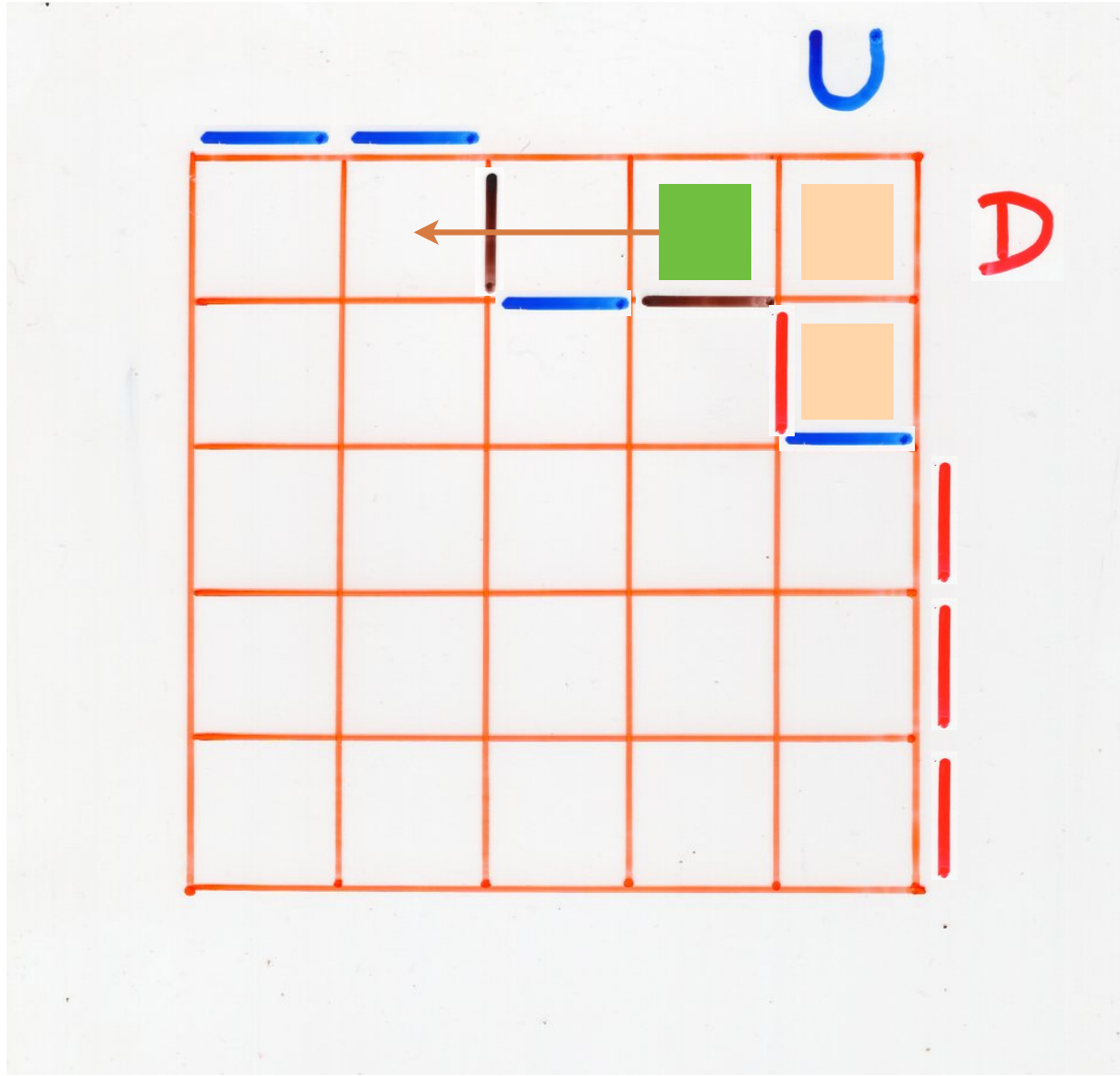


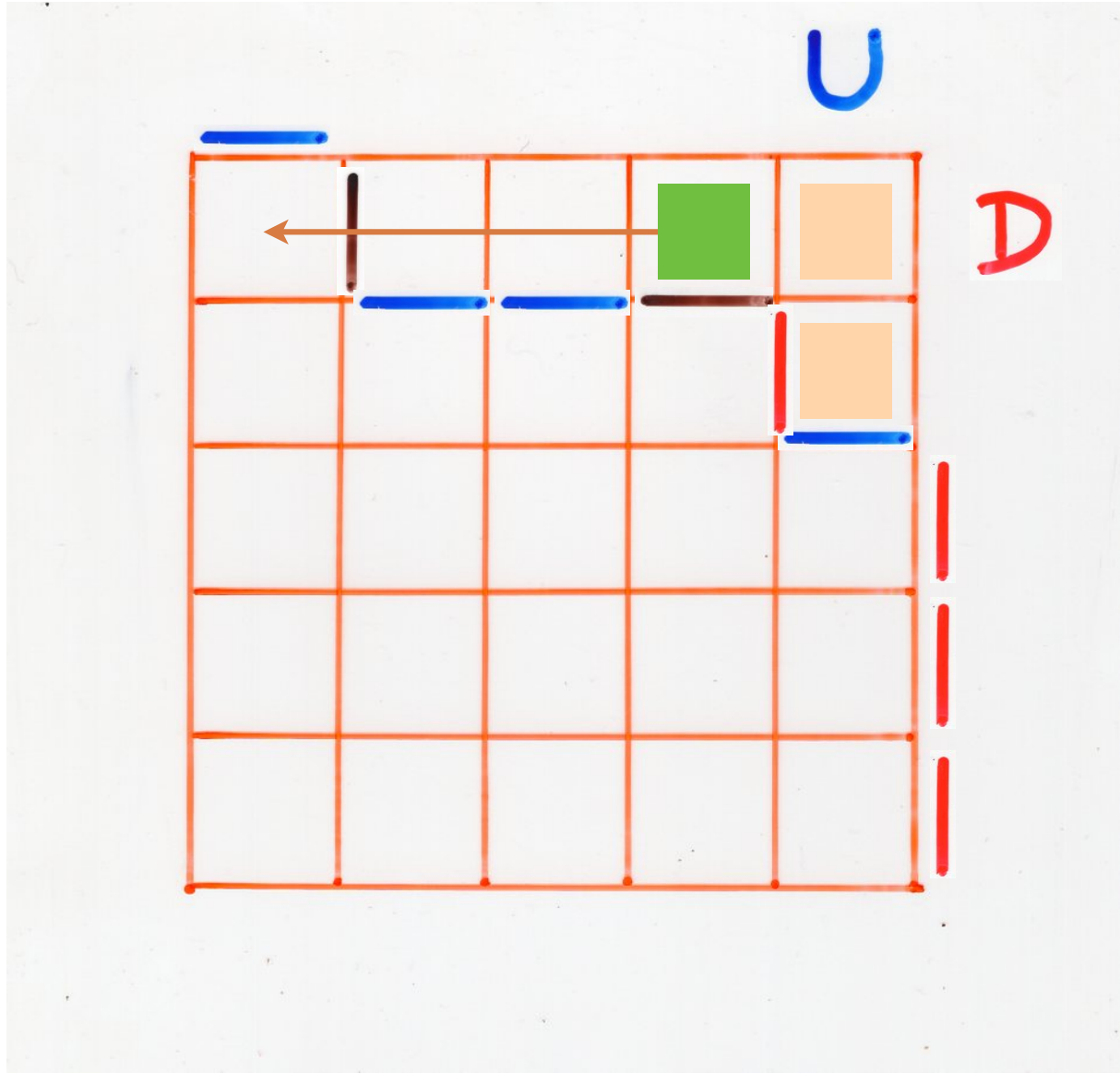


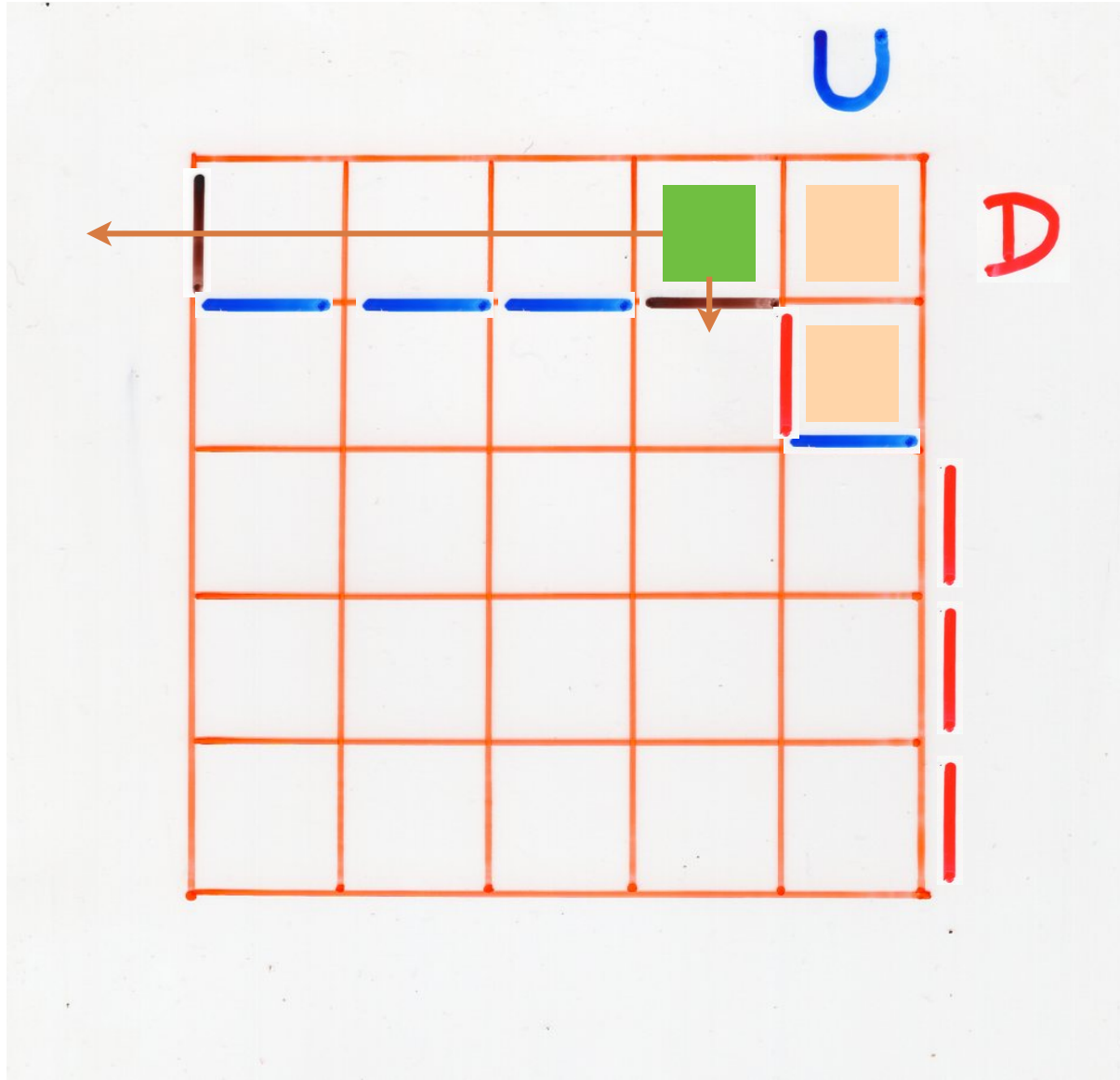
I_h

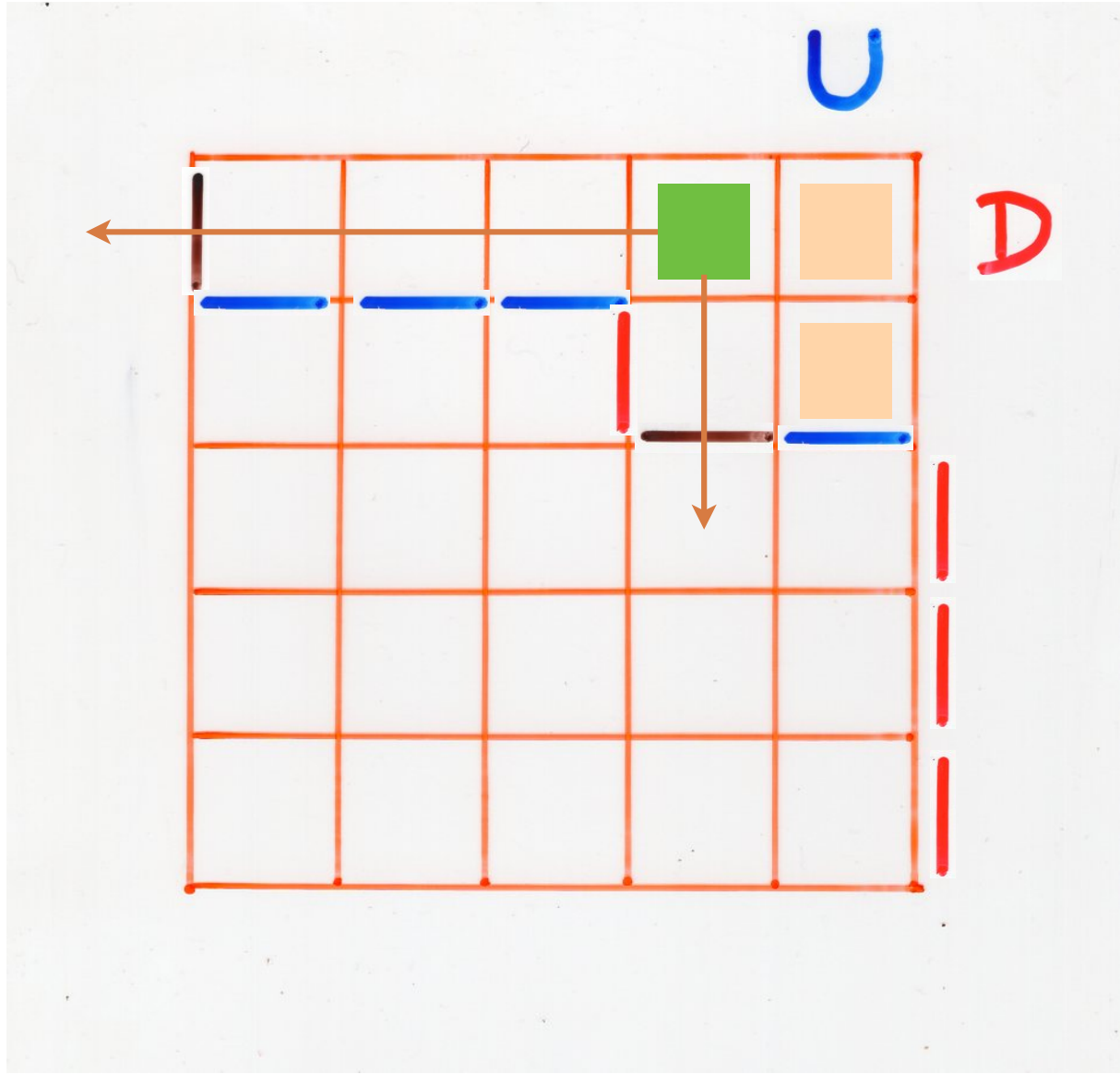


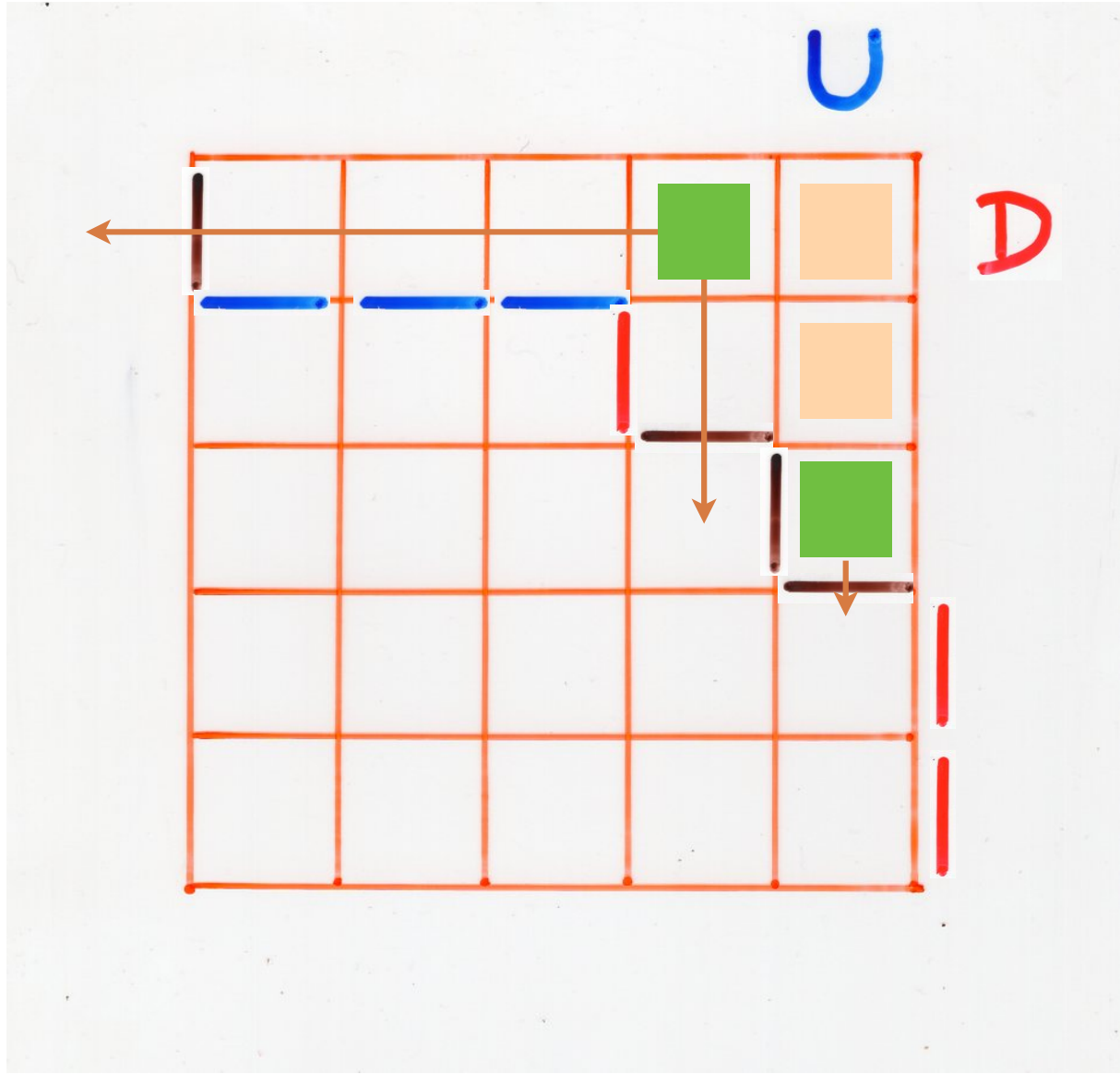
I_v

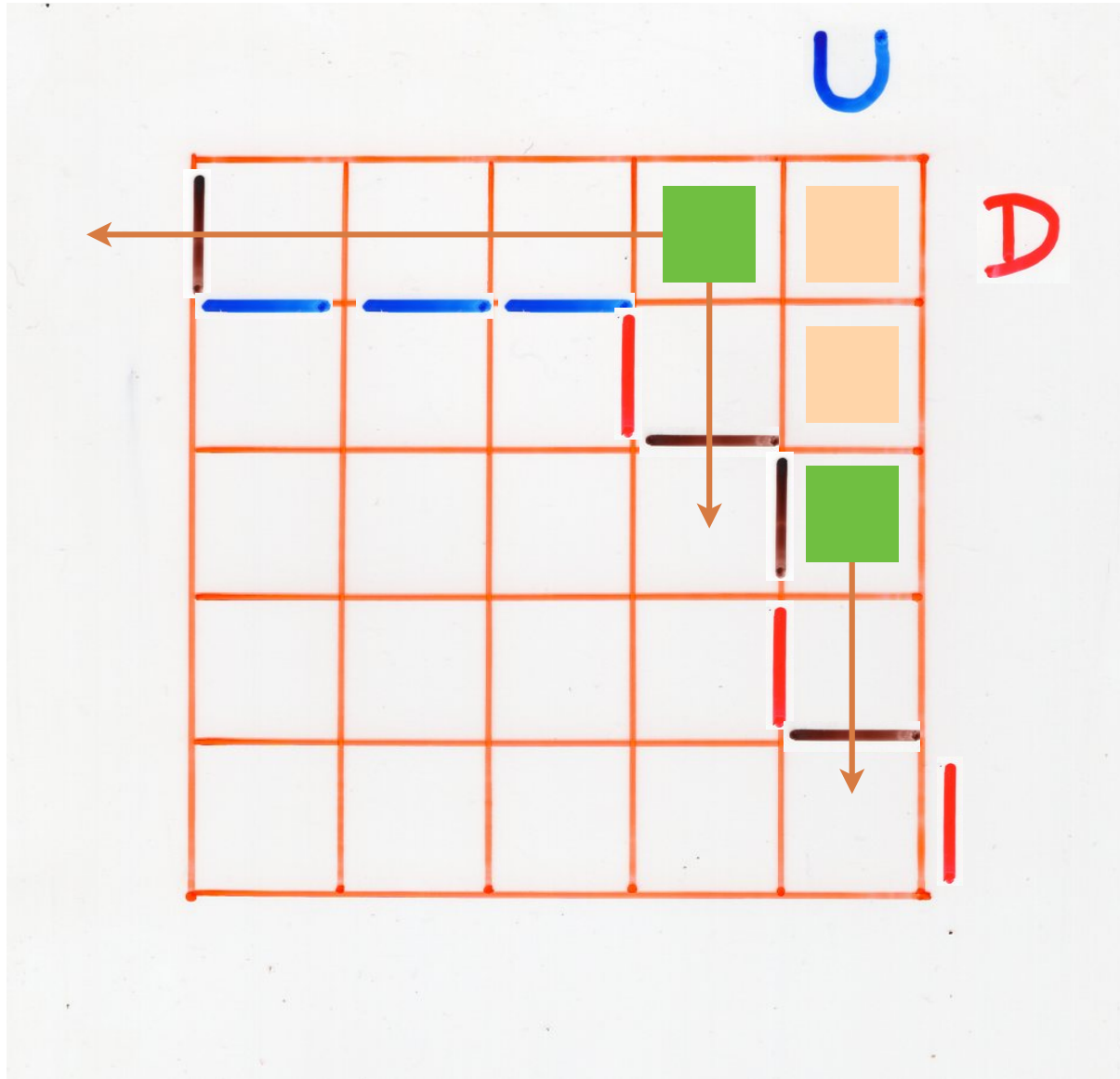


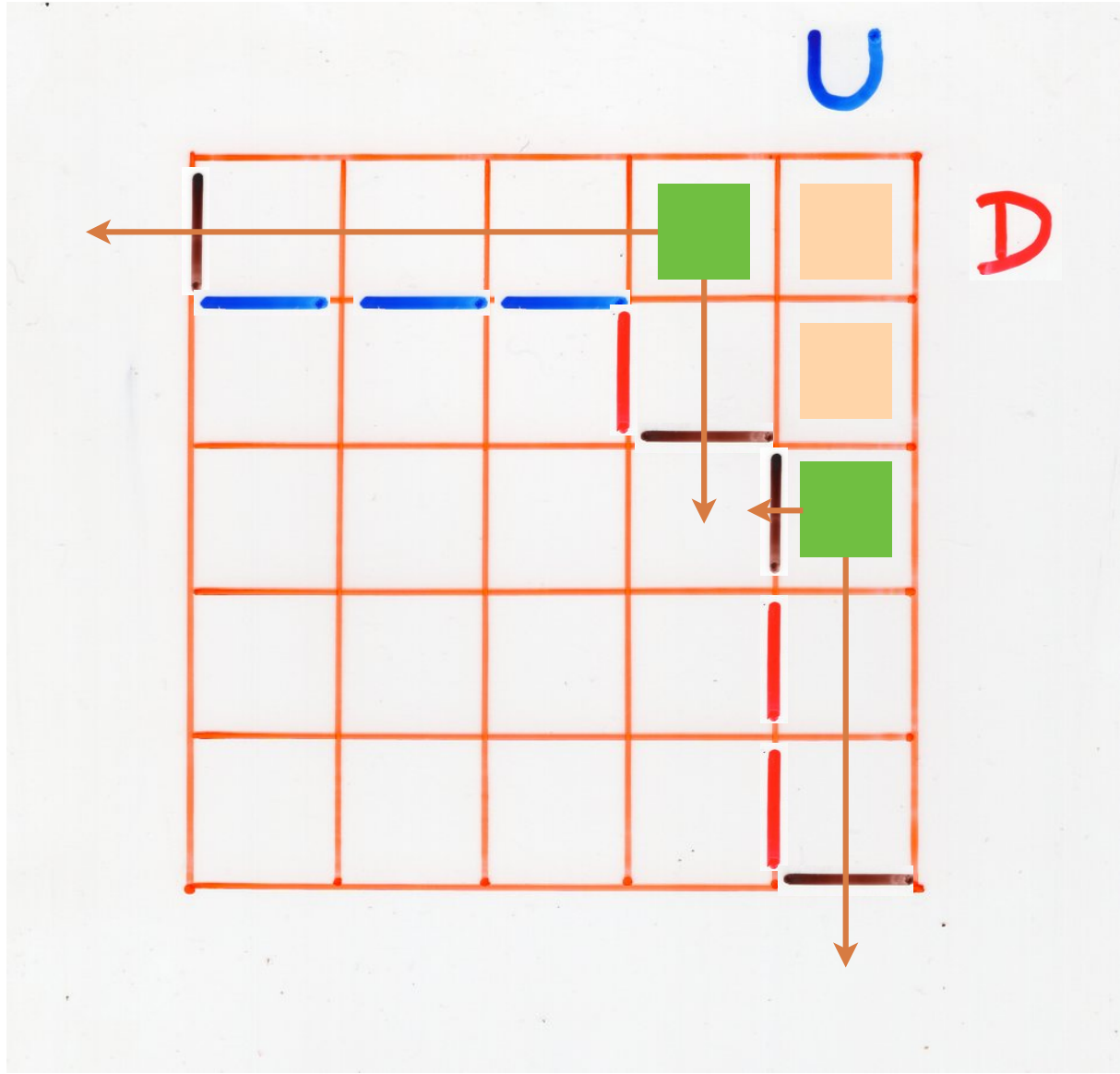


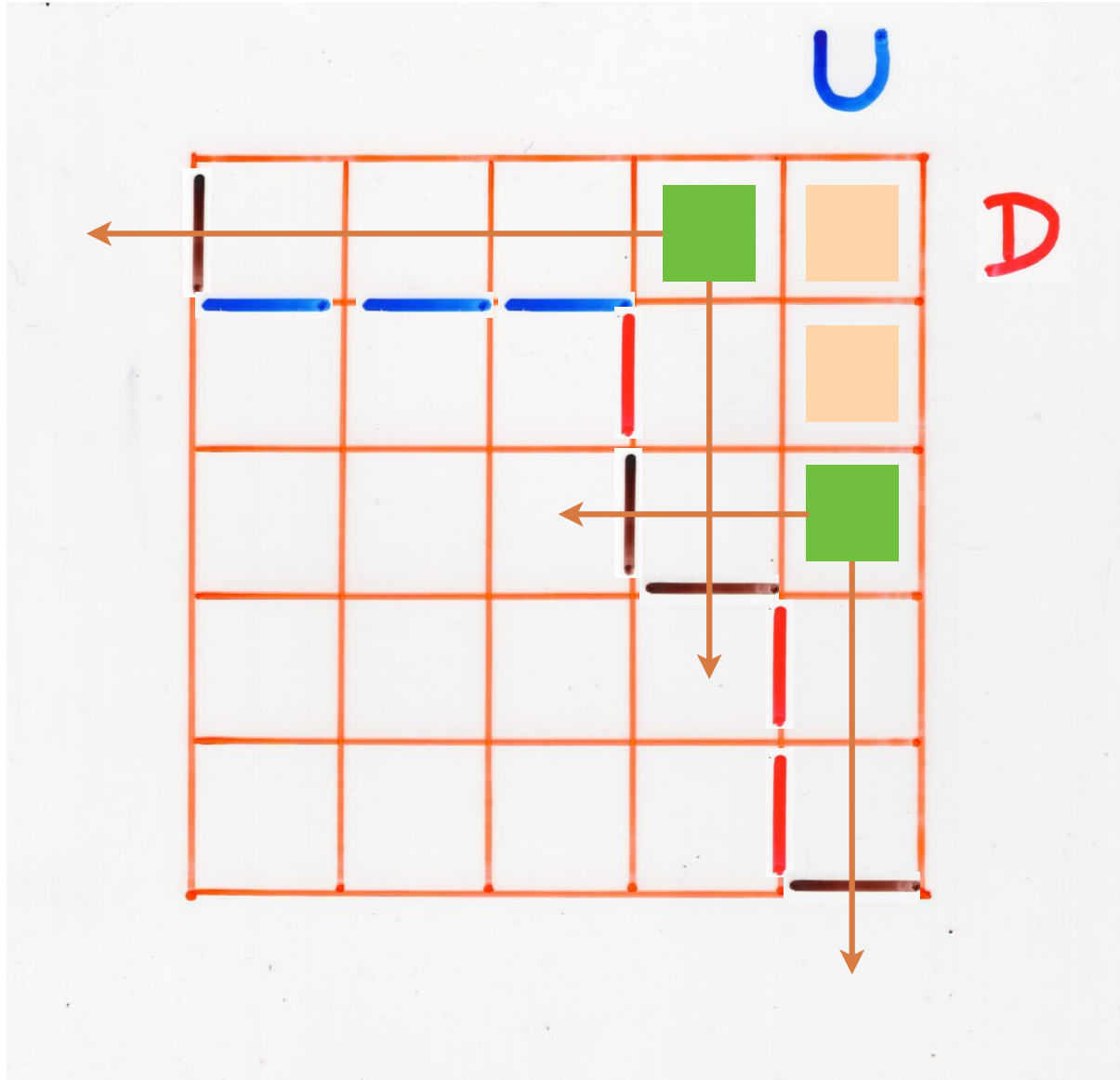


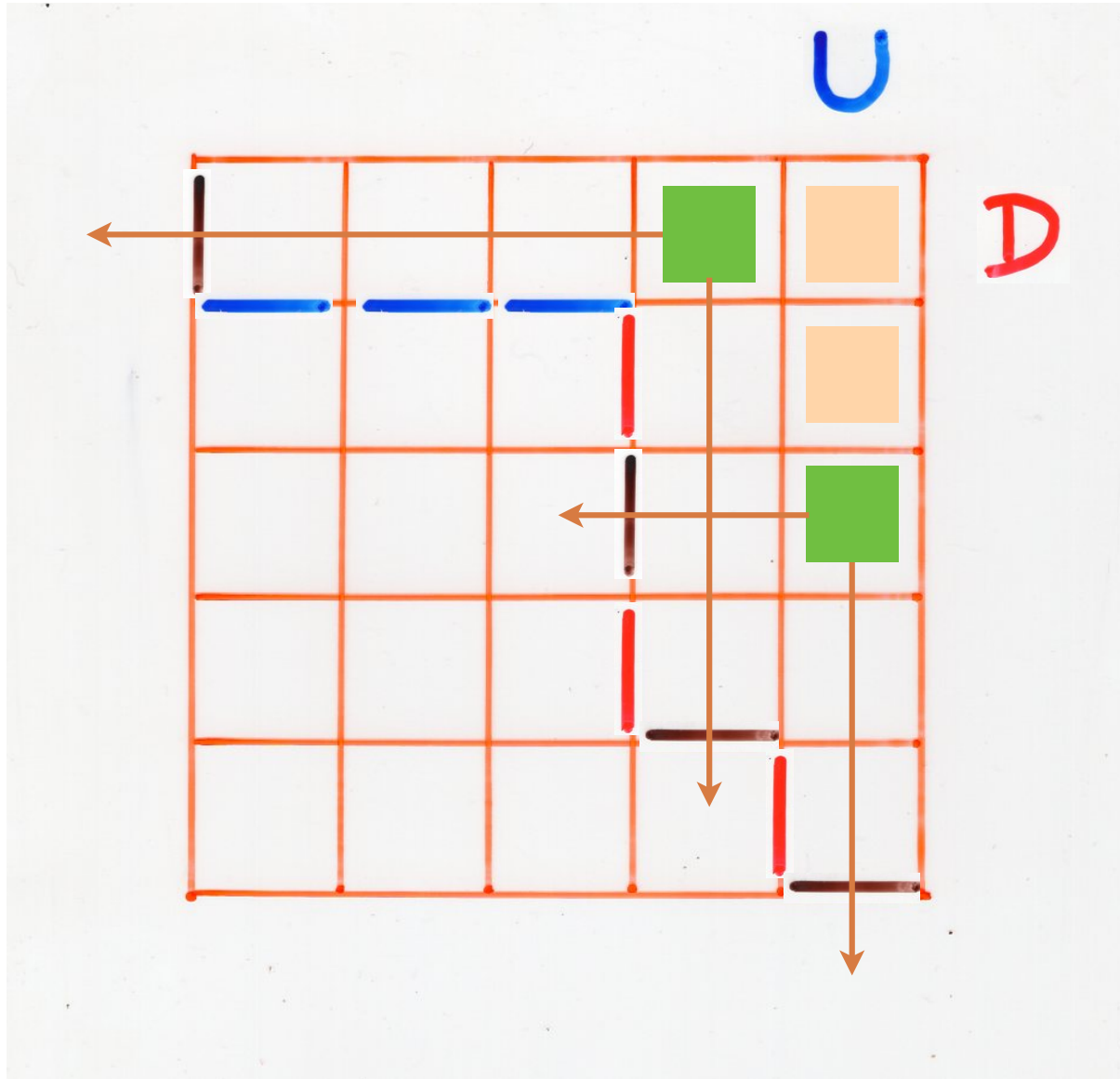


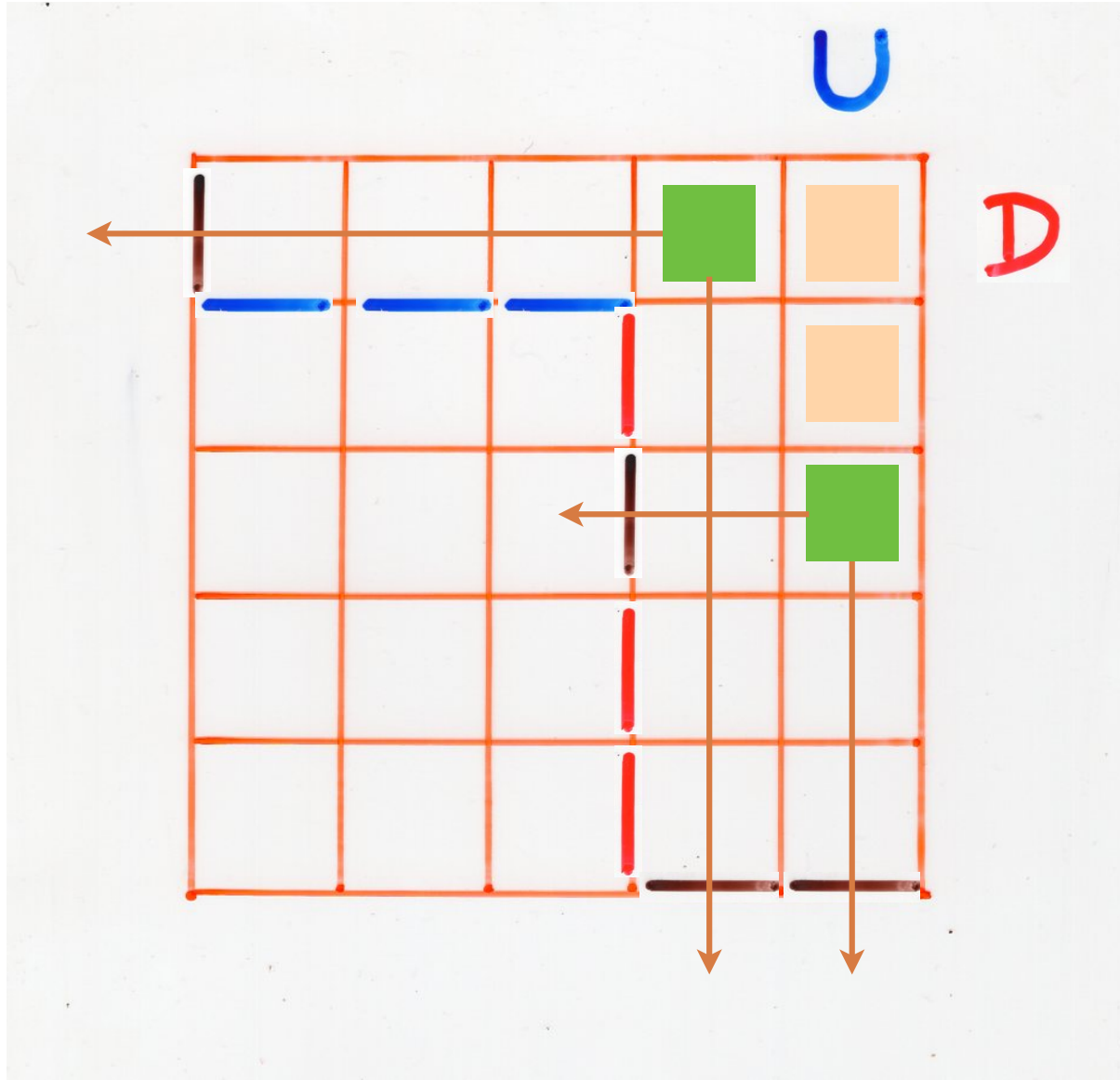


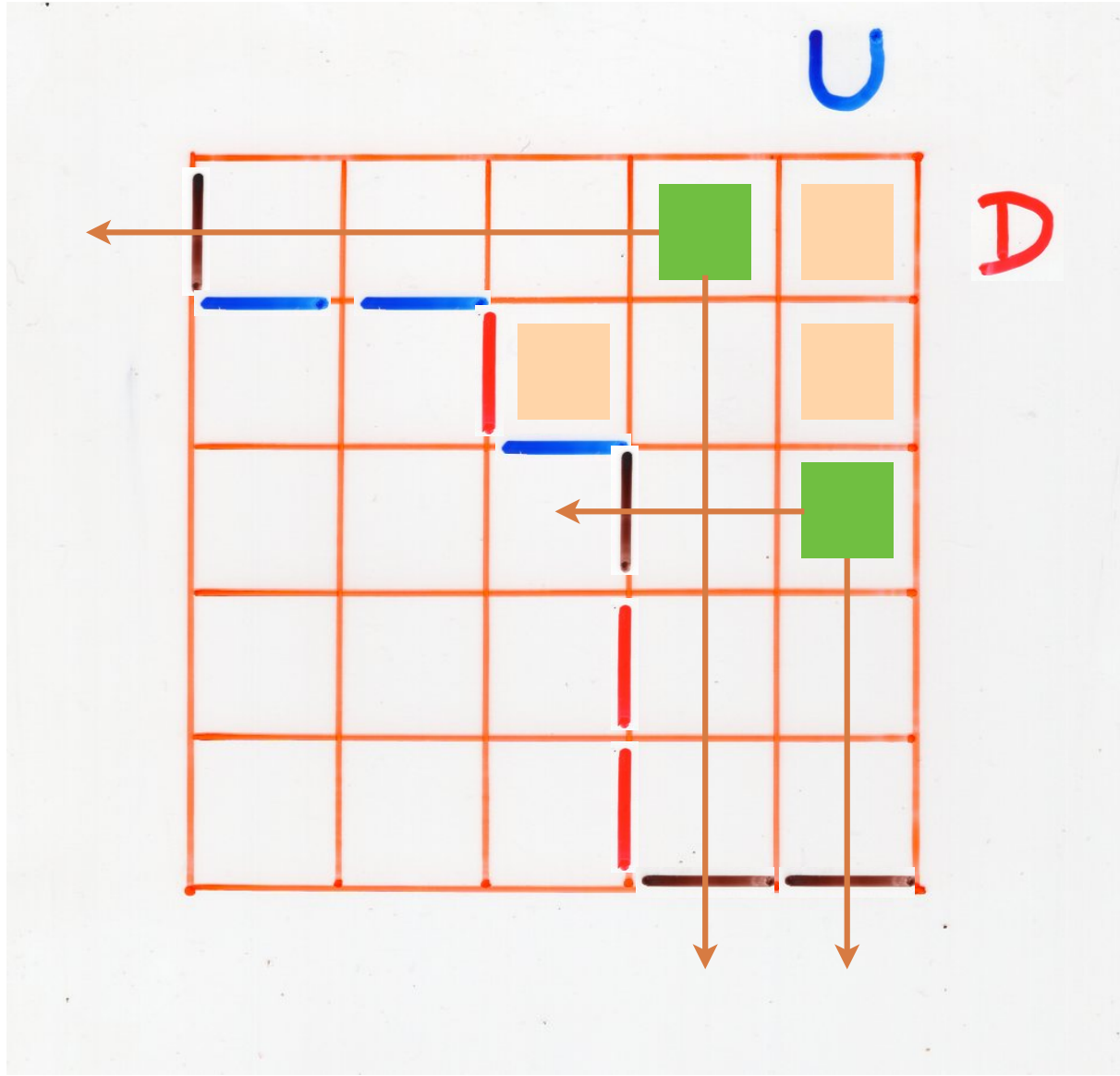


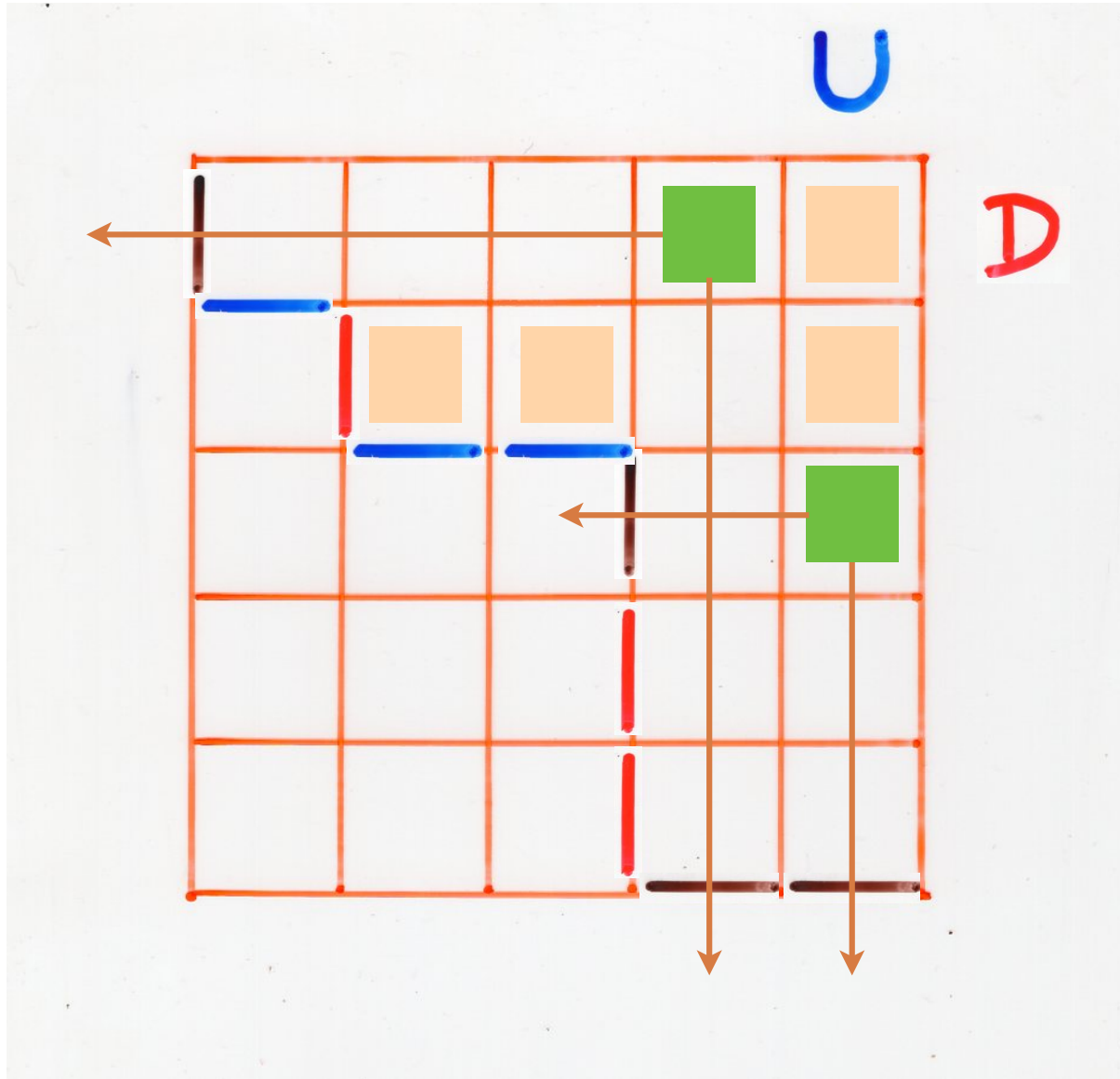


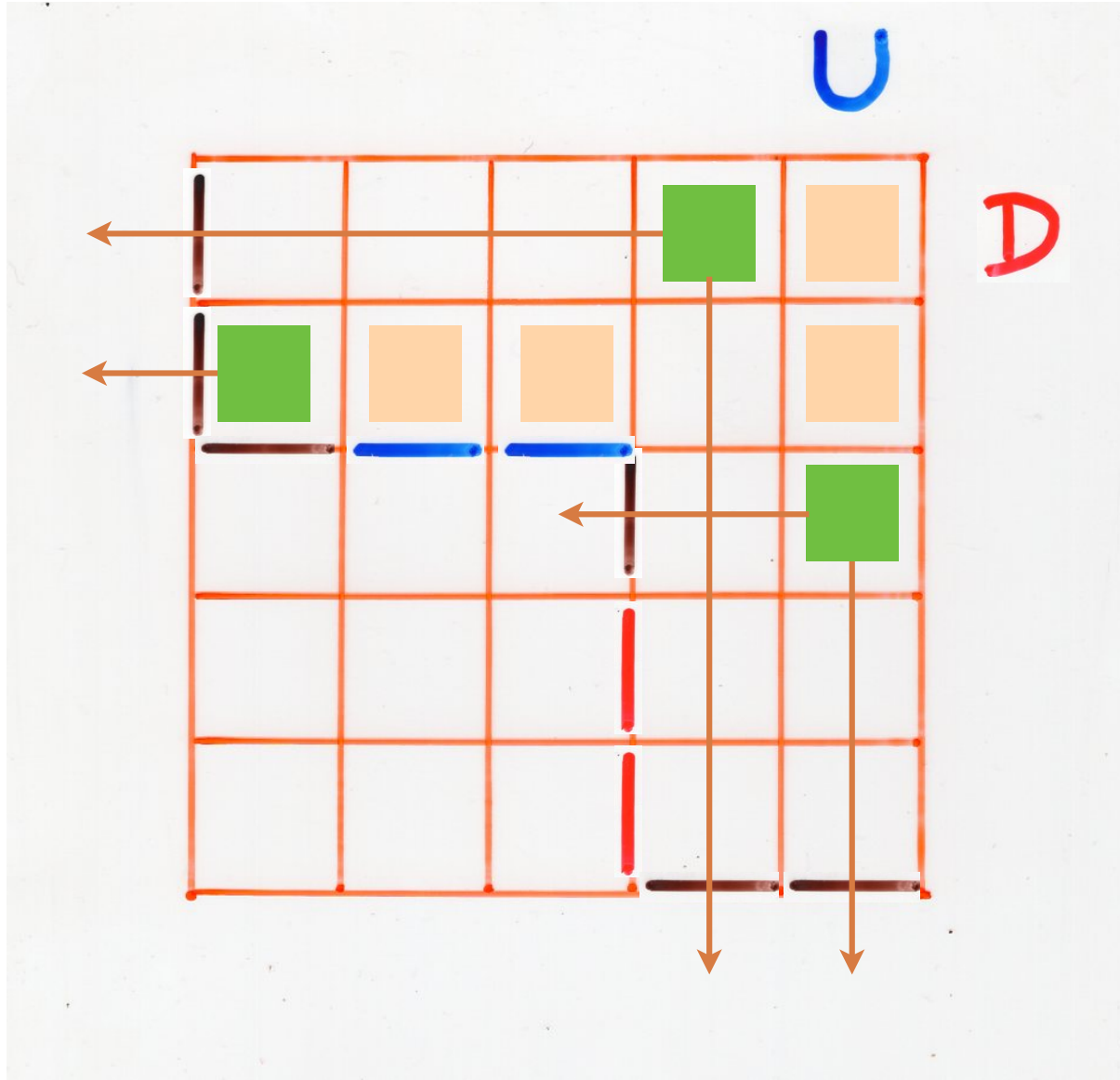


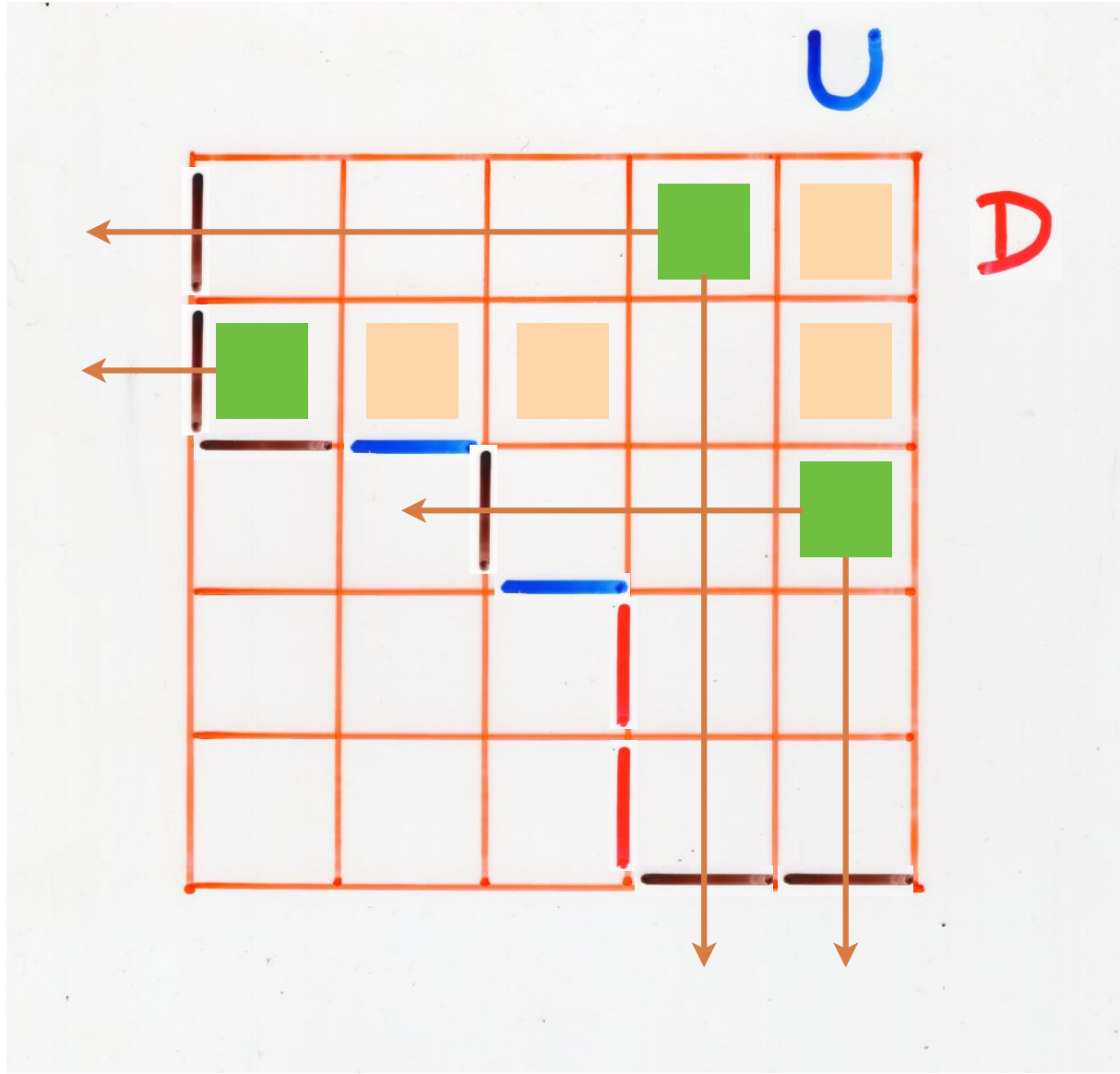


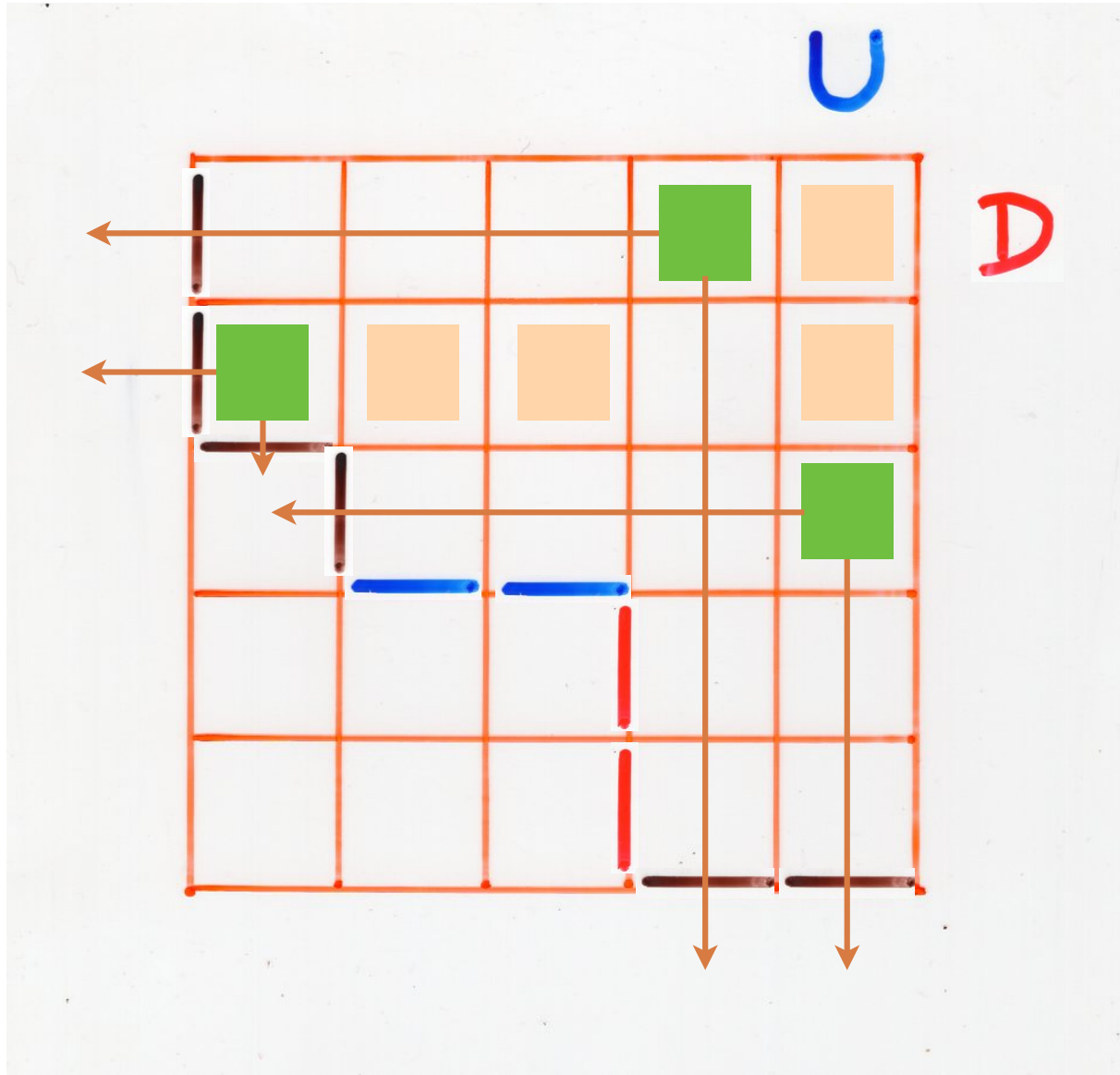


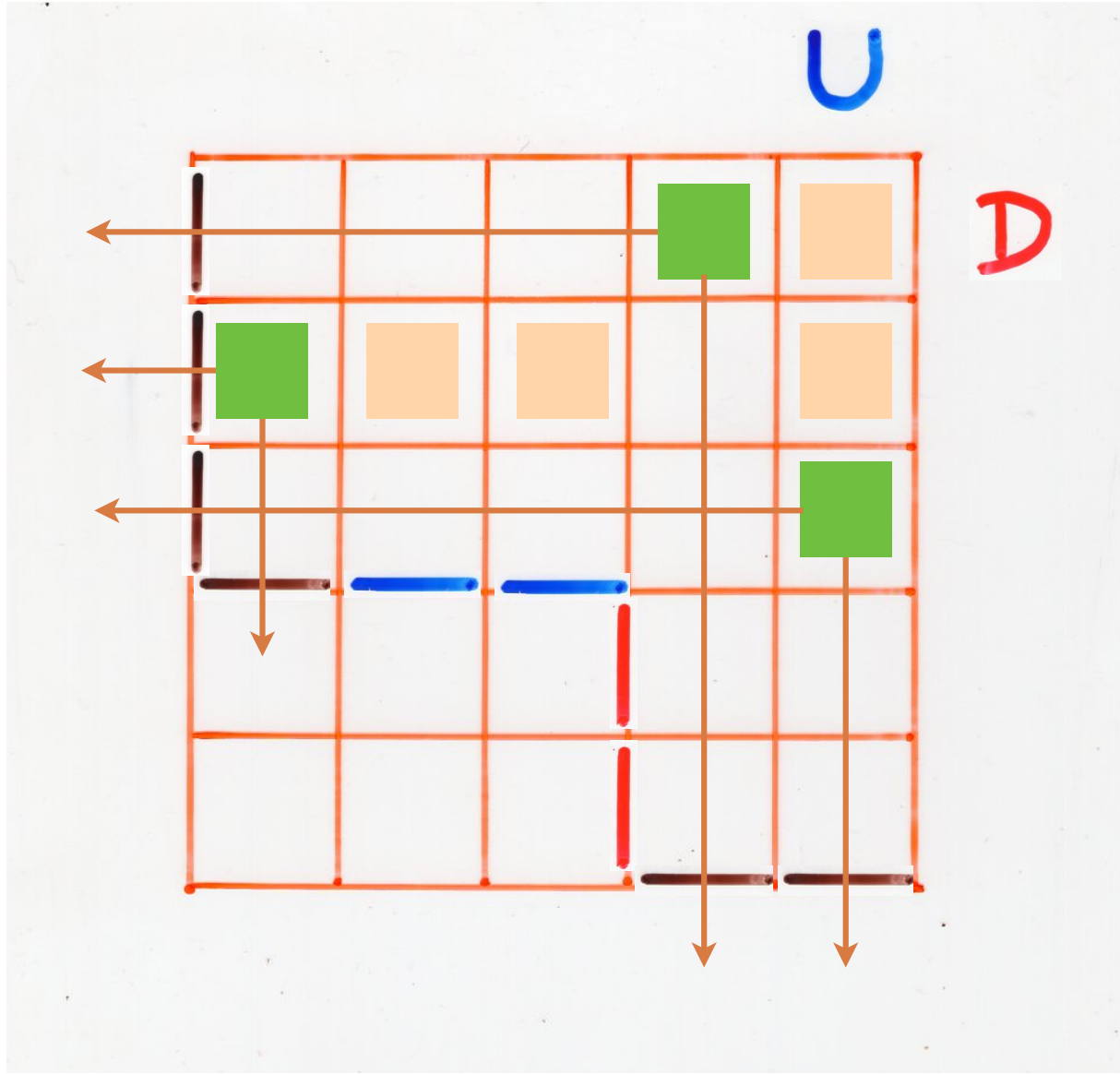


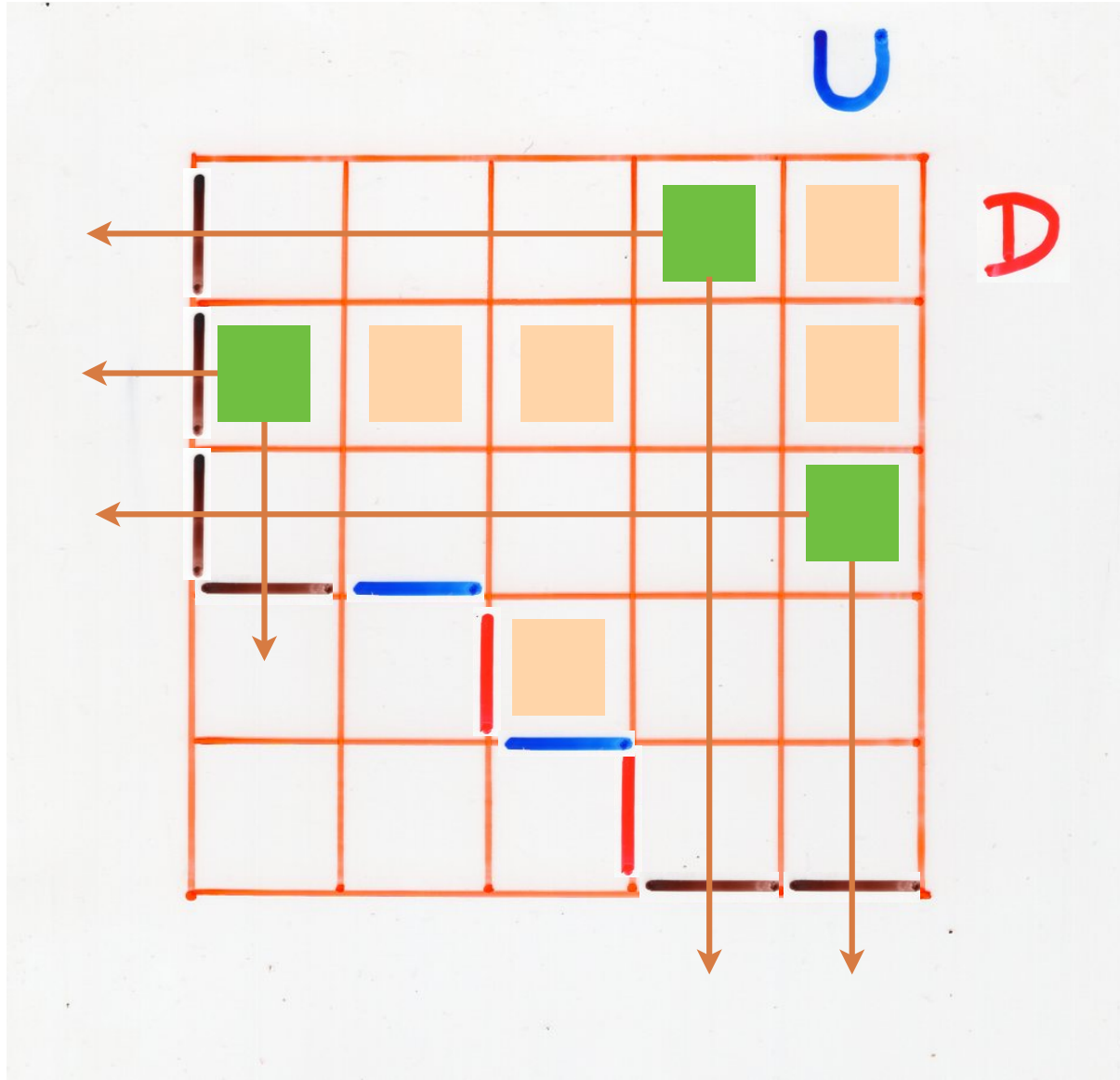


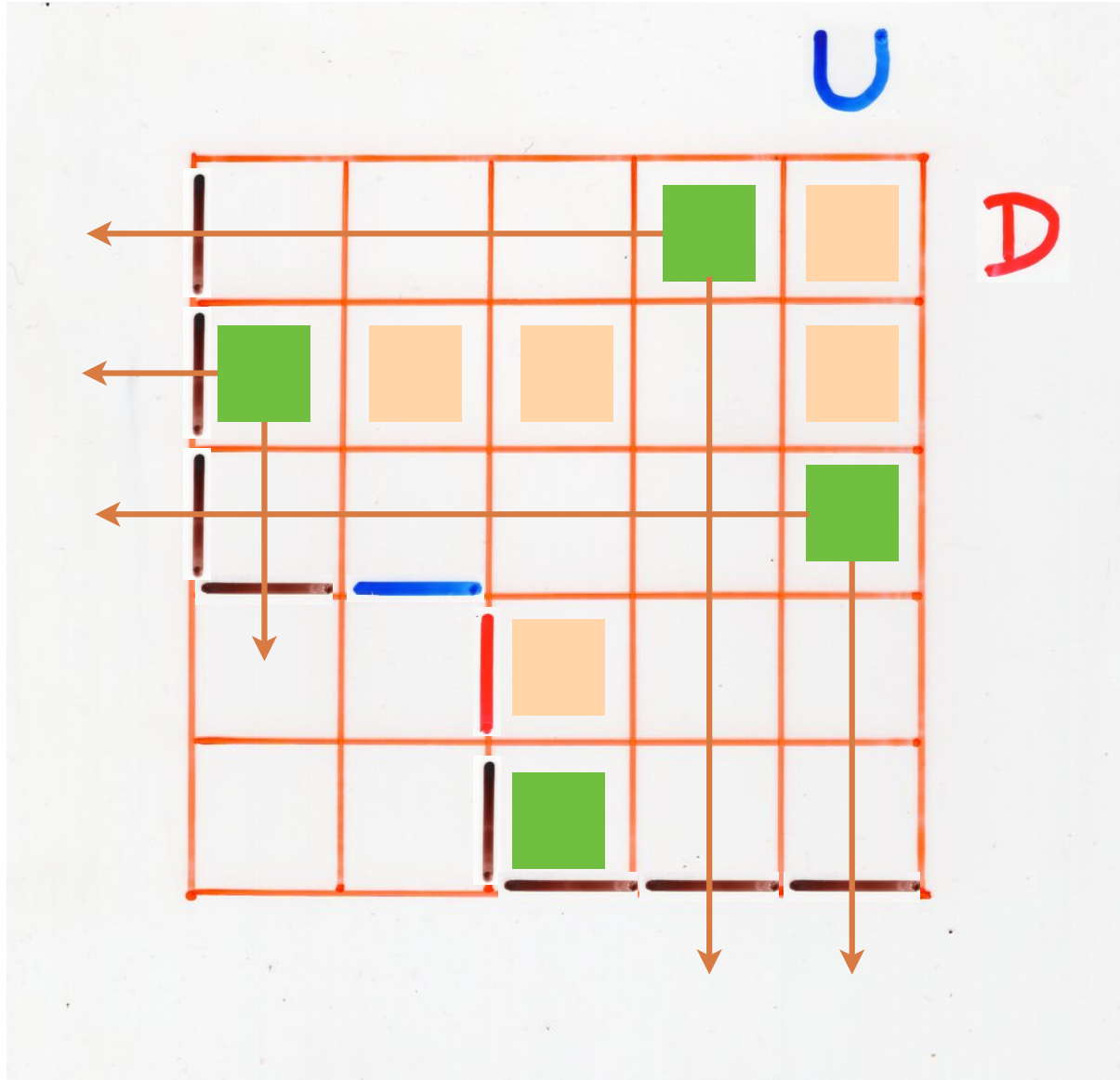


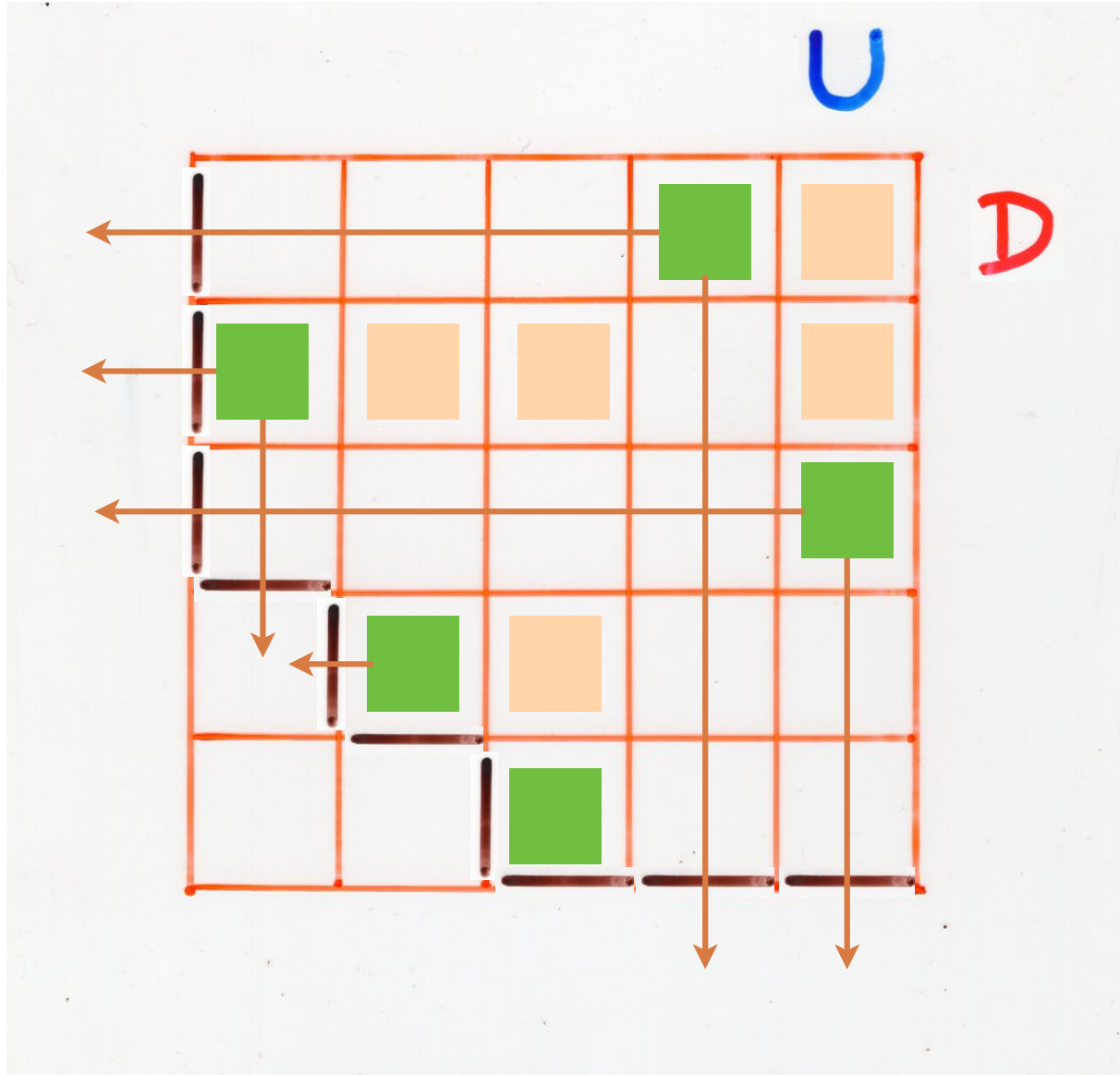


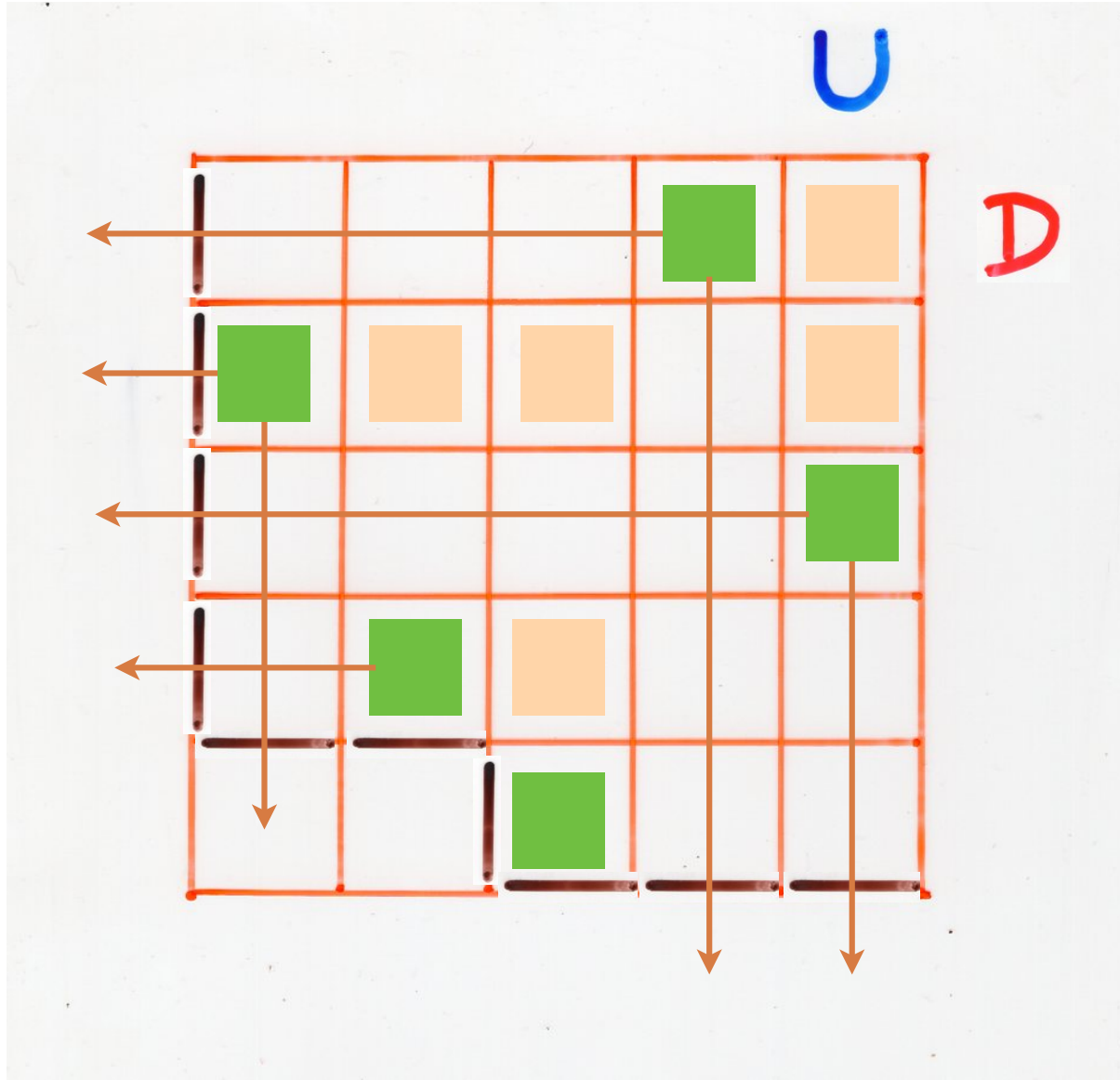


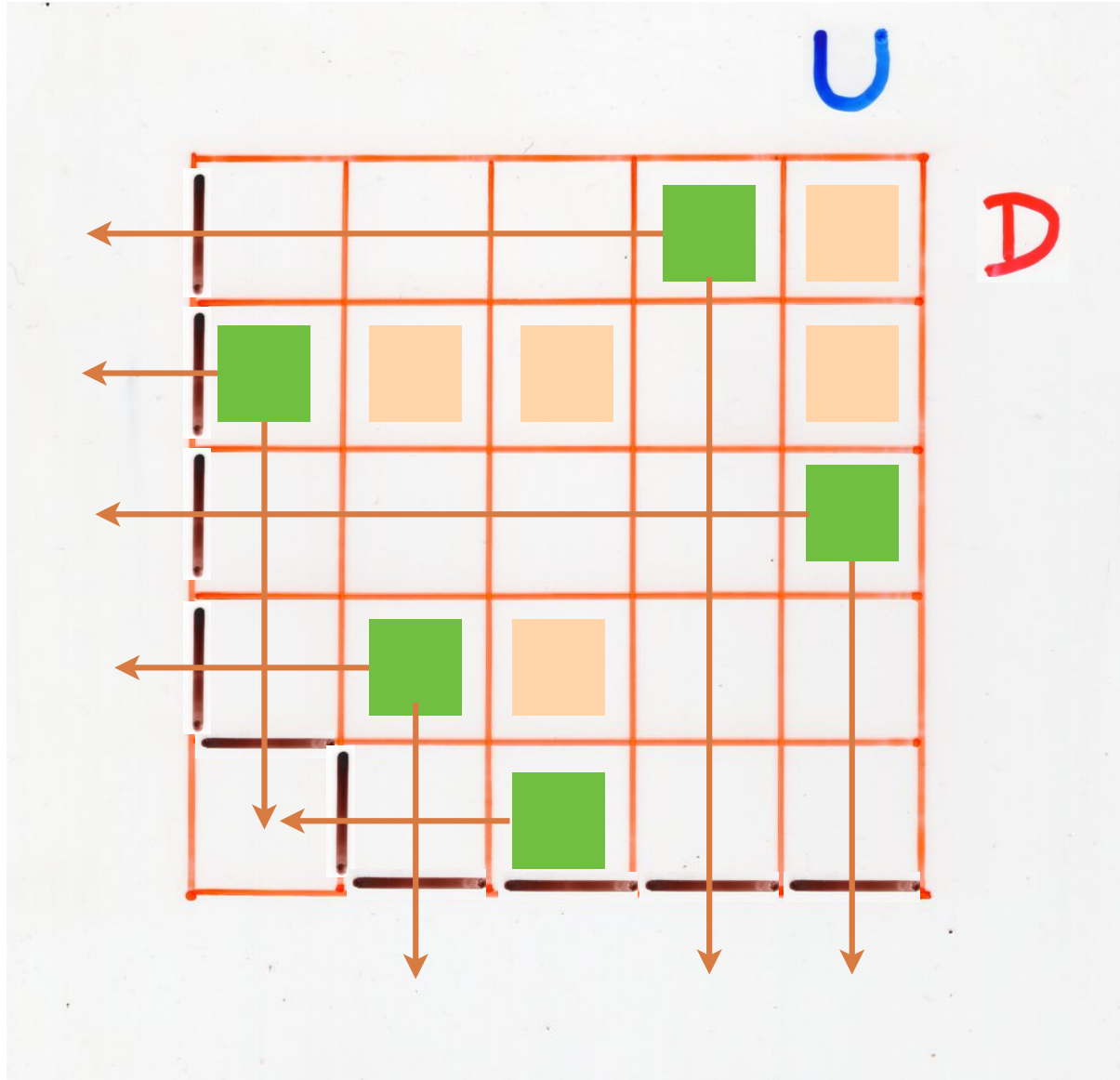


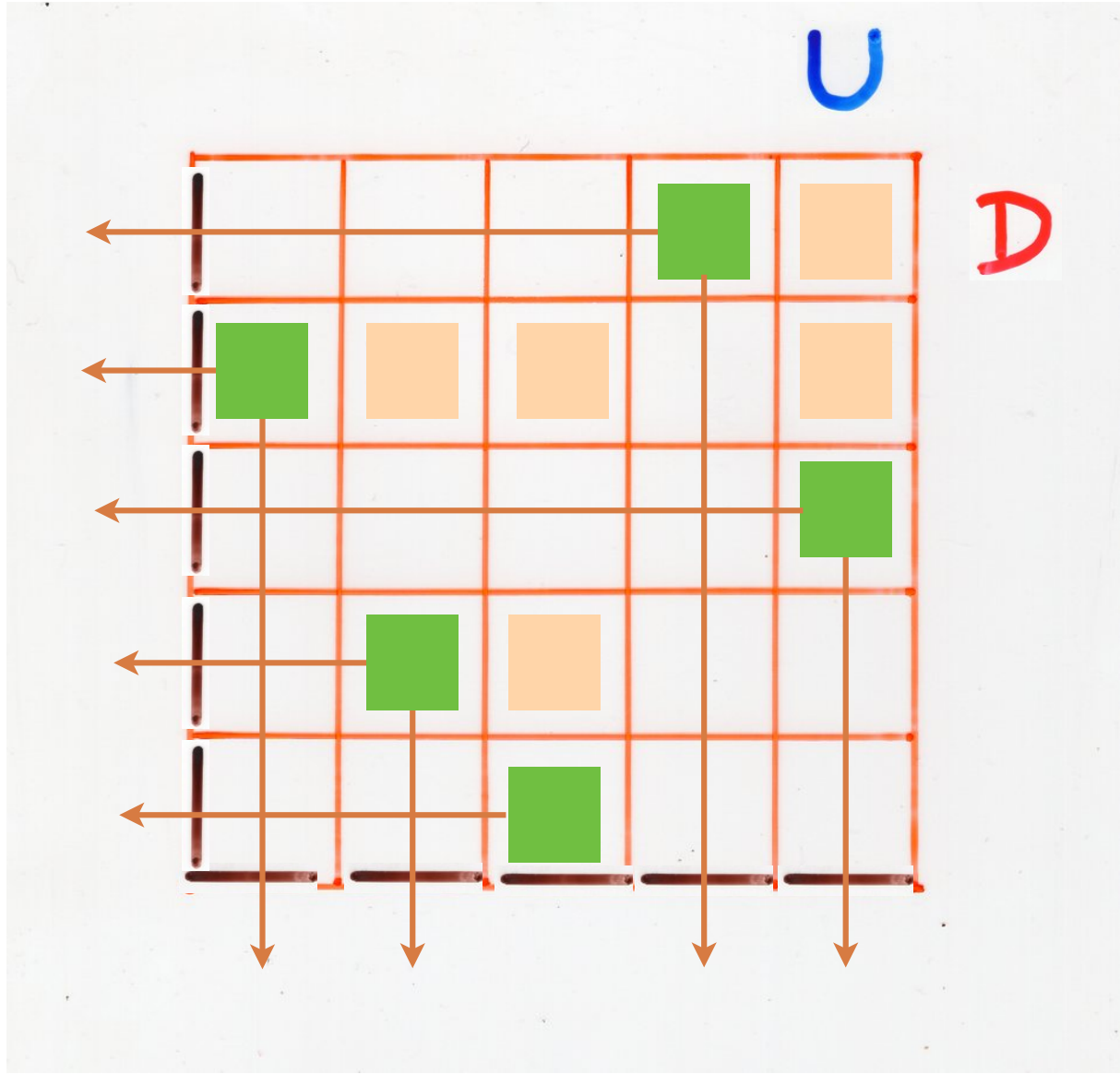


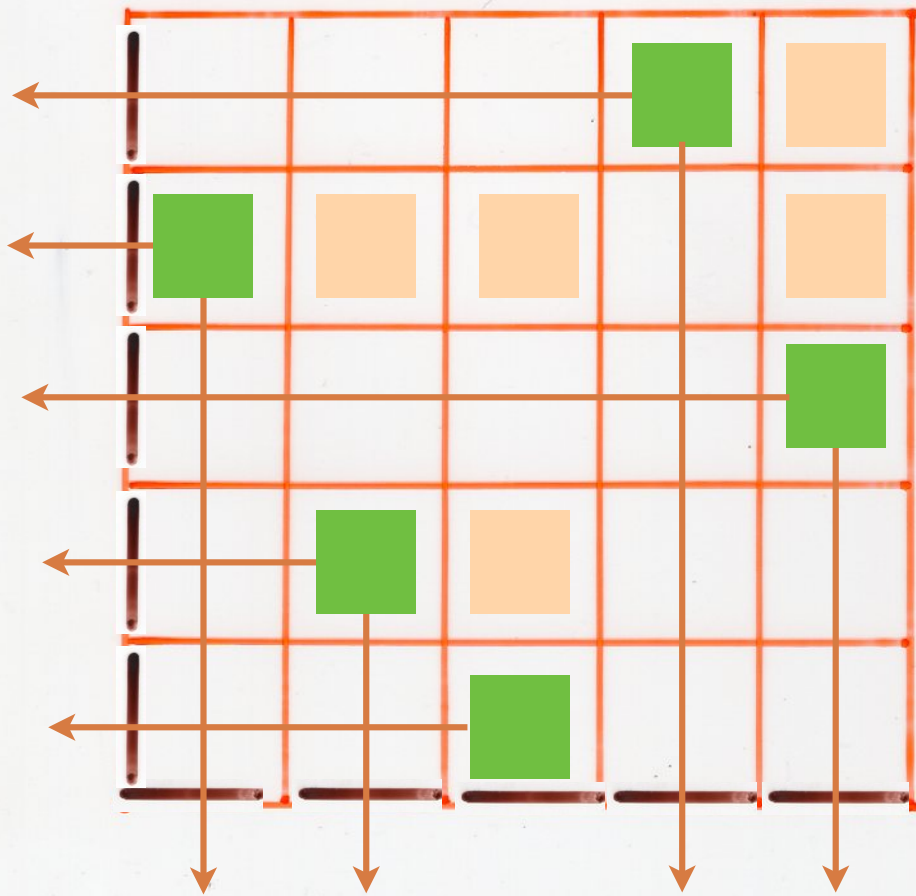










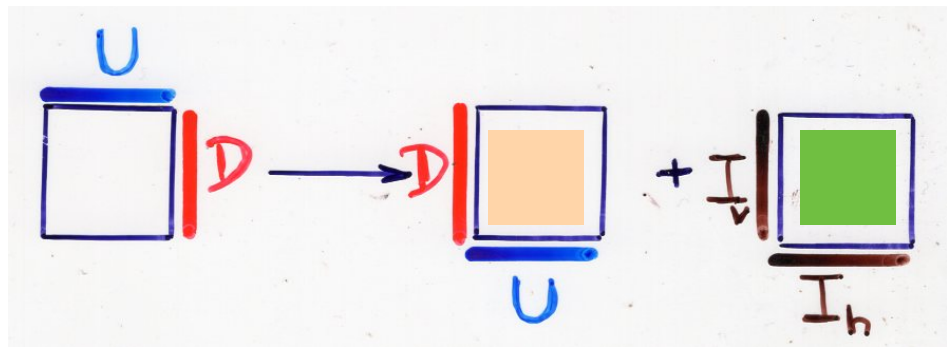


$$\begin{cases}
 U \mathcal{D} = \mathcal{D} U + I_v I_h \\
 U I_v = I_v U \\
 I_h \mathcal{D} = \mathcal{D} I_h \\
 I_h I_v = I_v I_h
 \end{cases}$$

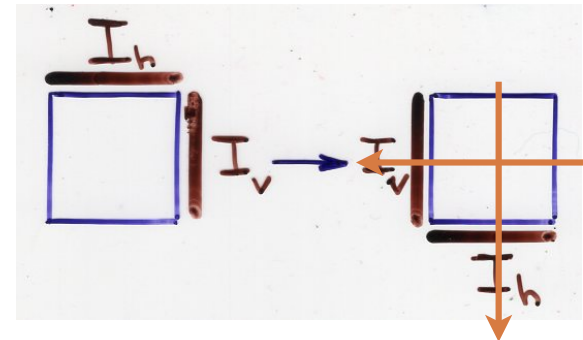
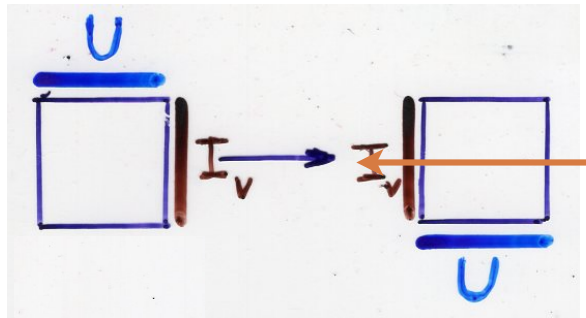
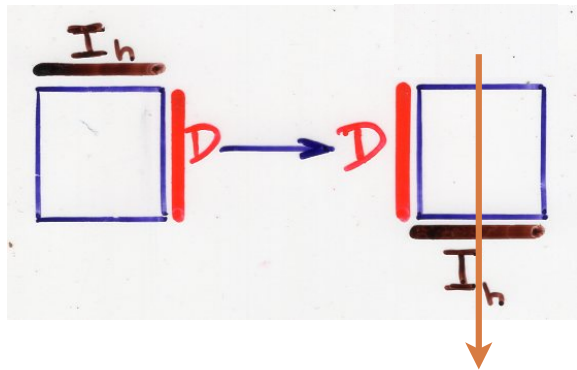
"complete"
Q-tableau

$$\begin{cases} U D = D U + I_v I_h \\ U I_v = I_v U \\ I_h D = D I_h \\ I_h I_v = I_v I_h \end{cases}$$

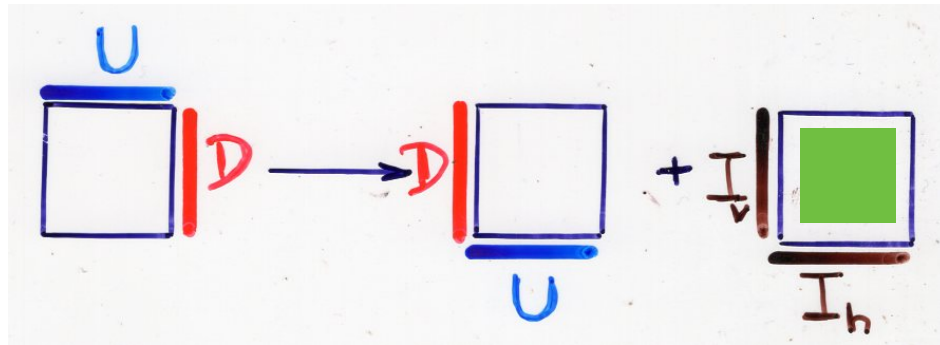
$$U D = q D U + I$$



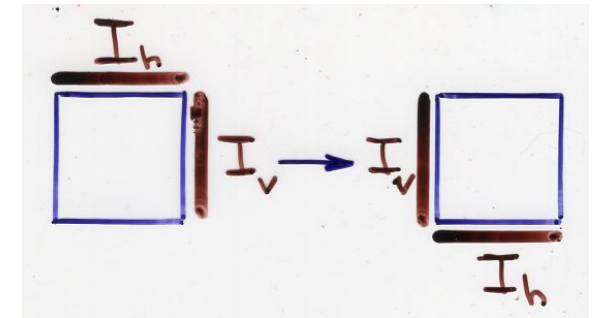
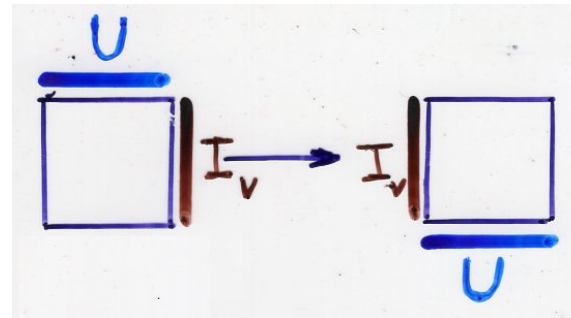
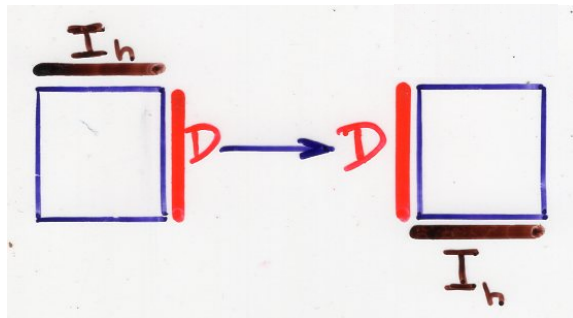
"complete"
Q-tableau

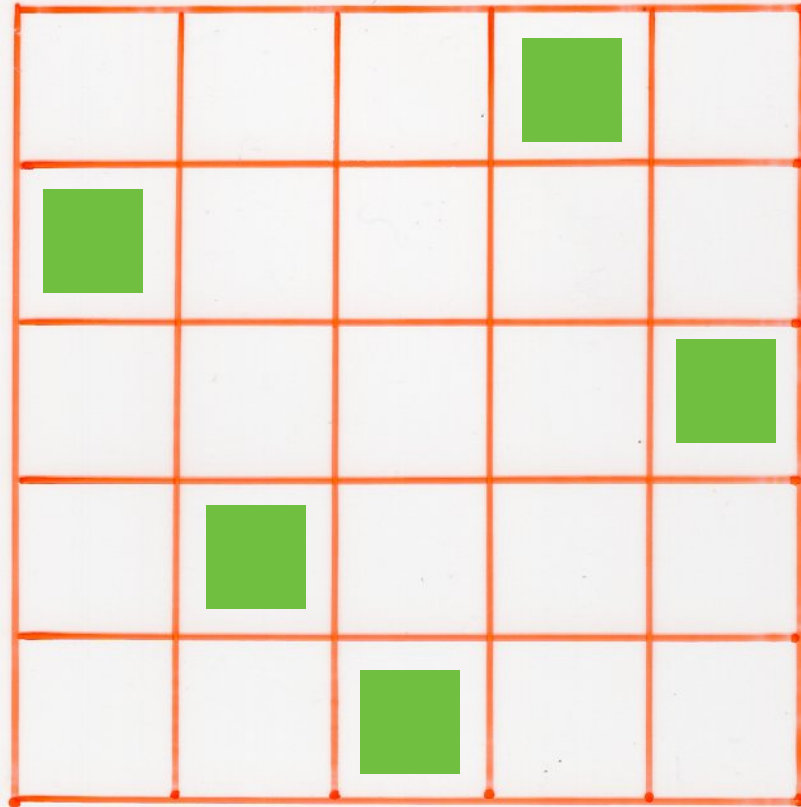


$$\begin{cases} U \mathcal{D} = \mathcal{D} U + I_v I_h \\ U I_v = I_v U \\ I_h \mathcal{D} = \mathcal{D} I_h \\ I_h I_v = I_v I_h \end{cases}$$



Q-tableau

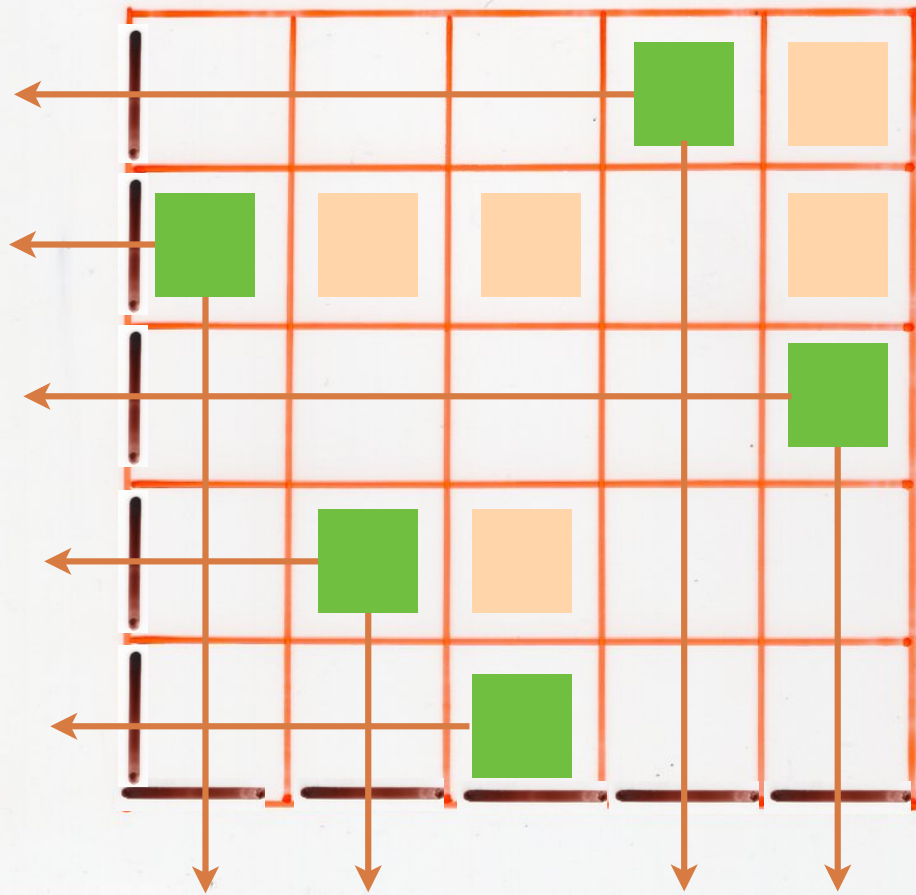




permutation
as a \mathbb{Q} -tableau

q

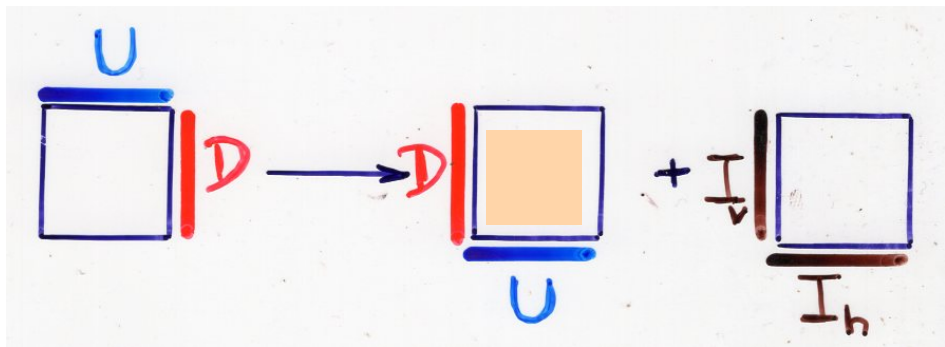
q-analog



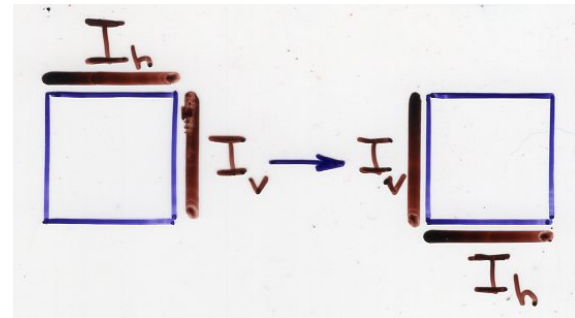
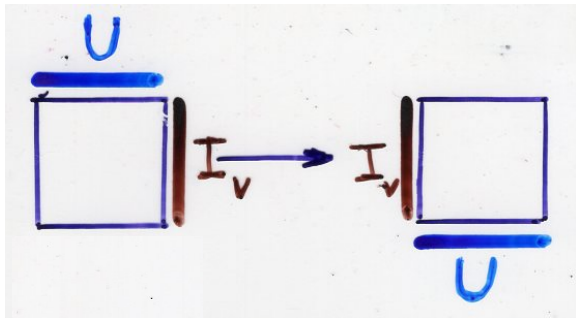
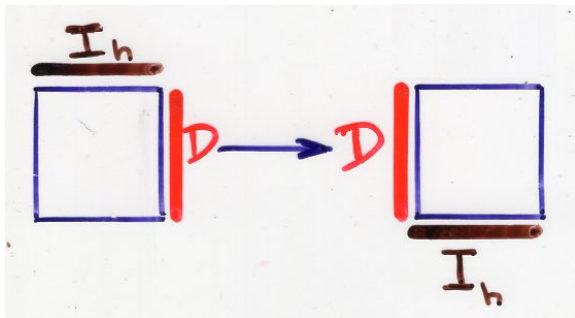
number of
inversions
of a permutation σ

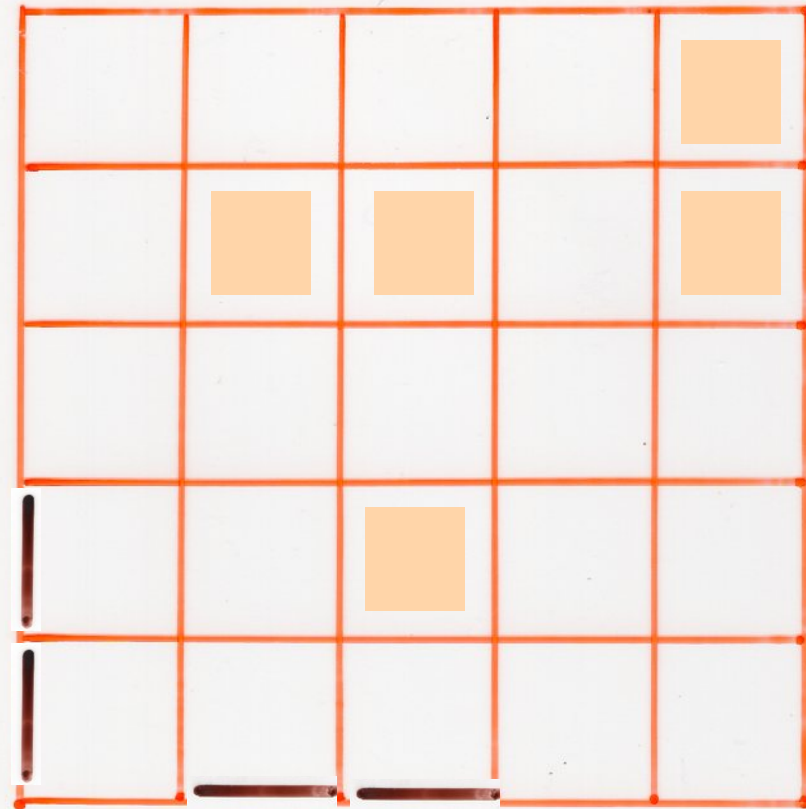
"complete"
Q-tableau

$$\left\{ \begin{array}{l} U \mathcal{D} = \mathcal{D} U + I_v I_h \\ U I_v = I_v U \\ I_h \mathcal{D} = \mathcal{D} I_h \\ I_h I_v = I_v I_h \end{array} \right.$$



Q-tableau





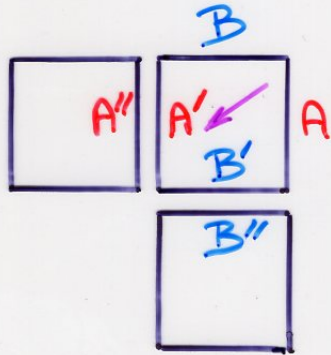
another Q-tableau
Rothe diagram
of a permutation

Definition $w(U, D)$ word of $\{U, D\}^*$

complete Q -tableau

- labeling of the cells of a Ferrers diagram $F = F(w)$ by the set R of rewriting rules
- with "compatibility" adjacent cells

i.e.



$A, A', B, B' \in \{U, D, I_v, I_h\}$
 A'', B''

$BA \rightarrow A'B'$

then $A' = A'', B' = B''$

- if the cell  is at the NE border of F , then $B = U, A = D$

$$\begin{array}{c} \varphi: R \longrightarrow L \\ \text{map} \end{array}$$

R = set of rewriting rules of the homogenous system associated to Q .

here 5 terms

L a set of "labels"
(for the cell of $[n] \times [n]$)

examples

$$L = \{ \square, \square \}$$

examples

$$L = \{ \square, \square \}$$

φ satisfies $(*)$:

$(*)$ if $\varphi(\alpha \rightarrow \beta) = \varphi(\alpha' \rightarrow \beta')$
then $\alpha \neq \alpha'$

i.e. in a single commutation equation

$$\alpha = \beta_1 + \dots + \beta_r$$

all elements $\varphi(\alpha \rightarrow \beta_i) \in L$ are \neq
set of labels

$$\left\{ \begin{array}{l} U \mathcal{D} = \mathcal{D} U + I_v I_h \\ U I_v = I_v U \\ I_h \mathcal{D} = \mathcal{D} I_h \\ I_h I_v = I_v I_h \end{array} \right.$$

if $\varphi: R \rightarrow L$
satisfies $(*)$
then

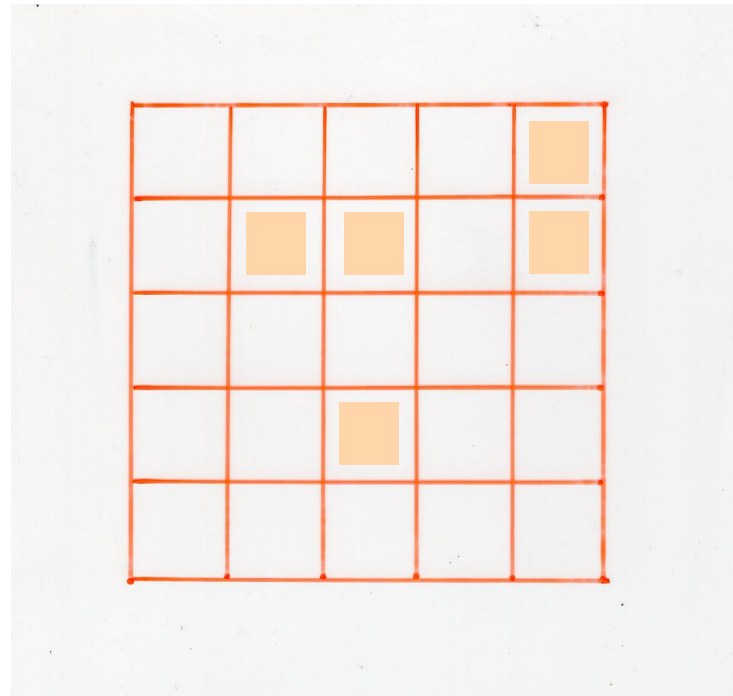
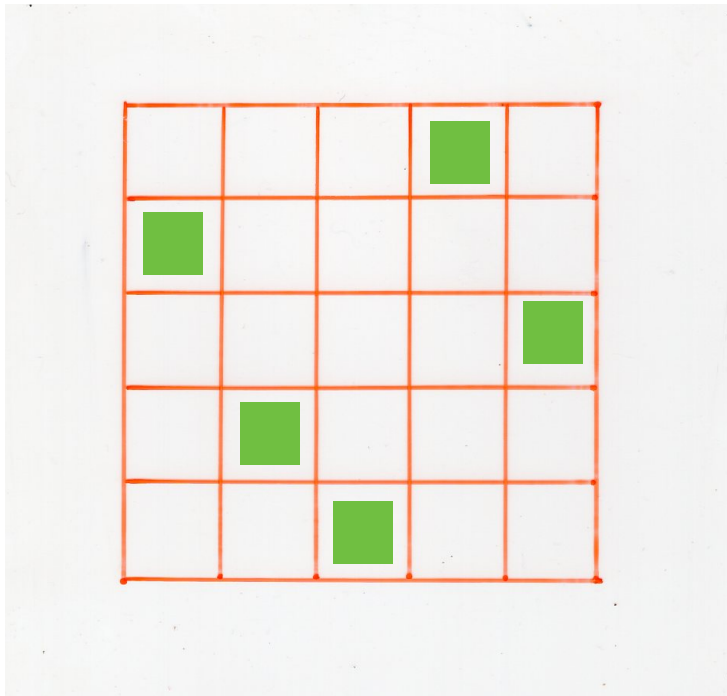
complete Q -tableaux $\xleftrightarrow{\text{bijection}}$ Q -tableaux

$$\varphi: \begin{array}{l} UD \rightarrow DU, UI_v \rightarrow I_v U, \\ I_h D \rightarrow DI_h, I_h I_v \rightarrow I_v I_h \end{array} \rightarrow \begin{array}{c} \square \\ \text{empty} \\ \text{cell} \end{array}$$

$$\varphi: \begin{array}{l} UD \rightarrow I_v I_h, UI_v \rightarrow I_v U, \\ I_h D \rightarrow DI_h, I_h I_v \rightarrow I_v I_h \end{array} \rightarrow \begin{array}{c} \square \\ \text{empty} \\ \text{cell} \end{array}$$

$$\varphi(UD \rightarrow I_v I_h) = \begin{array}{c} \square \\ \text{green} \end{array}$$

$$\varphi(UD \rightarrow DU) = \begin{array}{c} \square \\ \text{orange} \end{array}$$



"The cellular ansatz"

quadratic algebra Q

Q -tableaux

combinatorial objects
on a 2D lattice

$$UD = DU + Id$$

permutations

towers placements

(i) first step

commutations

rewriting rules

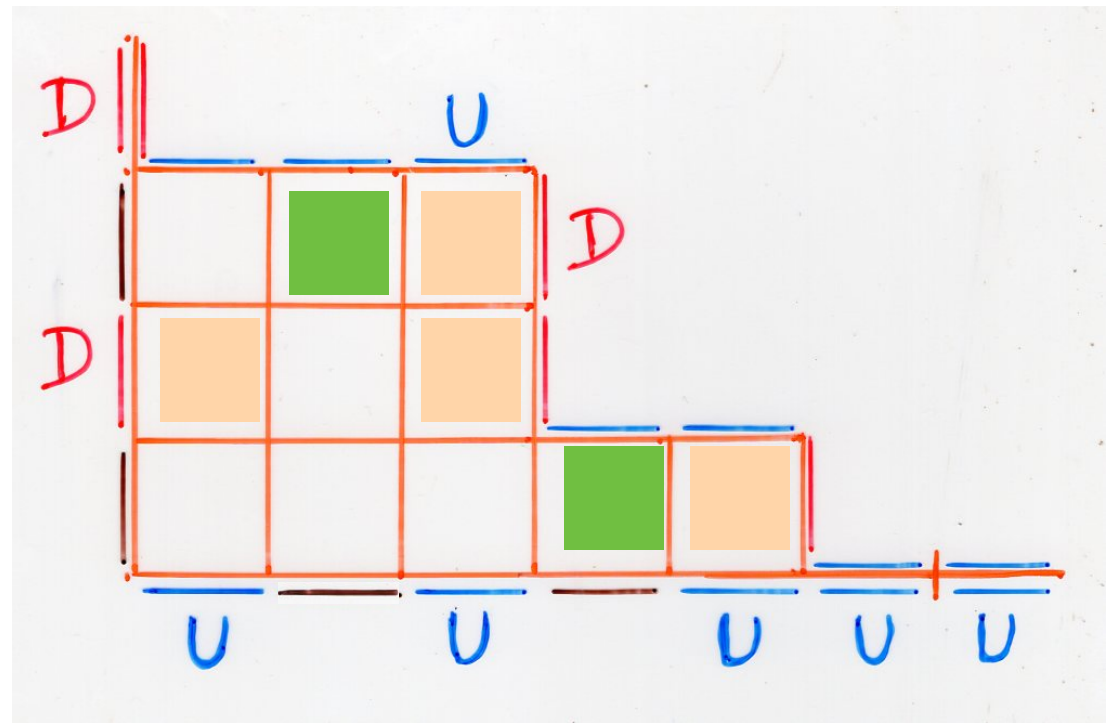
planarization

$$w = DU^3D^2U^2DU^2$$

$$w \rightarrow F = F(w)$$

F Ferrers diagram

Rooks placement



The cellular ansatz
second part:

guided construction
of a bijection
from the representation of U and D

$$UD = DU + I$$

"The cellular ansatz"

quadratic algebra Q

Q -tableaux

representation of Q
by combinatorial operators

$$UD = DU + Id$$

combinatorial objects
on a 2D lattice

bijections

permutations

RSK

pairs of
Young tableaux

towers placements



(i) first step

(ii) second step

commutations

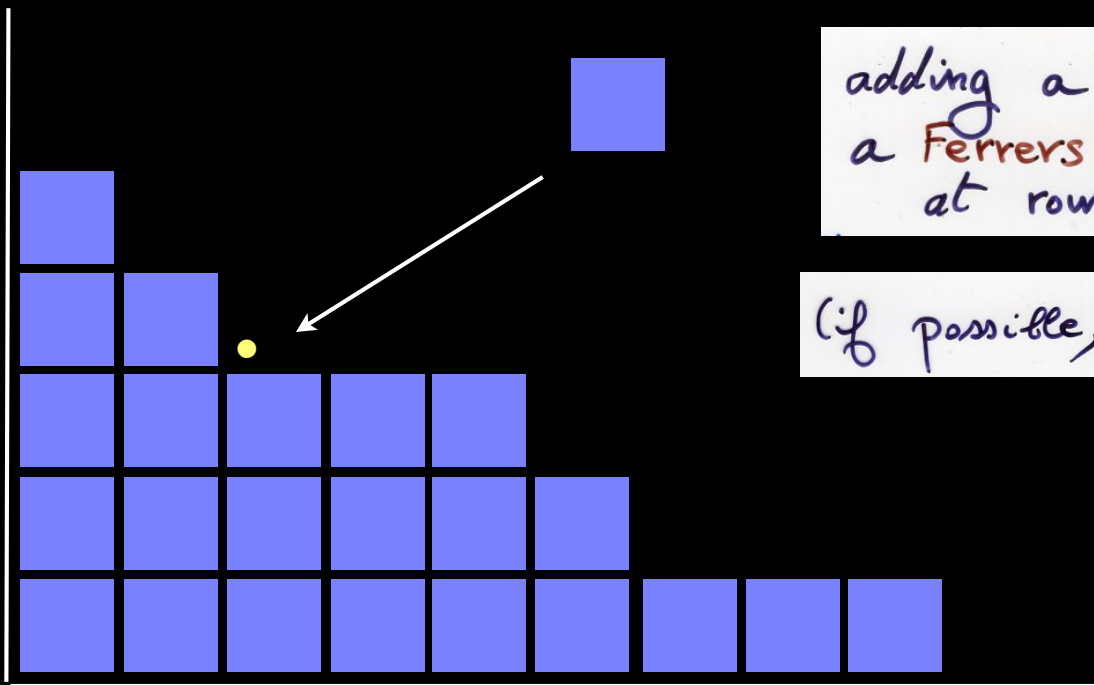
rewriting rules

planarization

notations

operator U_i

i



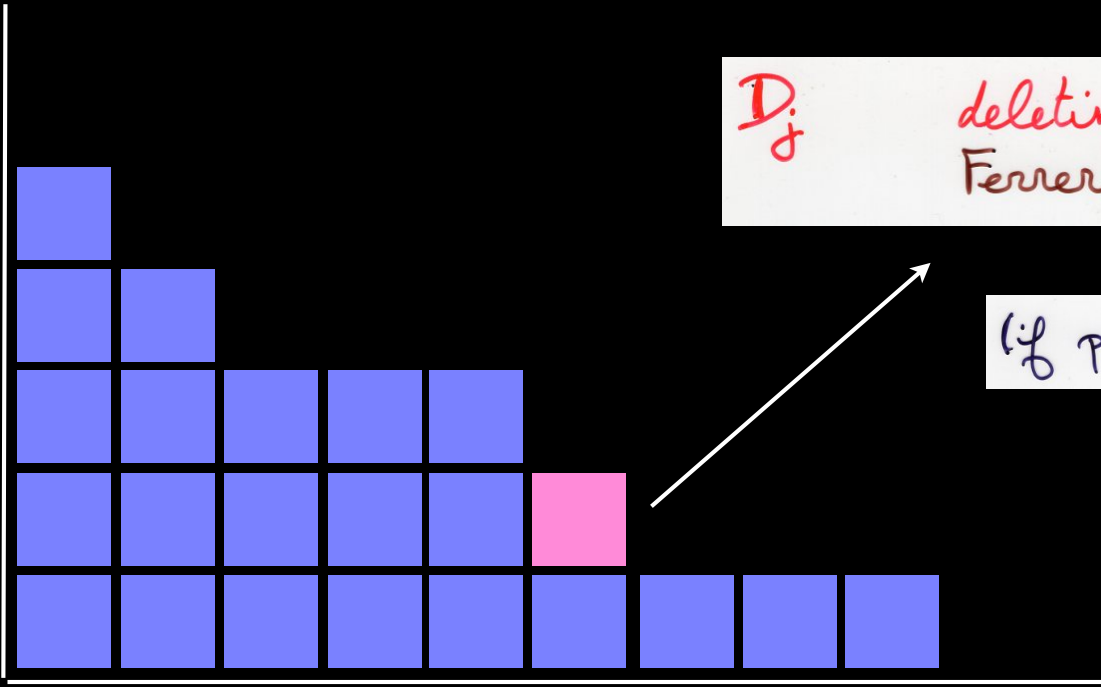
adding a cell in
a Ferrers diagram ρ
at row i

(if possible, else $U_i(\rho) = 0$)

$$U_i(\rho) = \rho + (i)$$

$$D_j(\rho) = \rho - (j)$$

j



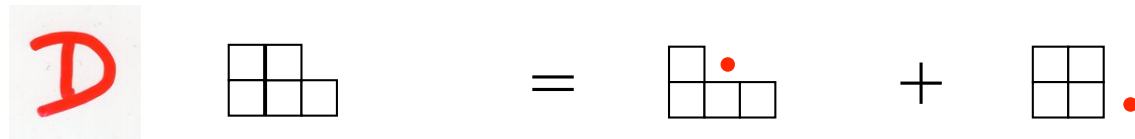
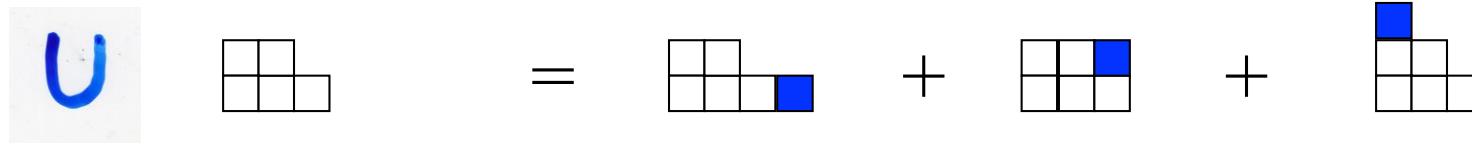
D_j deleting a cell in a Ferrers diagram ρ at row j

(if possible, else $D_j(\rho) = 0$)

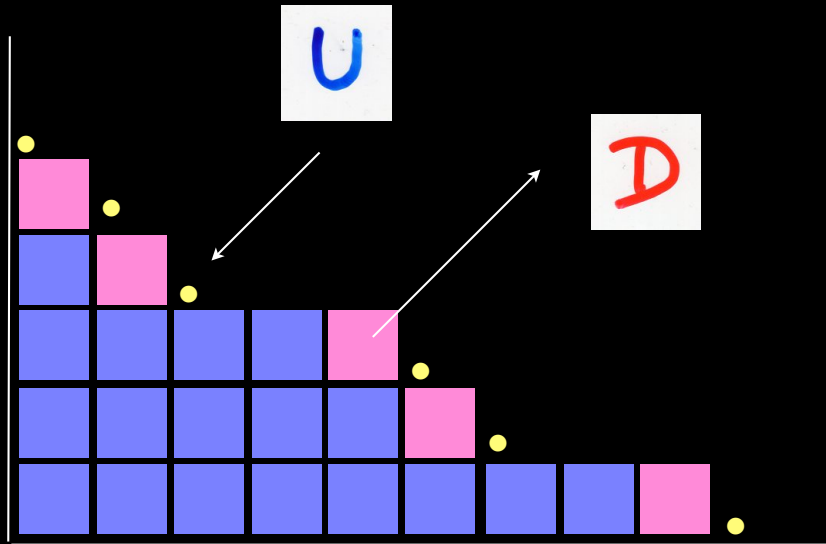
$$U = \sum_{i \geq 1} U_i$$

$$D = \sum_{i \geq 1} D_i$$

U and D are operators acting on the vector space generated by Ferrers diagrams



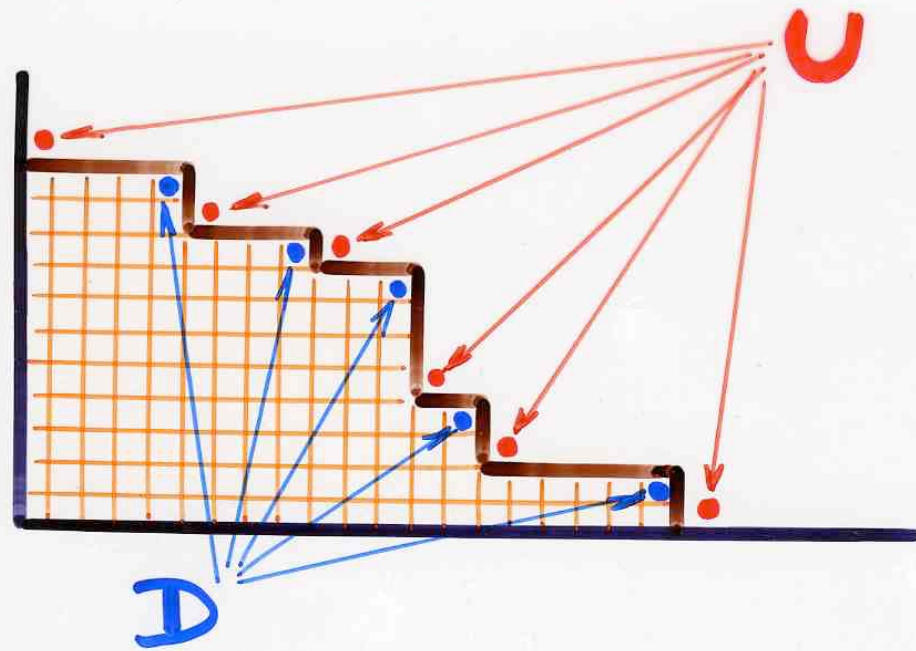
operators
U and D



Young lattice

{ U adding a cell in a Ferrers diagram
D deleting

$$UD = DU + I$$



In this course, product of operators are written from left to right

$$\underline{A} B (\mu) = B (A (\mu))$$

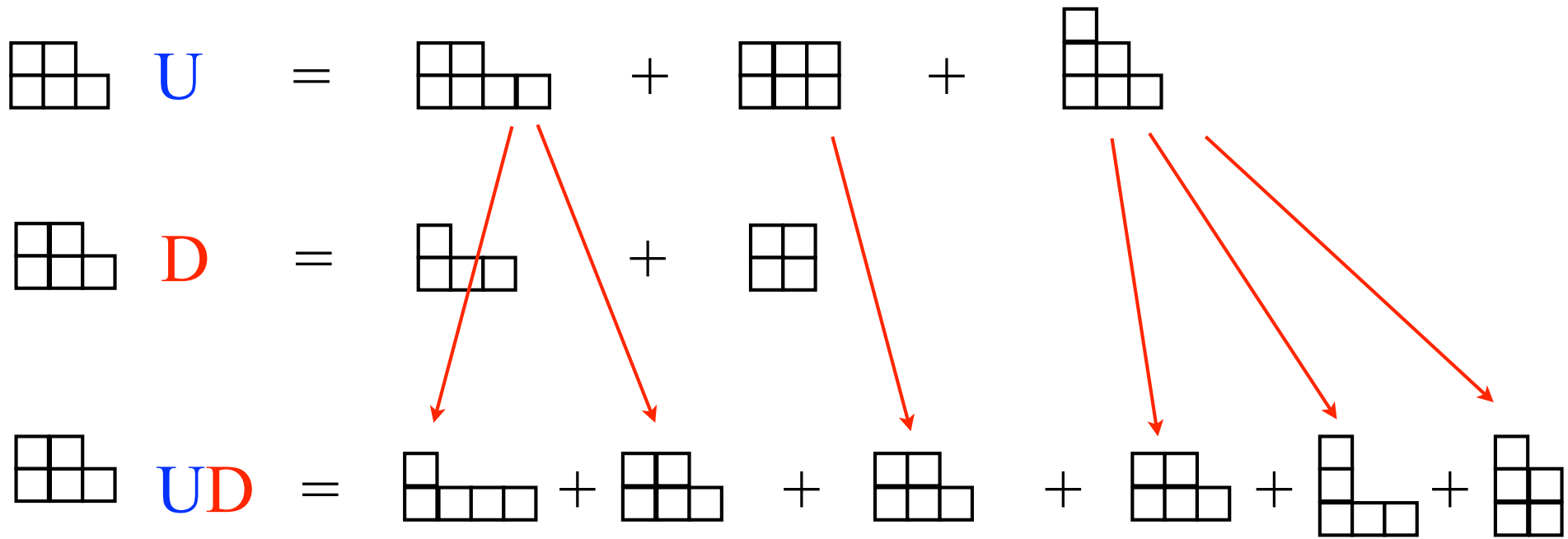
should be written $(\mu) A B$
or $\langle \mu | A B$

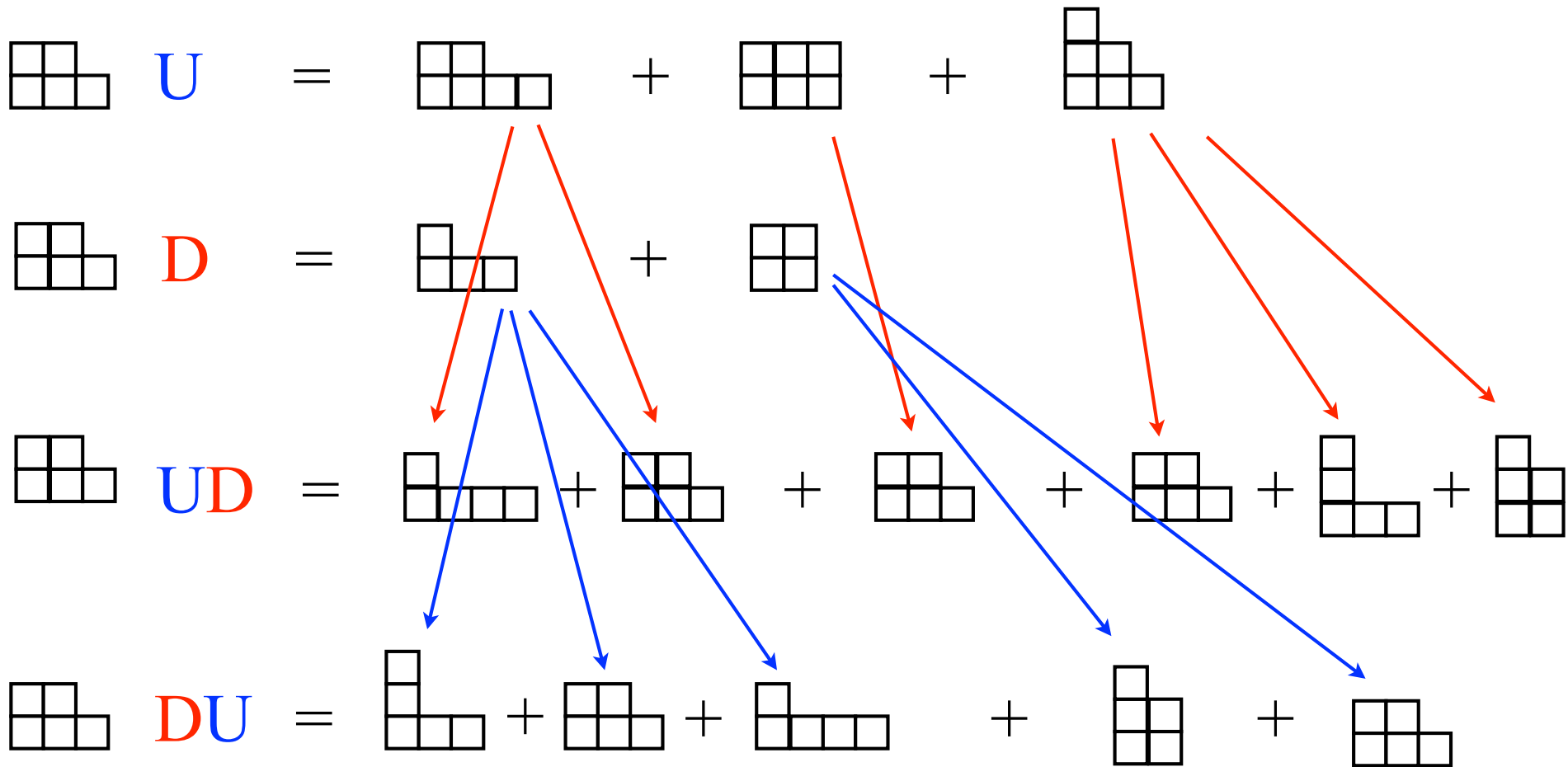
with operators written from right to left

$$B A (\mu) = B (A (\mu))$$

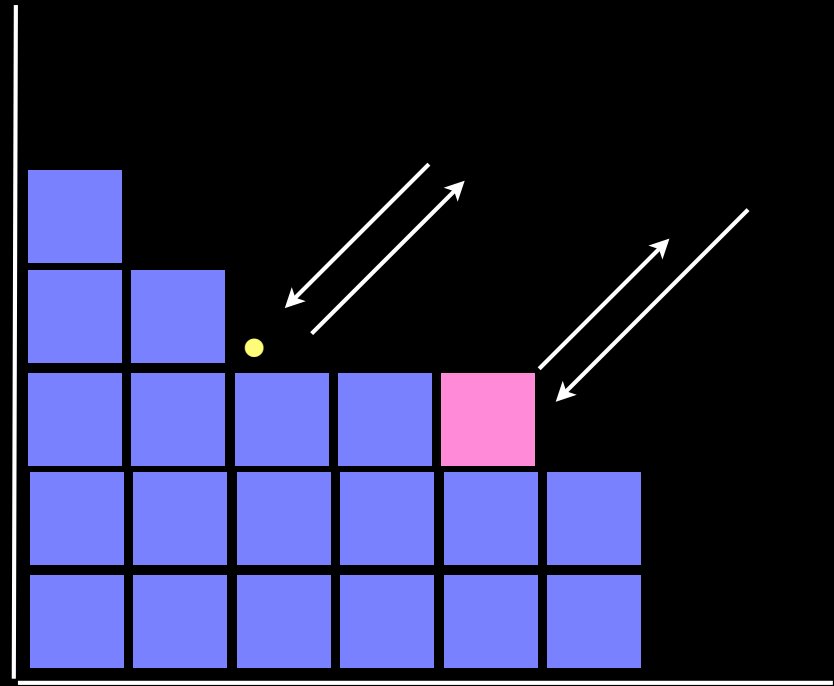
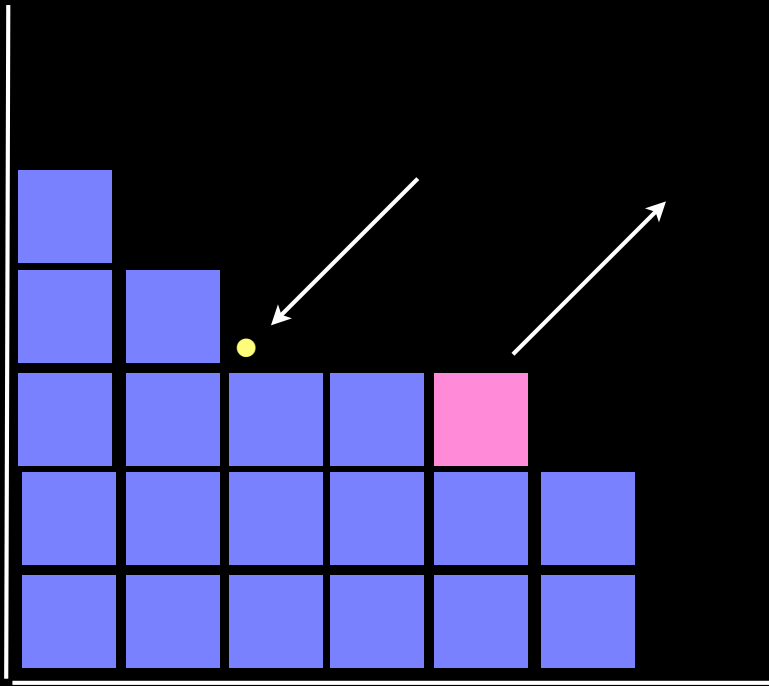
$$U D - D U = I \quad \text{becomes}$$

$$D U - U D = I$$

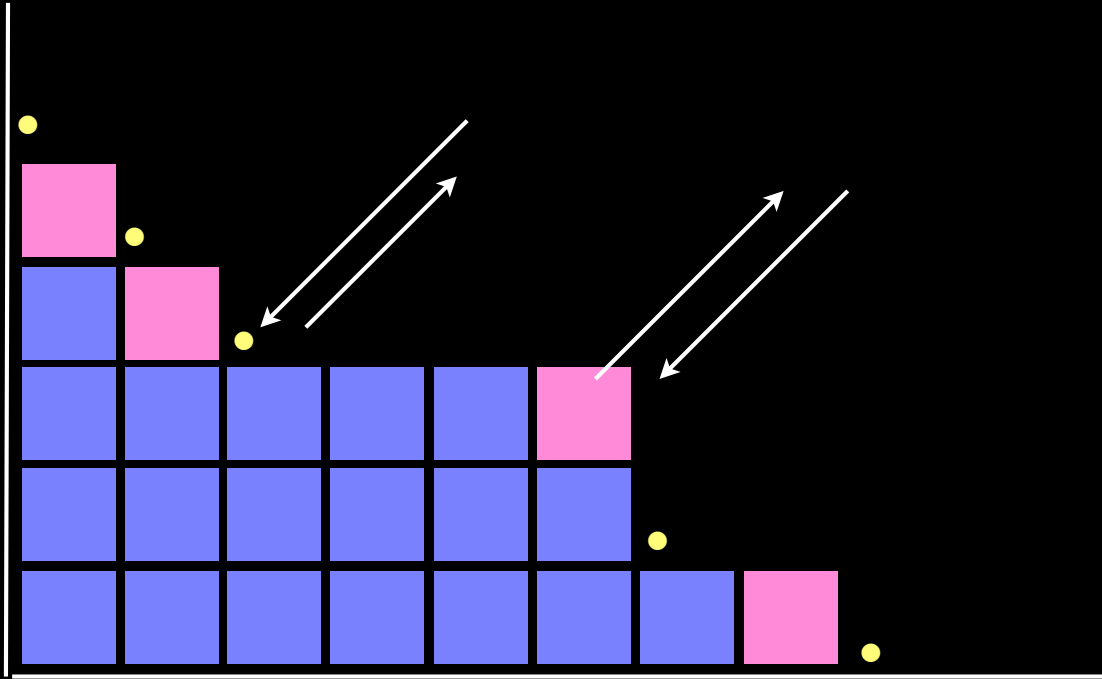




$$UD = DU + I$$

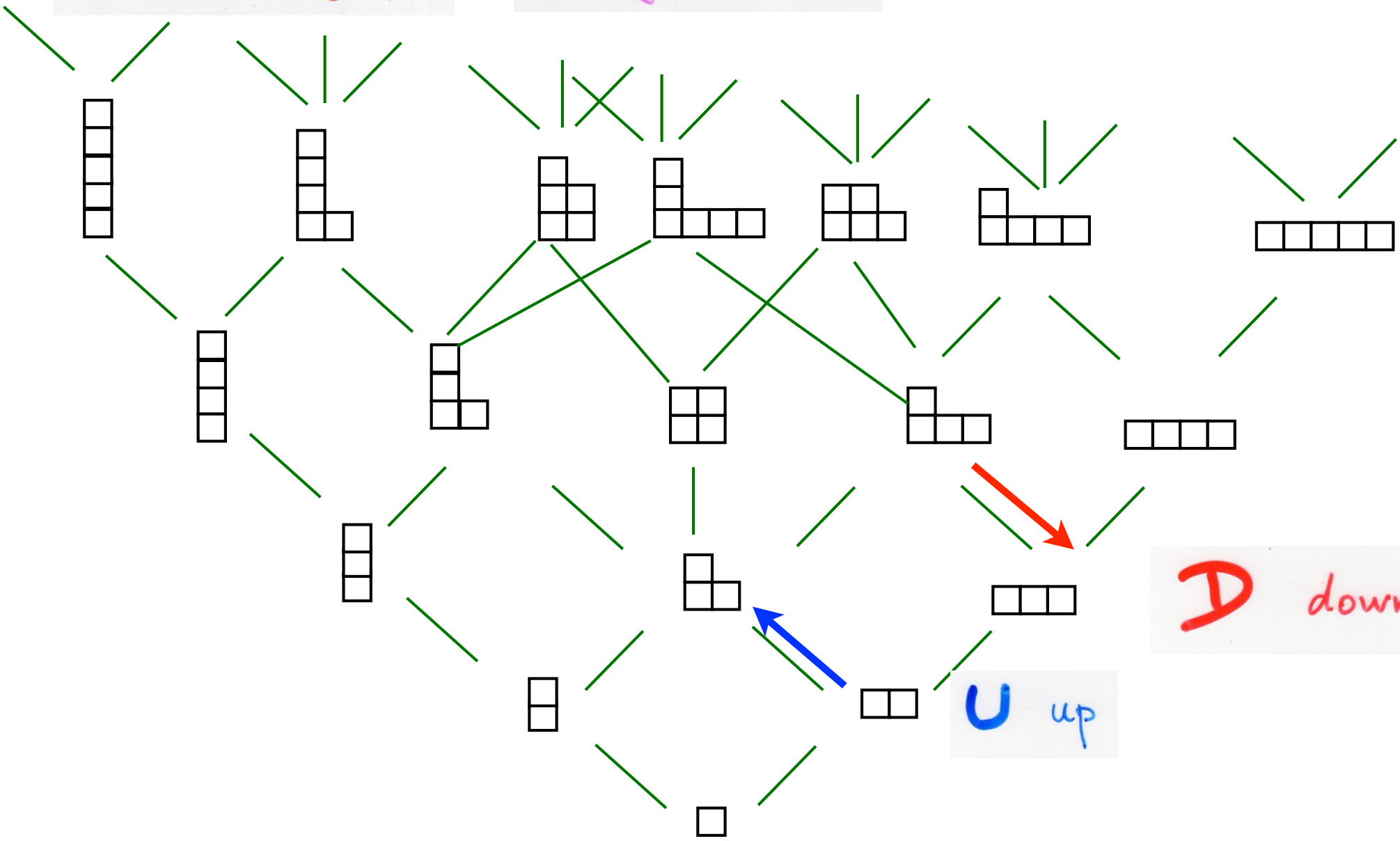


$$UD = DU + I$$



Hasse diagram

Young lattice



D down

U up

differential poset

Fomin (1992, 1995)

Stanley (1988, 1990)

Roby (1991)

$$UD = DU + I$$

U up

D down

direct proof of the identity

permutations pairs of Young tableaux,
same shape

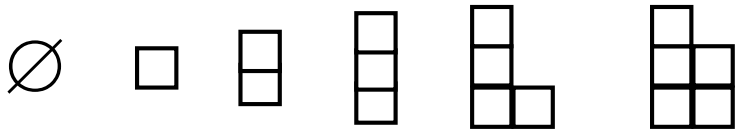
$$n! = \sum_{\lambda} (\mathfrak{f}_{\lambda})^2$$

partition
of n

$$UD = DU + I$$

$$\langle \emptyset | U^n D^n | \emptyset \rangle$$

\emptyset empty Ferrers diagram



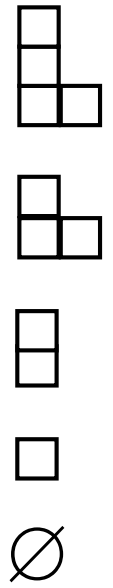
$$= \sum_{\lambda} \binom{n}{\lambda}^2$$

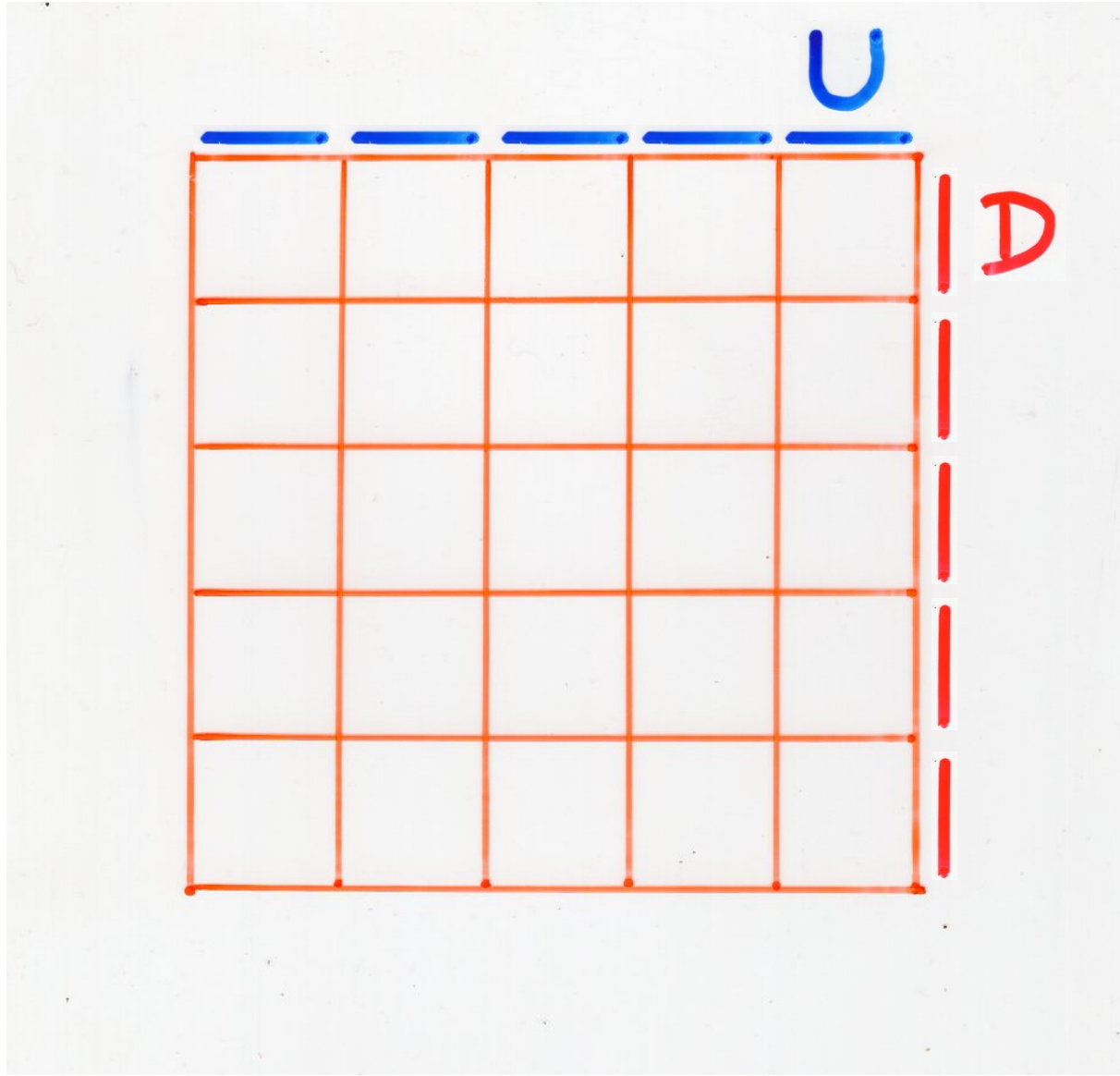
partition of n

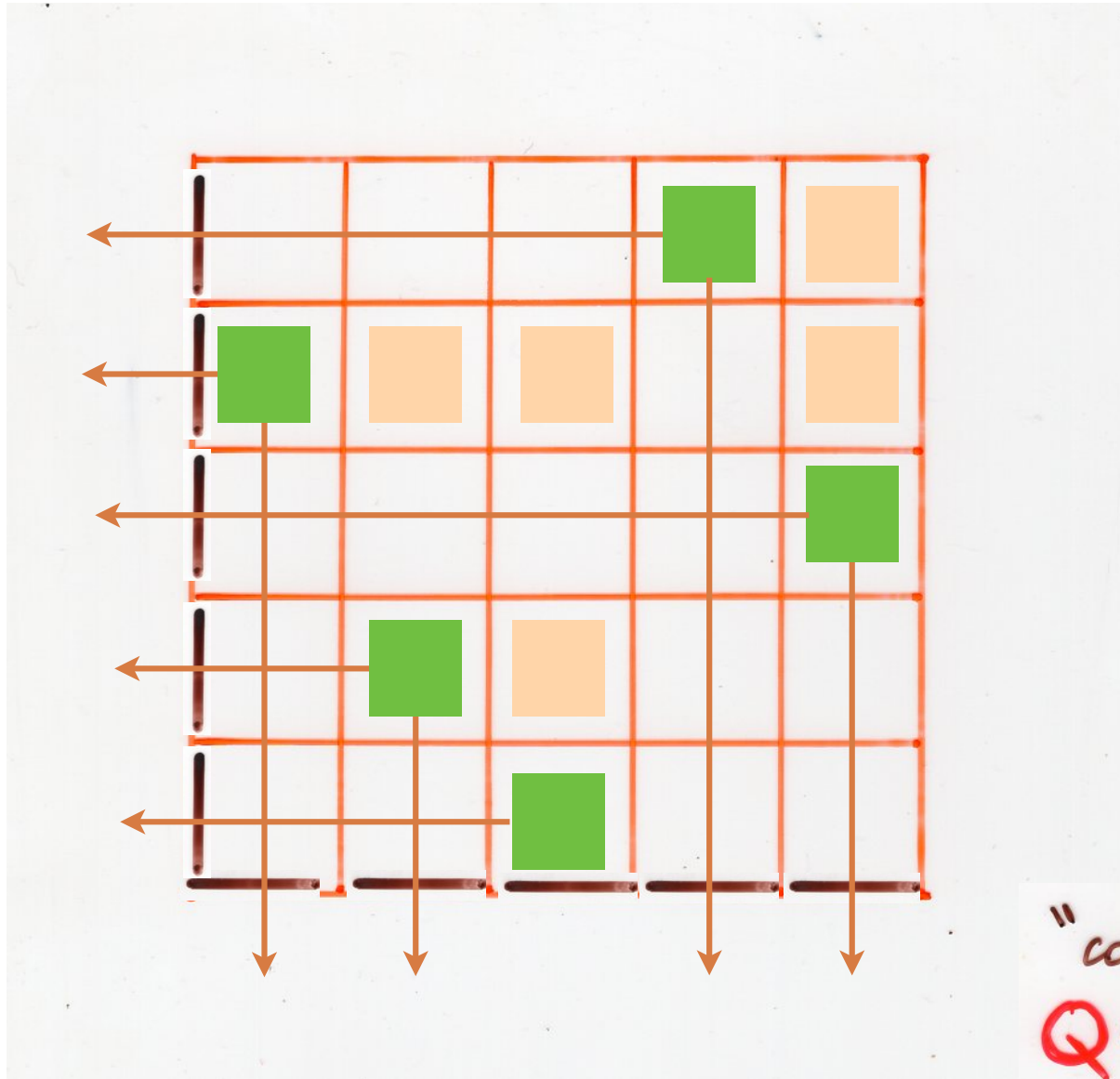
$$= \sum_{i \geq 0} c_{n,i} \langle \emptyset | D^i U^i | \emptyset \rangle$$

$$= c_{n,0}$$

$$= n!$$





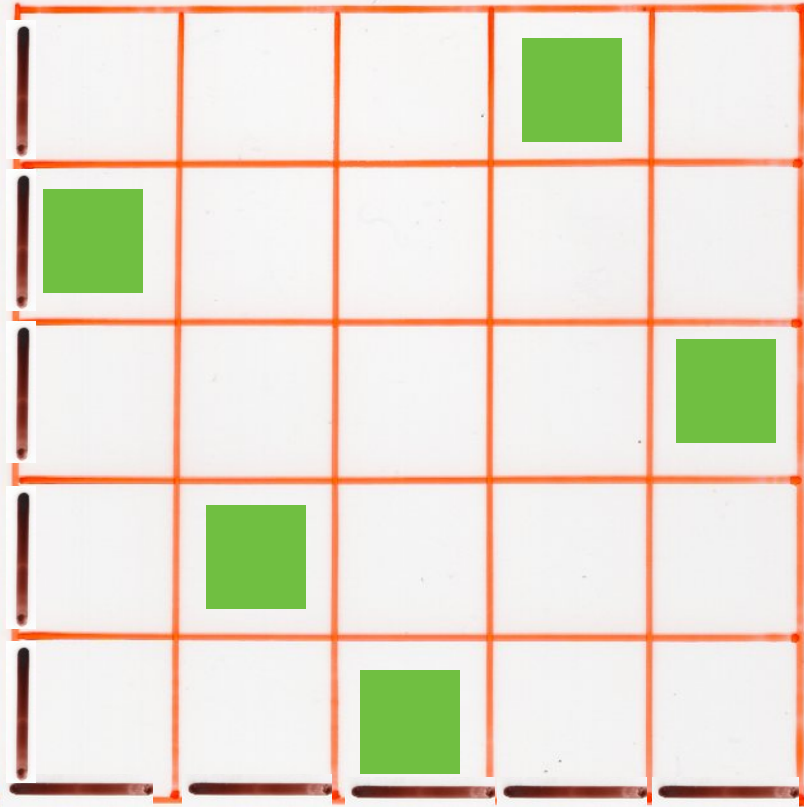


"complete"
Q-tableau

$$c_{n,0} = n!$$

$$= \sum_{\lambda} \left(f_{\lambda} \right)^2$$

partition
of n



permutation
as a Q -tableau

construction of the RSK correspondence
by «propagation» on the grid
of an elementary «diagram bijection»
related to each cell of the grid

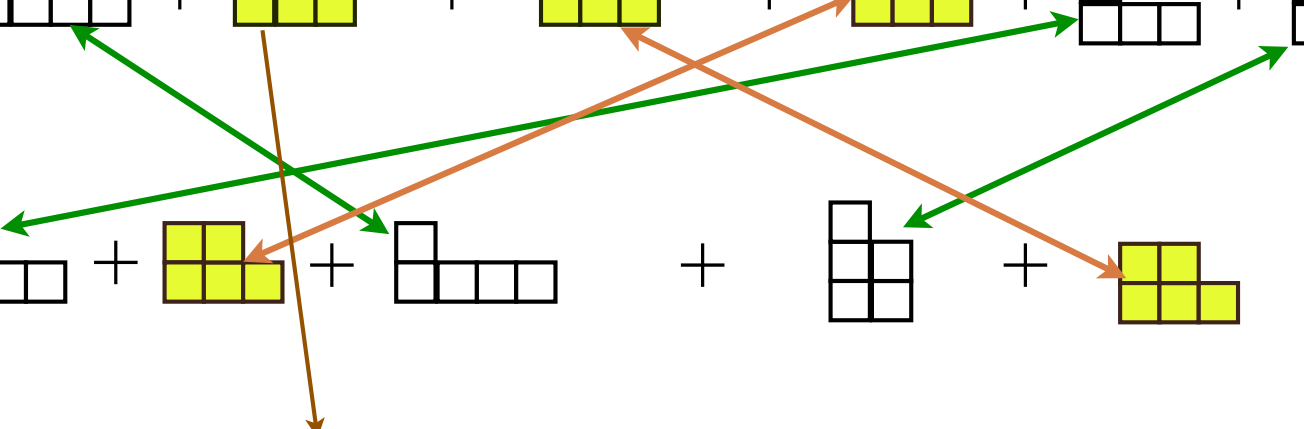
$$\begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 \quad \mathbf{U}
 =
 \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}$$

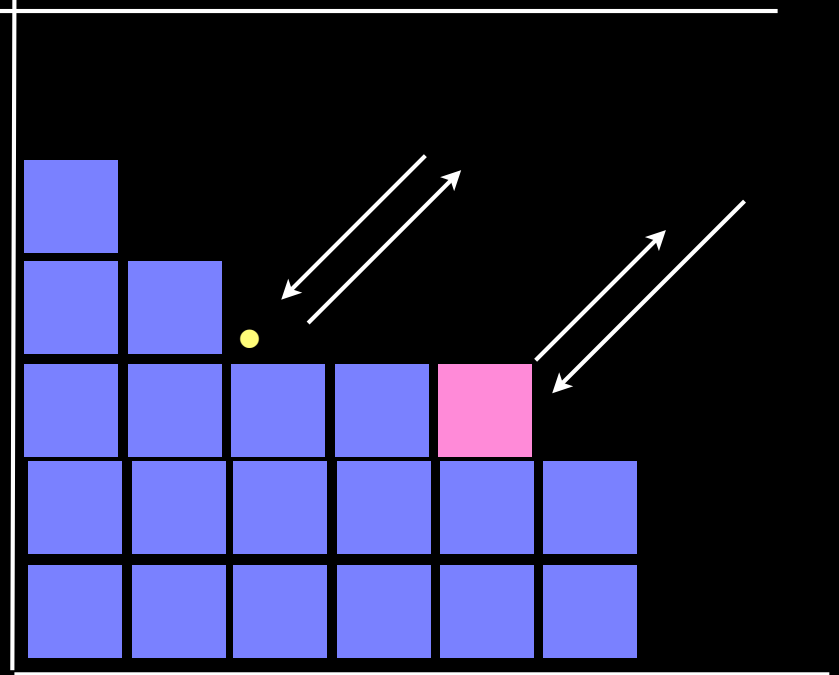
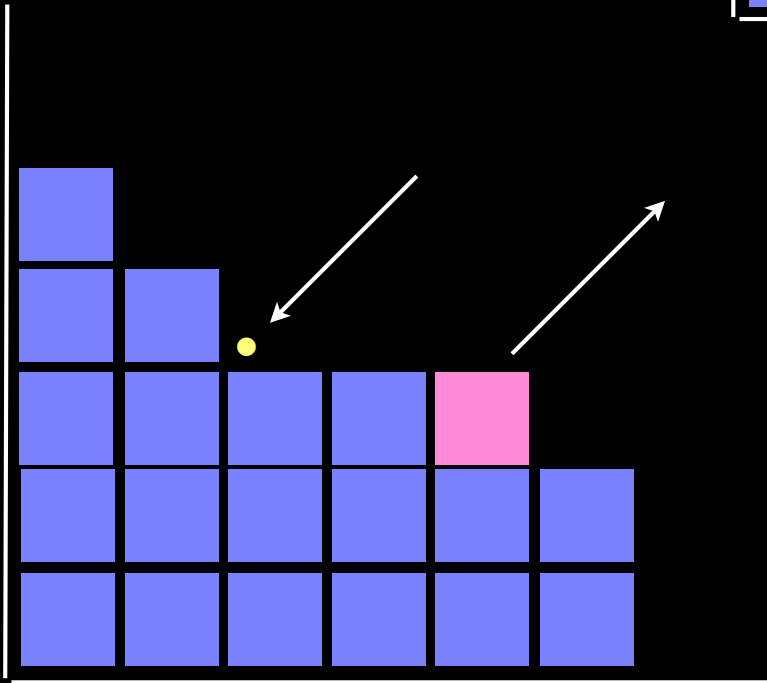
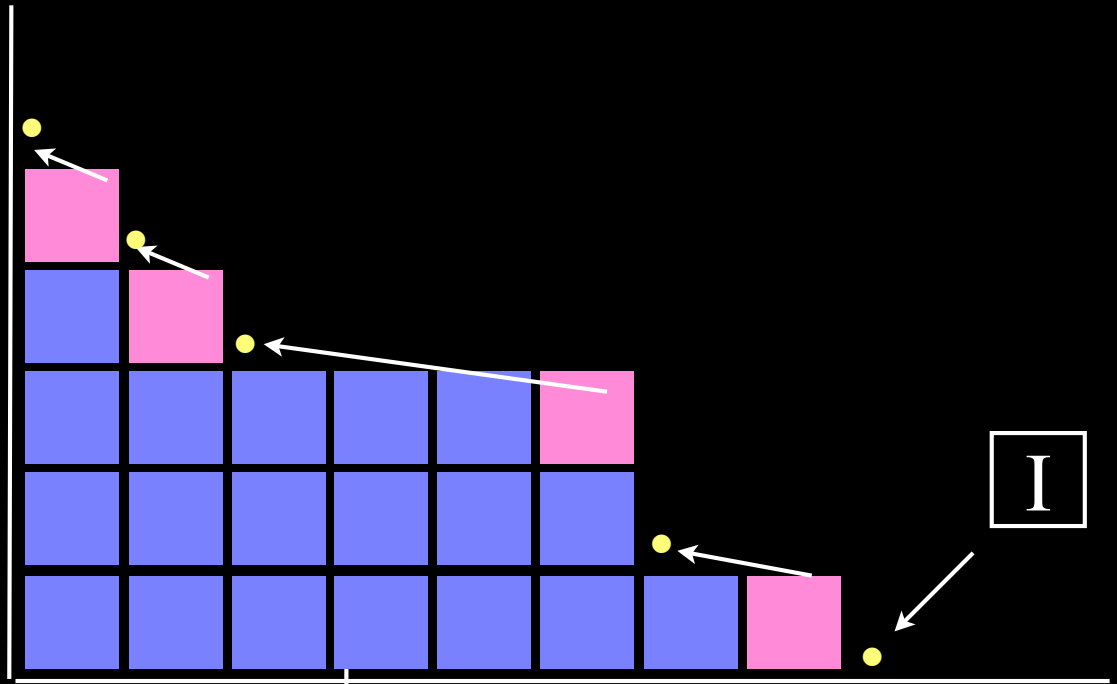
$$\begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 \quad \mathbf{D}
 =
 \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}$$

$$\begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 \quad \mathbf{UD}
 =
 \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \square & \square & \square & \square \\ \hline \square & \square & \square & \square \\ \hline \square & \square & \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}$$

$$\begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 \quad \mathbf{DU}
 =
 \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \square & \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 +
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}$$

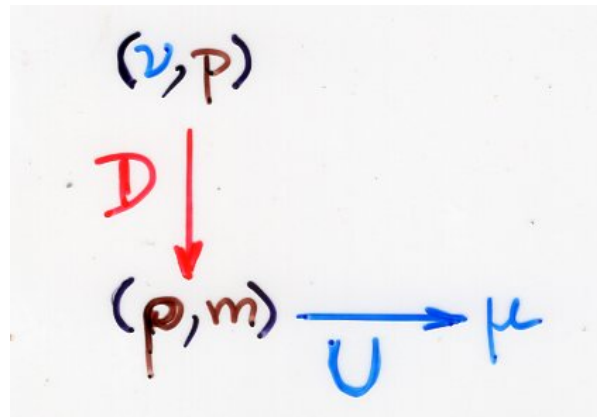
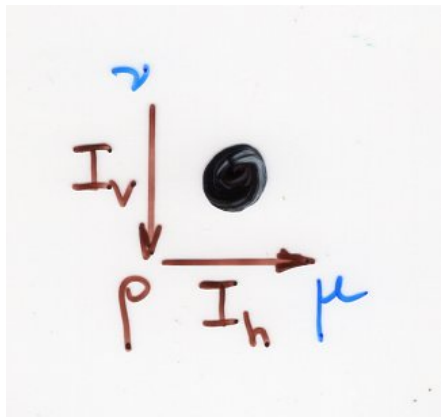
$$\begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}
 \quad (\mathbf{UD-DU})
 =
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}$$



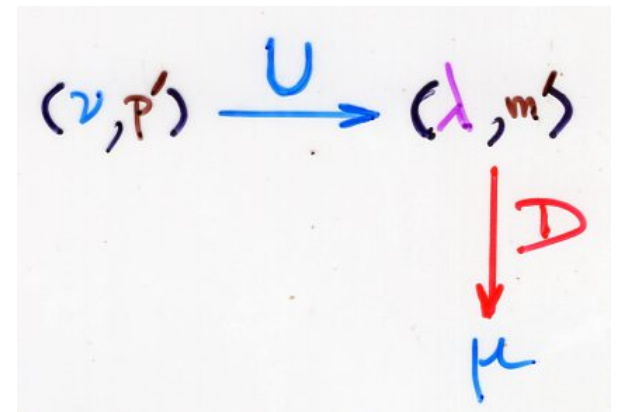


$$UD = DU + I_v I_h$$

"commutation diagrams"



bijection



p, m, p', m' are "positions"

in ν, ρ, ν, λ respectively

$$(v, p') \xrightarrow{U} (\lambda, m')$$

$$(v, p)$$

 $D \downarrow$

$$(p, m)$$

$$\xrightarrow{U} \mu$$

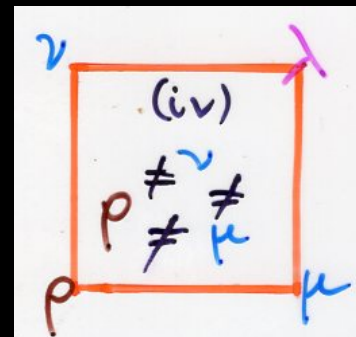
 $D \downarrow$
 μ


$$p = j$$

$$m = i$$

$$p' = i$$

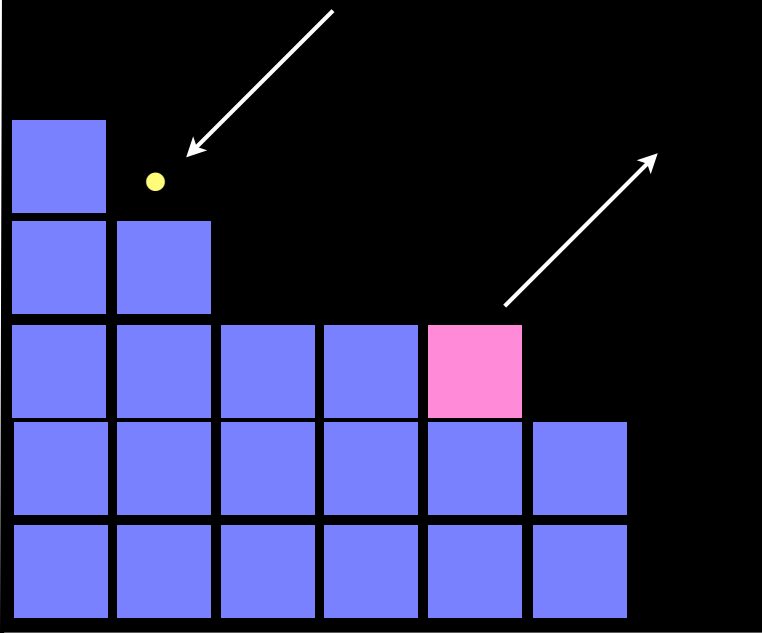
$$m' = j$$

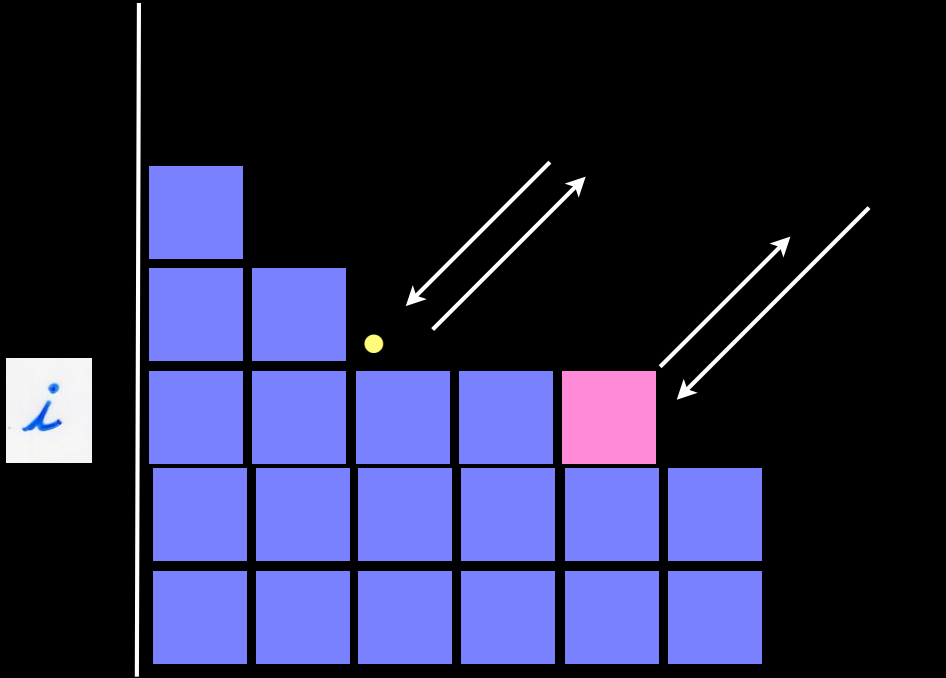
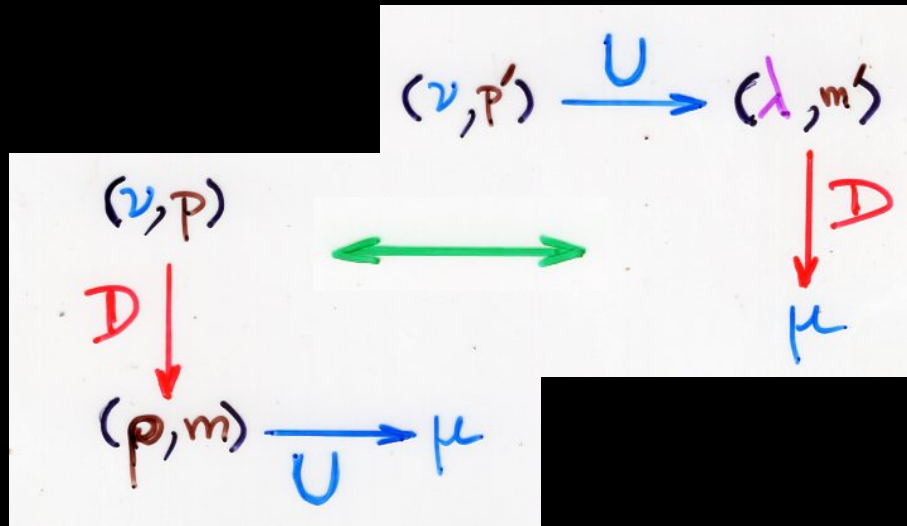


$$v = p + (j)$$

$$\mu = p + (i)$$

$$\lambda = p + (i) + (j)$$

 i
 j


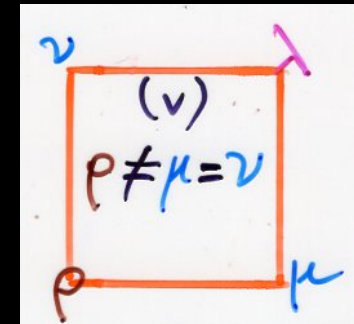


$$p = i$$

$$m = i$$

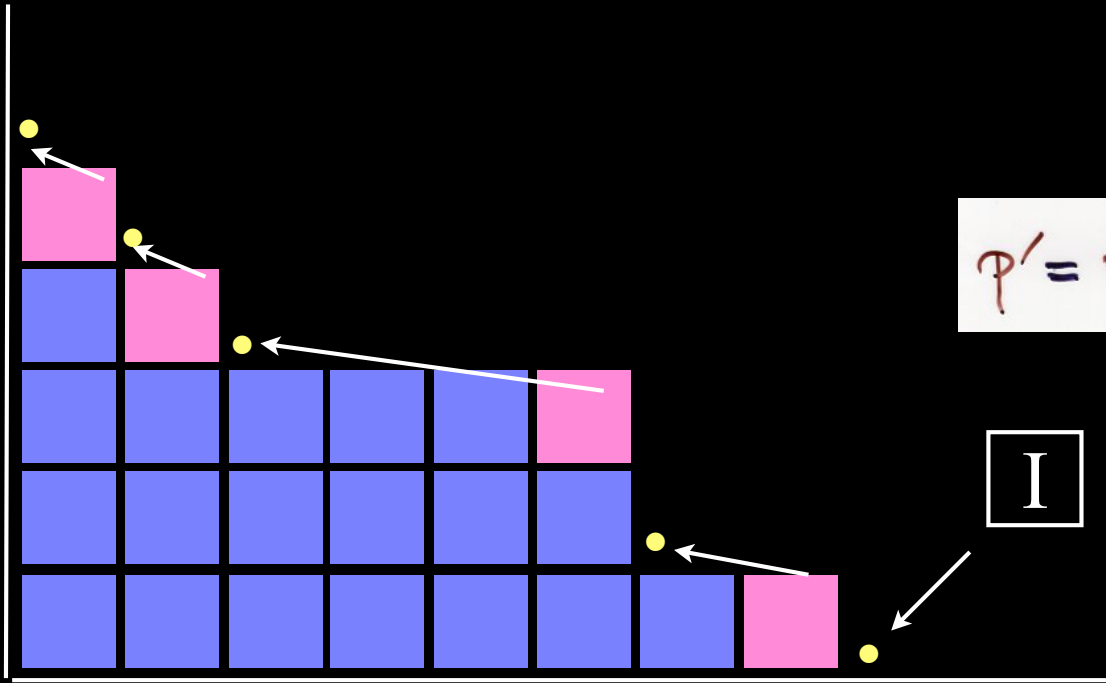
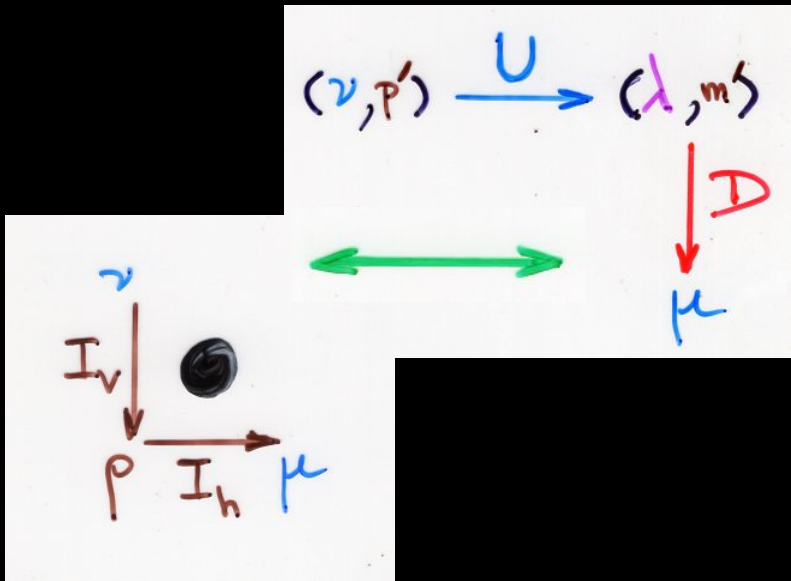
$$p' = i+1$$

$$m' = i+1$$

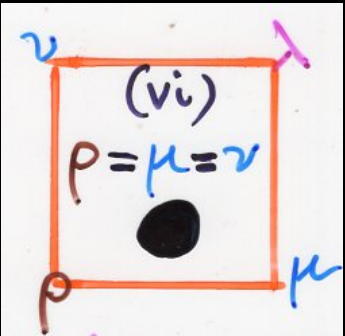


$$\mu = \nu = \rho + (i)$$

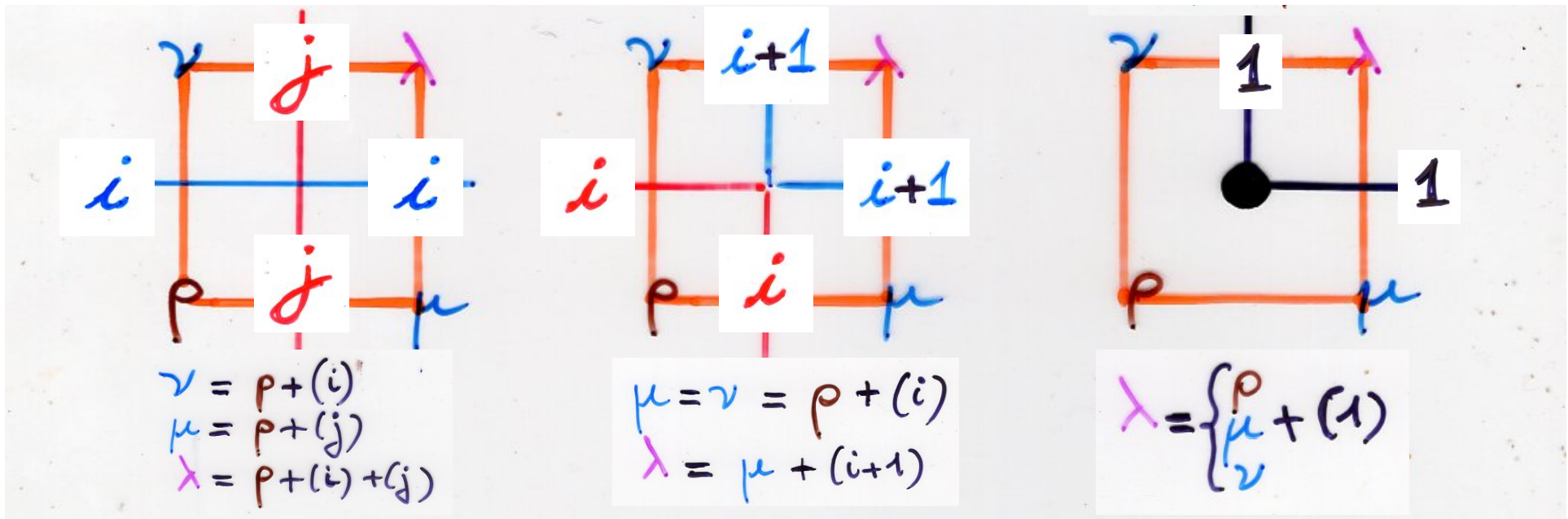
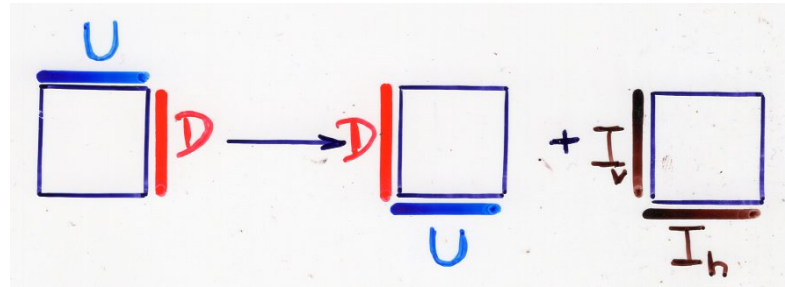
$$\lambda = \mu + (i+1)$$



$$\rho' = m' = 1$$



$$\lambda = \begin{cases} \rho \\ \mu + (1) \\ \nu \end{cases}$$



$$\begin{aligned} \nu &= \rho + (i) \\ \mu &= \rho + (j) \\ \lambda &= \rho + (i) + (j) \end{aligned}$$

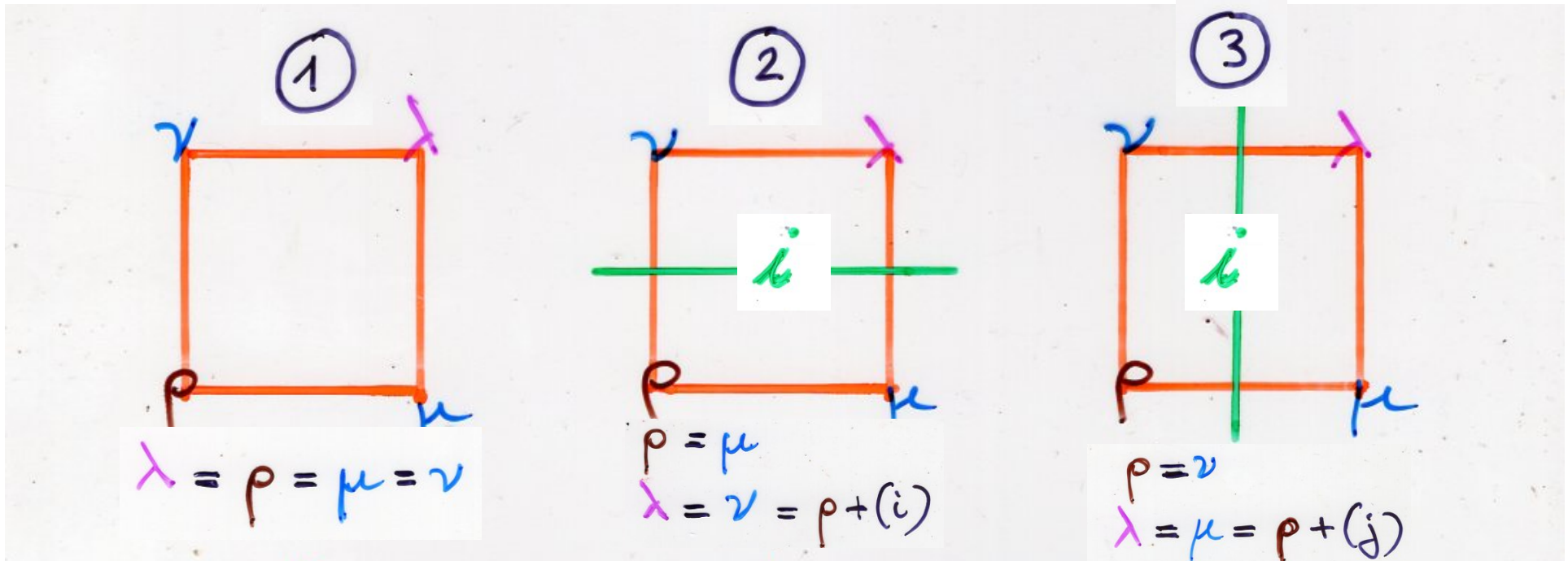
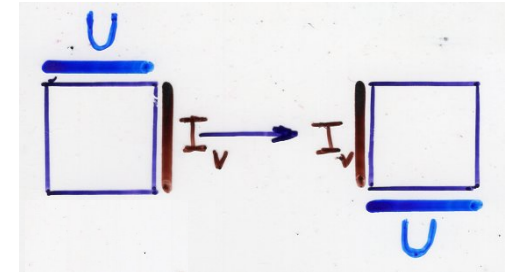
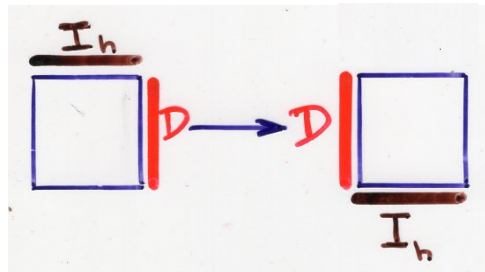
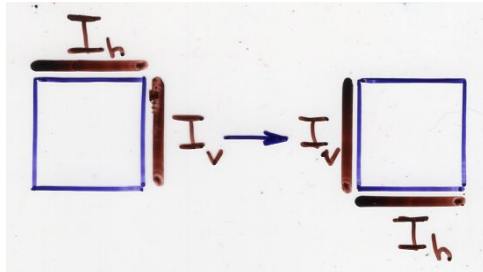
$$\begin{aligned} \mu &= \nu = \rho + (i) \\ \lambda &= \mu + (i+1) \end{aligned}$$

$$\lambda = \begin{cases} \rho \\ \mu + (1) \\ \nu \end{cases}$$

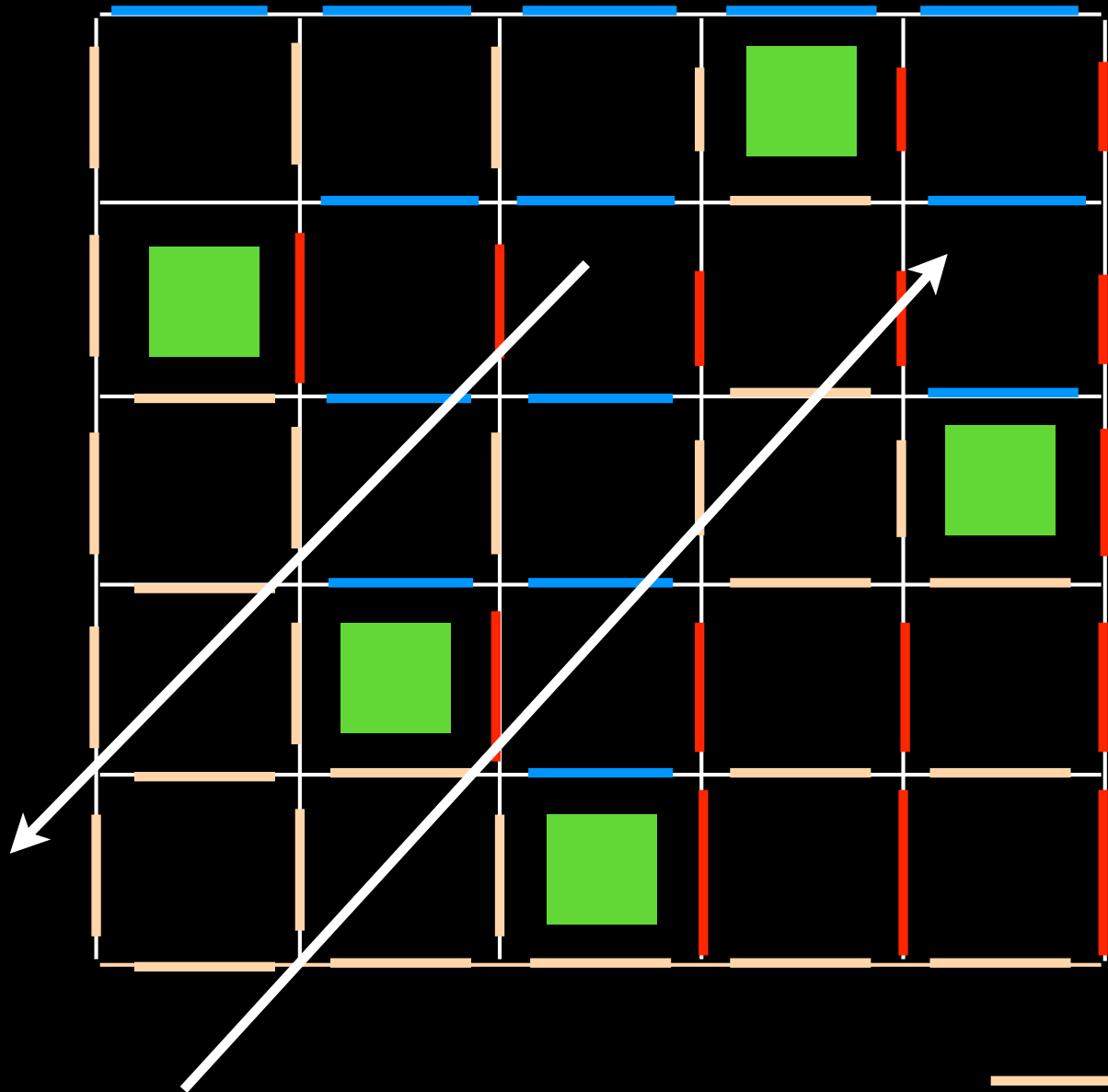
④

⑤

⑥



I



U

D

propagation
of the
diagrams
bijection
related
to one cell

formalisation



I

For $w = w_1 \dots w_n$ word of $\{U, D, I_h, I_v\}^*$
we consider sequences h

$$h = ((\mu_1, p_1), \dots, (\mu_n, p_n), \mu_{n+1})$$

where $\mu_i, i=1, \dots, n+1$ are partitions (Ferrers diagrams)

and for $i=1, \dots, n$ μ_{i+1} is obtained from μ_i
by applying the operator w_i at
position p_i

$$w = w(h)$$

if $w_i = I_h$ or I_v , then $(\mu_{i+1}, p_{i+1}) = (\mu_i, p_i)$

h "histories"

admissible sequence

2-colored vacillating tableaux

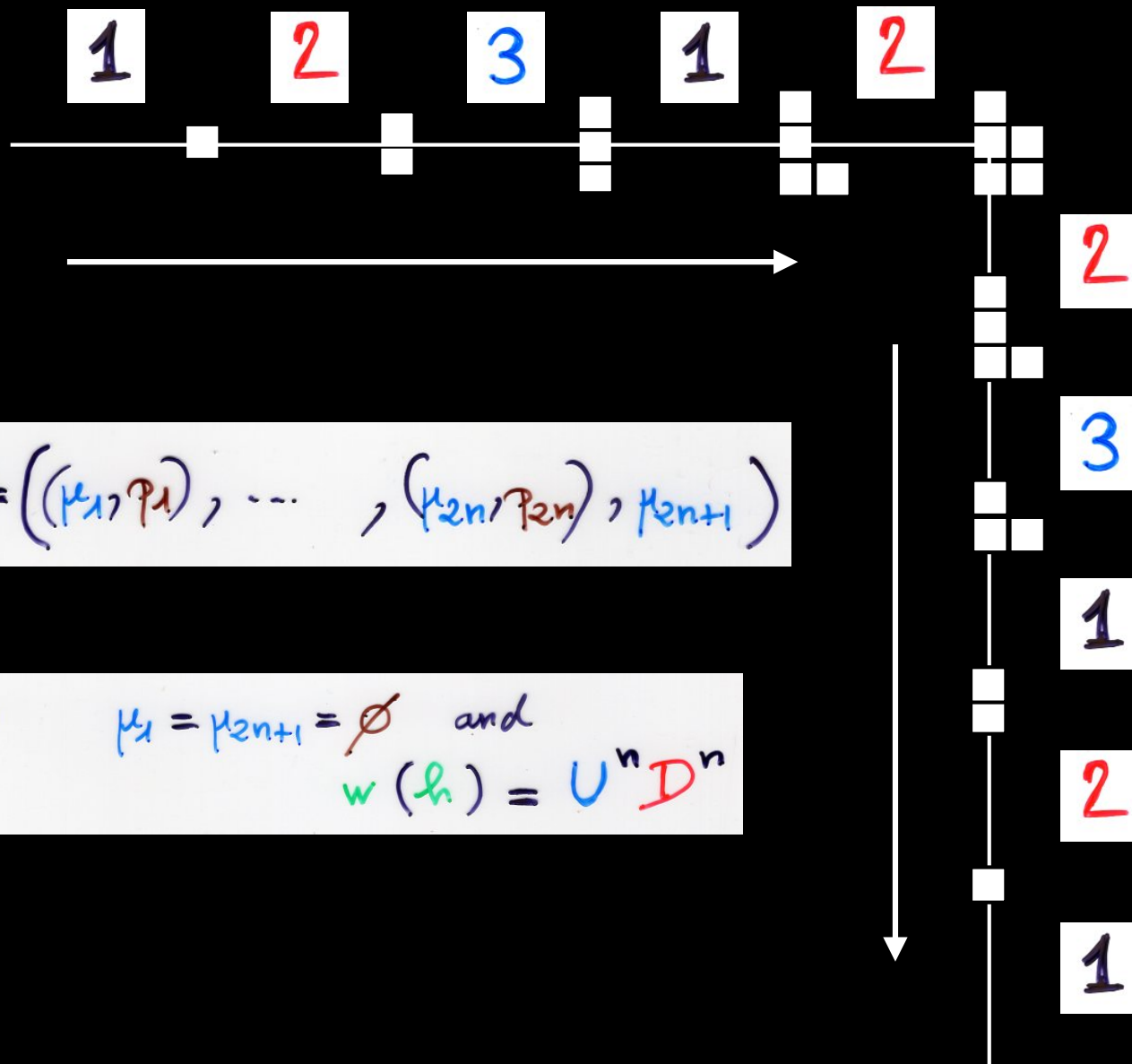
$(P, Q) \rightarrow (\alpha, \beta) \rightarrow \text{sequence } h$
 Young tableaux same shape λ pair of maximal chains
 $\emptyset \rightarrow \lambda$

$$h = ((\mu_1, \rho_1), \dots, (\mu_n, \rho_n), \mu_{n+1})$$

with $\mu_1 = \mu_{n+1} = \emptyset$ and

$$w(h) = U^n D^n$$

3	
2	5
1	4



$$h = ((\mu_1, p_1), \dots, (\mu_n, p_n), \mu_{n+1})$$

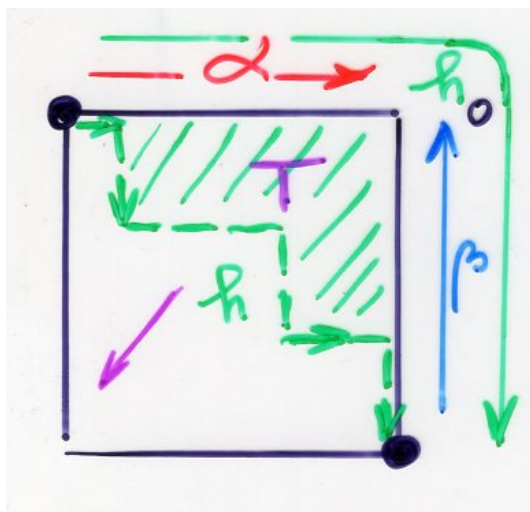
with $\mu_1 = \mu_{n+1} = \emptyset$ and $w(h) = U^n D^n$

4	
2	5
1	3

Starting from $h_0(\alpha, \beta) = h_0(P, Q)$ $T = \emptyset$

we "propagate" the "commutation diagrams" through the lattice $[n] \times [n]$.

At any step, we have a pair



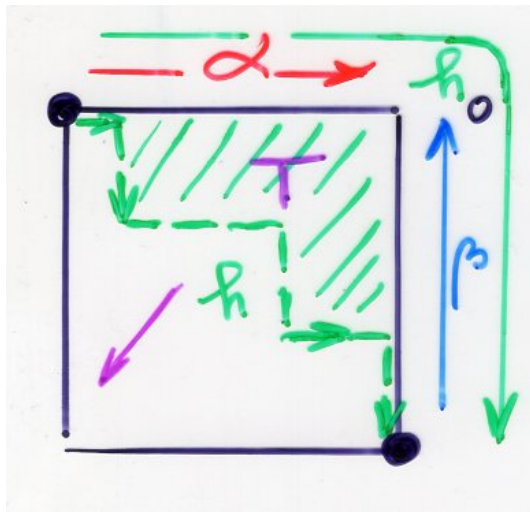
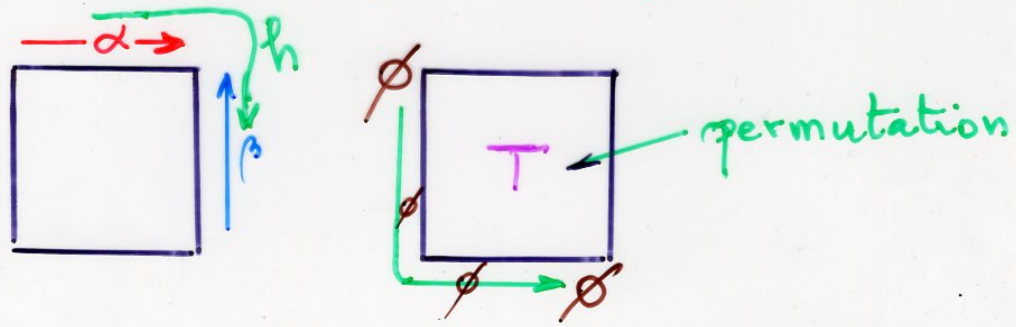
(h, T)

T tableau above the path associated to $w(h)$ with cells labeled by \square \blacksquare

$(h, T) \longleftrightarrow h_0 = h(\alpha, \beta)$
are in bijection

By recurrence

Thus $h(\alpha, \beta)$ in bijection with



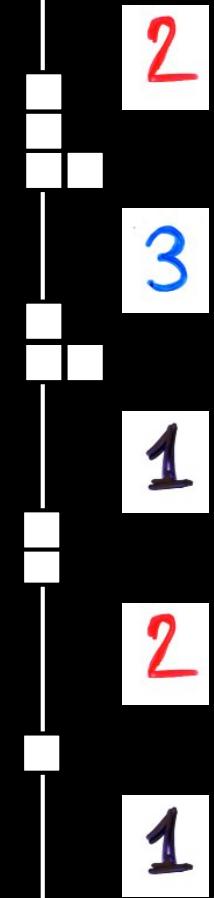
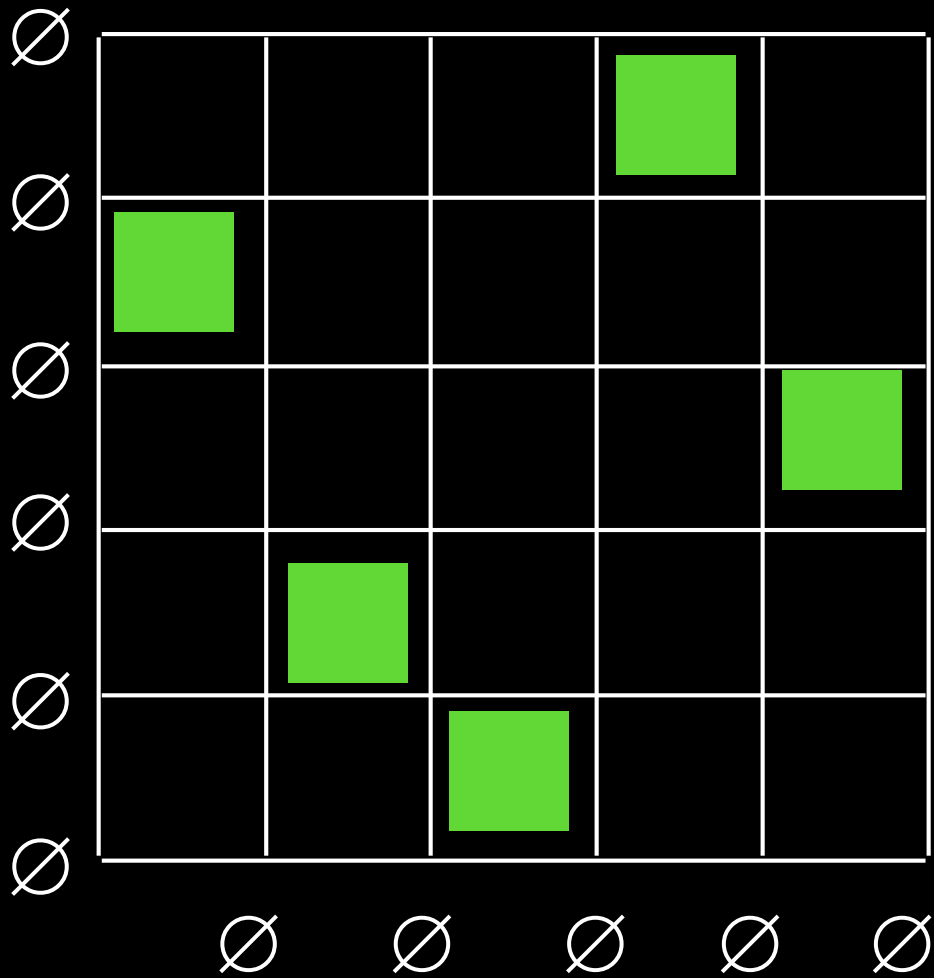
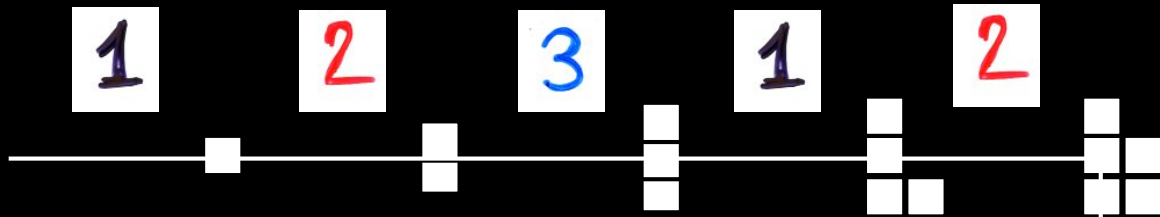
(h, T)

T tableau above the path associated to $w(h)$ with cells labeled by \square \blacksquare

$(h, T) \longleftrightarrow h_0 = h(\alpha, \beta)$
are in bijection

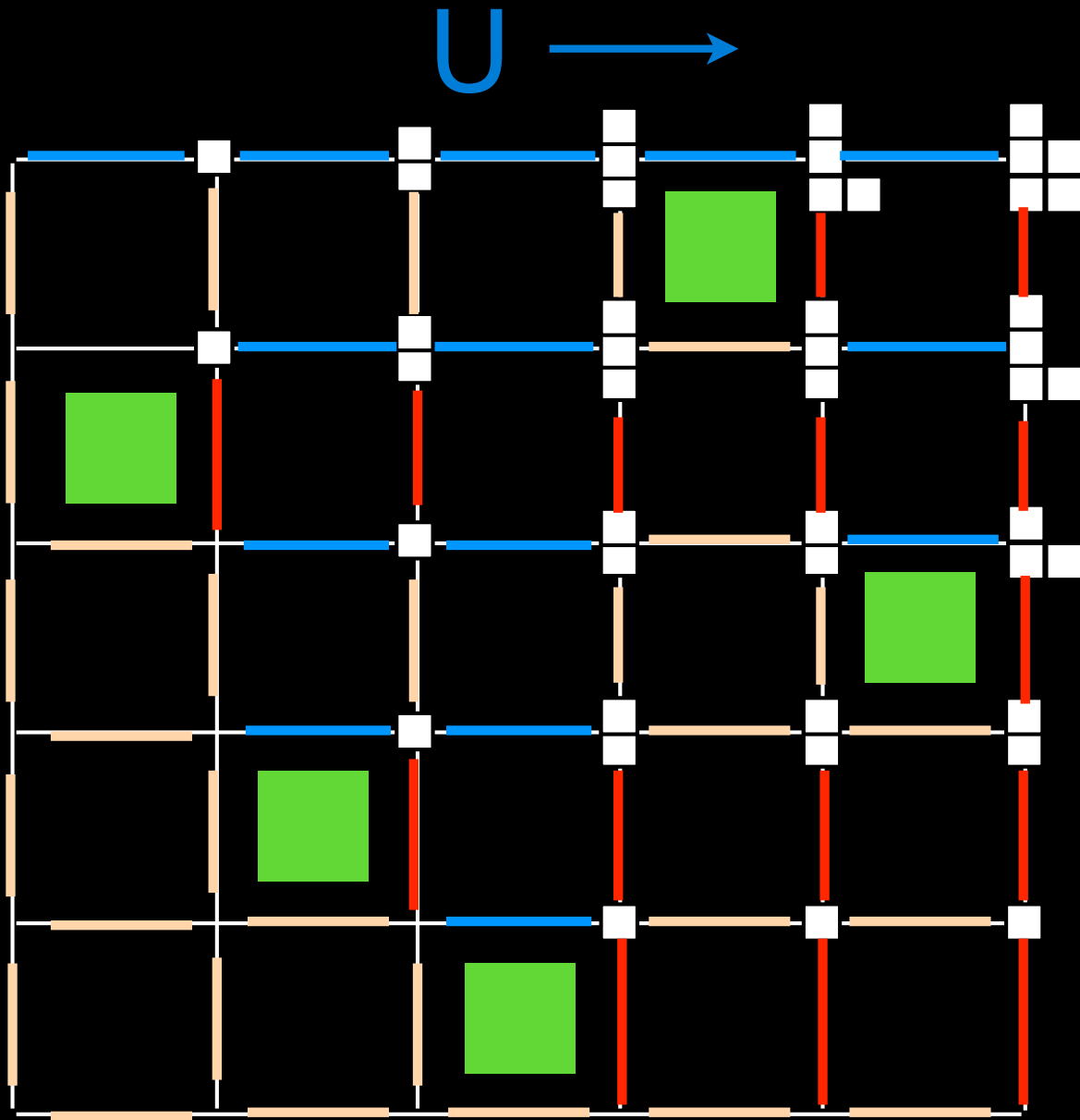
By recurrence

3	
2	5
1	4



4	
2	5
1	3

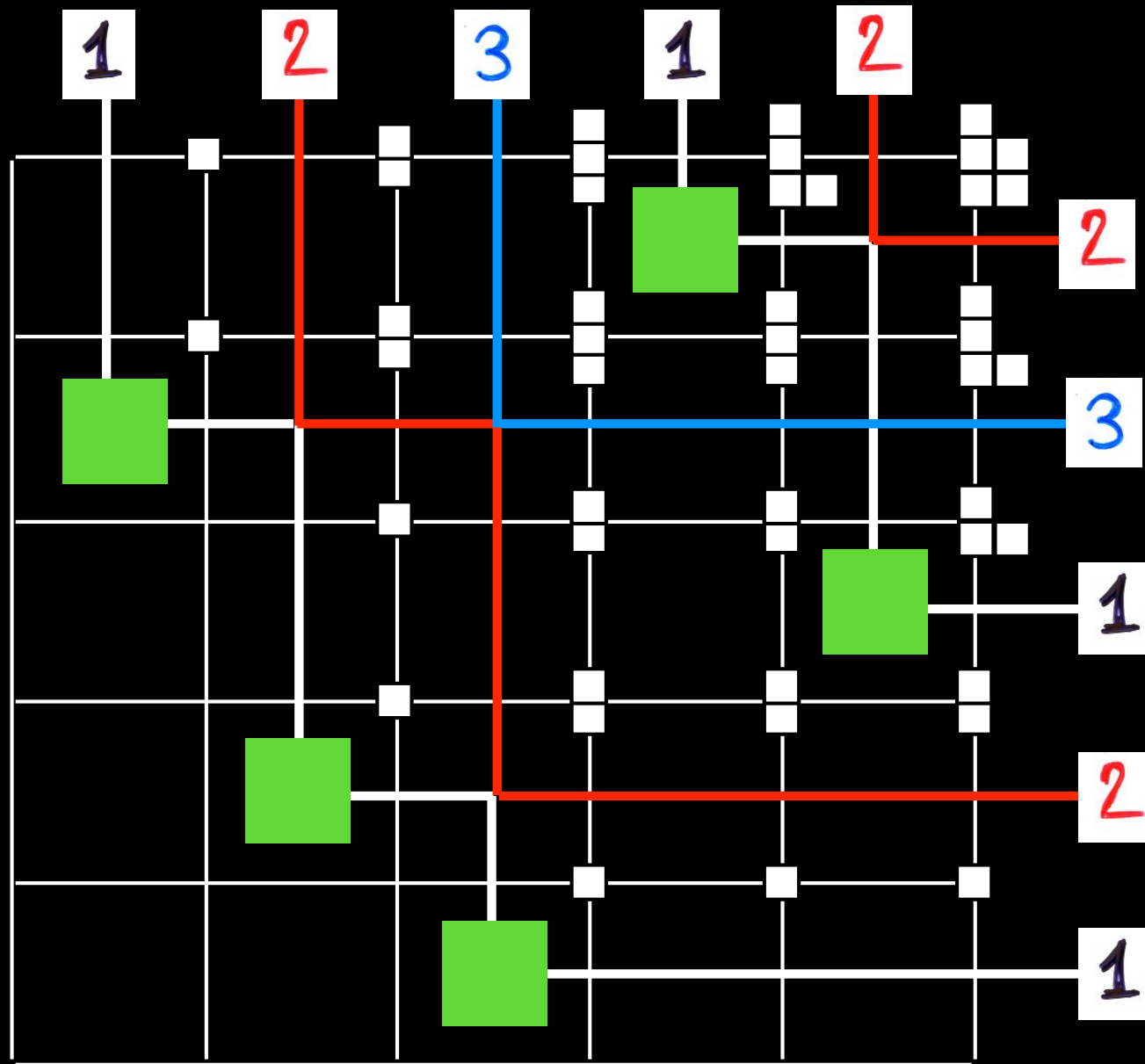
I



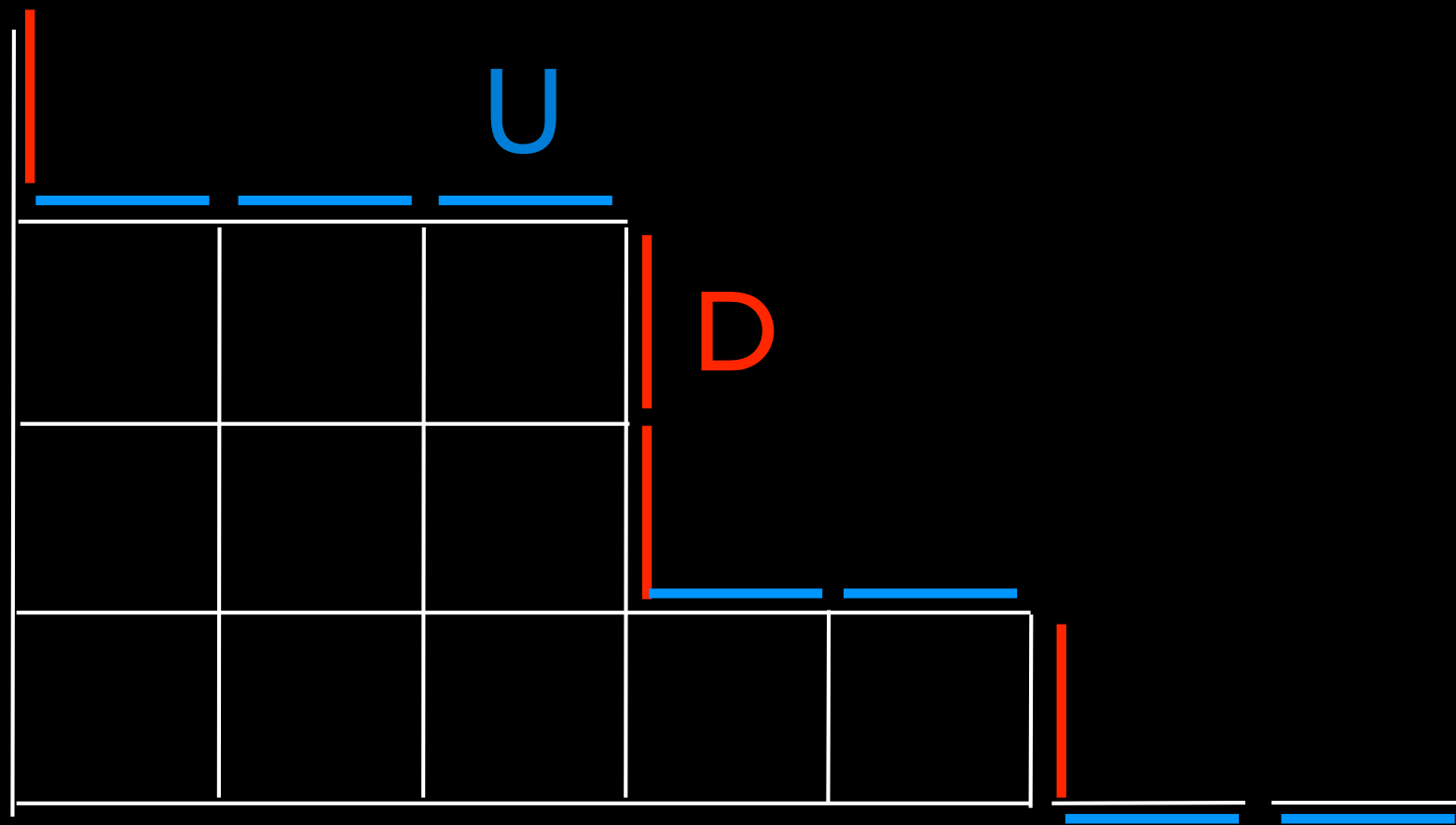
D

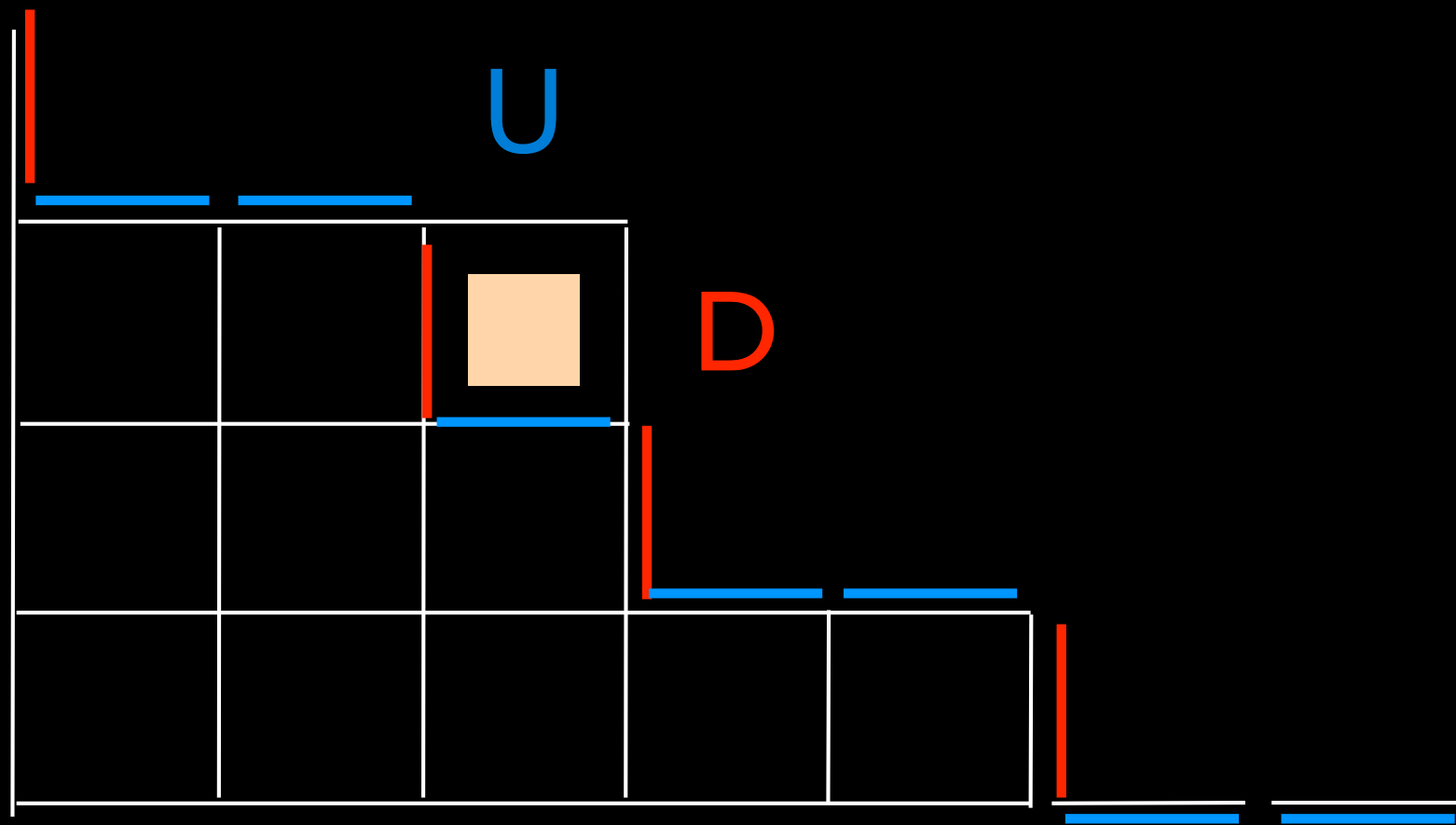
This "propagation" algorithm is exactly the reverse of Fomin's "growth diagrams"

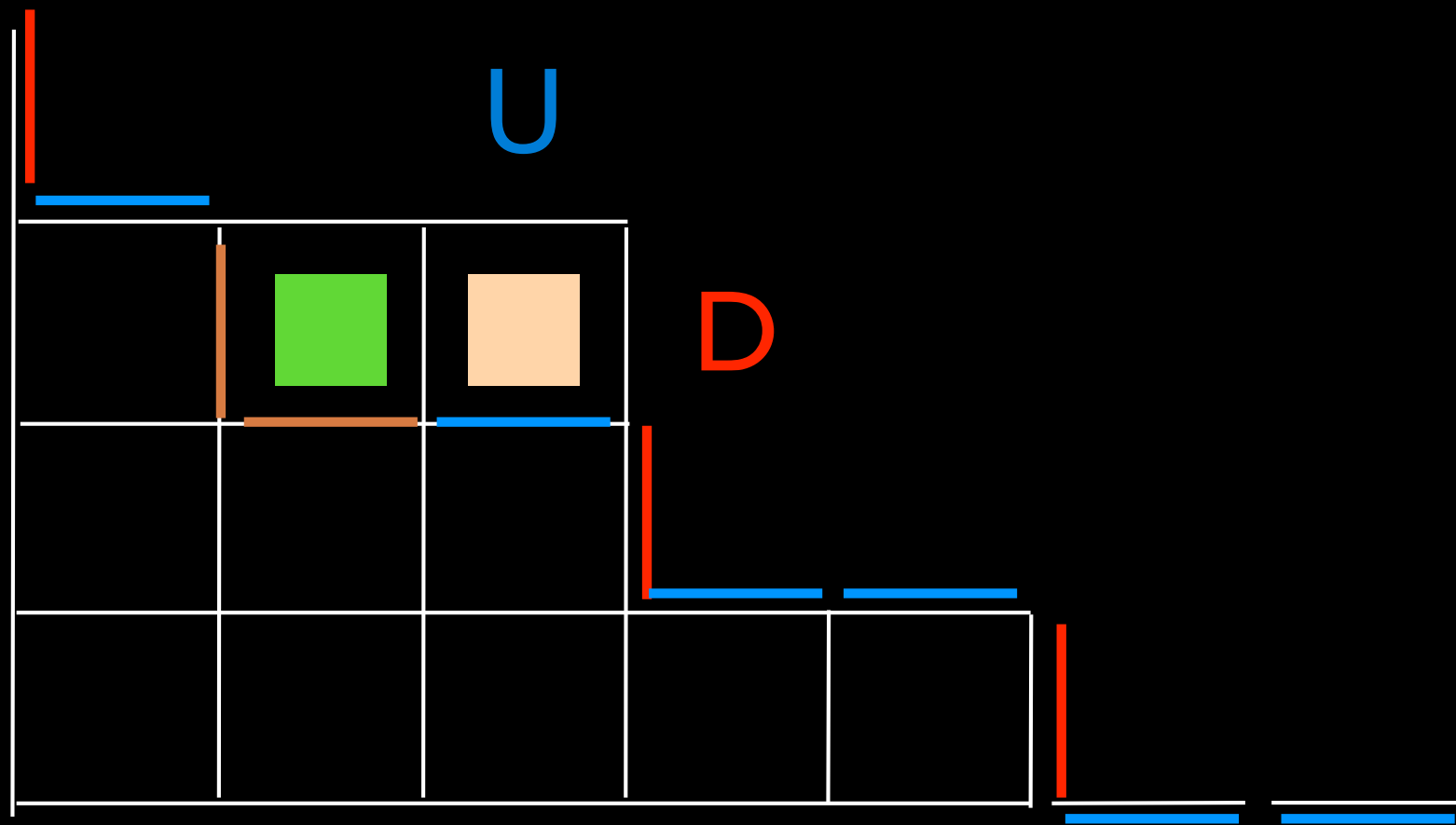
I

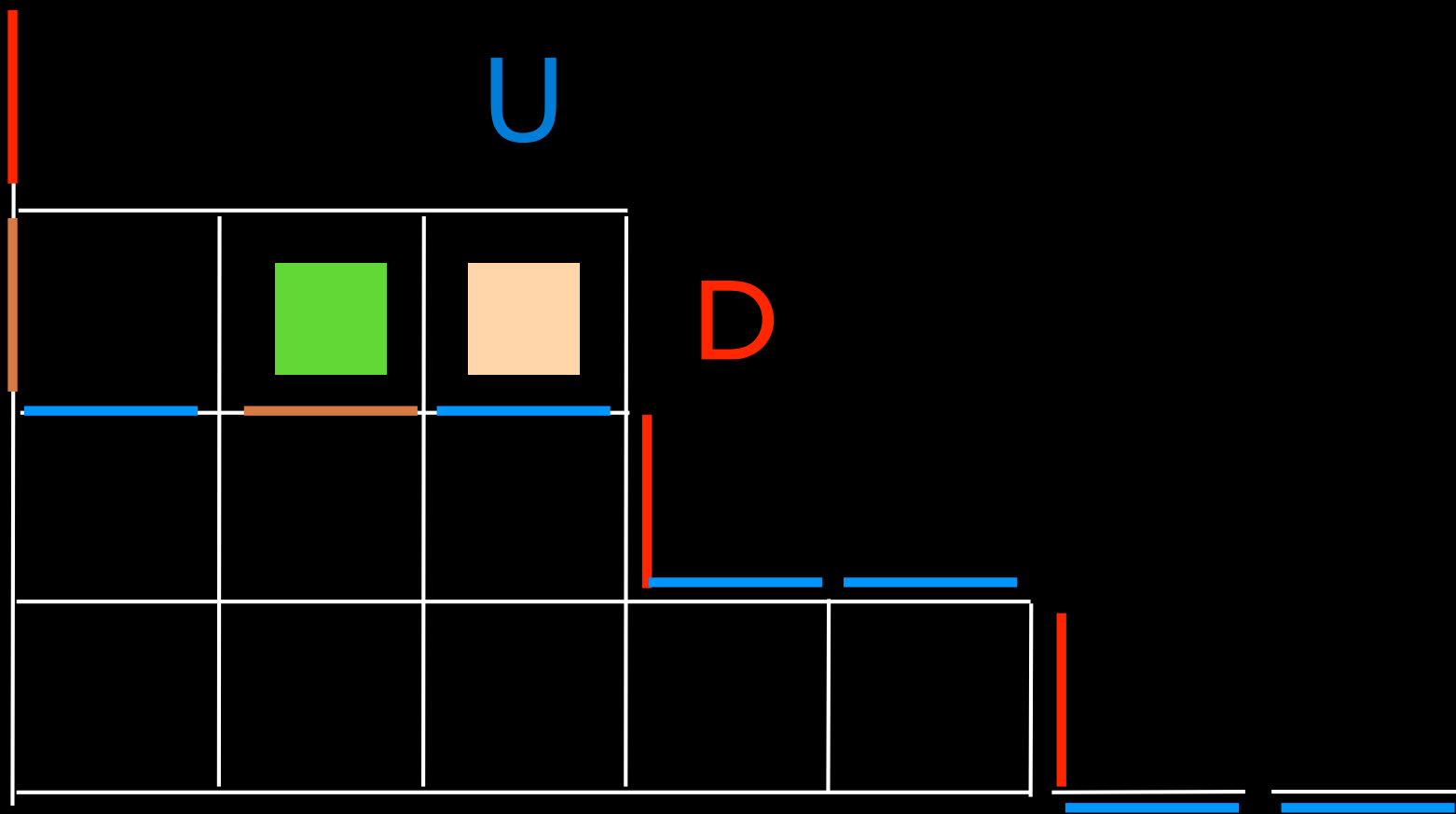


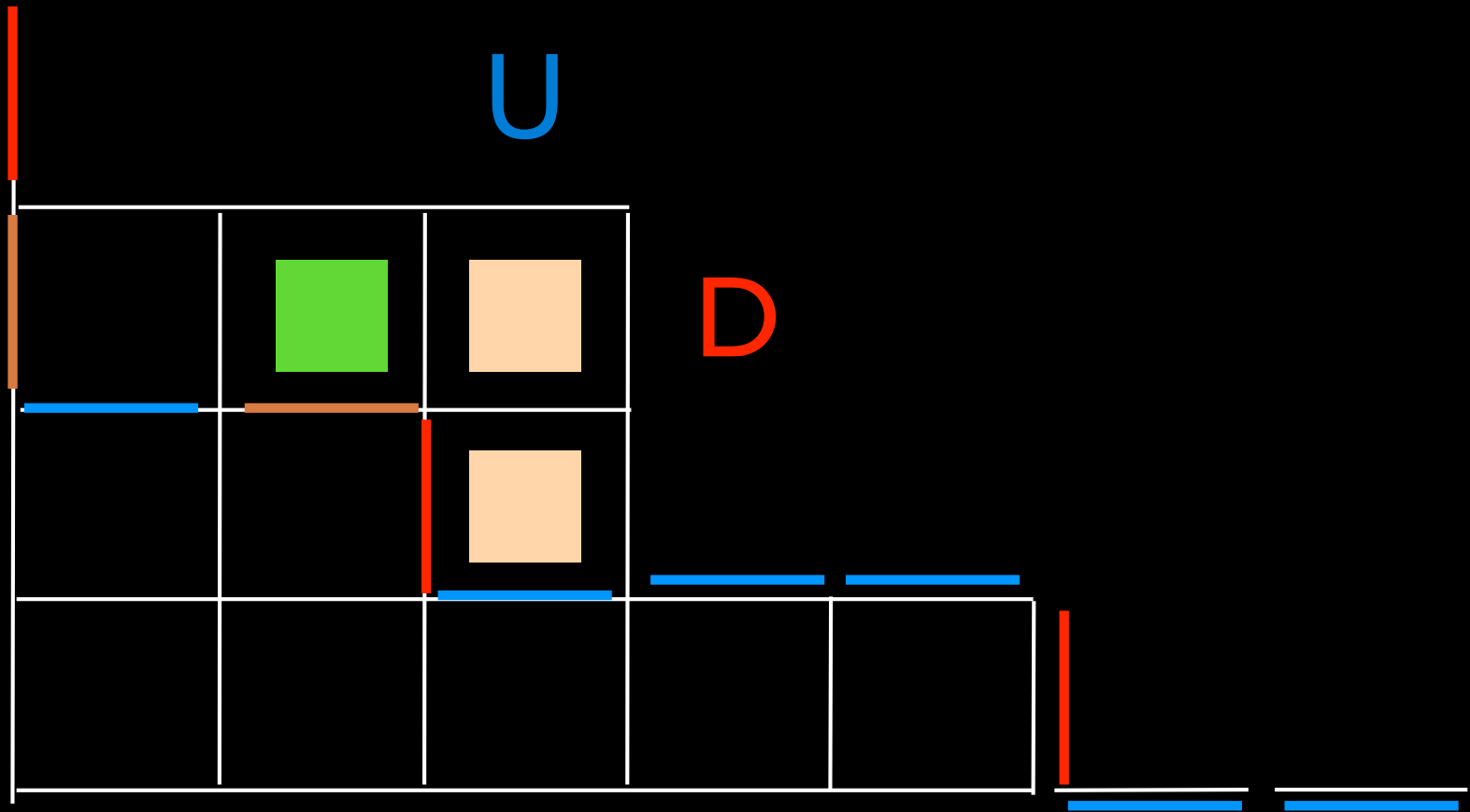
rook placements

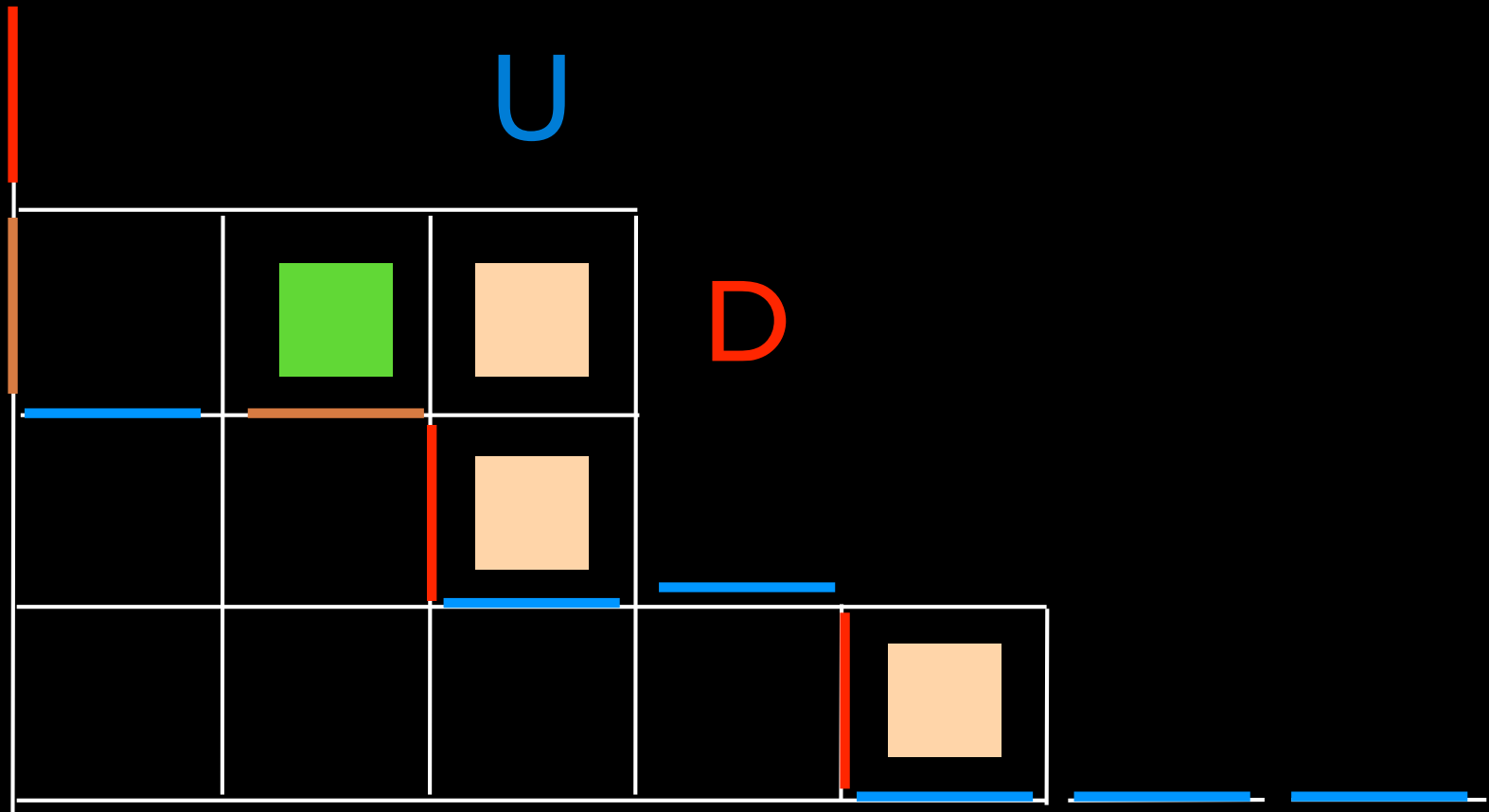


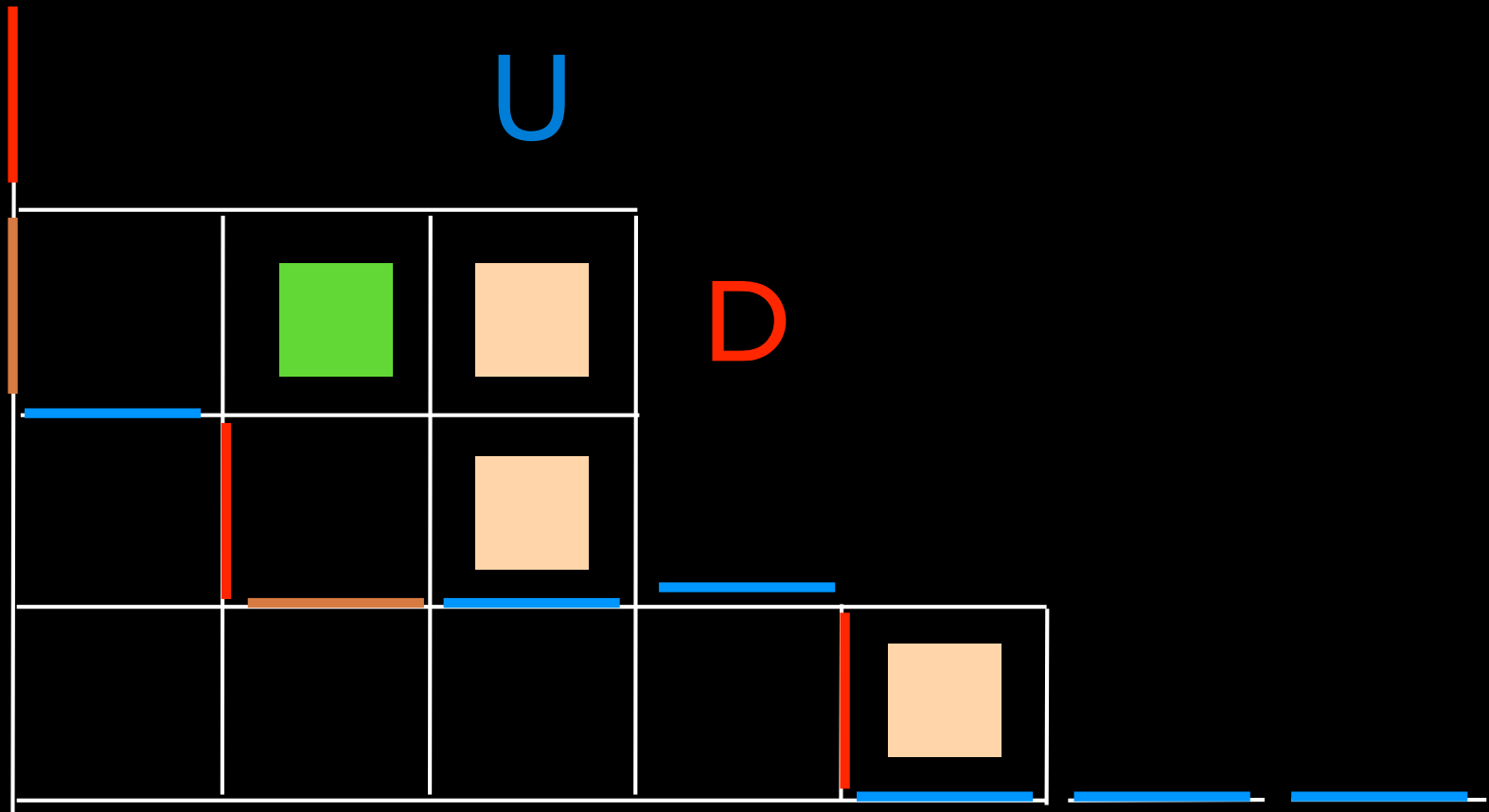


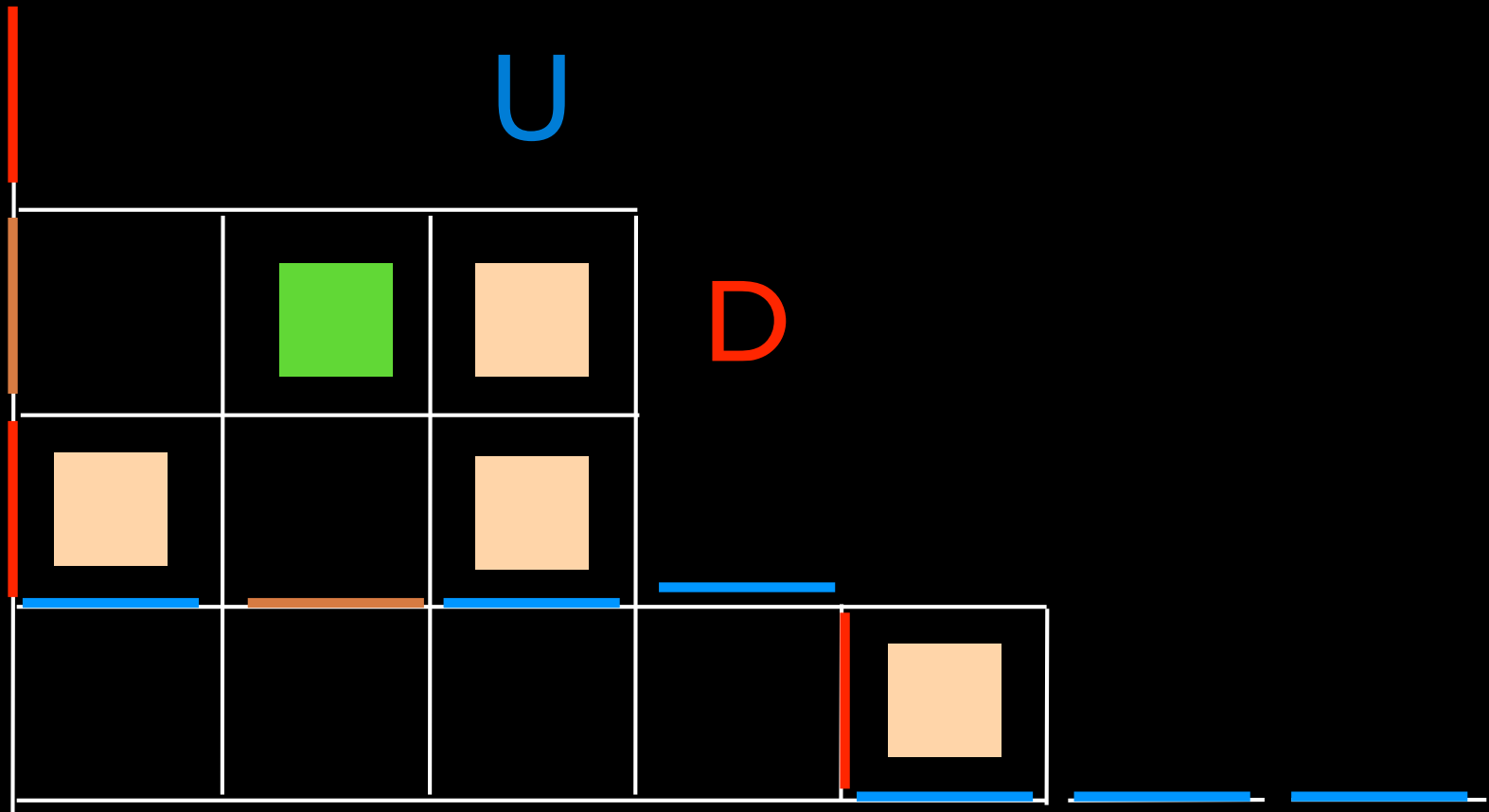


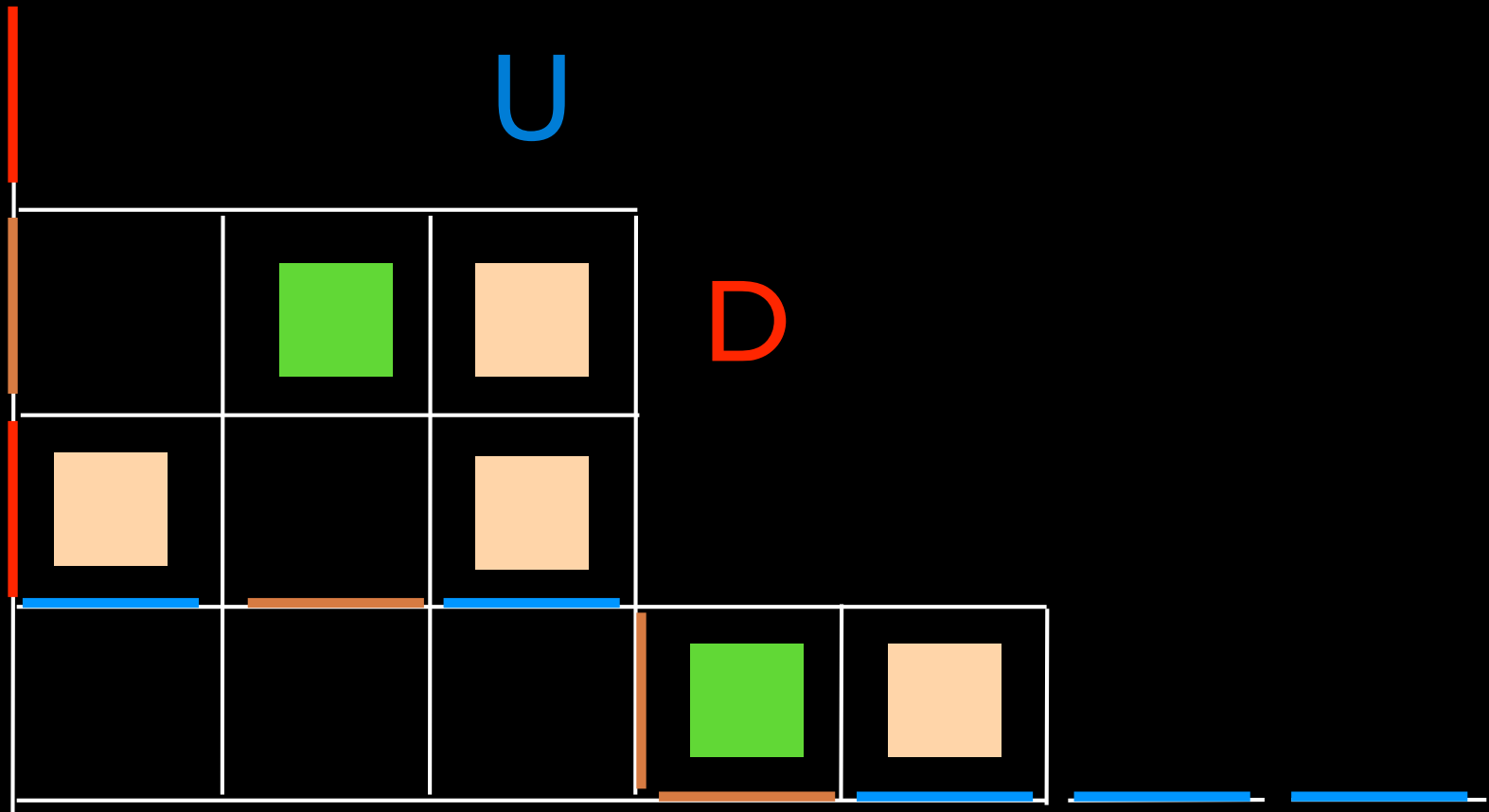


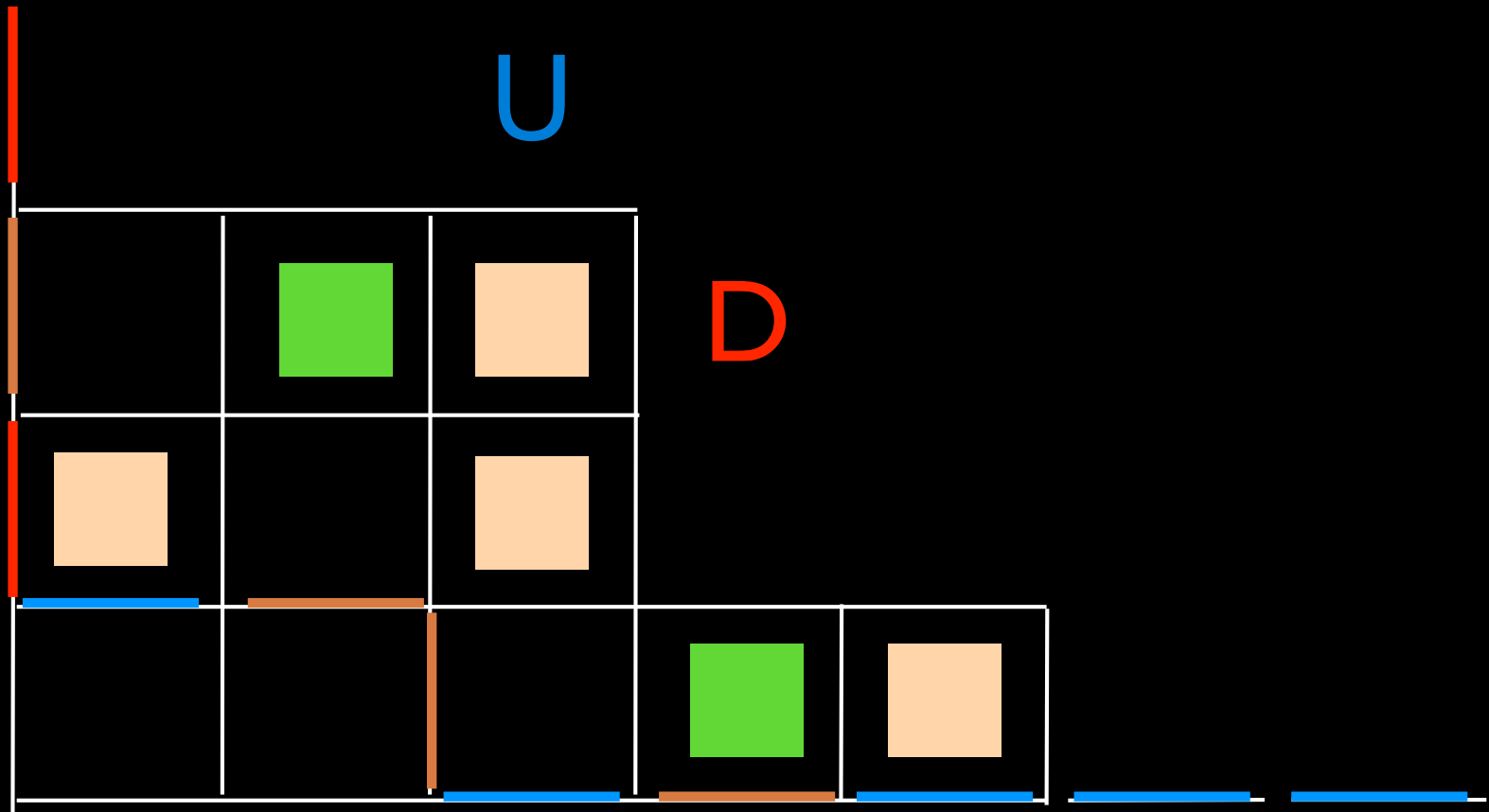


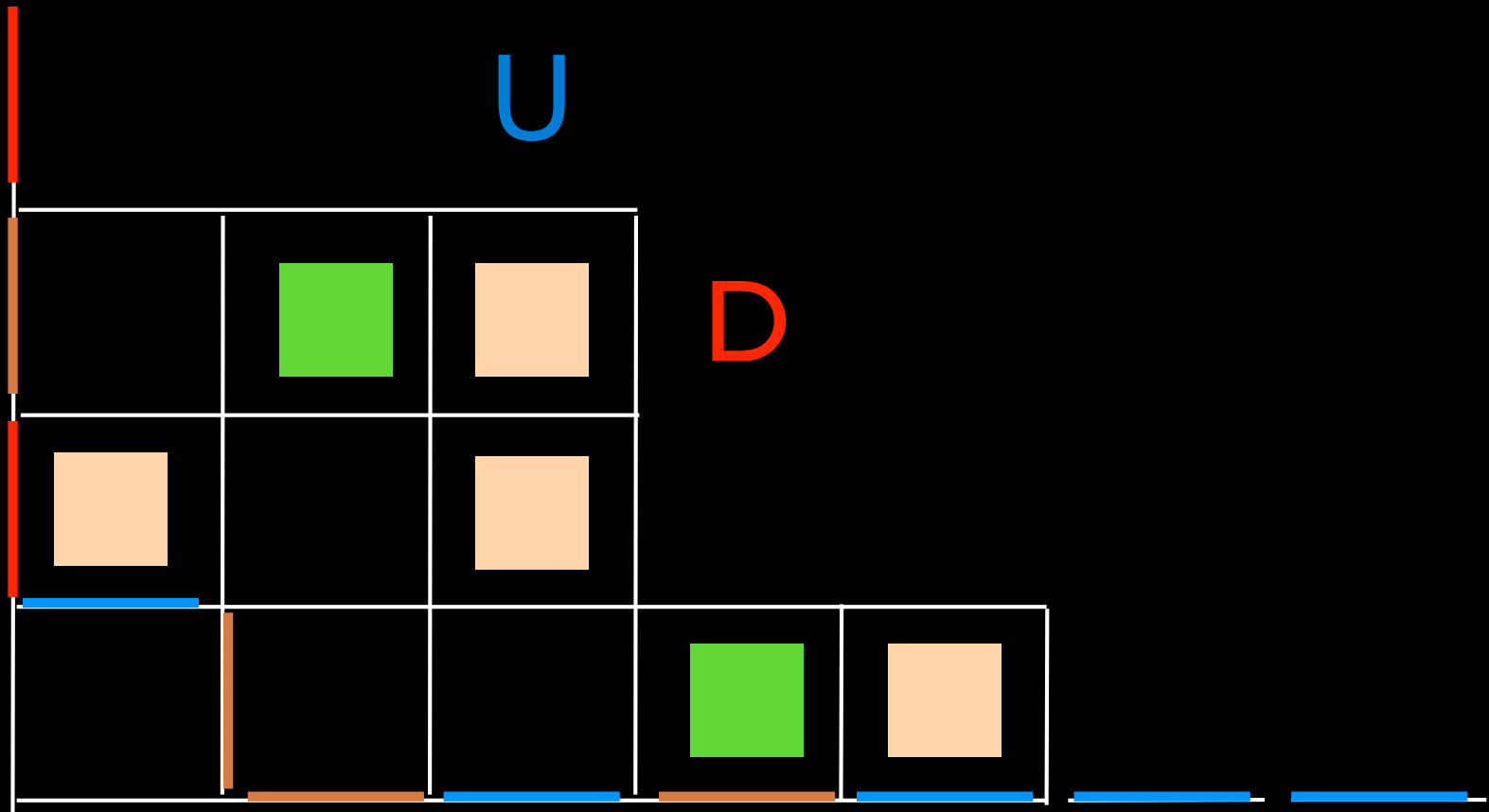




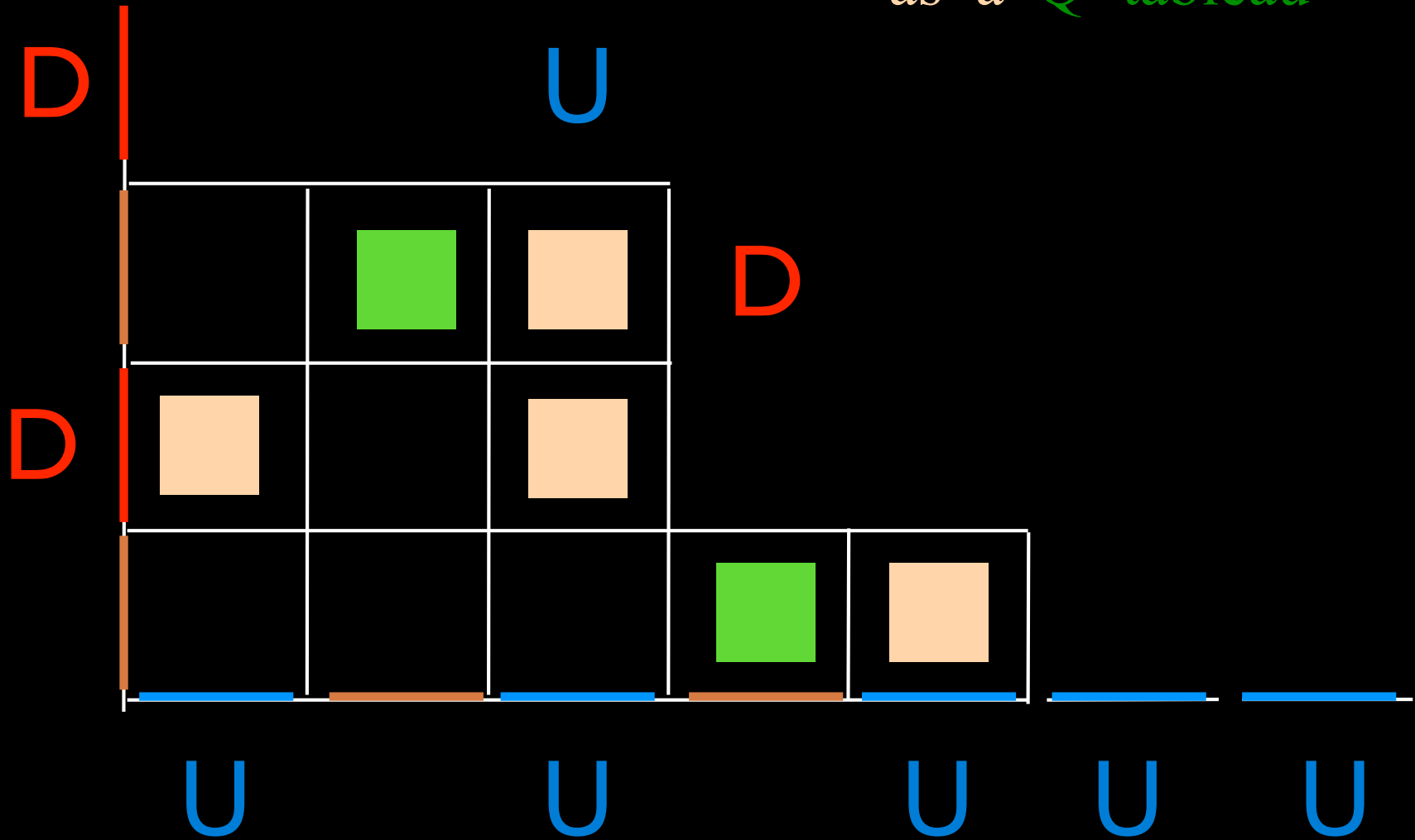




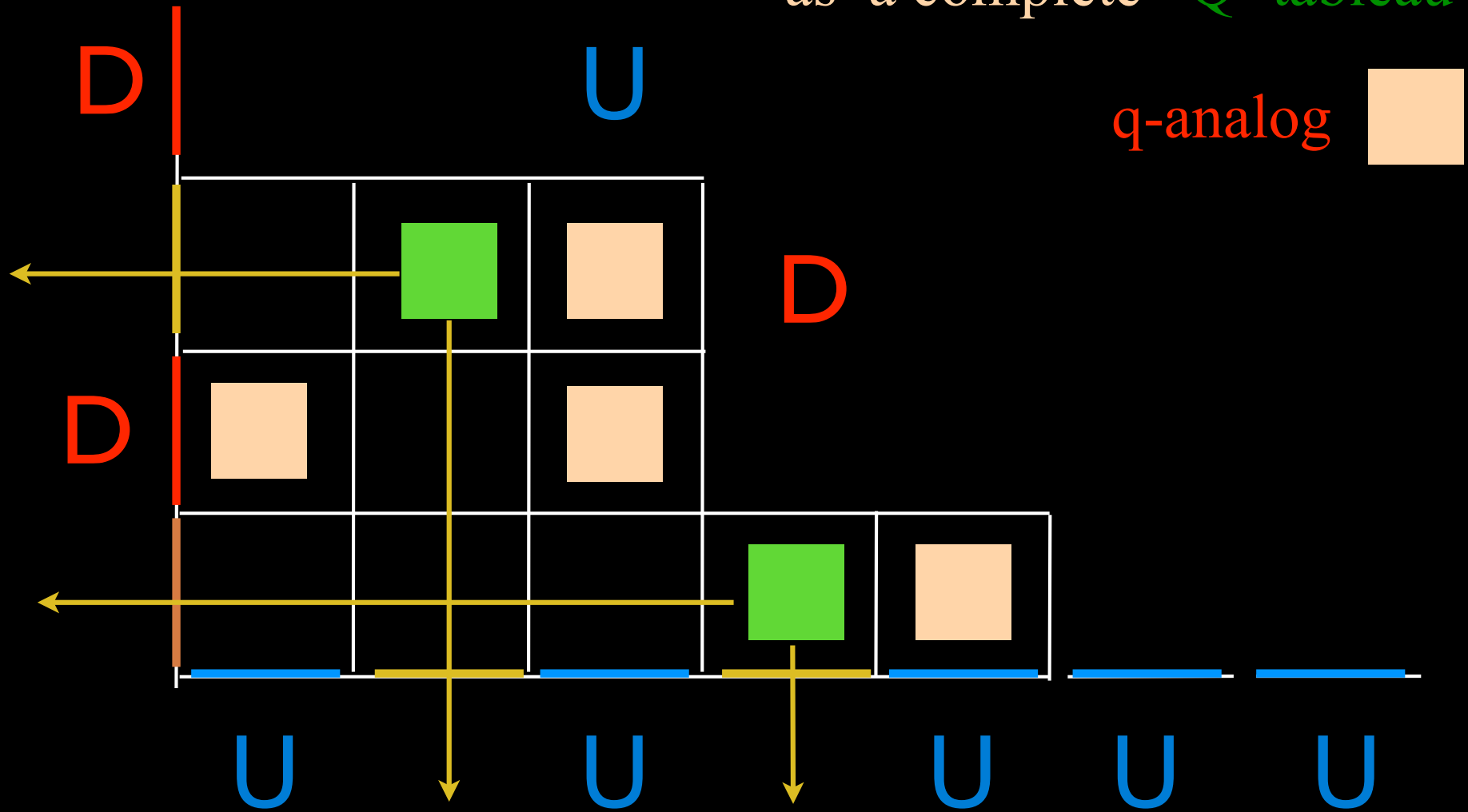




rook placement
as a Q-tableau



rook placement
as a complete Q-tableau

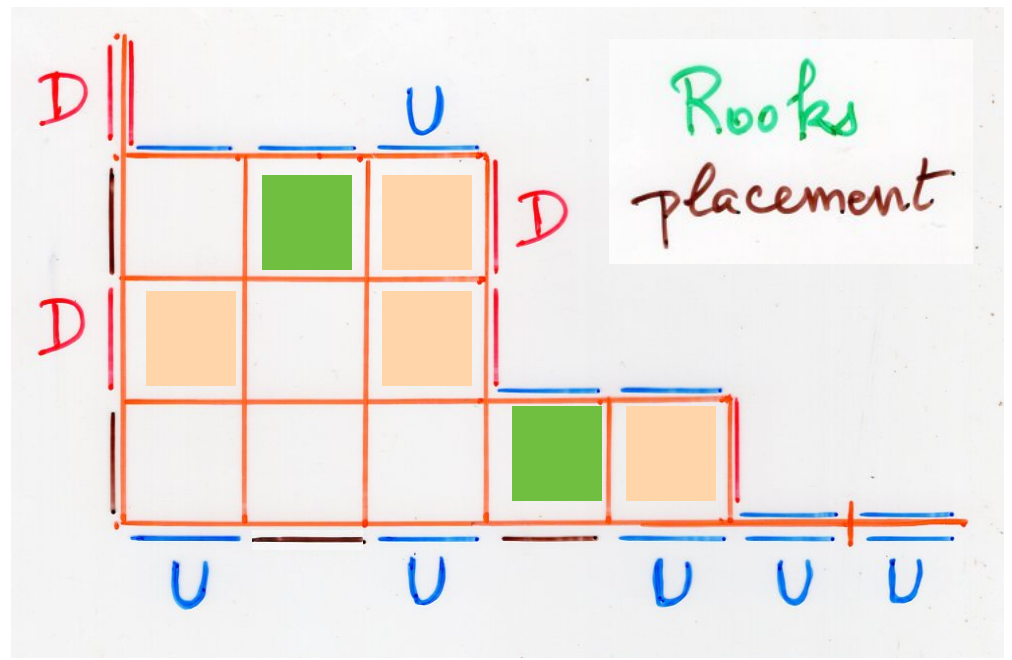


$$w = D U^3 D^2 U^2 D U^2$$

$$w \rightarrow F = F(w) \quad F \text{ Ferrers diagram}$$

Proposition

$$w(U, D) = \sum_T D^{i(T)} U^{j(T)}$$



$i(T)$ = number of rows with no cell labeled
 $j(T)$ = number of columns with no cell labeled
 $UD \rightarrow I_v I_h$

$$q^{k(T)}$$

Lemma Every word w with letters U and D can be written in a unique way

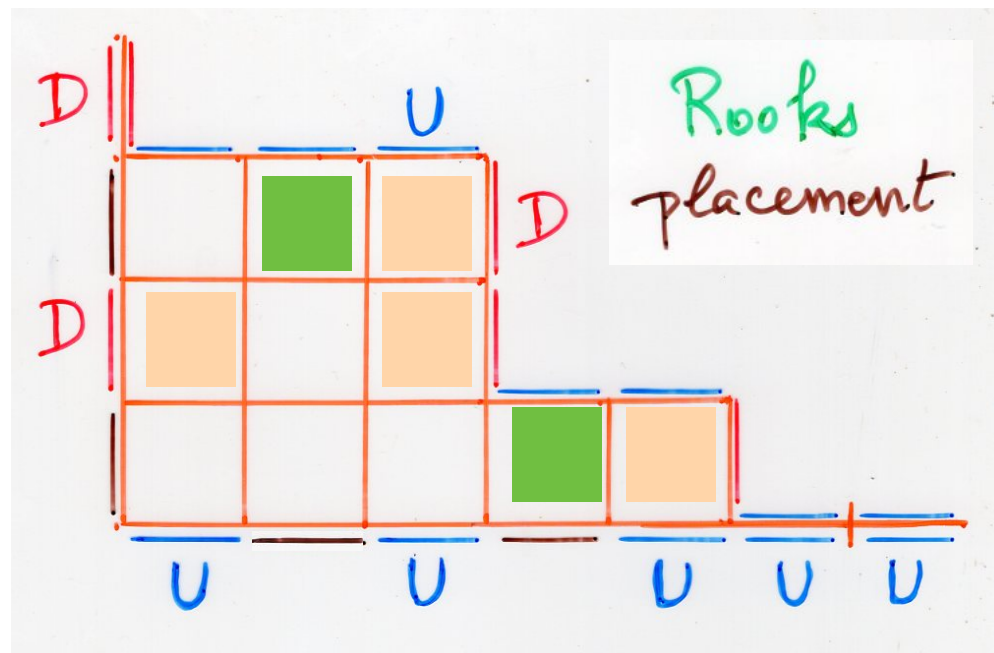
$$w = \sum_{i,j \geq 0} c_{ij}(w) D^i U^j$$

Proposition

$c_{ij}(w)$ = number of placements of k rooks on the Ferrers "board" F

with $i = |w|_D - k$

$j = |w|_U - k$



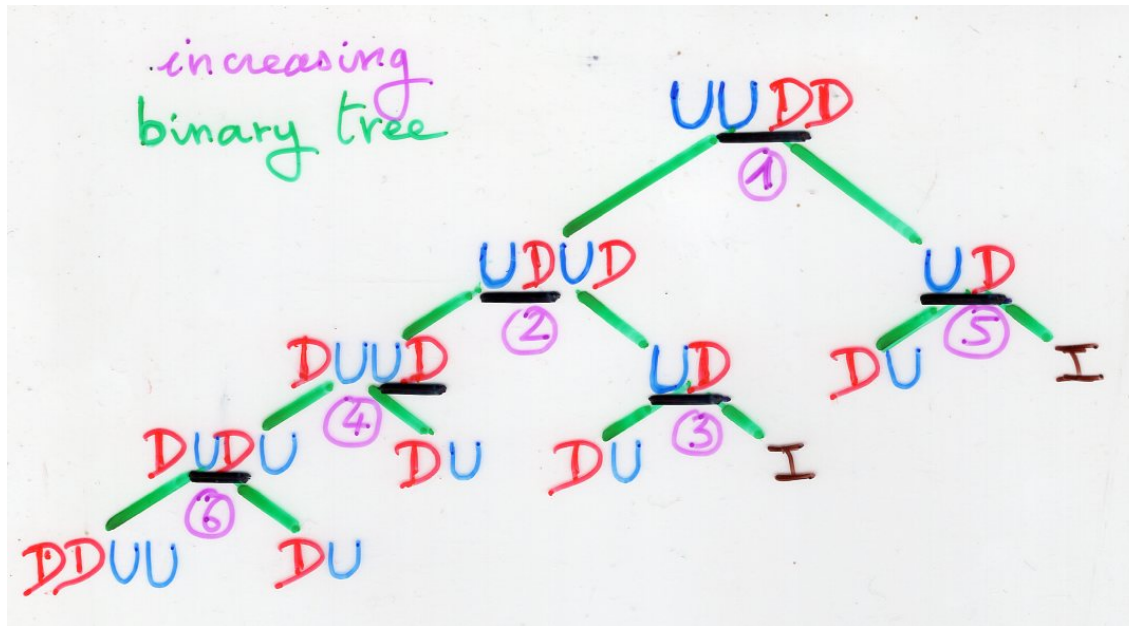
$$U^n D^n = \sum_{0 \leq i \leq n} c_{n,i} D^i U^i$$

$$c_{n,0} = n!$$

permutations

$$c_{n,i} = \binom{n}{i}^2 (n-i)!$$

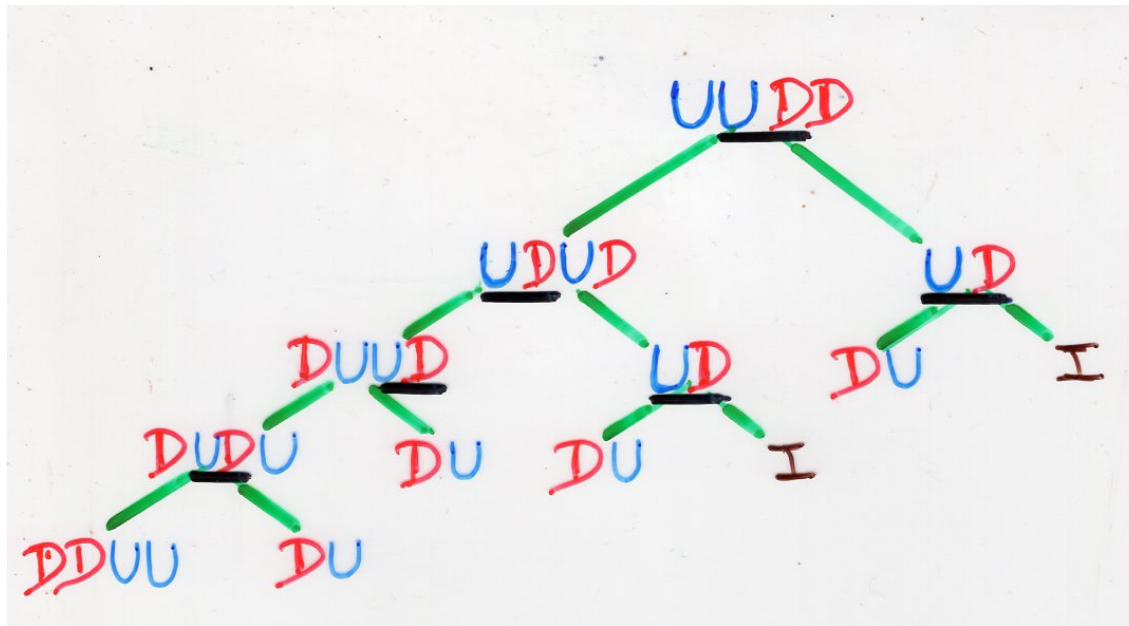
choice of the underlying grid $\binom{n}{i}^2$



binary tree T
 associated
 to a possible
 rewriting process

$$UD^2 = D^2U^2 + 4DU + 2I$$

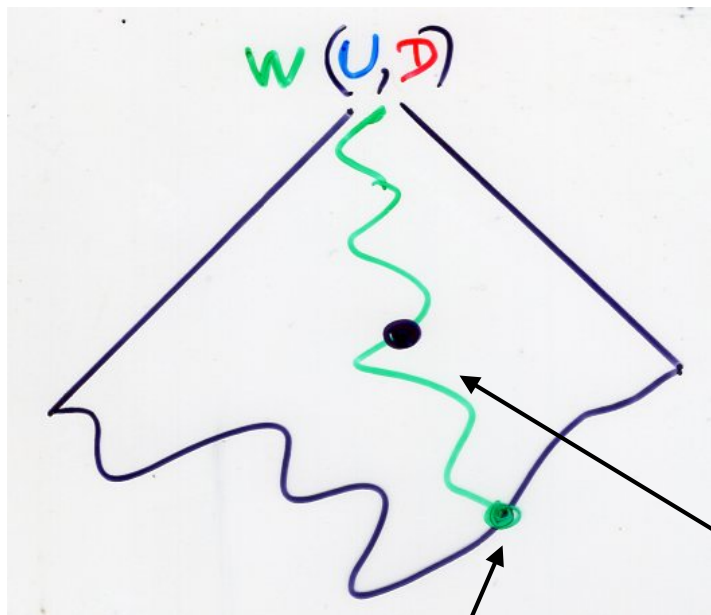
this polynomial is independent
 of the order of the substitutions



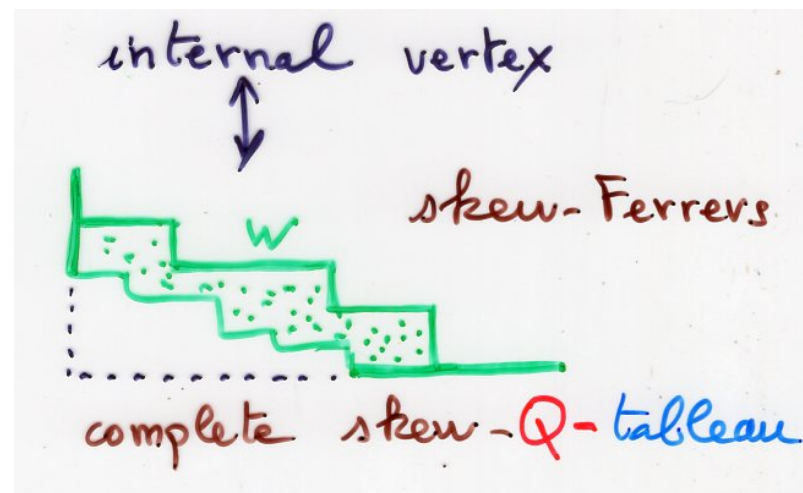
binary tree T
 associated
 to a possible
 rewriting process

$$U^2 D^2 = D^2 U^2 + 4DU + 2I$$

this polynomial is independent
 of the order of the substitutions



binary tree T
 associated
 to a possible
 rewriting process



leaves of T

bijection
 \longleftrightarrow

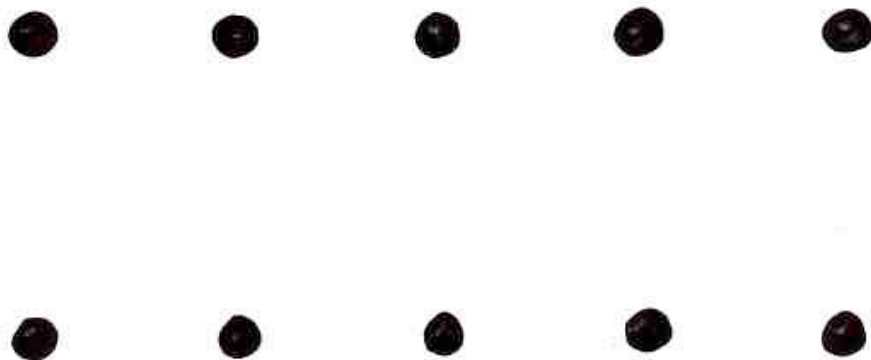
complete
 Q -tableaux
 shape λ

$$\lambda = F(w)$$

Another representation of the algebra

$$UD = DU + 1$$

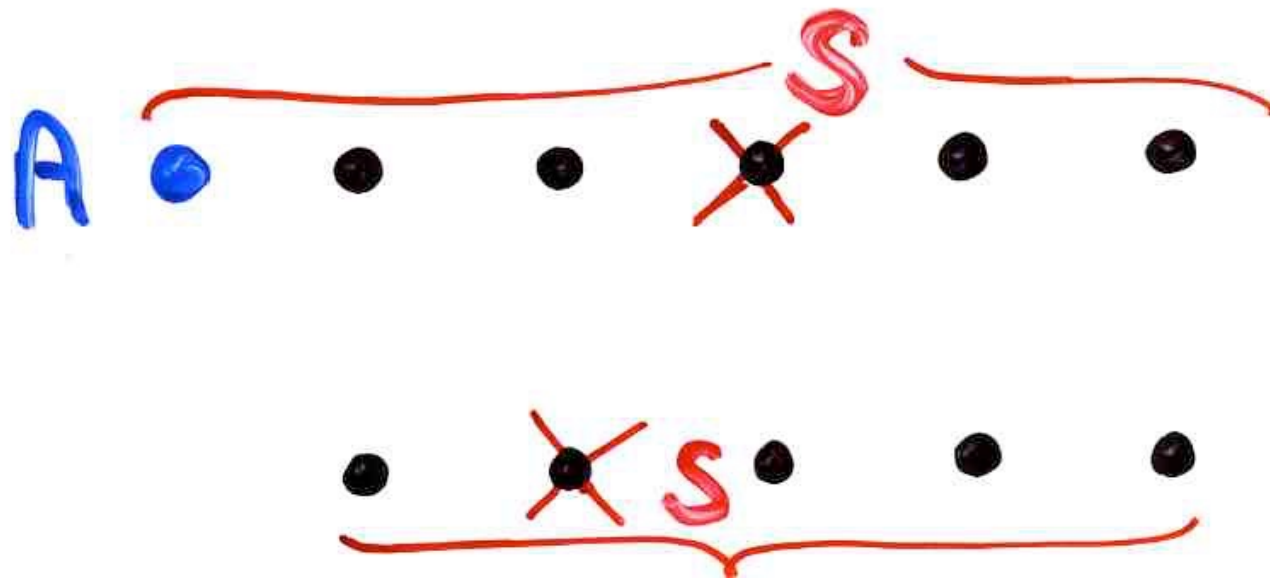
Polya urn



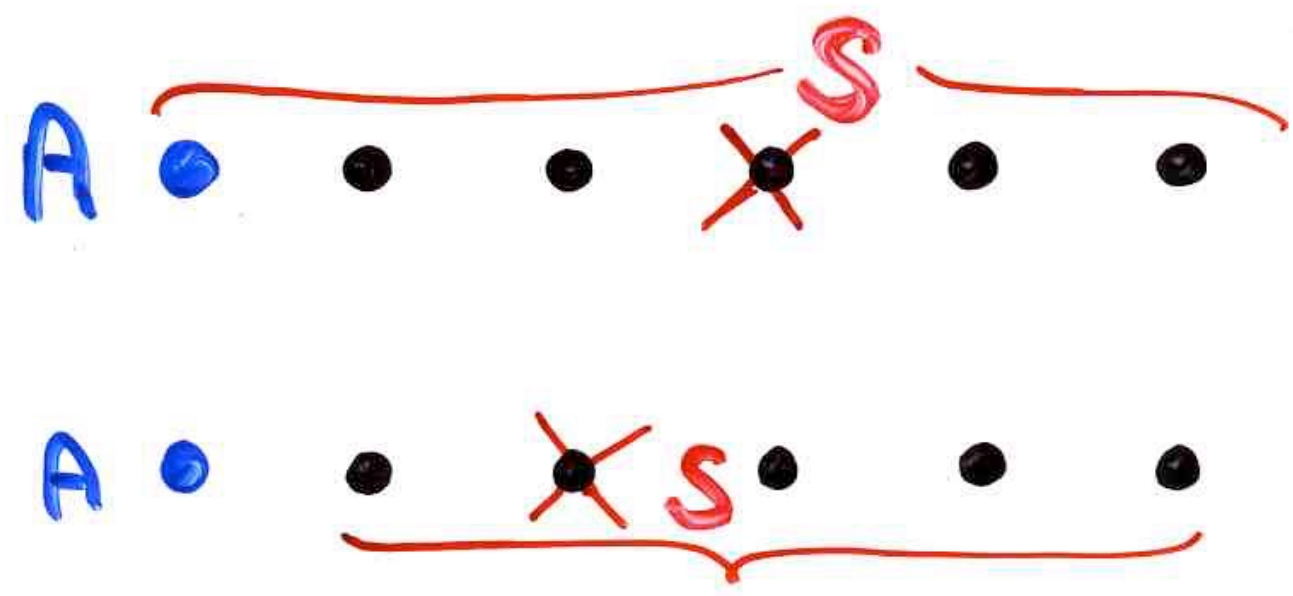
A

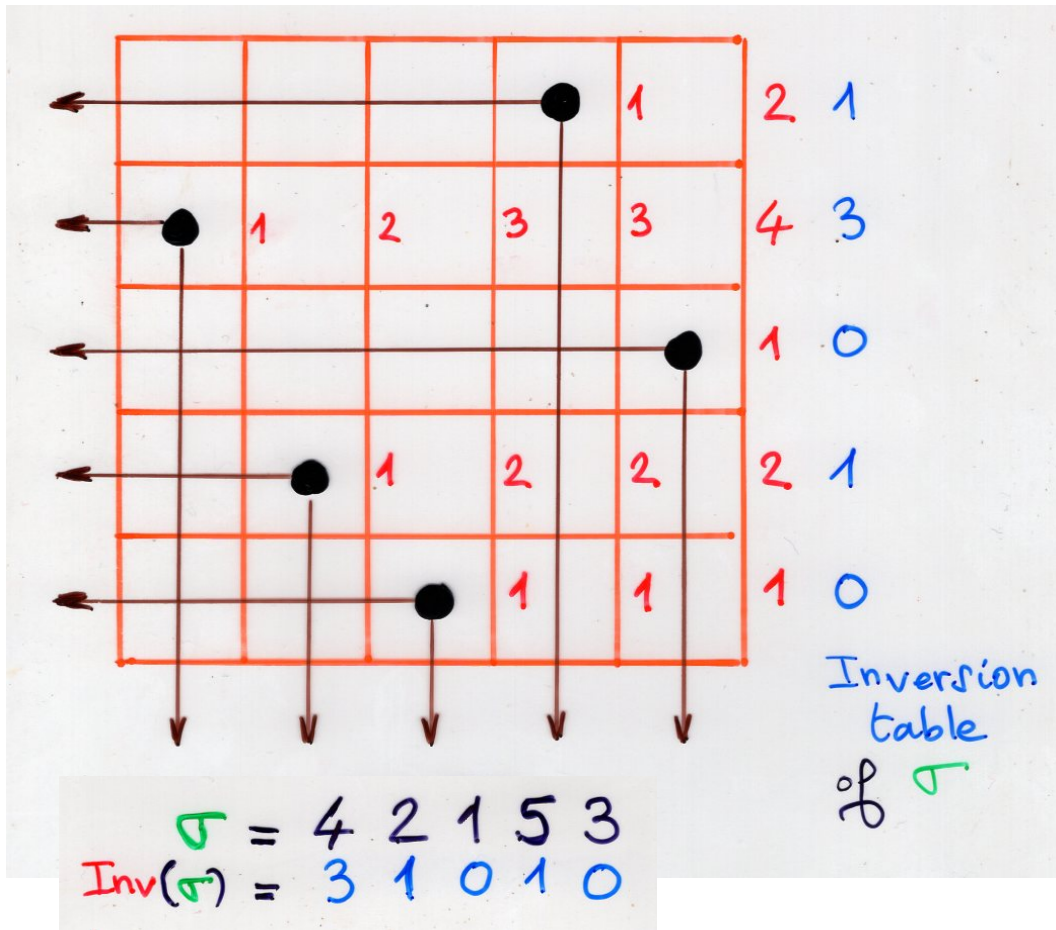
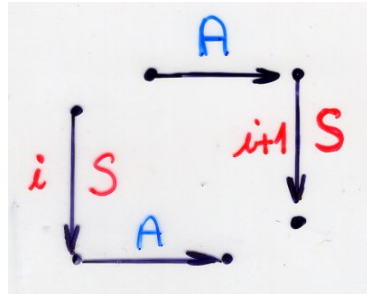
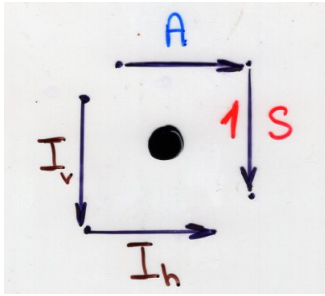






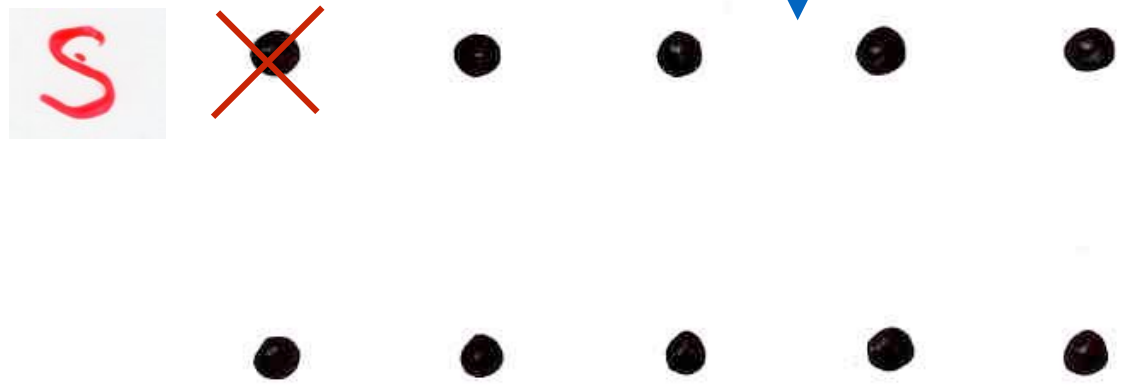
$$A S - S A = I$$





Priority queue

$$AS-SA = I$$



data structures

Computer Science

