

Course IMSc Chennai, India

January-March 2017

Enumerative and algebraic combinatorics,
a bijective approach:

commutations and heaps of pieces

(with interactions in physics, mathematics and computer science)

Monday and Thursday 14h-15h30

www.xavierviennot.org/coursIMSc2017



IMSc

January-March 2017

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Chapter 1
Commutation monoids
and
heaps of pieces:

basic definitions
(2)

IMSc, Chennai

9 January 2017

from the previous lecture

commutation

relation

C

antireflexive
symmetric

\equiv_C

congruence of A^* generated
by the commutations

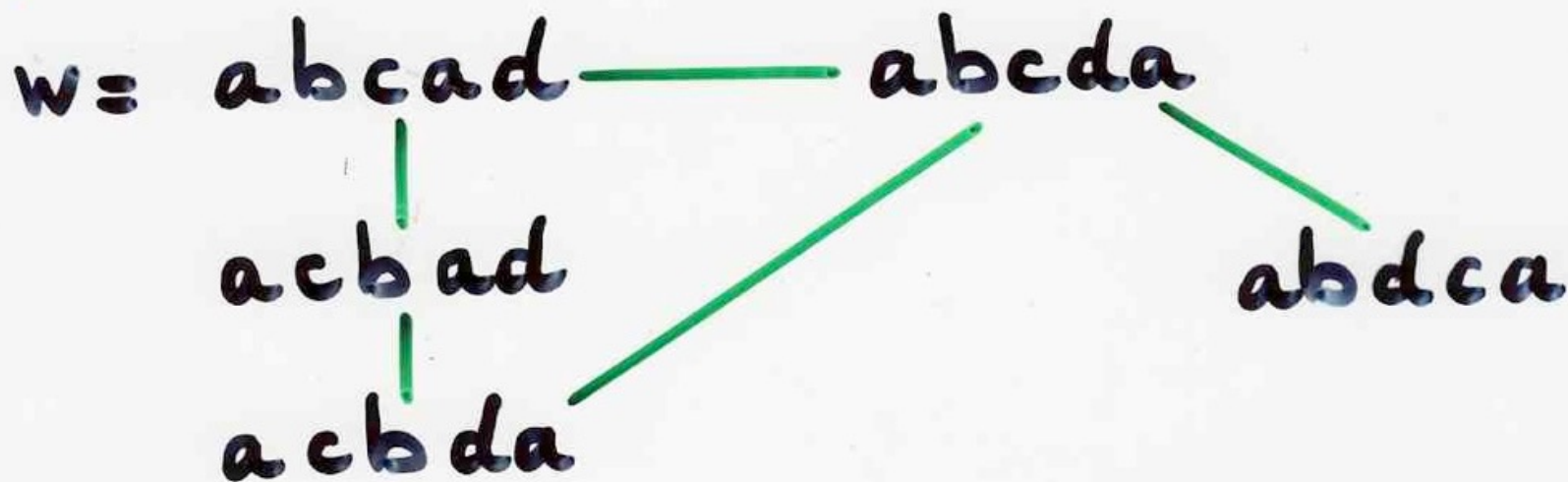
$$ab \equiv ba \text{ iff } aCb$$

- $aCb \Leftrightarrow bCa$
- ~~aCa~~

ex: $A = \{a, b, c, d\}$

$$C \begin{cases} ad = da \\ bc = cb \\ cd = dc \end{cases}$$

equivalence class



commutation
monoid

$$A^* / \equiv C$$

$[w]$

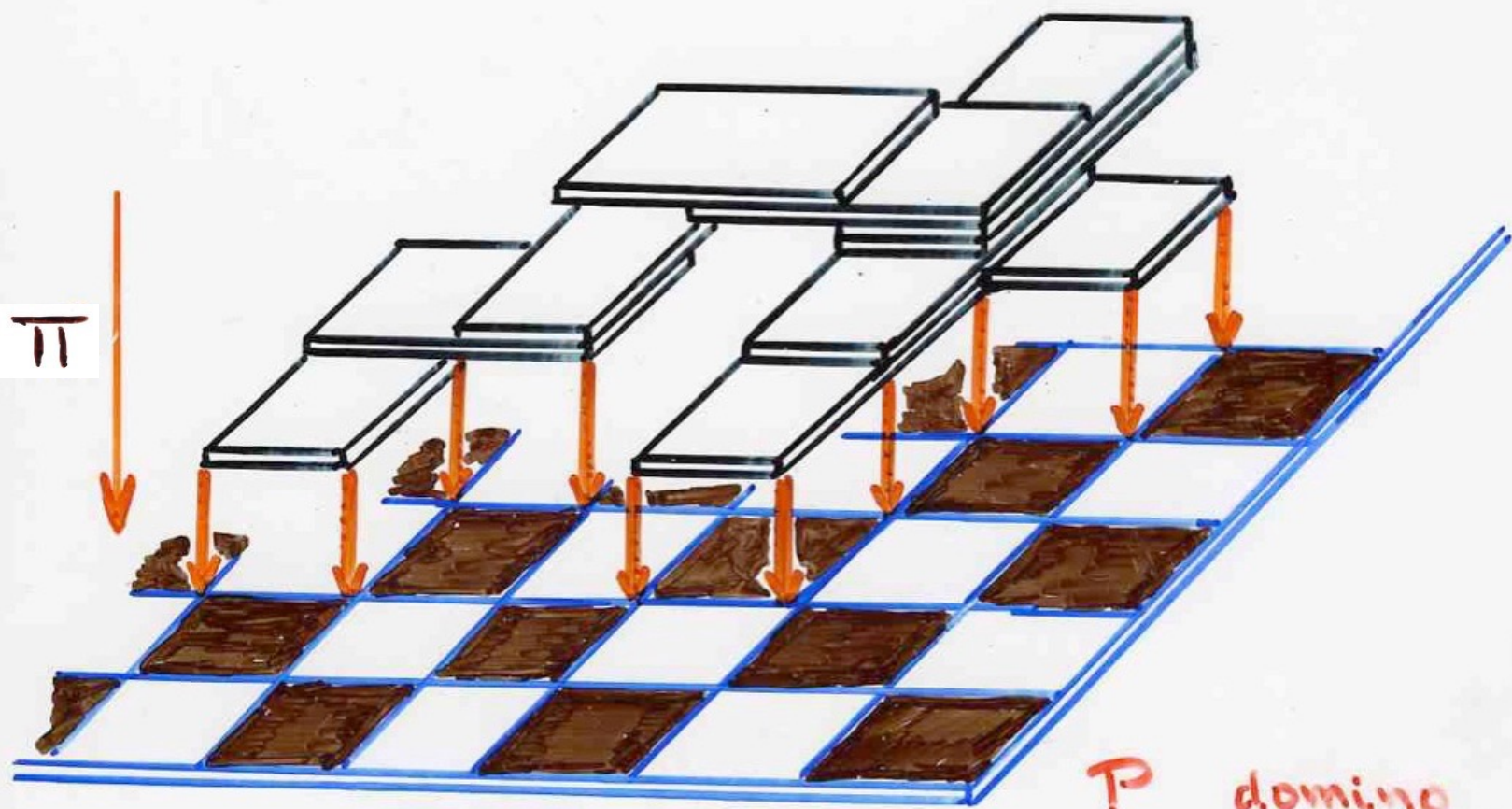
equivalence class
of the word $w \in A^*$

$$A^* / \equiv C$$

- product in the
commutation monoid

$$[u] \cdot [v] = [uv]$$

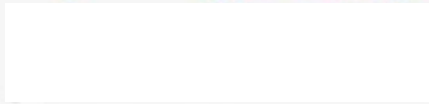
independent of the choices
of representants u and v



π

$$B = R \times R$$

P domino



heap

definition

- \mathcal{P} set (of basic pieces)
- \mathcal{C} binary relation on \mathcal{P} $\left\{ \begin{array}{l} \text{symmetric} \\ \text{reflexive} \end{array} \right.$
(dependency relation)
- heap E , finite set of pairs
 (α, i) $\alpha \in \mathcal{P}, i \in \mathbb{N}$ (called pieces)

projection

level

$$(i) \quad (\alpha, i), (\beta, j) \in E, \alpha \mathcal{C} \beta \implies i \neq j$$

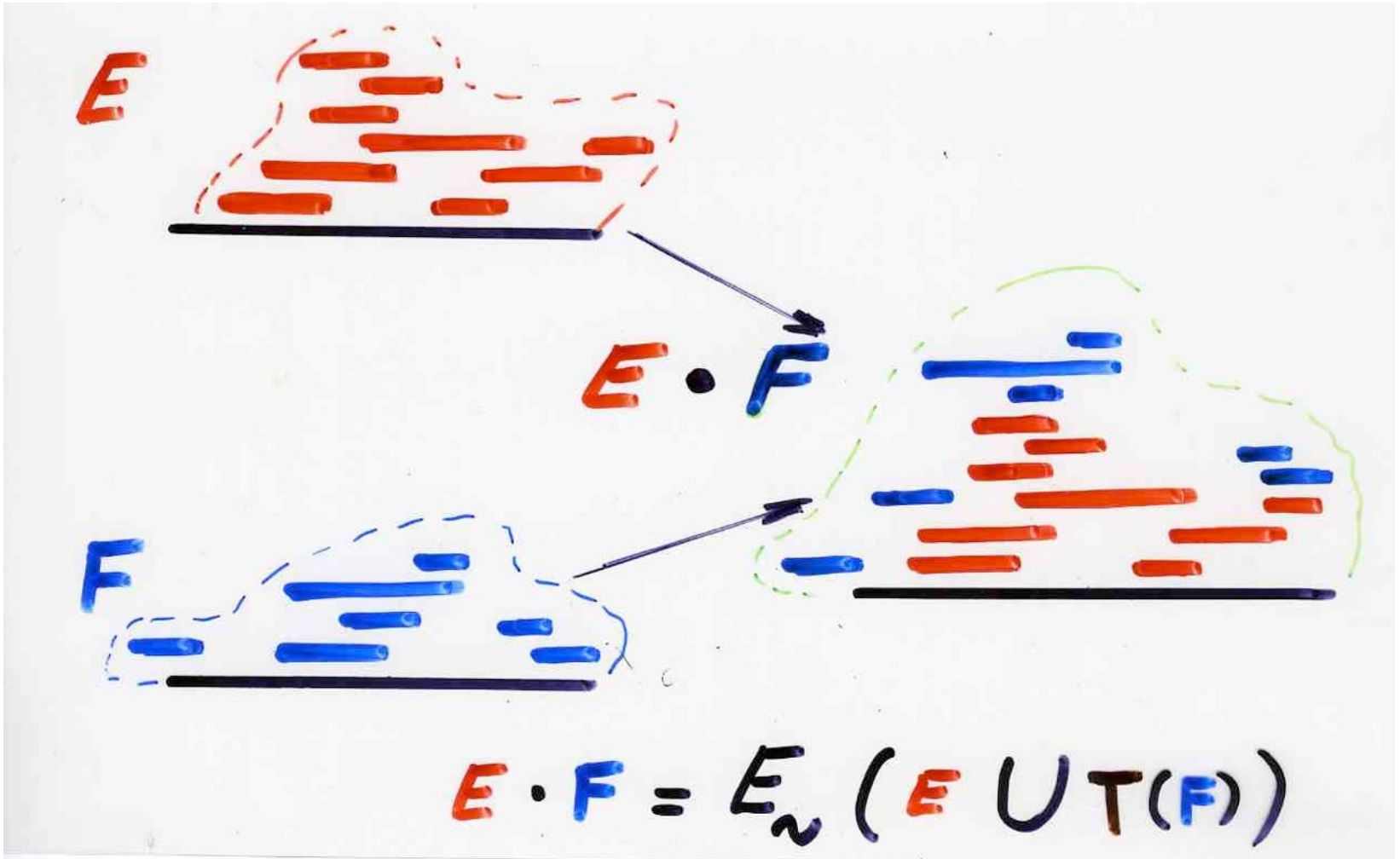
$$(ii) \quad (\alpha, i) \in E, i > 0 \implies \exists \beta \in \mathcal{P}, \alpha \mathcal{C} \beta, \\ (\beta, i-1) \in E$$

Heaps monoid

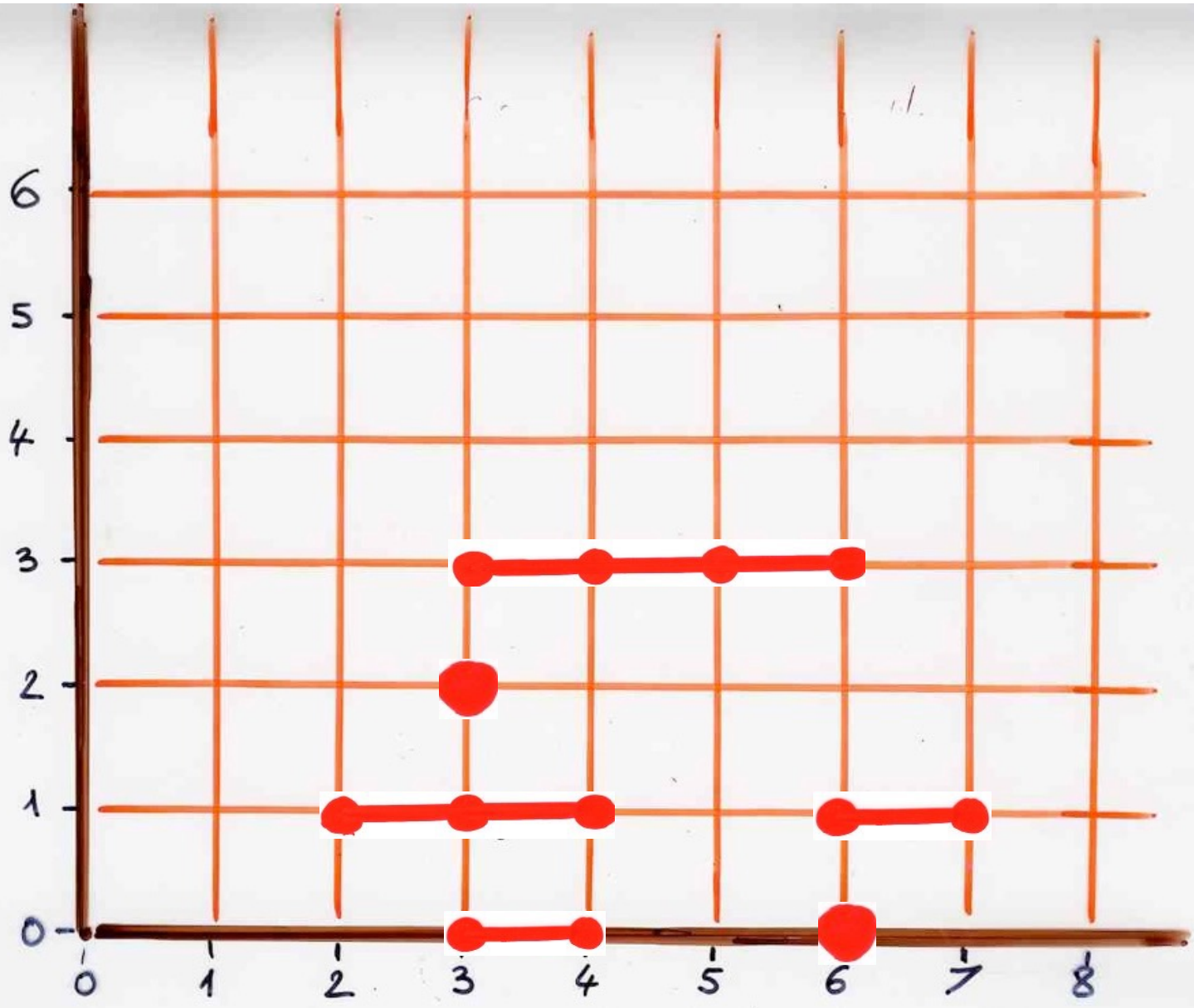
$H(P, \mathcal{E})$

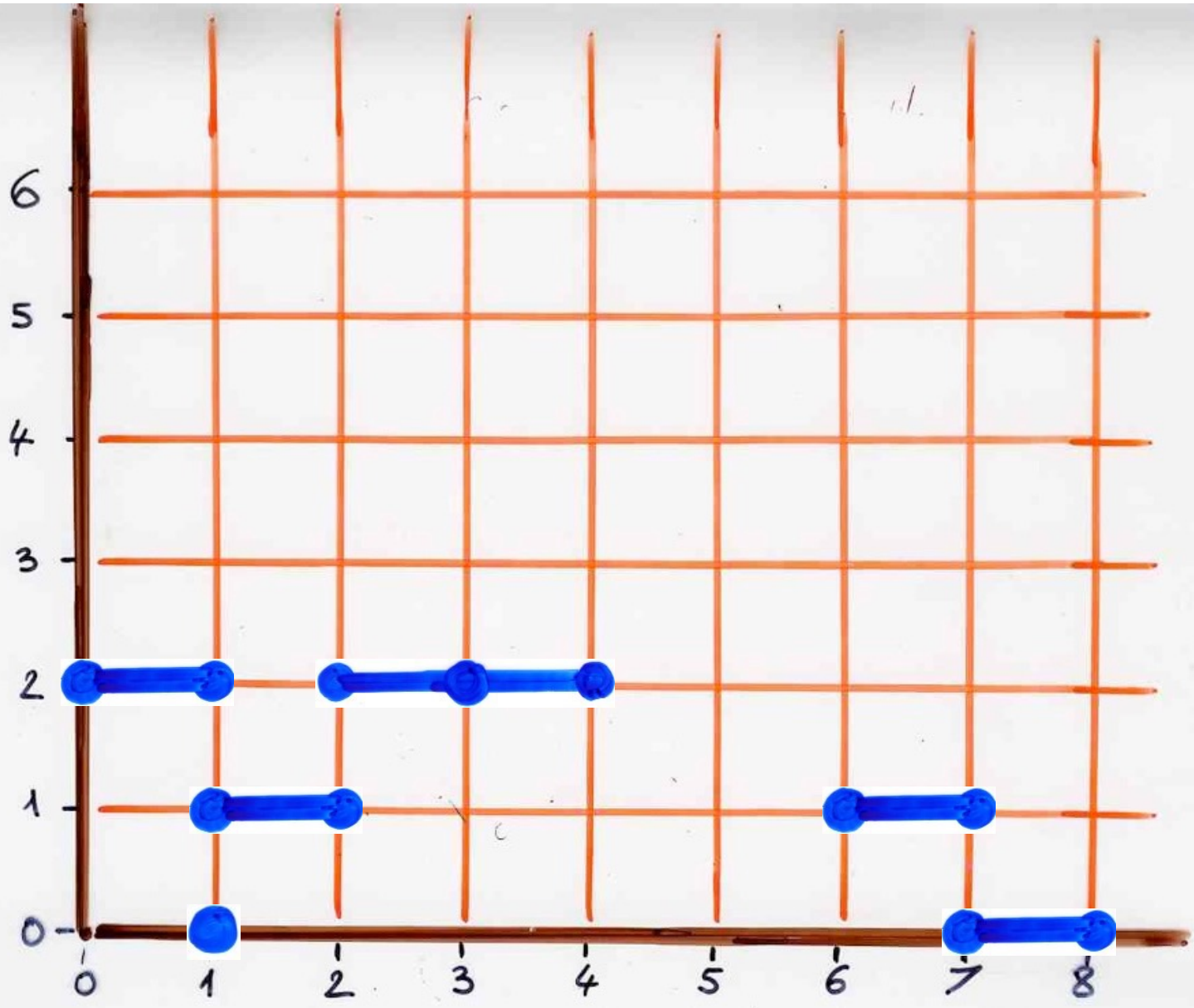
product of two heaps

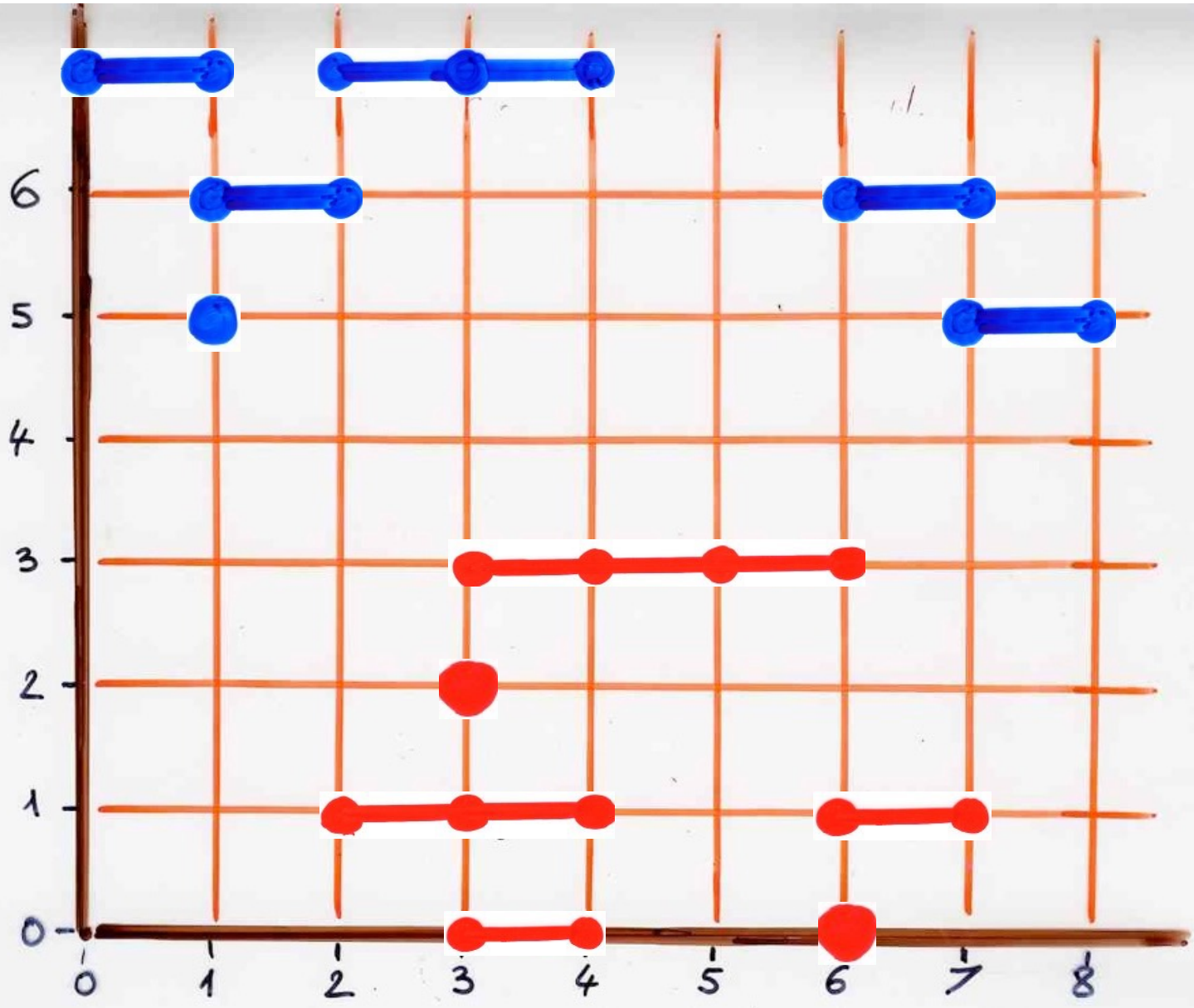
$E \circ F$

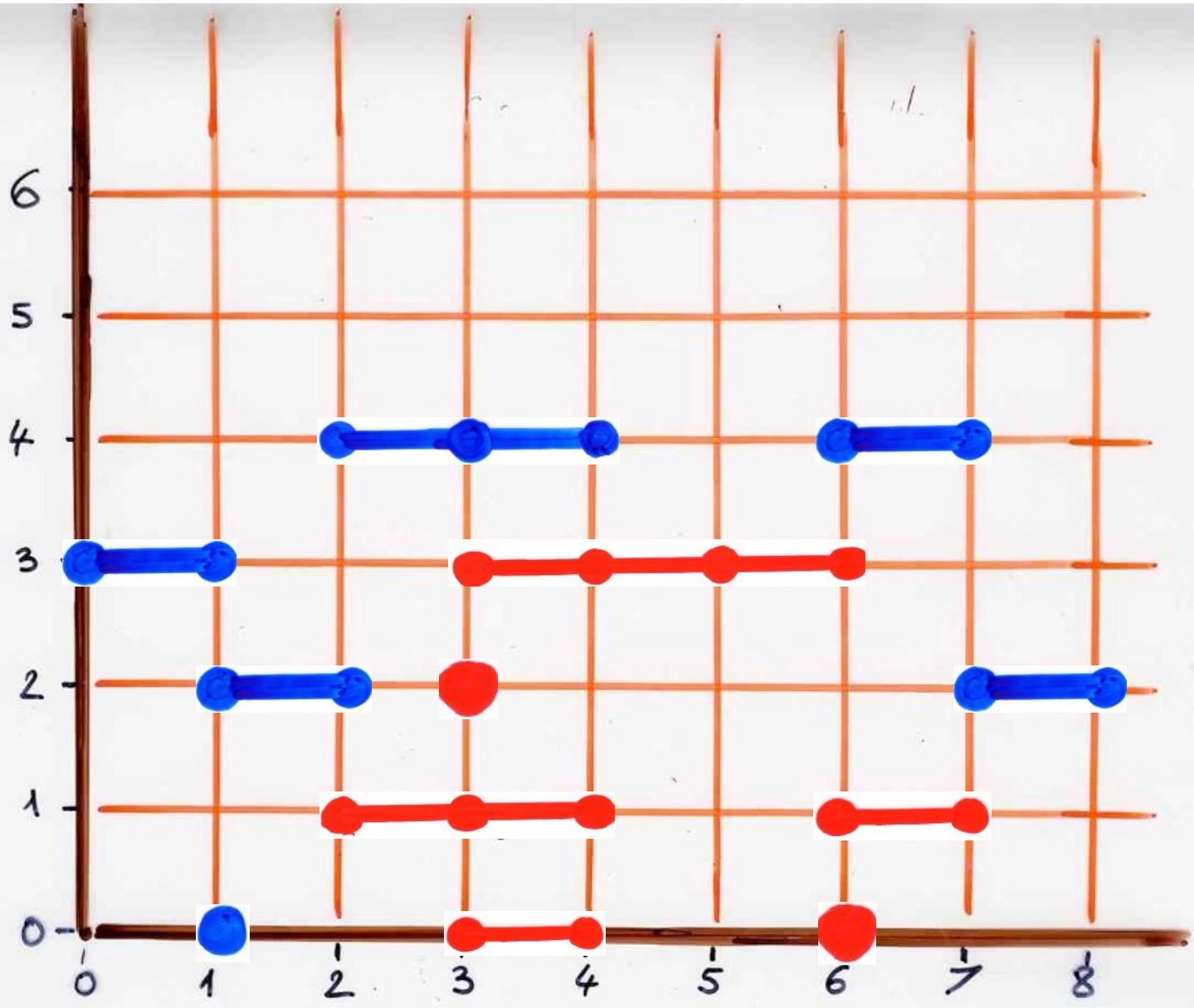


$$E \cap F = E \cap (E \cup T(F))$$









Equivalence
commutation monoids
and heaps monoids

$$\mathcal{P} \subseteq \text{Heap}(\mathcal{P}, \mathcal{E})$$

$$\alpha \longleftrightarrow \{(\alpha, 0)\}$$

$$\varphi : \mathcal{P}^* \longrightarrow \text{Heap}(\mathcal{P}, \mathcal{E})$$

$$w = \alpha_1 \alpha_2 \dots \alpha_n \longrightarrow \alpha_1 \odot \alpha_2 \odot \dots \odot \alpha_n$$

word heap

$$\mathcal{C} = \overline{\mathcal{E}}$$

commutation
relation

complementary
of the
dependency
relation

ex: heaps of dimers on \mathbb{N}

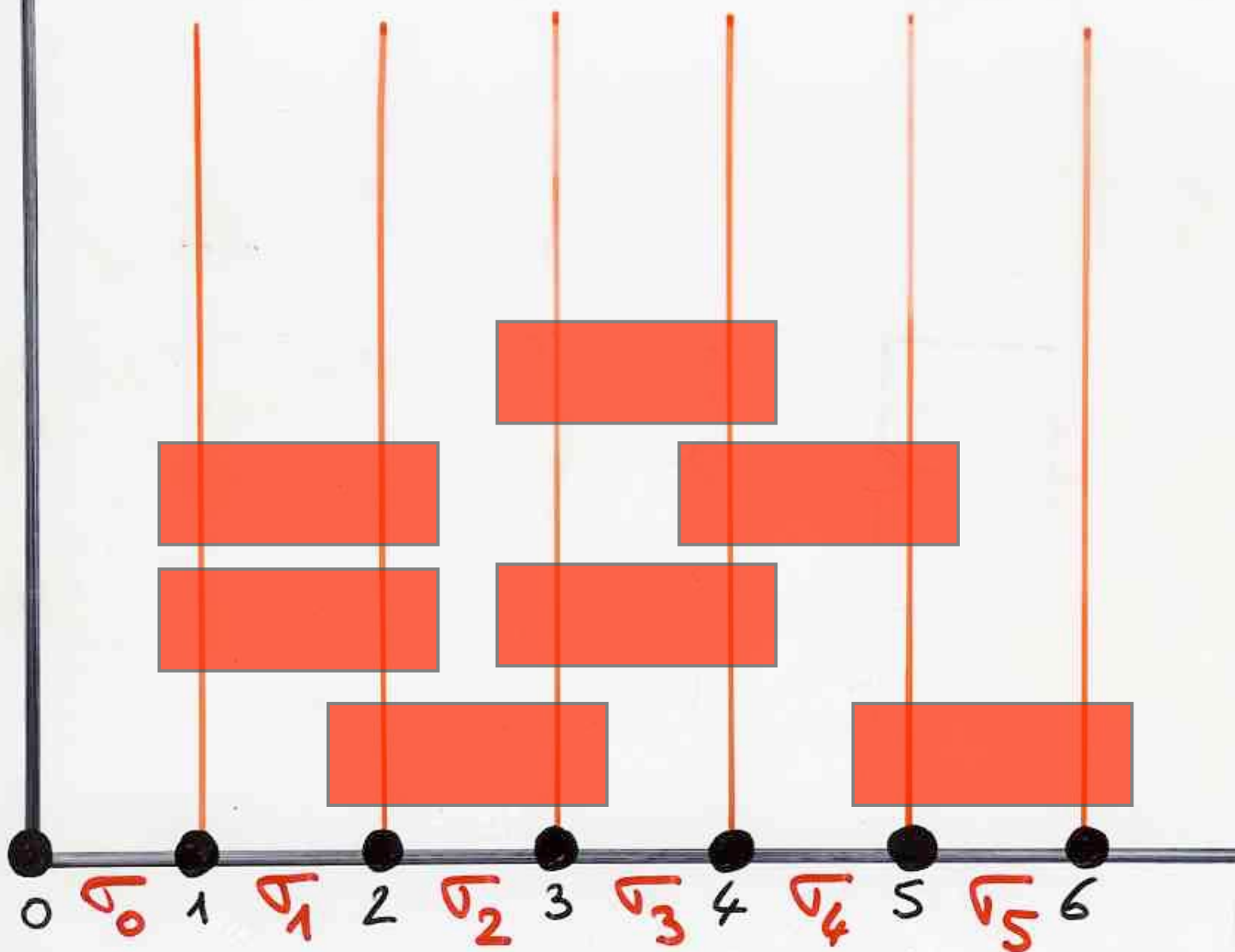
$$P = \{ [i, i+1] = \sigma_i, i \geq 0 \}$$

\mathcal{C}

\subset commutations

$$\sigma_i \sigma_j = \sigma_j \sigma_i \text{ iff } |i-j| \geq 2$$

$$W = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



$$P \subseteq \text{Heap}(P, \mathcal{E})$$

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$$w = \alpha_1 \alpha_2 \dots \alpha_n \longrightarrow \alpha_1 \odot \alpha_2 \odot \dots \odot \alpha_n$$

word heap

Lemma 1

$$u \equiv_C v \Rightarrow \varphi(u) = \varphi(v)$$

Lemma 2

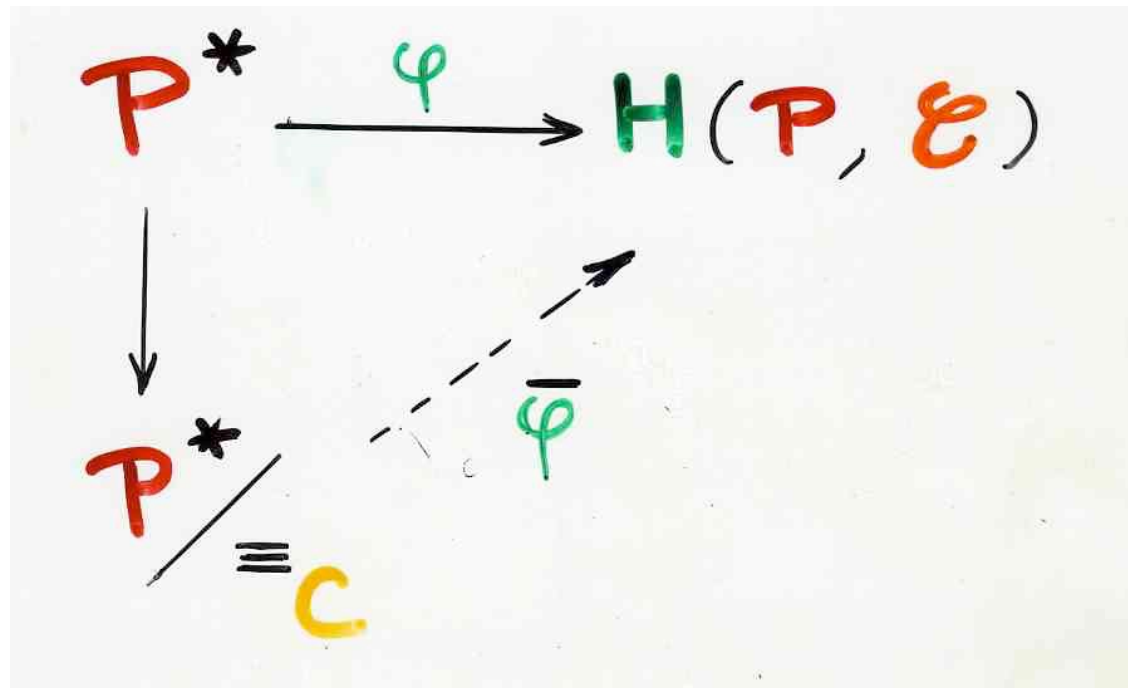
$$\varphi(u) = \varphi(v) \Rightarrow u \equiv_C v$$

C
commutation
relation

$\overline{\mathcal{E}}$
complementary
of the
dependency
relation

Definition

$$\overline{\varphi}([u]) = \varphi(u)$$



Proposition

$\overline{\varphi}$

is an isomorphism
of monoids

Heap (P, \mathcal{E})

\cong

$P^* / \equiv C$

heaps
monoid

commutation
monoid

$C = \overline{\mathcal{E}}$

complementary
relation

Proofs of Lemma 1, 2
and Proposition

$$P \subseteq \text{Heap}(P, \mathcal{E})$$

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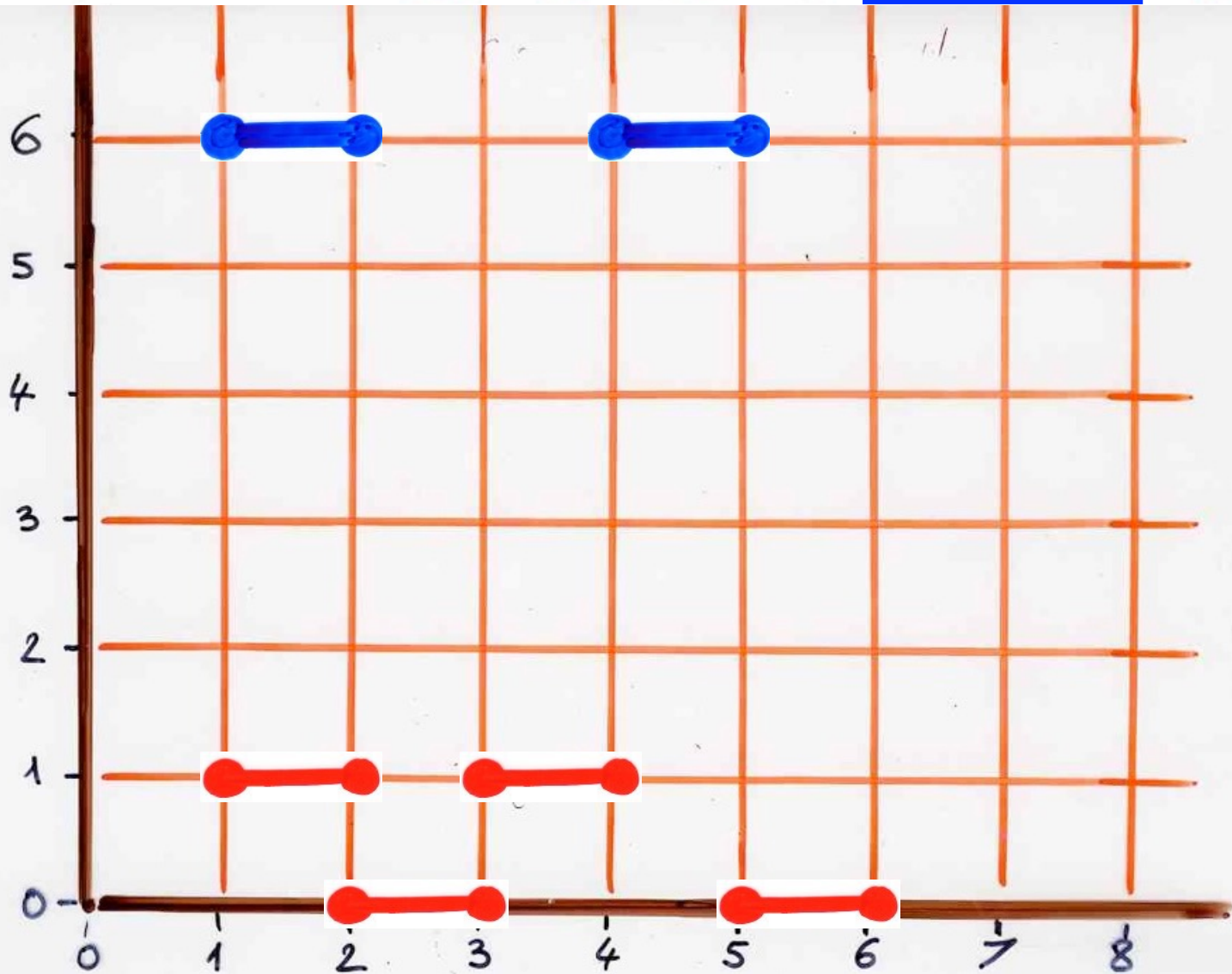
Lemma 1

$$u \equiv_C v \implies \varphi(u) = \varphi(v)$$

Proof: obvious

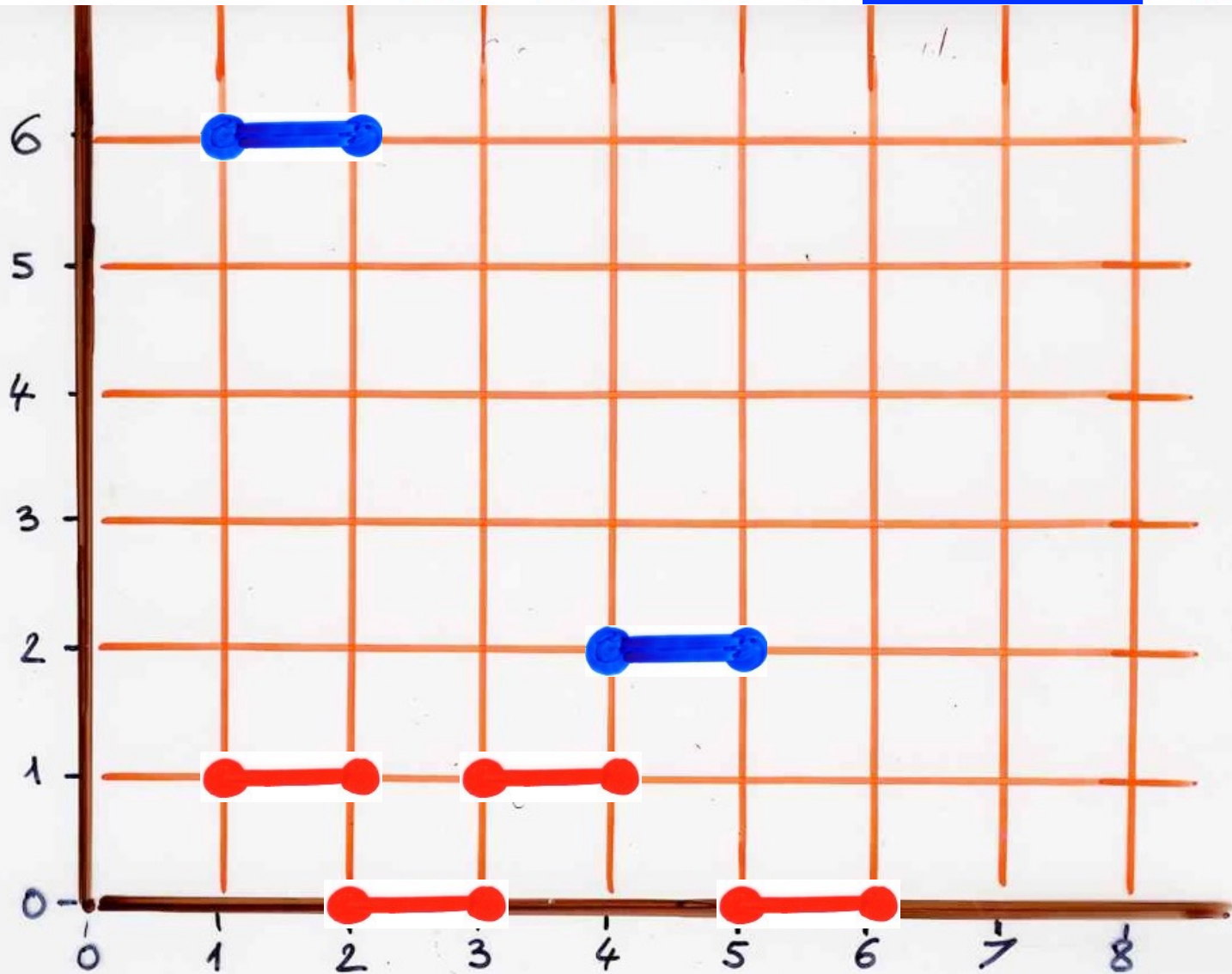
example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



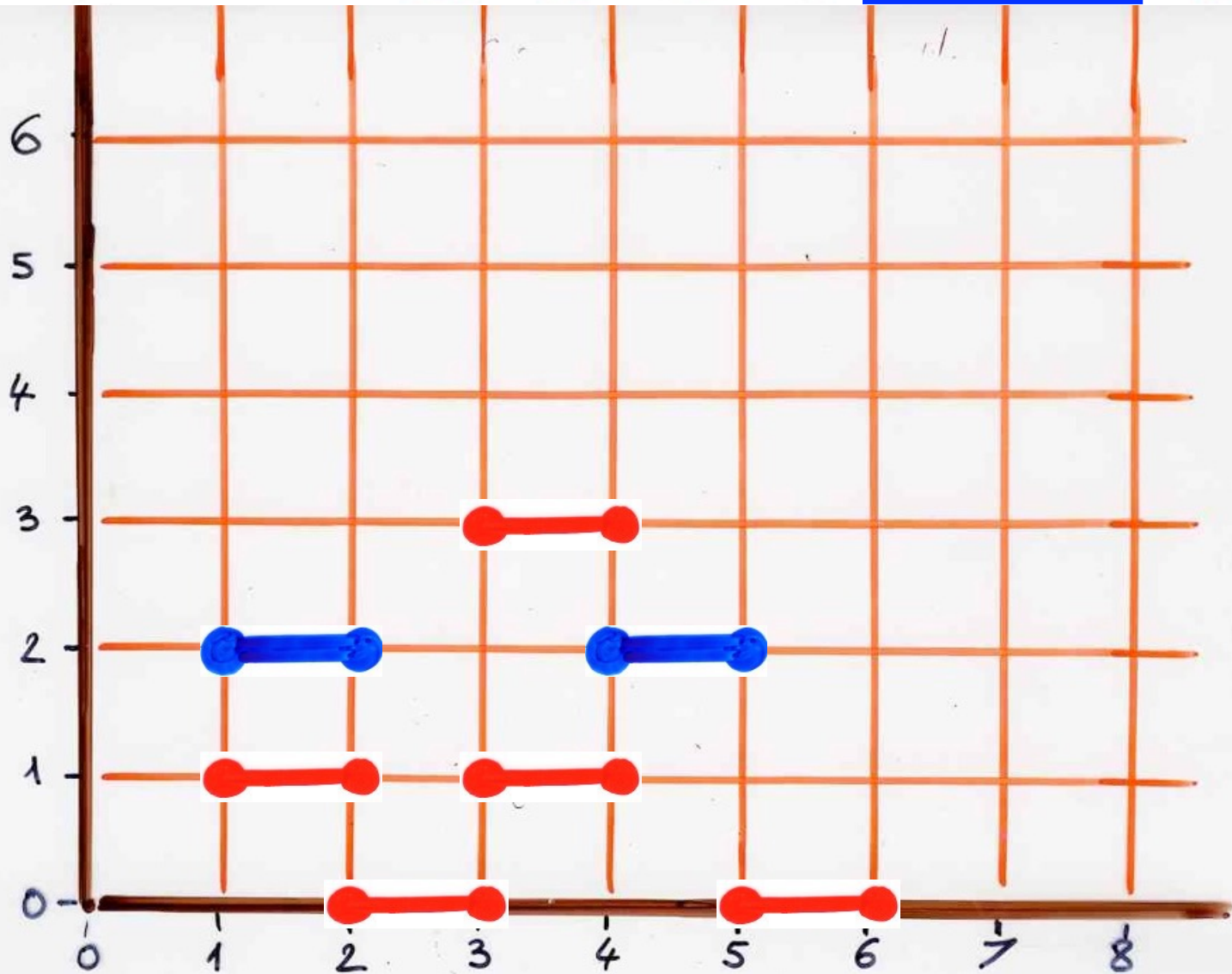
example

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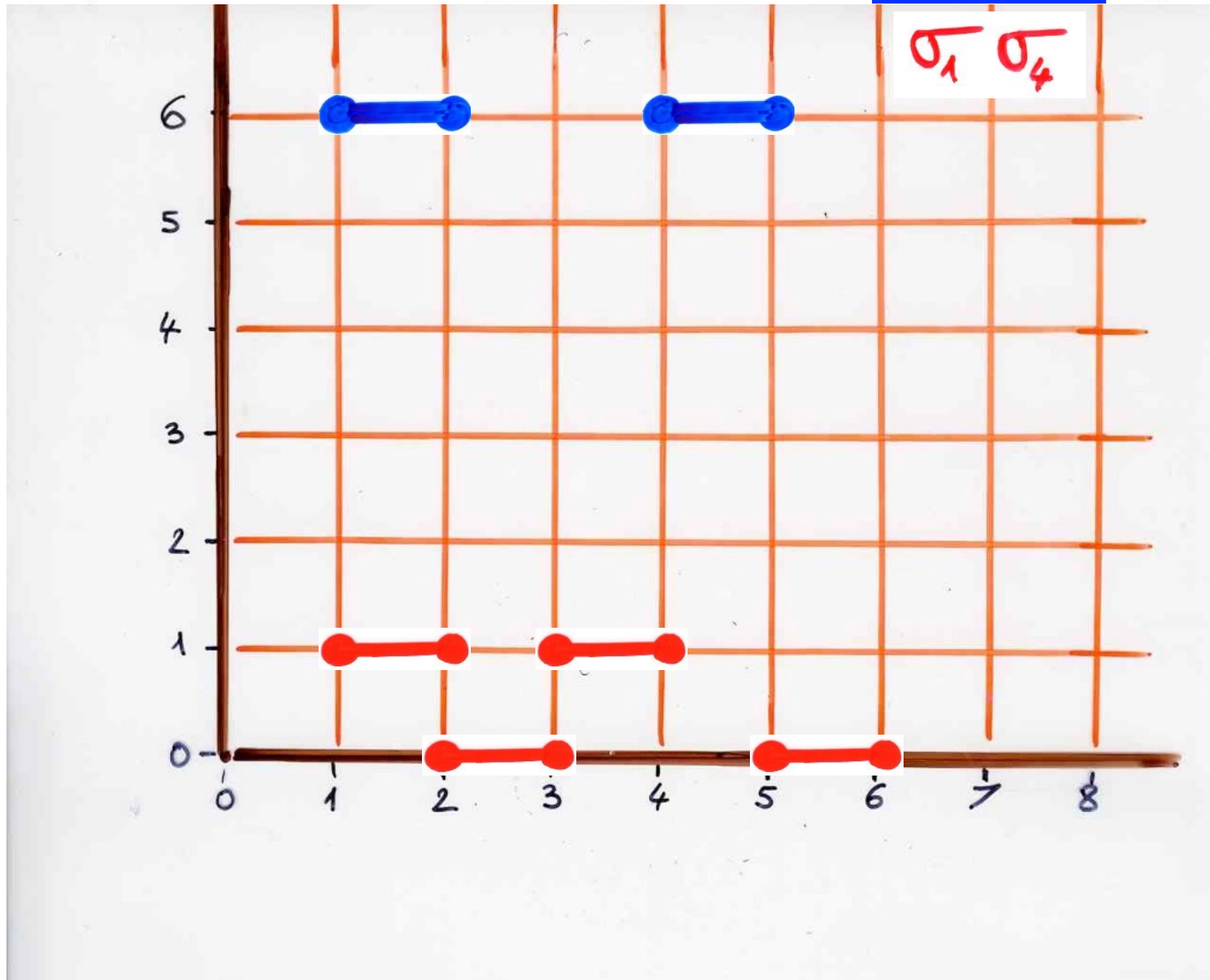
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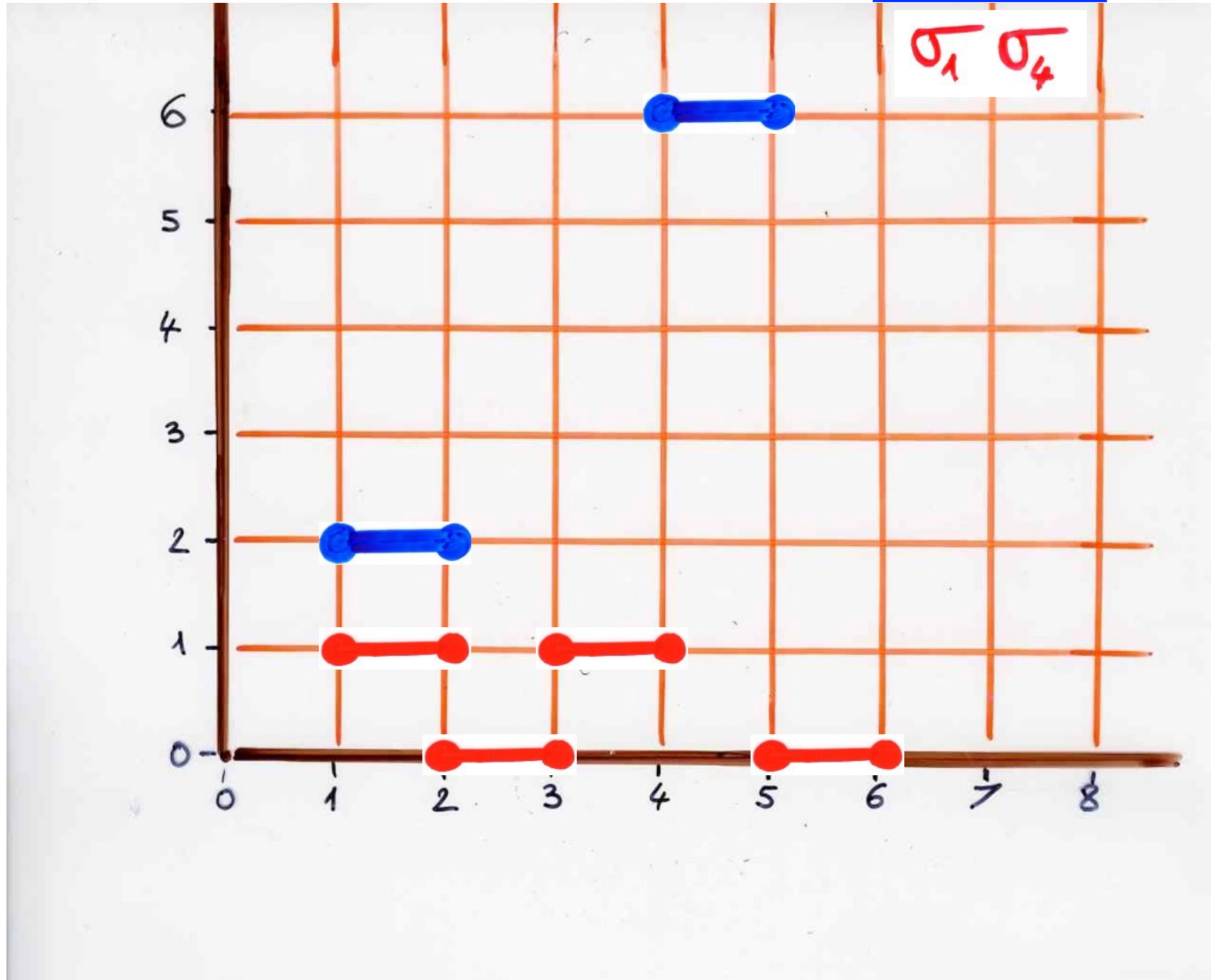
example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



example

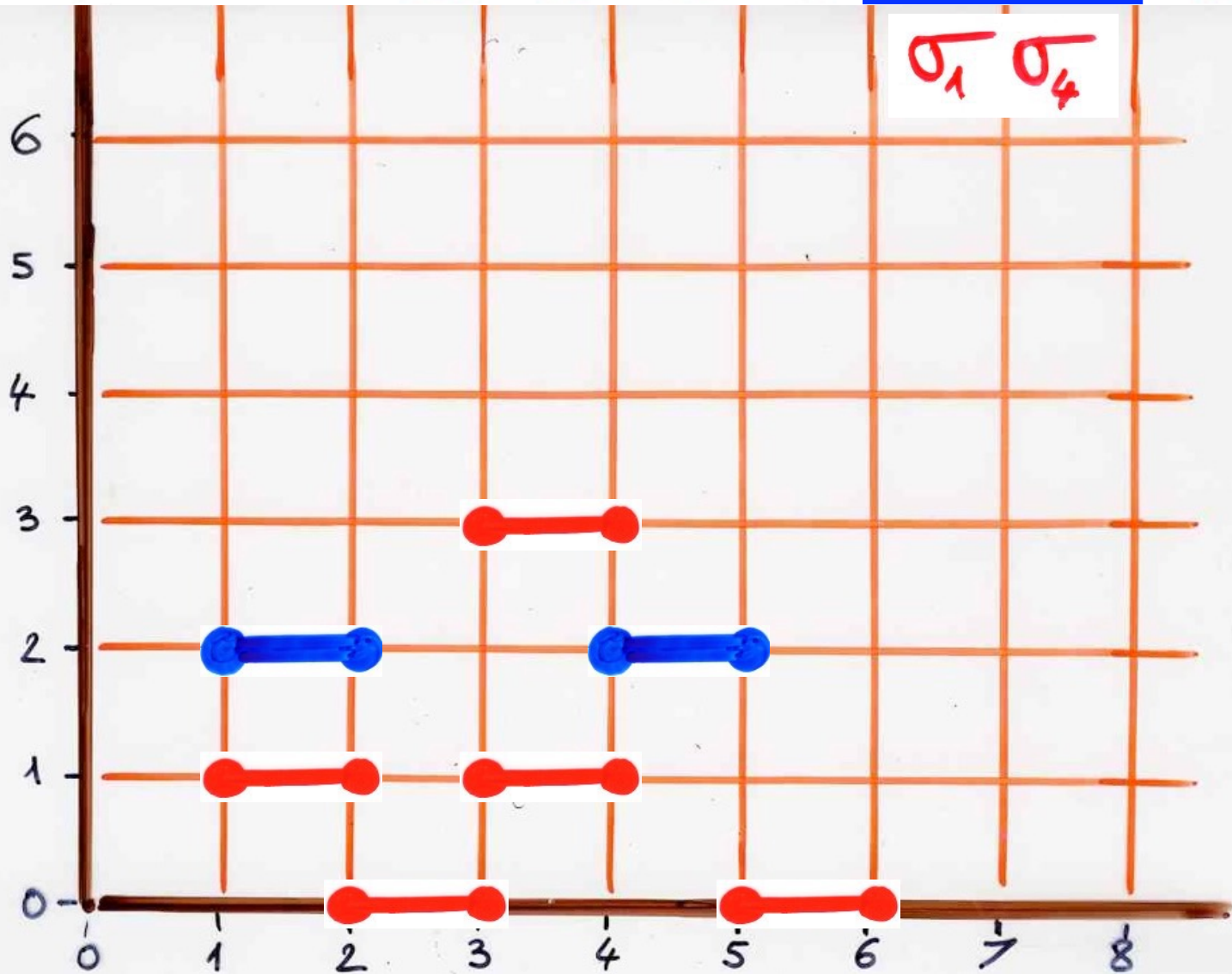
$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$\sigma_1 \sigma_4$$



$$\mathcal{P} \subseteq \text{Heap}(\mathcal{P}, \mathcal{E})$$

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$$\mathcal{C} = \overline{\mathcal{E}}$$

commutation
relation

complementary
of the
dependency
relation

Lemma 2

$$\varphi(u) = \varphi(v) \Rightarrow u \equiv_{\mathcal{C}} v$$

proof of Lemma 2 with

Cartier - Foata normal form

Cartier-Foata normal form

Lemma Every element $[w] \in L(A, C)$ has a unique decomposition into blocks

$$[w] = [w_1] [w_2] \dots [w_r]$$

- where each w_i is a word where the letters "commute" two by two $(y C z)$
- for every letter z of the $(j+1)^{\text{th}}$ block, there exist a letter y of the j^{th} block such that ~~$y C z$~~ (i.e. does not "commute")

in particular all the letters of each w_i are distinct (C antireflexive)

Proof Let F_1 the set of letters y of w
such that $w \equiv_C y v_1$

(i.e. applying commutations y can be put
as the first letter)

- the letters of F_1 "commute" 2 by 2

$$w = \dots y \dots z \dots \quad y, z \in F_1$$

$[w_1]$ is the equivalence class of the product
of the letters of F_1

- Let F_2 be the set of letters y of w
such that $w \equiv_C w_1 y v_2$

$[w_2]$ is the equivalence class of the product
of the letters of F_2 and these
letters "commute" 2 by 2

- ... etc... for F_2, \dots, F_r

example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$\equiv \sigma_2 \sigma_5 / \sigma_3 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$\equiv \sigma_2 \sigma_5 / \sigma_3 \sigma_1 / \sigma_4 \sigma_1 \sigma_3$$

$$\equiv \sigma_2 \sigma_5 / \sigma_1 \sigma_3 / \sigma_1 \sigma_4 / \sigma_3$$

unicity

... obvious

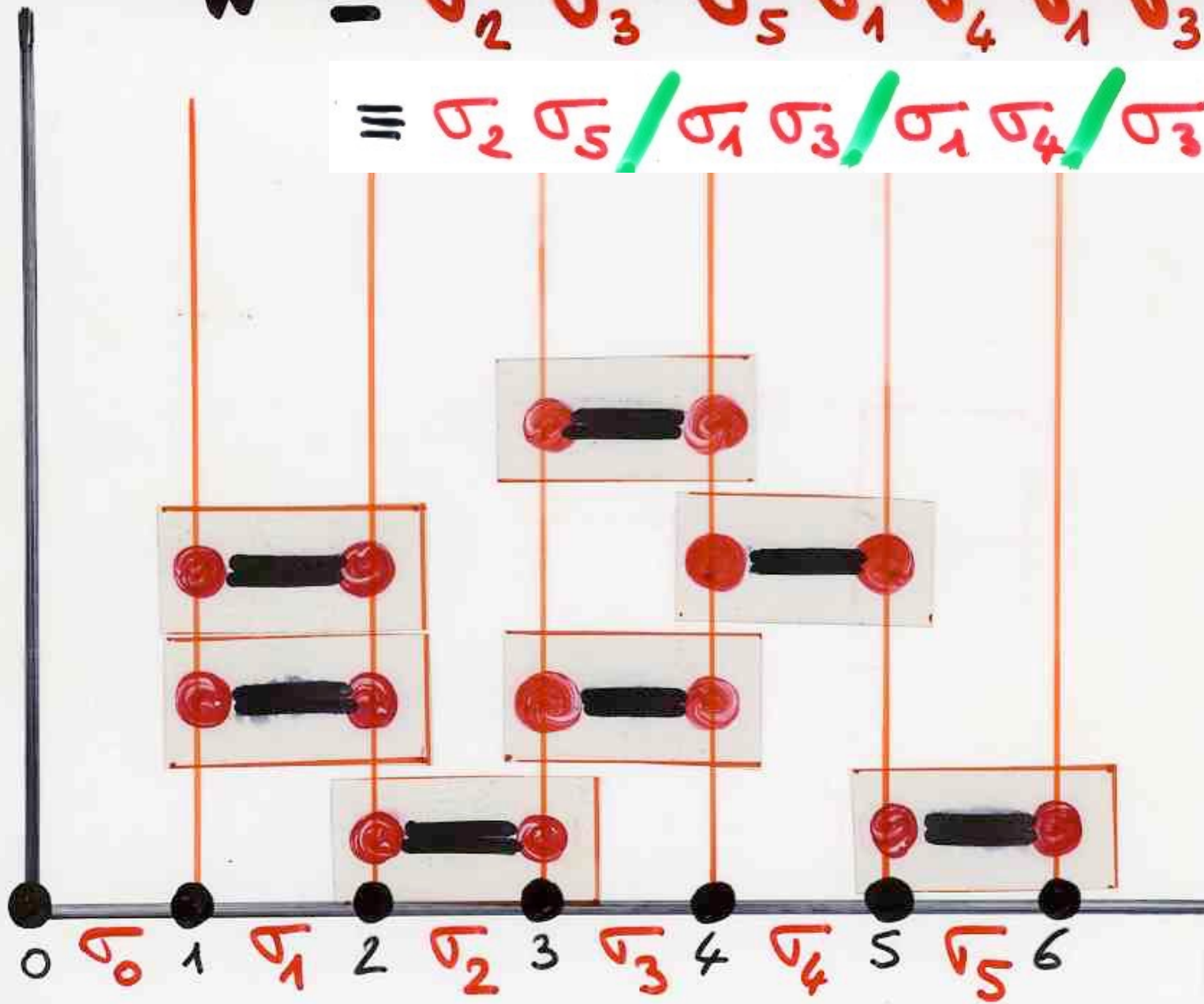
Lemma

Let $[w] = [w_1] \dots [w_r]$
be the Cartier-Foata normal form
of $[w] \in L(A, C)$.

Each block $[w_i]$ corresponds to the
elements of the heap $\varphi(w)$ located
at level $i-1$.

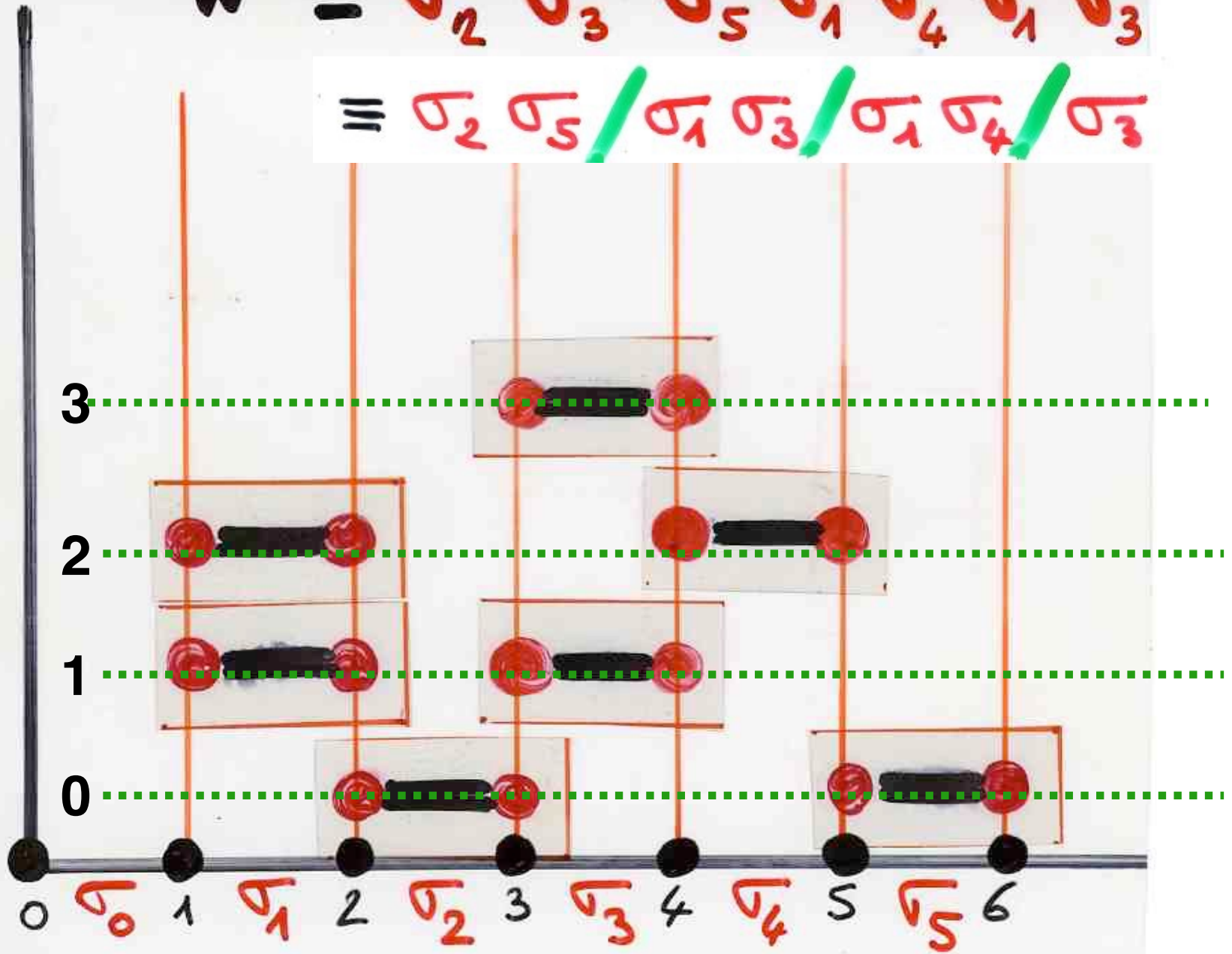
$$W = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$\equiv \sigma_2 \sigma_5 / \sigma_1 \sigma_3 / \sigma_1 \sigma_4 / \sigma_3$$



$$W = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$\equiv \sigma_2 \sigma_5 / \sigma_1 \sigma_3 / \sigma_1 \sigma_4 / \sigma_3$$



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word heap

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commutation
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Lemma 2

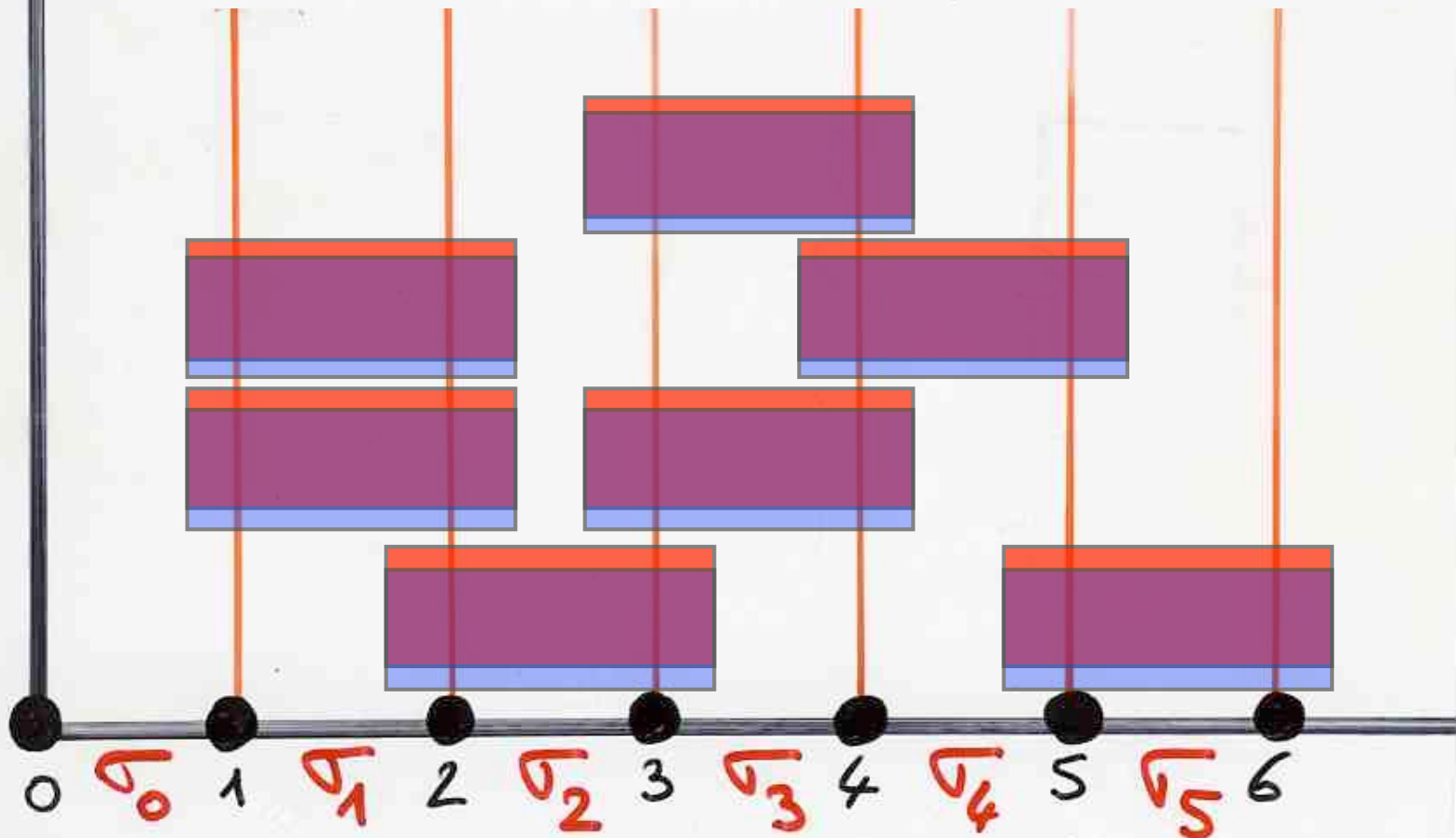
$$\varphi(u) = \varphi(v) \Rightarrow u \equiv_{\mathcal{C}} v$$

Proof of Lemma 2

$$W = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$= \sigma_2 \sigma_5 / \sigma_1 \sigma_3 / \sigma_1 \sigma_4 / \sigma_3$$

$$W = \sigma_5 \sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_4 \sigma_3$$



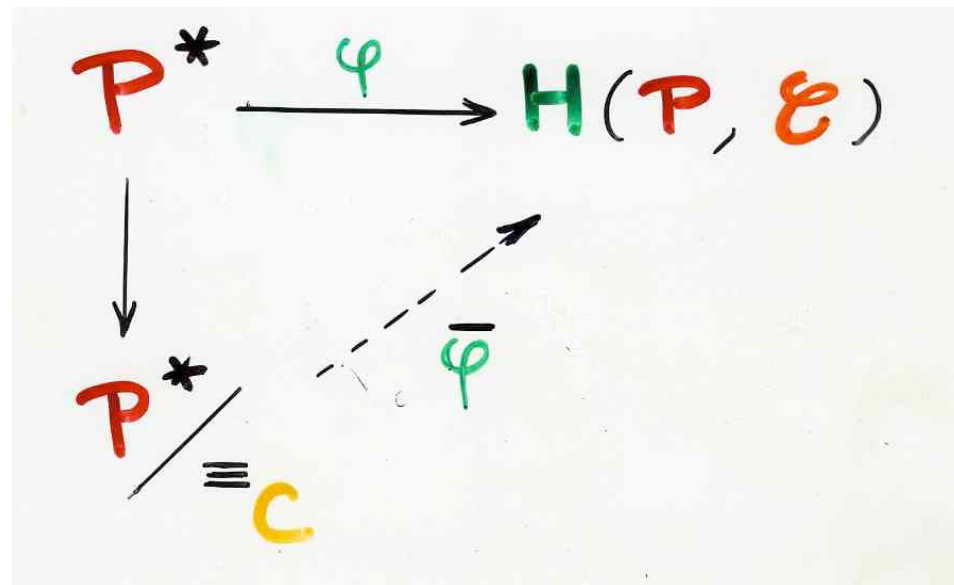
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Lemma 2

$$\varphi(u) = \varphi(v) \Rightarrow u \equiv_C v$$

Definition $\bar{\varphi}([u]) = \varphi(u)$



Proposition

$\overline{\varphi}$

is

an isomorphism
of monoids

Heap $(P, \mathcal{E}) \cong$

heaps
monoid

$P^* / \equiv C$

commutation
monoid

$C = \overline{\mathcal{E}}$

complementary
relation

$[w]$

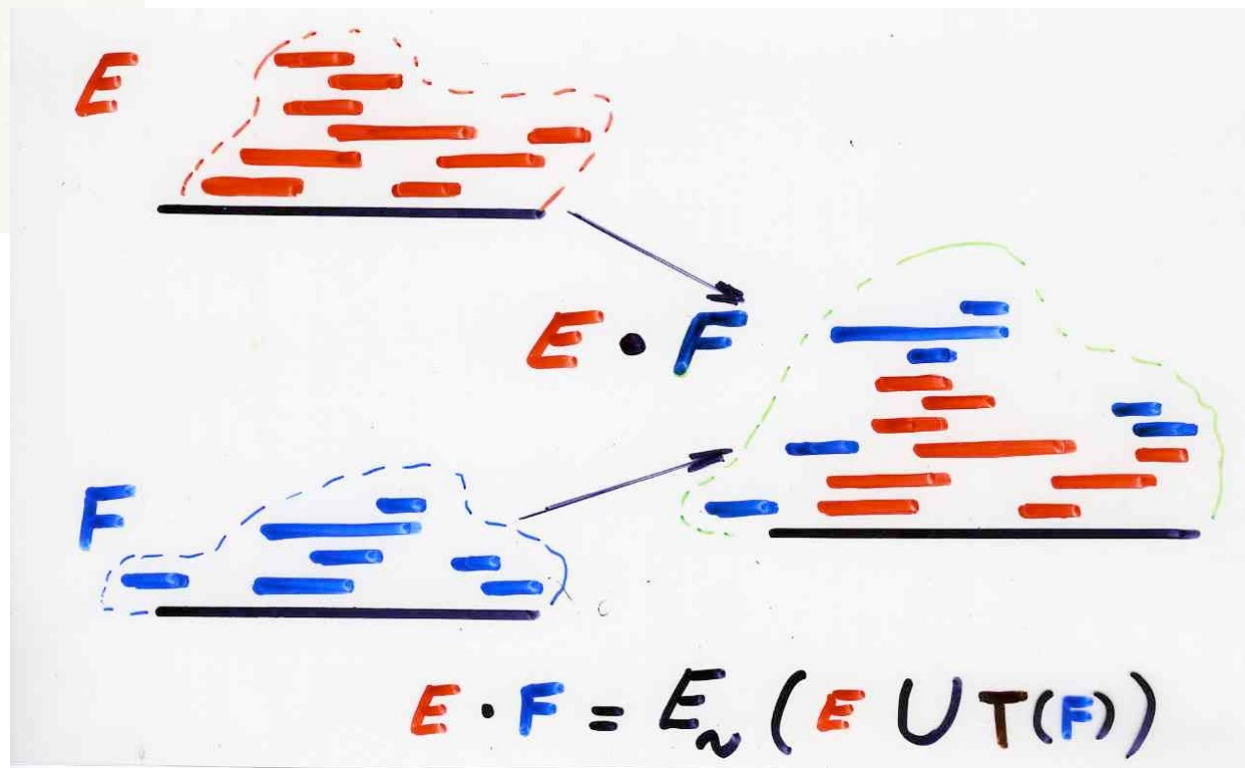
equivalence class
of the word $w \in A^*$

A^*/\equiv

- product in the commutation monoid

$$[u] \cdot [v] = [uv]$$

independent of the choices
of representants u and v



exercise using Cartier-Foata normal form

prove that the commutation monoid $L(A, C)$

is simplifiable, i.e.

$$uv = uv' \Rightarrow v = v'$$

$$uv = u'v \Rightarrow u = u'$$

lexicographic normal form

(« Knuth »)

lexicographic

normal

form

(Knuth)

$H(P, \mathcal{E})$

total order on
set

P
of

basic pieces

minimal

letter of a class $[w]$

$$[w] = [y x_1]$$

Cartier-Foata normal form

$$[w] = [w_1] [w_2] \dots [w_r]$$

minimal

letter : any letter of w_1

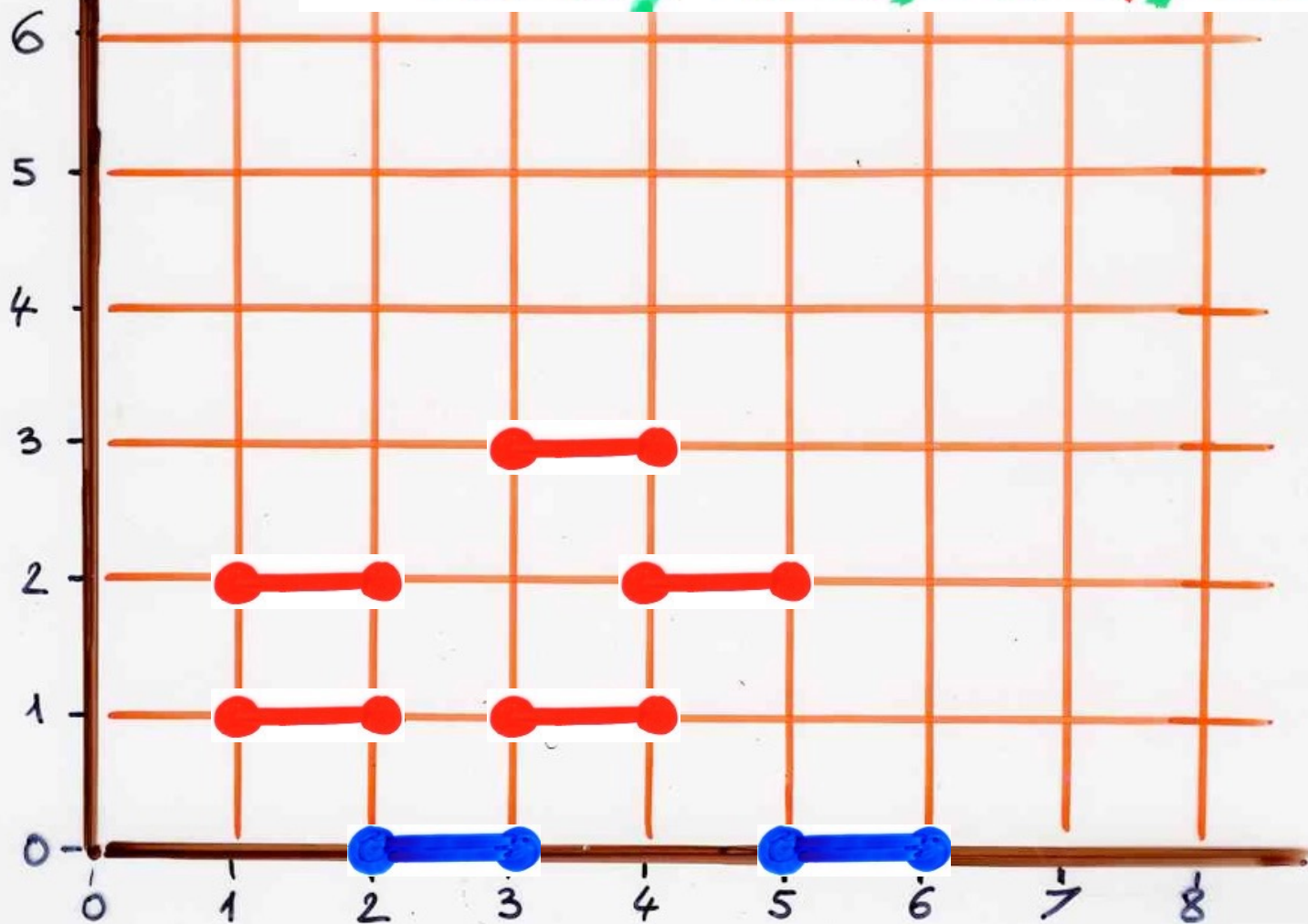


piece of the associated heap
at level 0

example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

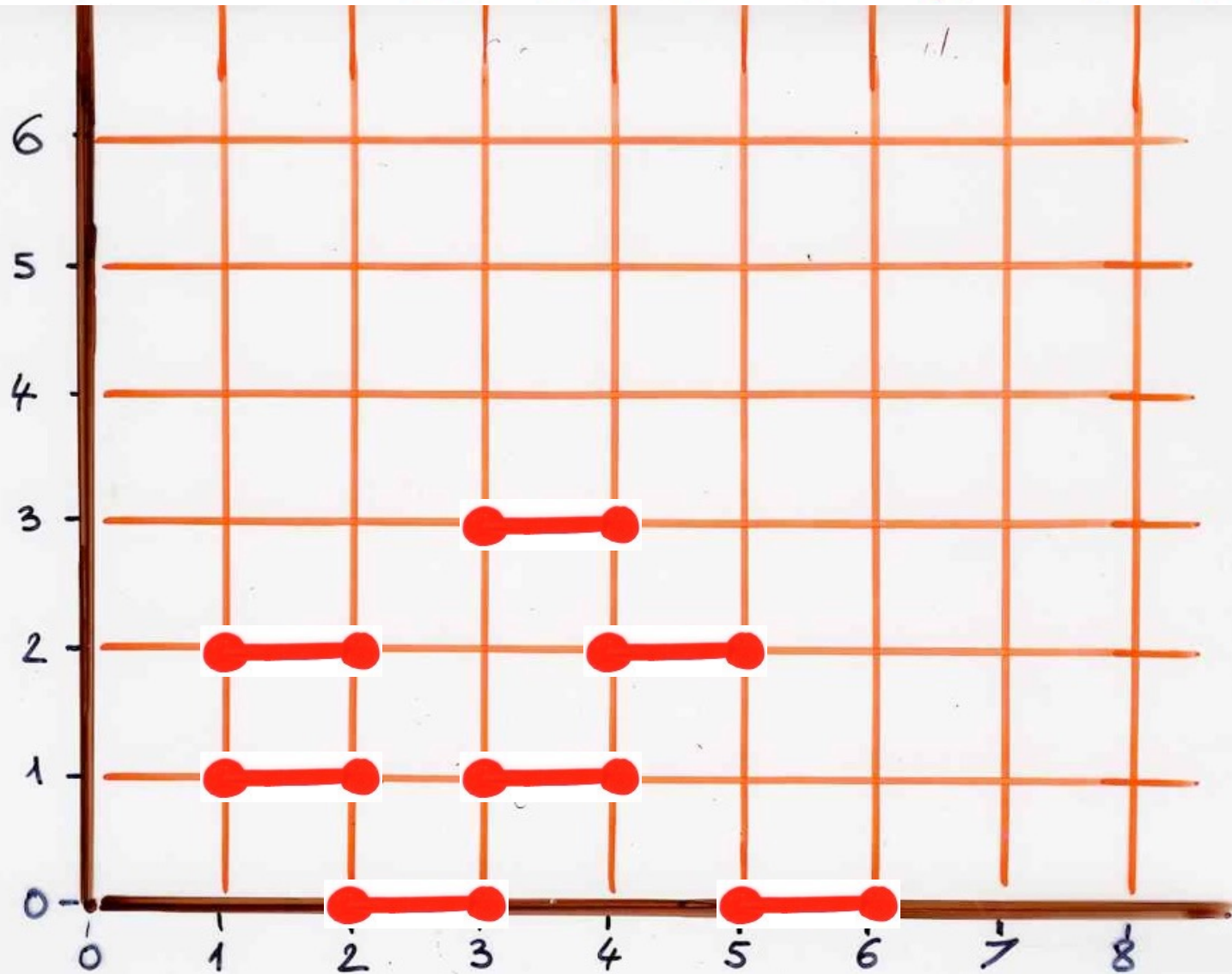
$$\equiv \sigma_2 \sigma_5 / \sigma_1 \sigma_3 / \sigma_1 \sigma_4 / \sigma_3$$



Lemma In the commutation class $[w]$, the smallest word $v = v_1 \dots v_n$ for the lexicographic order is obtained by taking v_1 the smallest minimal letter of $[w]$, $w = v_1 w_2$, then v_2 the smallest minimal letter of $[w_2]$, then \dots

example

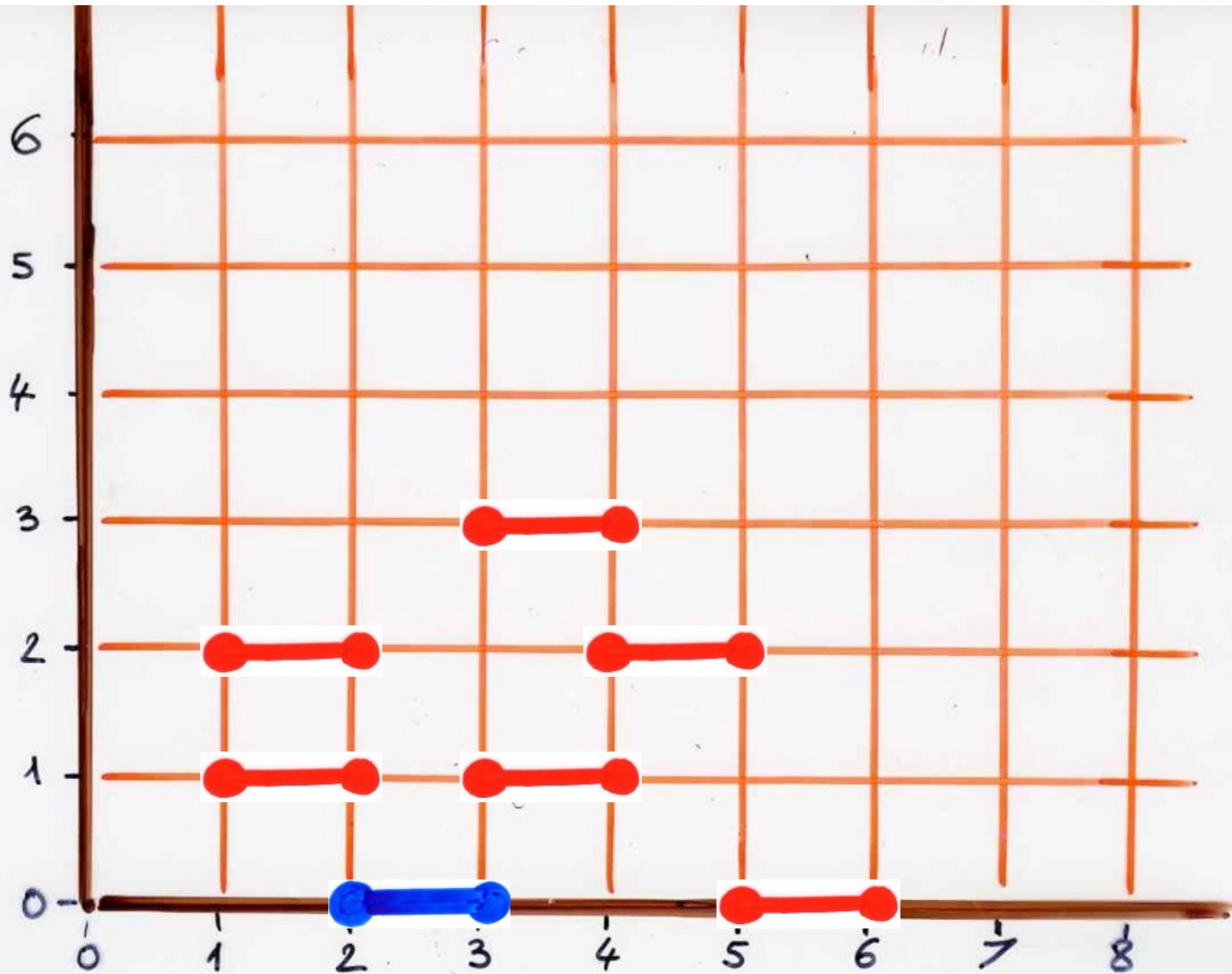
$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



$$\sigma_0 < \sigma_1 < \sigma_2 < \dots < \sigma_5$$

example

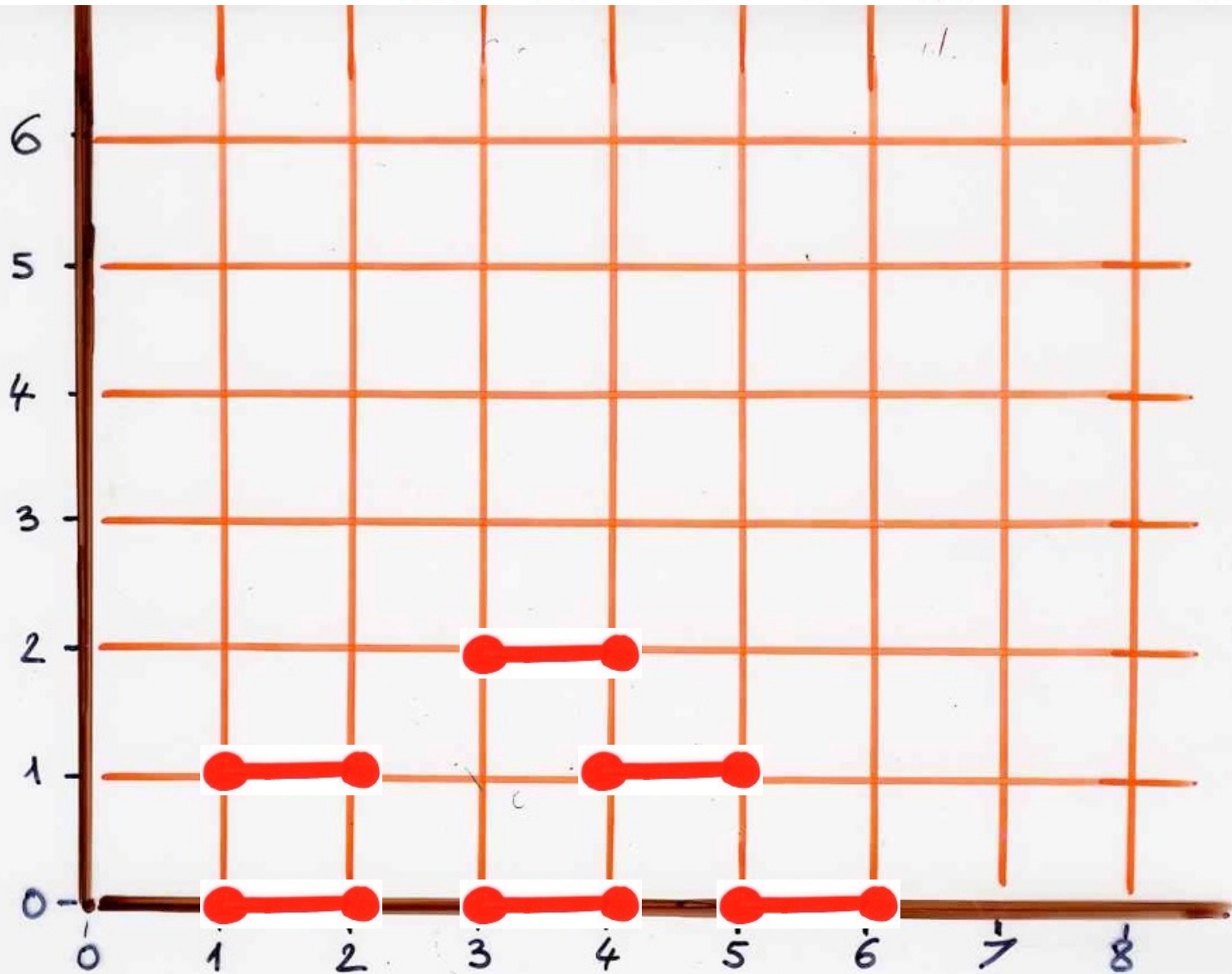
$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



σ_2

example

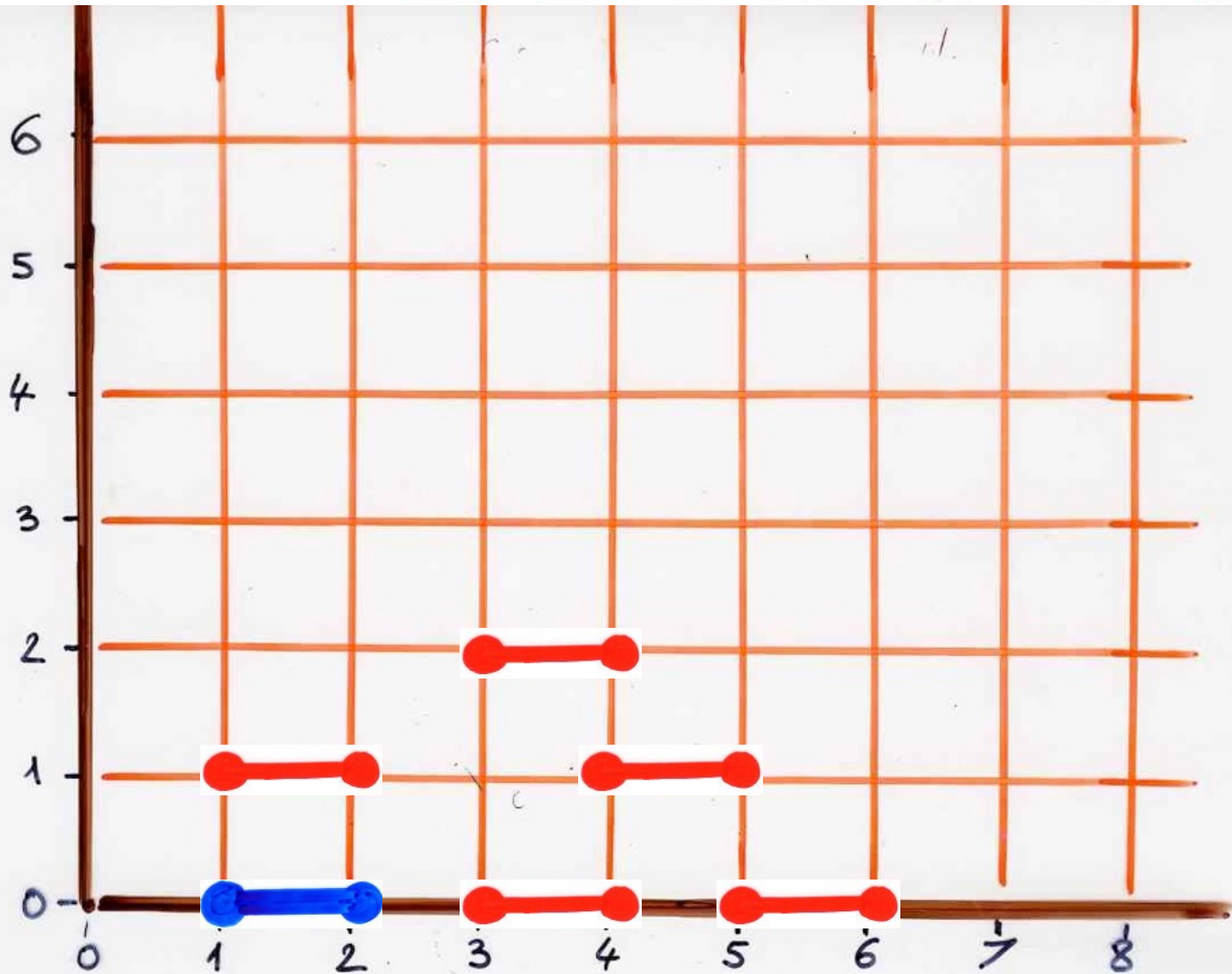
$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



σ_2

example

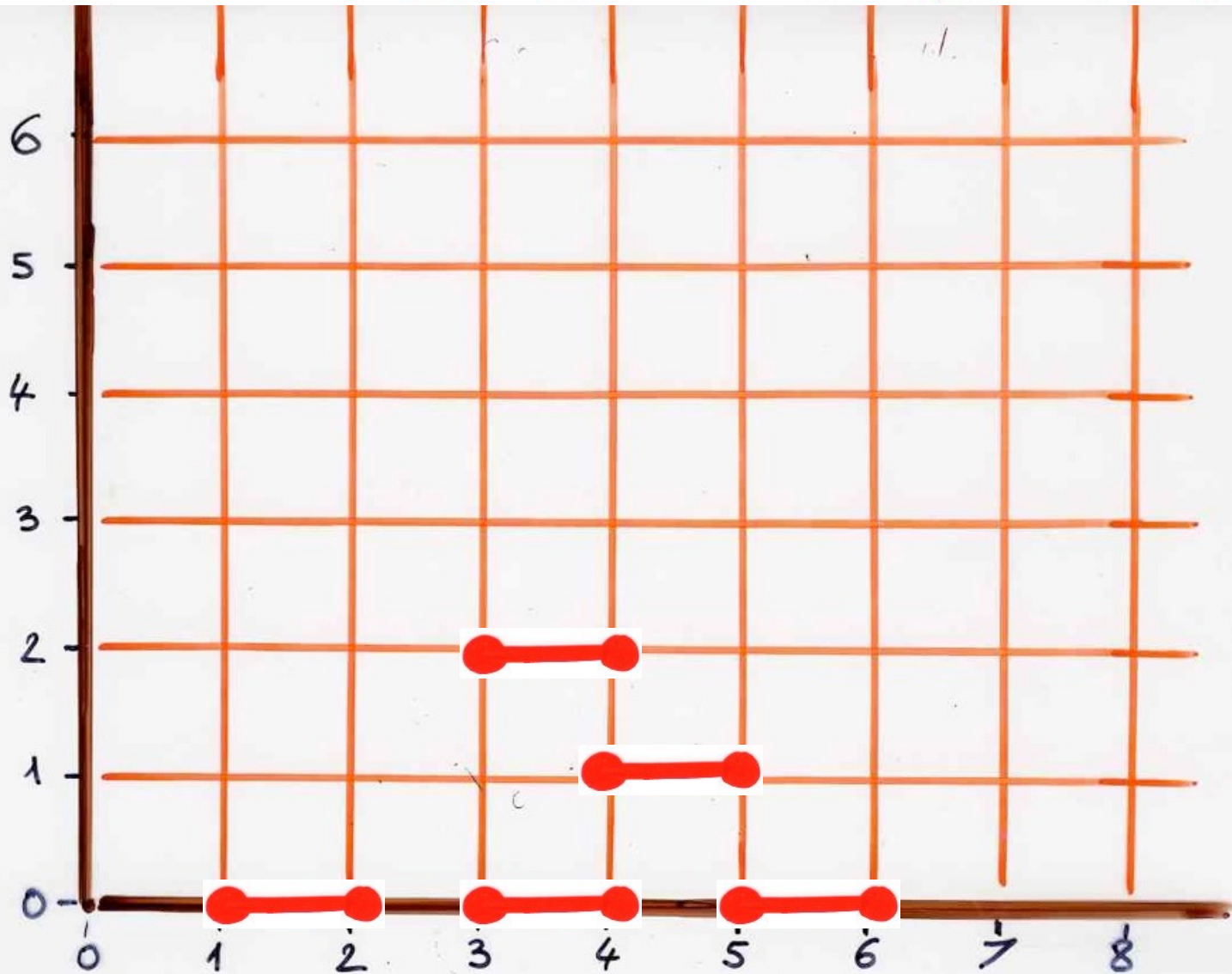
$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



$$\sigma_2 \sigma_1$$

example

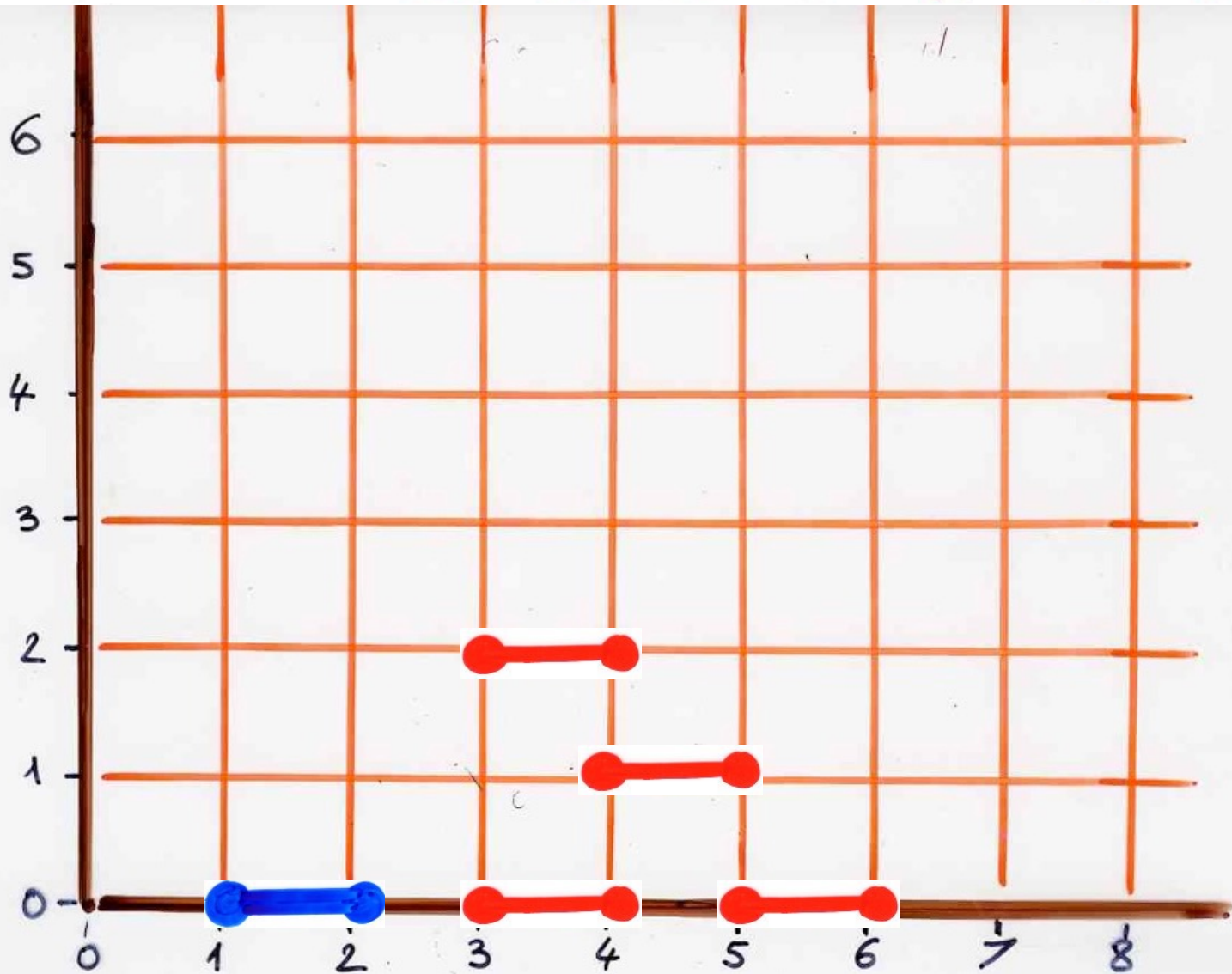
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$\sigma_2 \sigma_1$

example

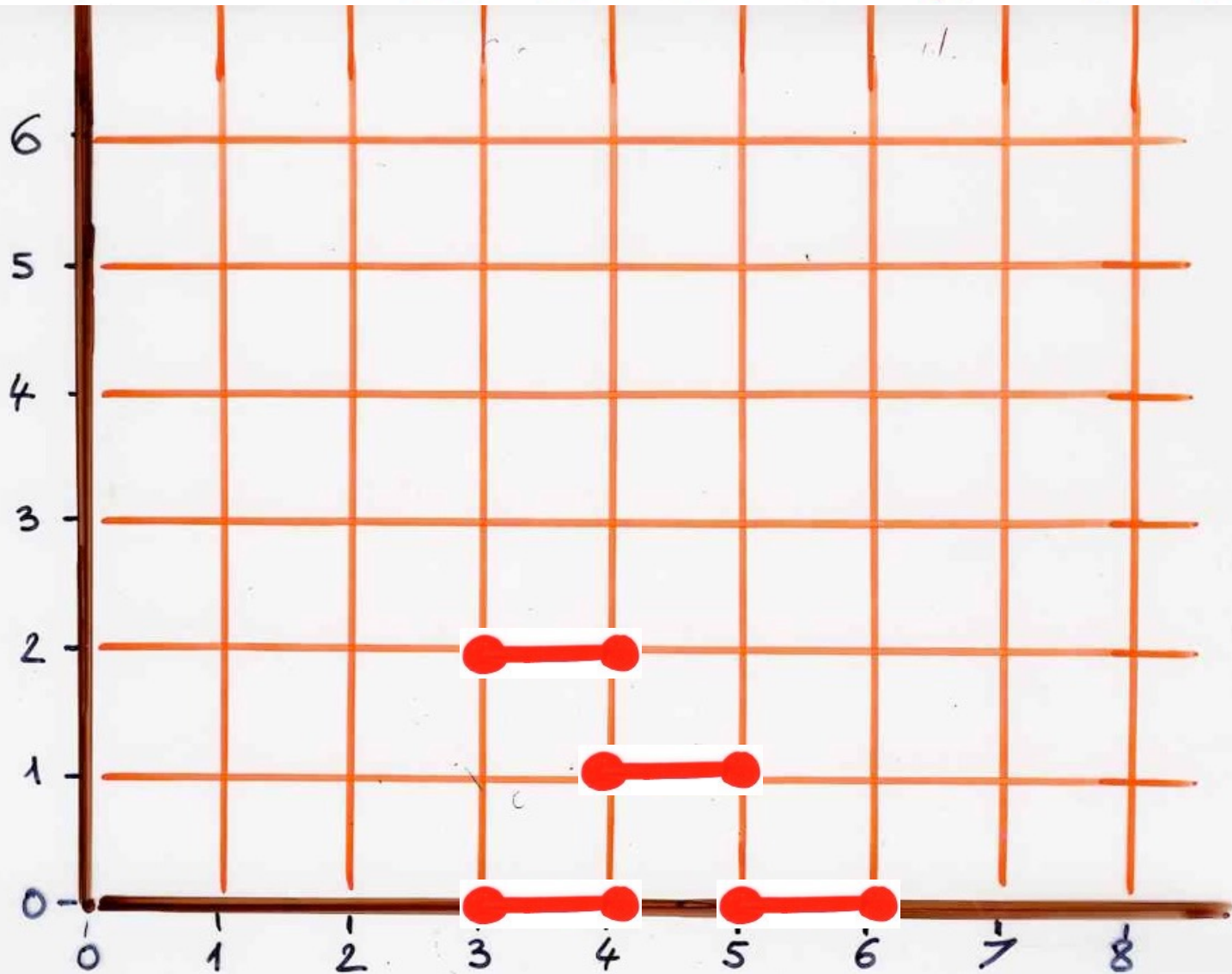
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$$\sigma_2 \sigma_1 \sigma_1$$

example

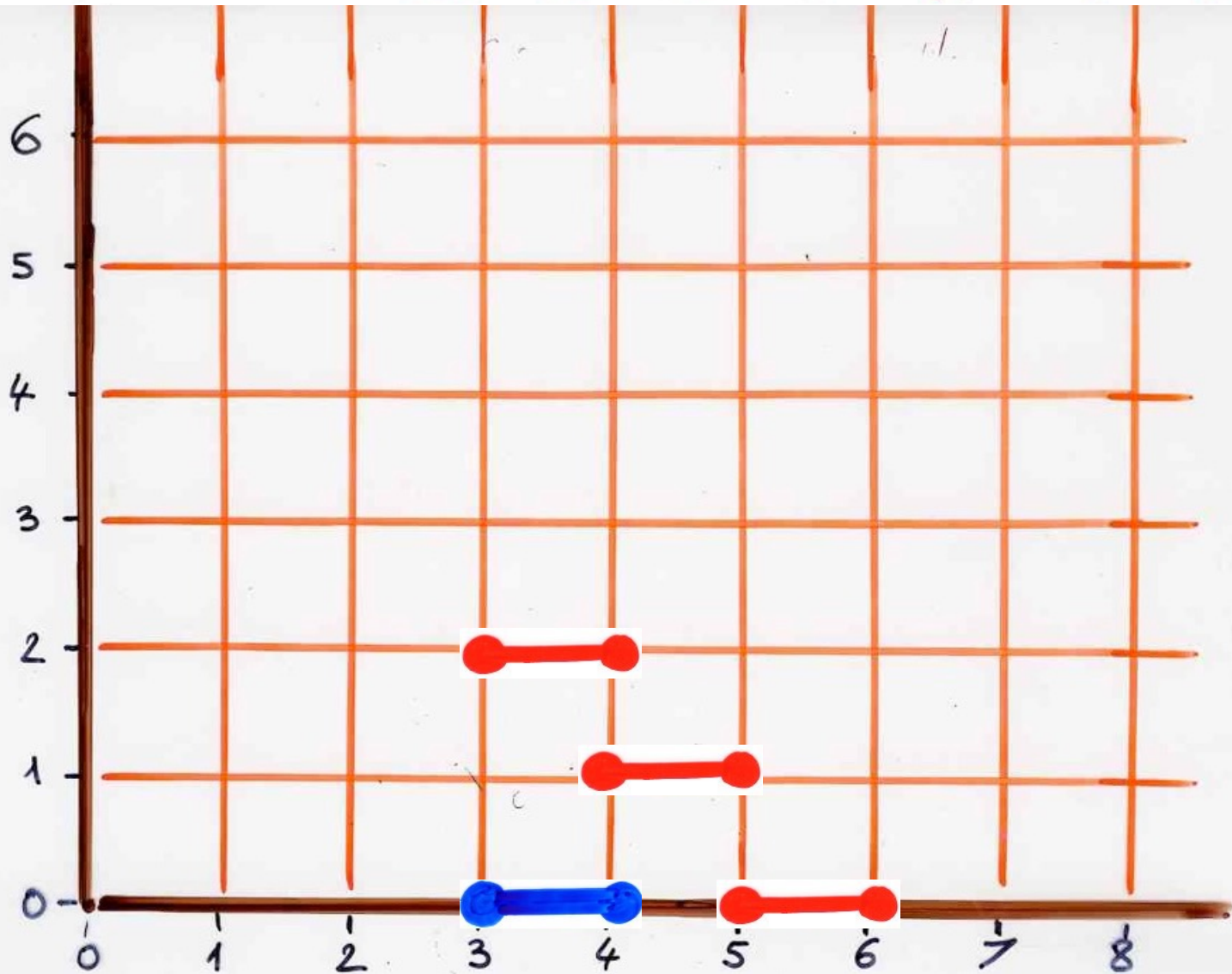
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$$\sigma_2 \sigma_1 \sigma_1$$

example

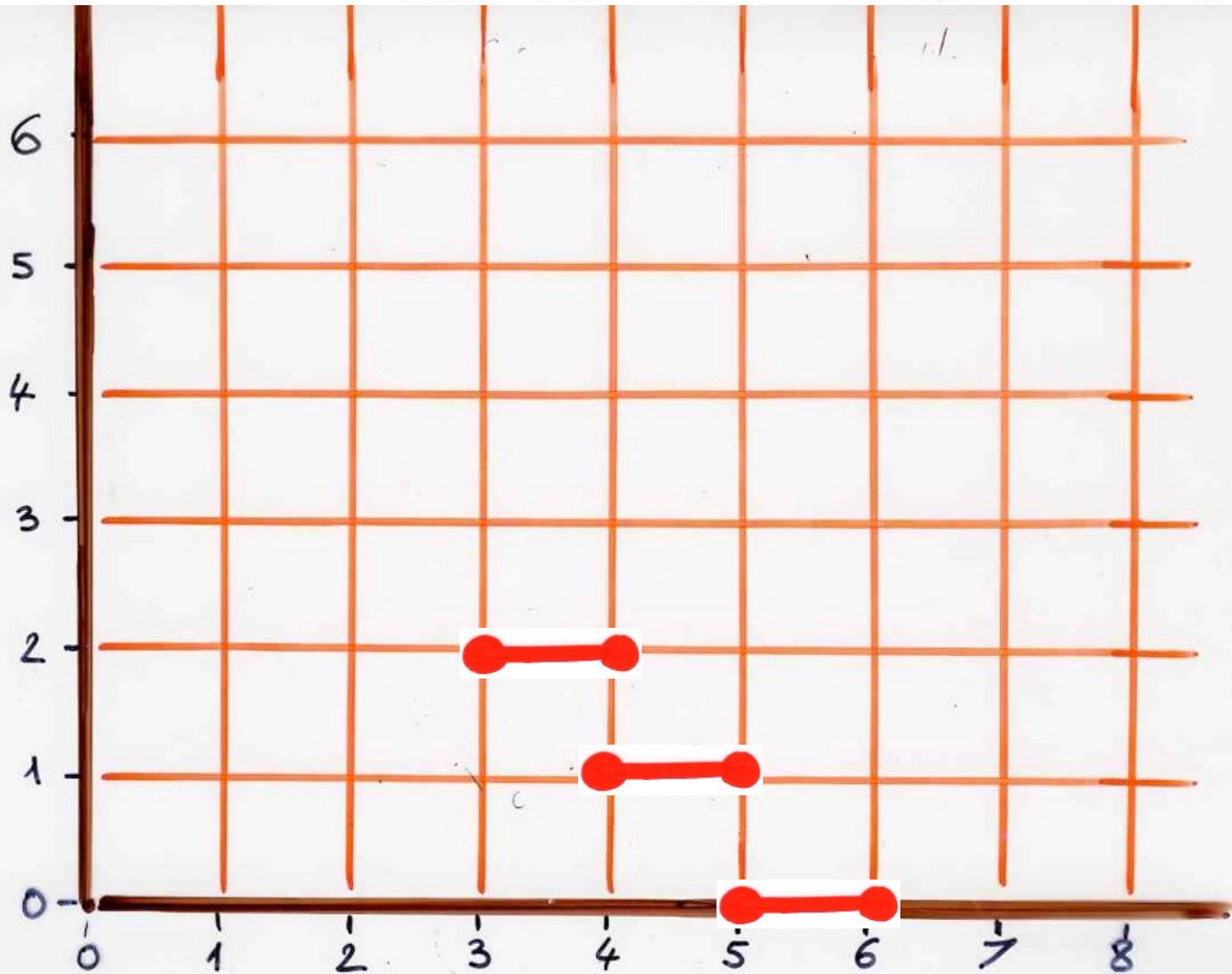
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3$$

example

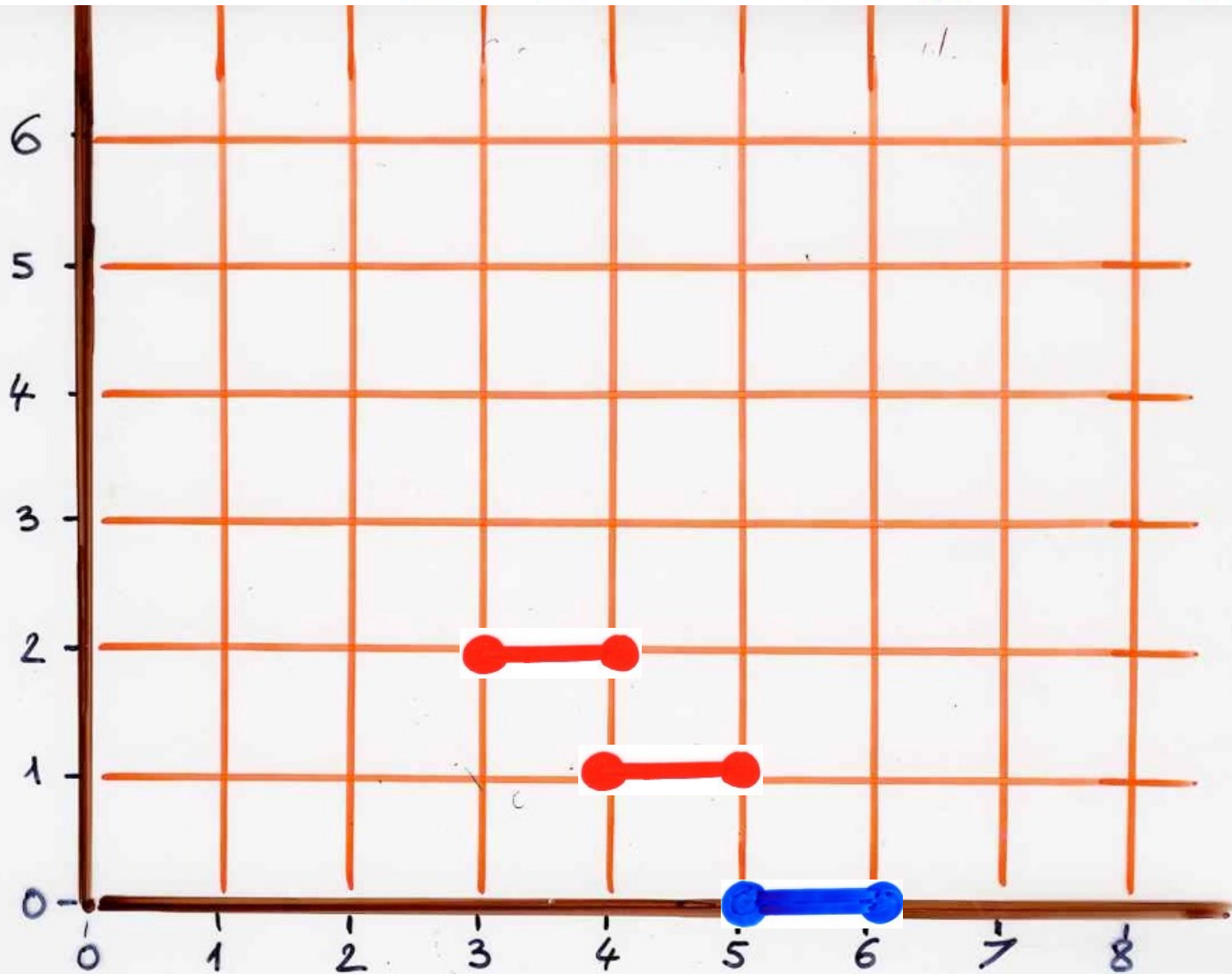
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3$$

example

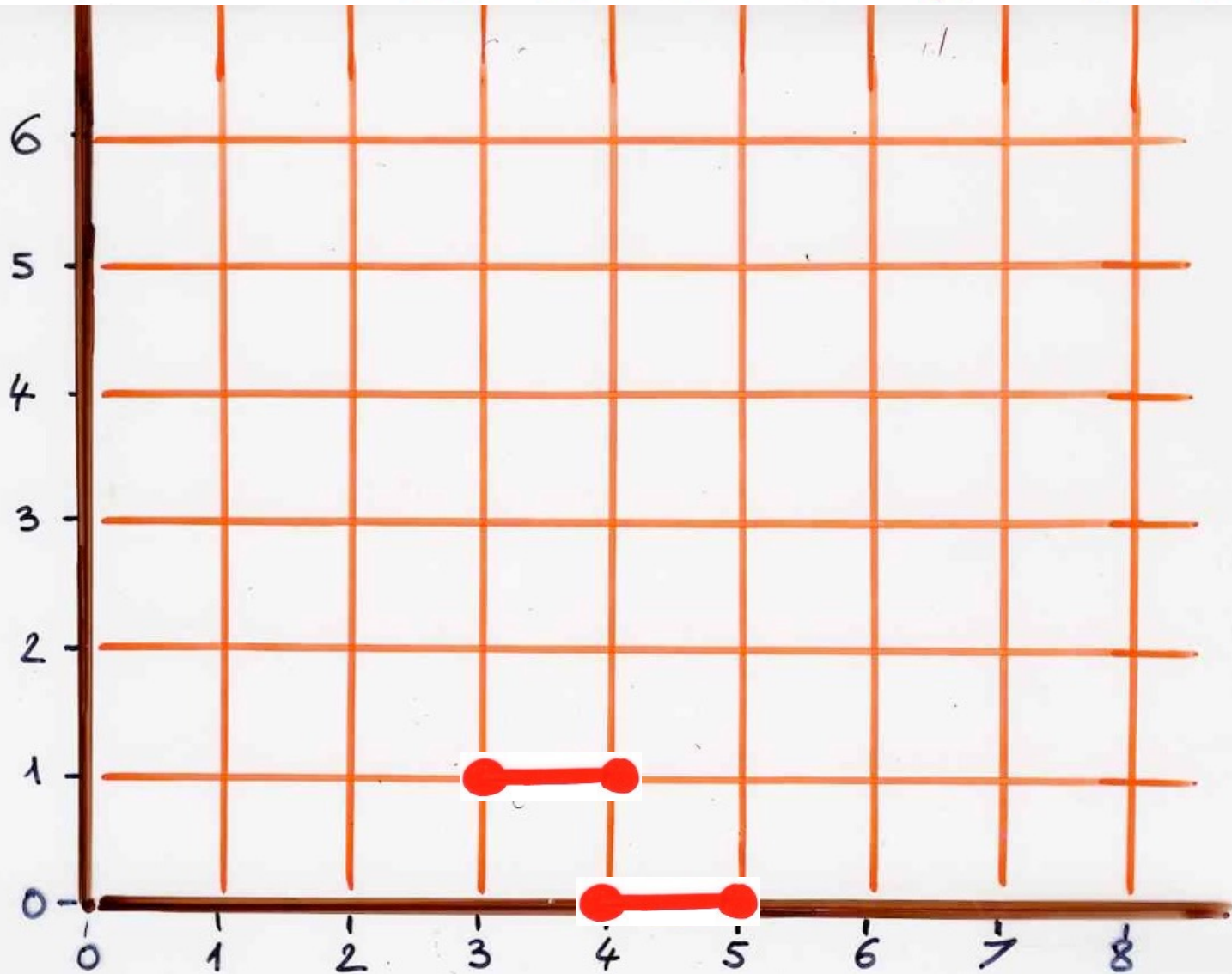
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_5$$

example

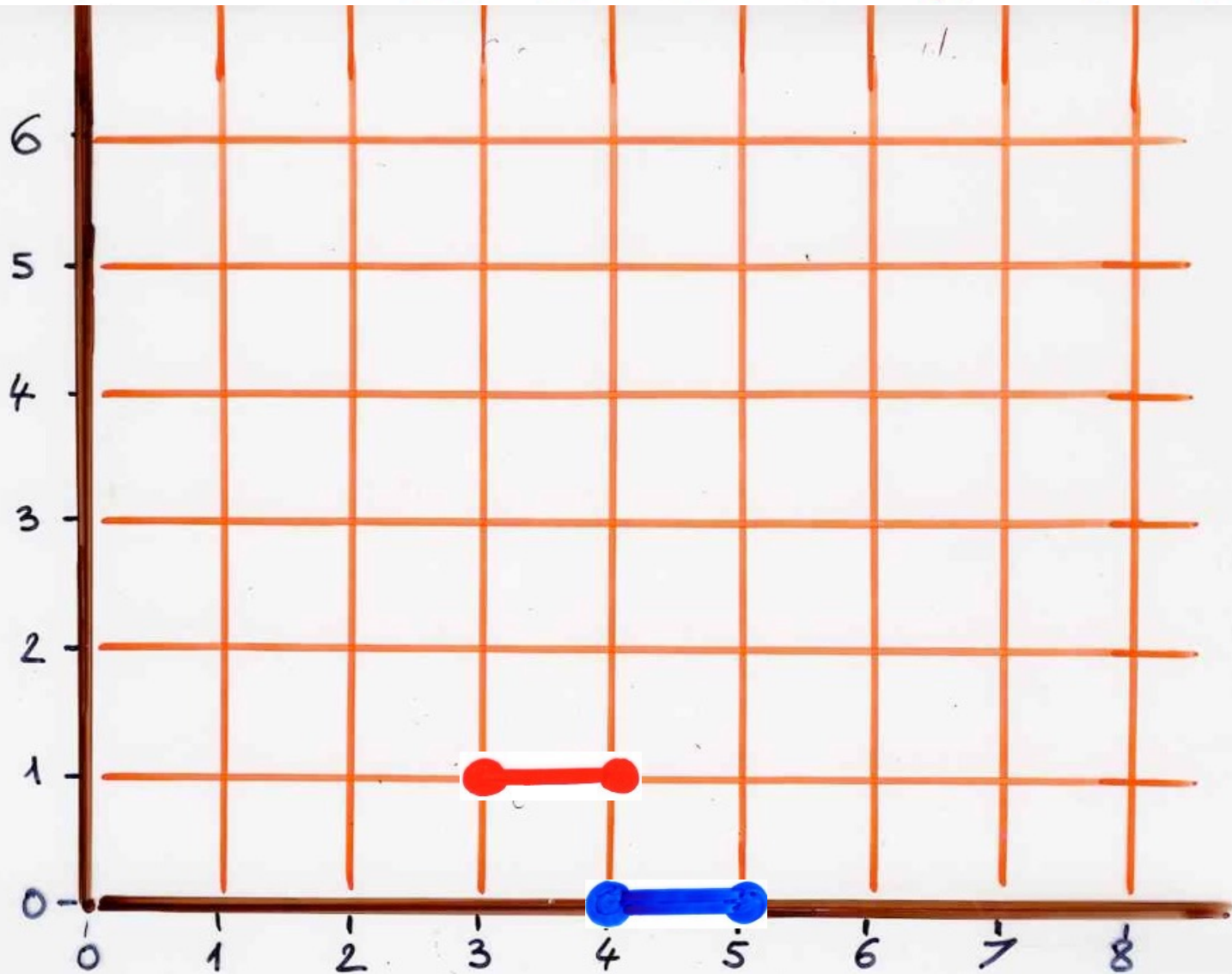
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_5$$

example

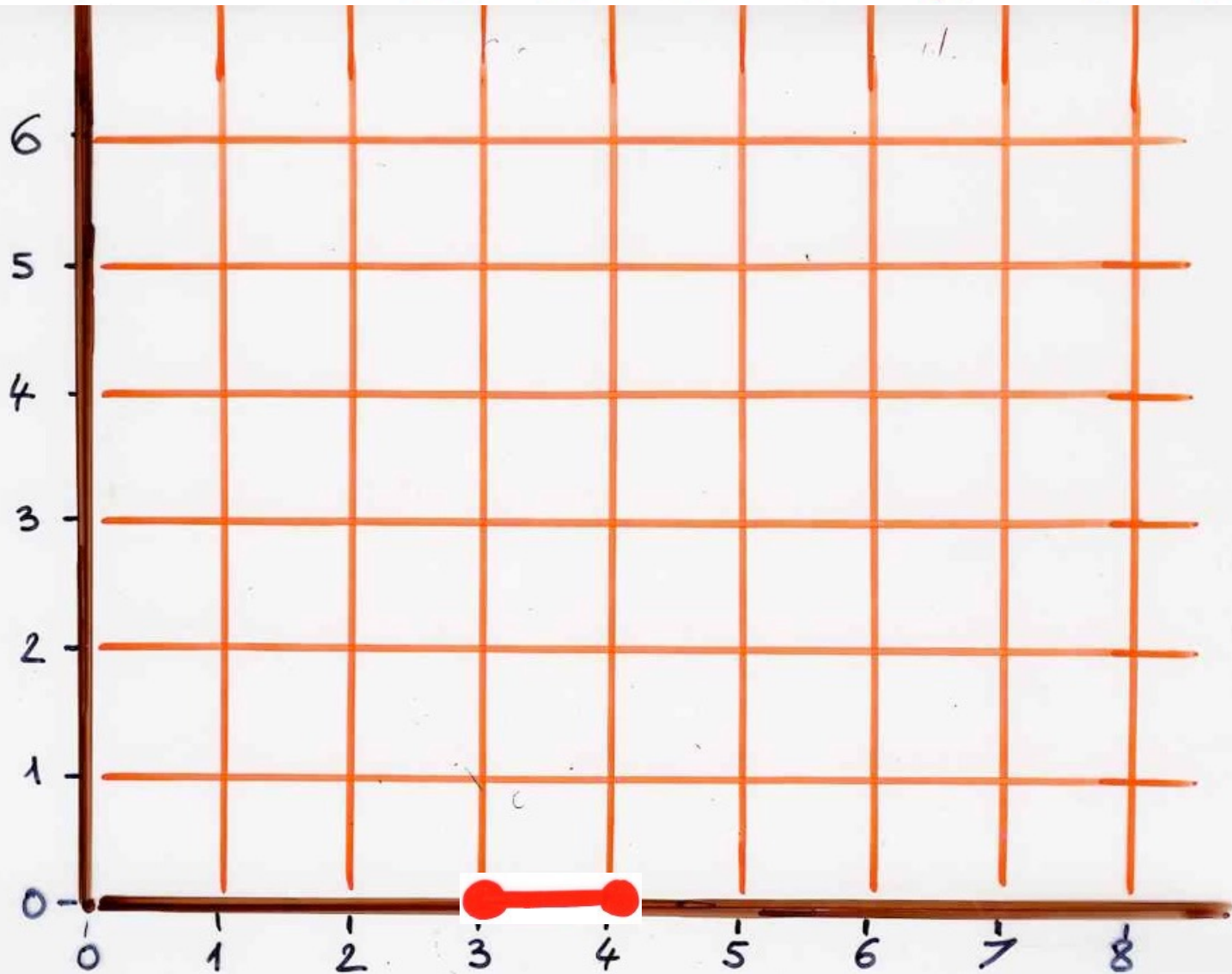
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_5 \sigma_4$$

example

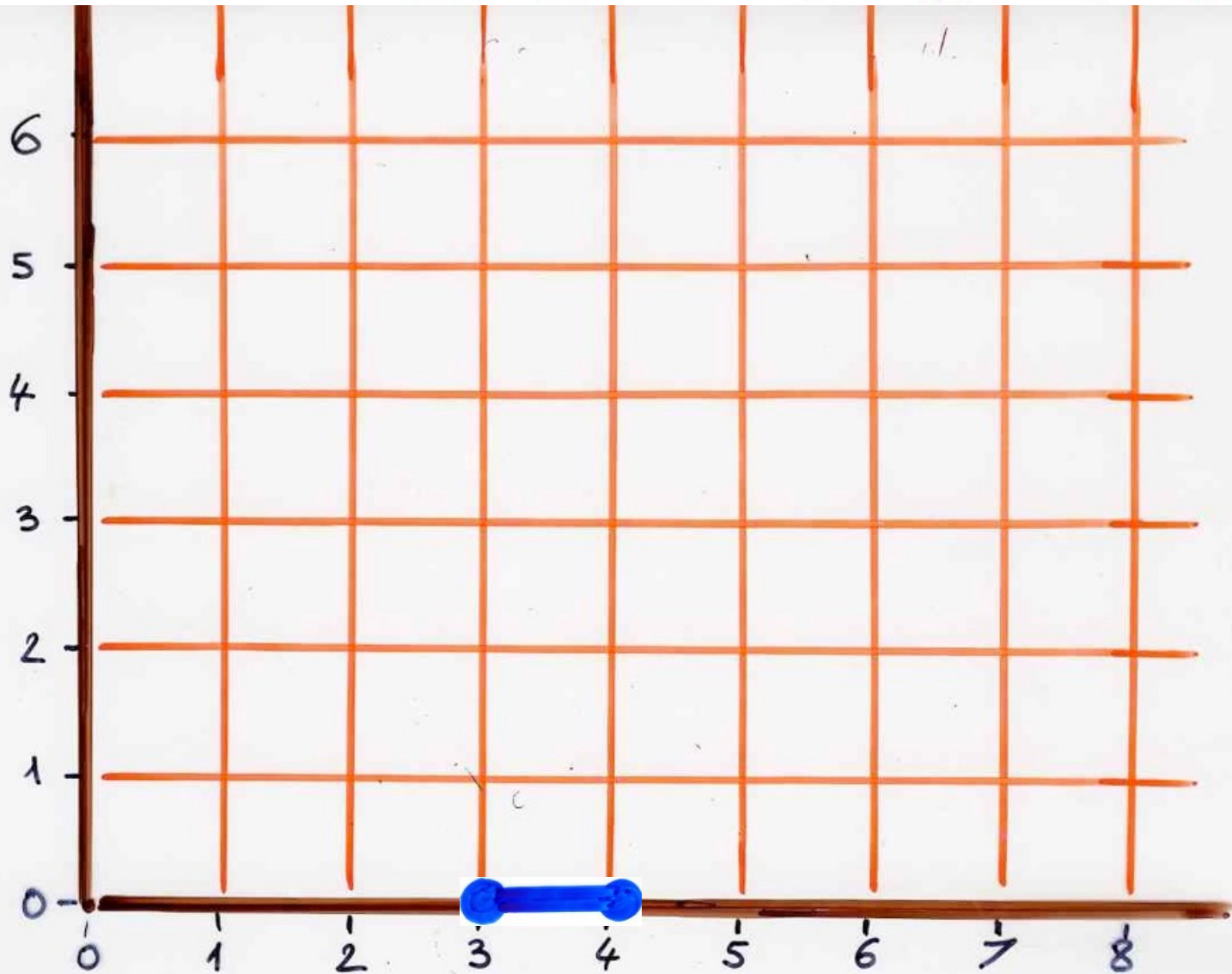
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example

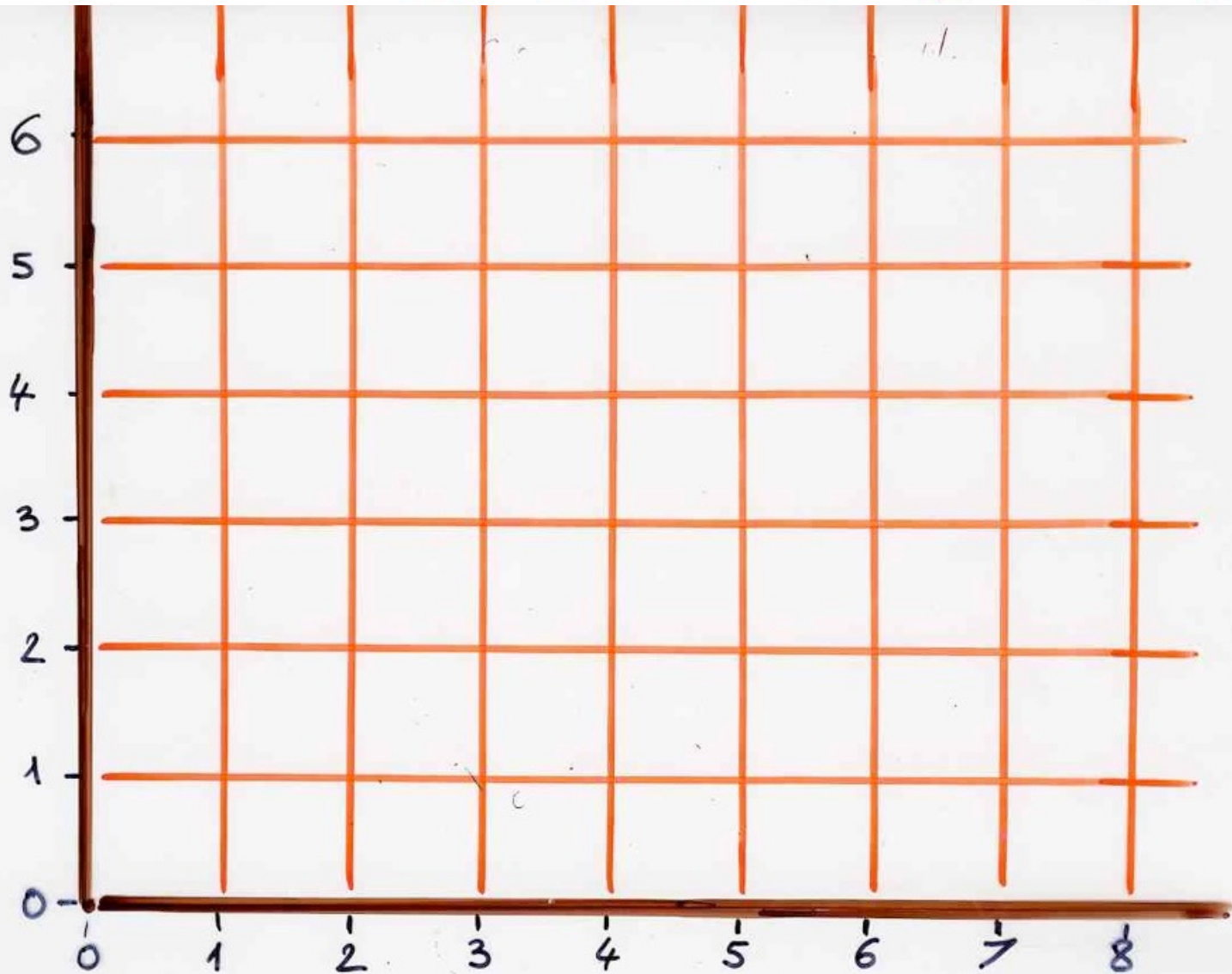
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_5 \sigma_4 \sigma_3$$

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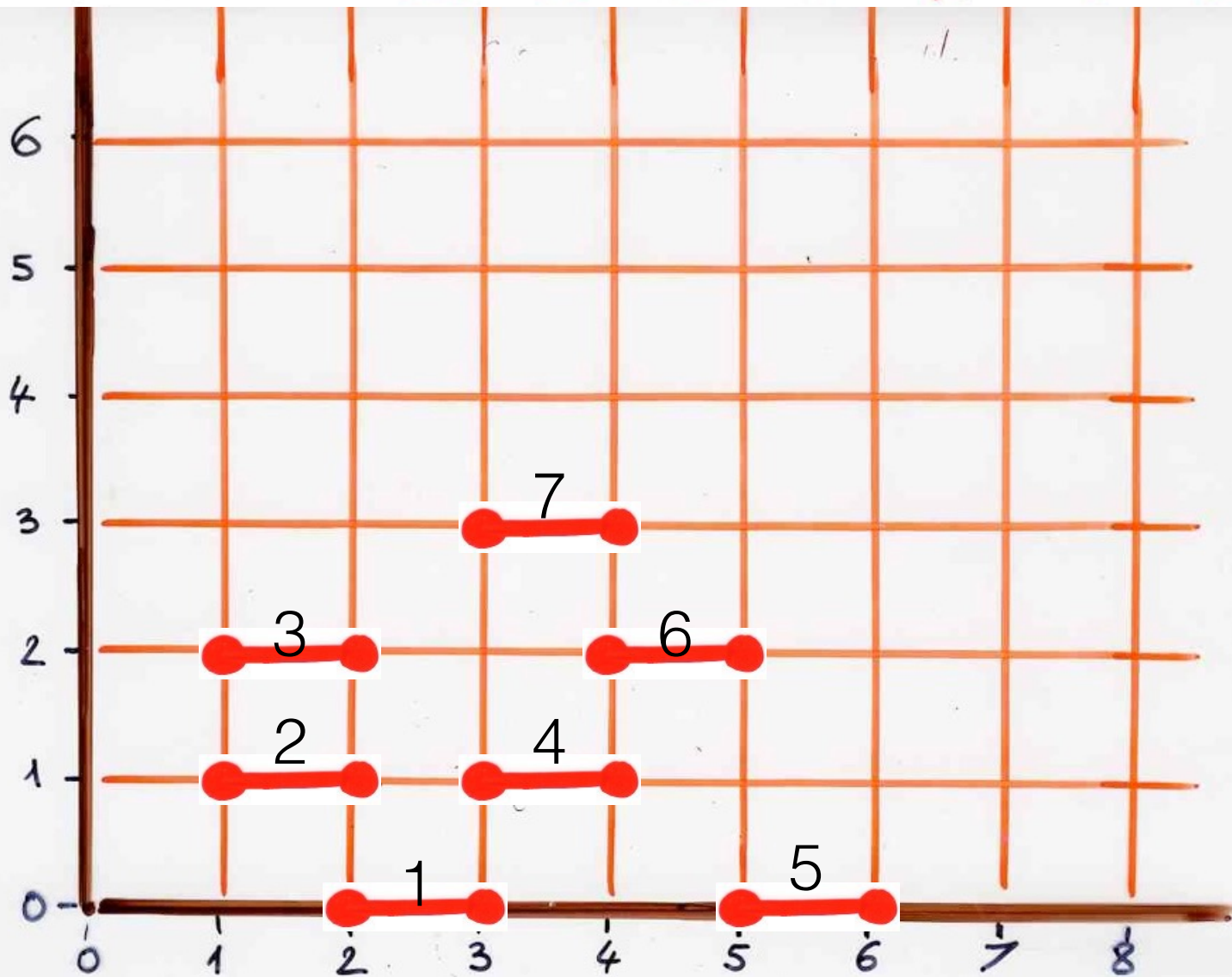
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$$\sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_5 \sigma_4 \sigma_3$$

example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$



$$\sigma_2 \sigma_1 \sigma_1 \sigma_3 \sigma_5 \sigma_4 \sigma_3$$

minimal

letter of a class $[w]$

$$[w] = [y, y_1]$$

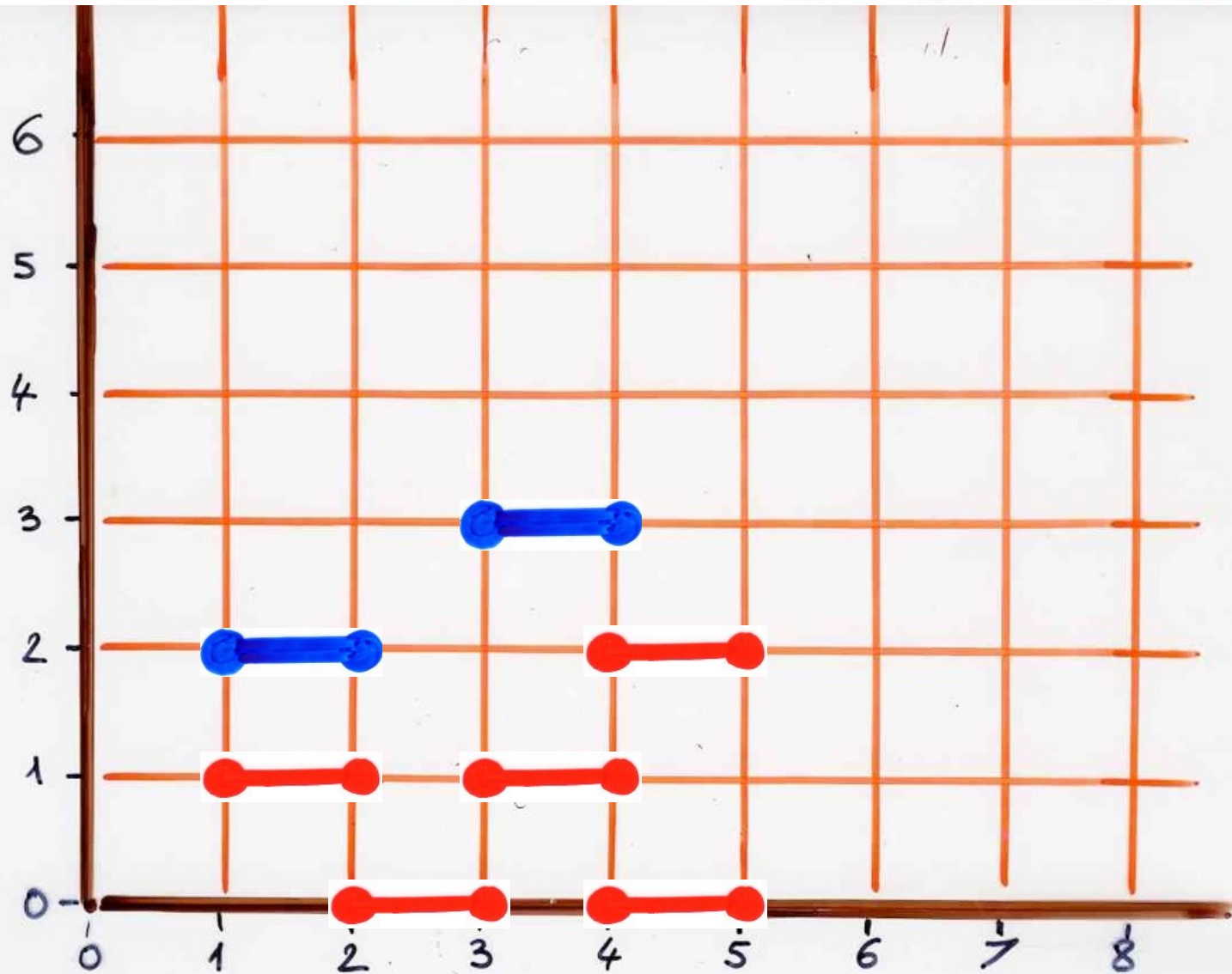
maximal

letter of a class $[w]$

$$[w] = [u_1, z]$$

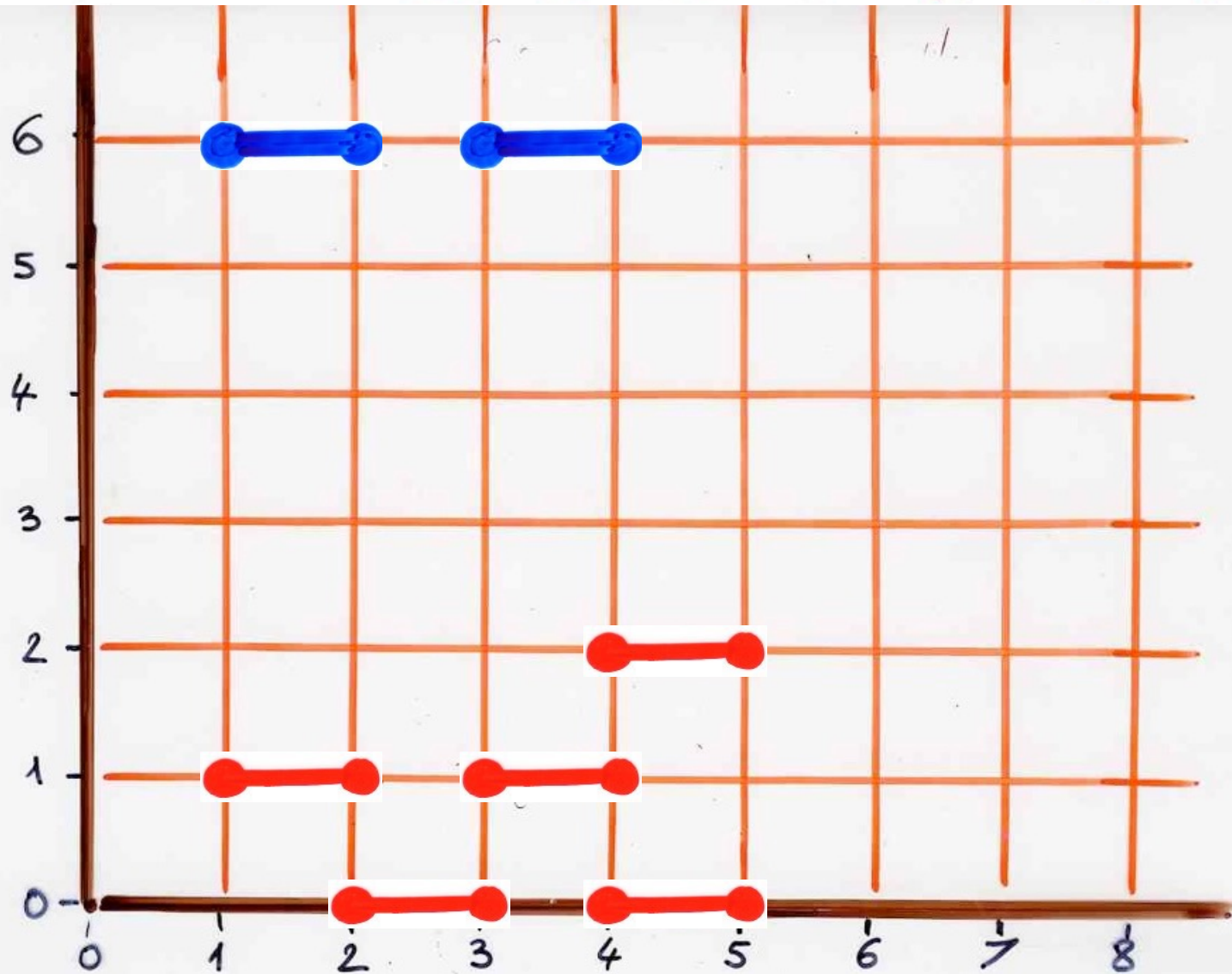
example

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example

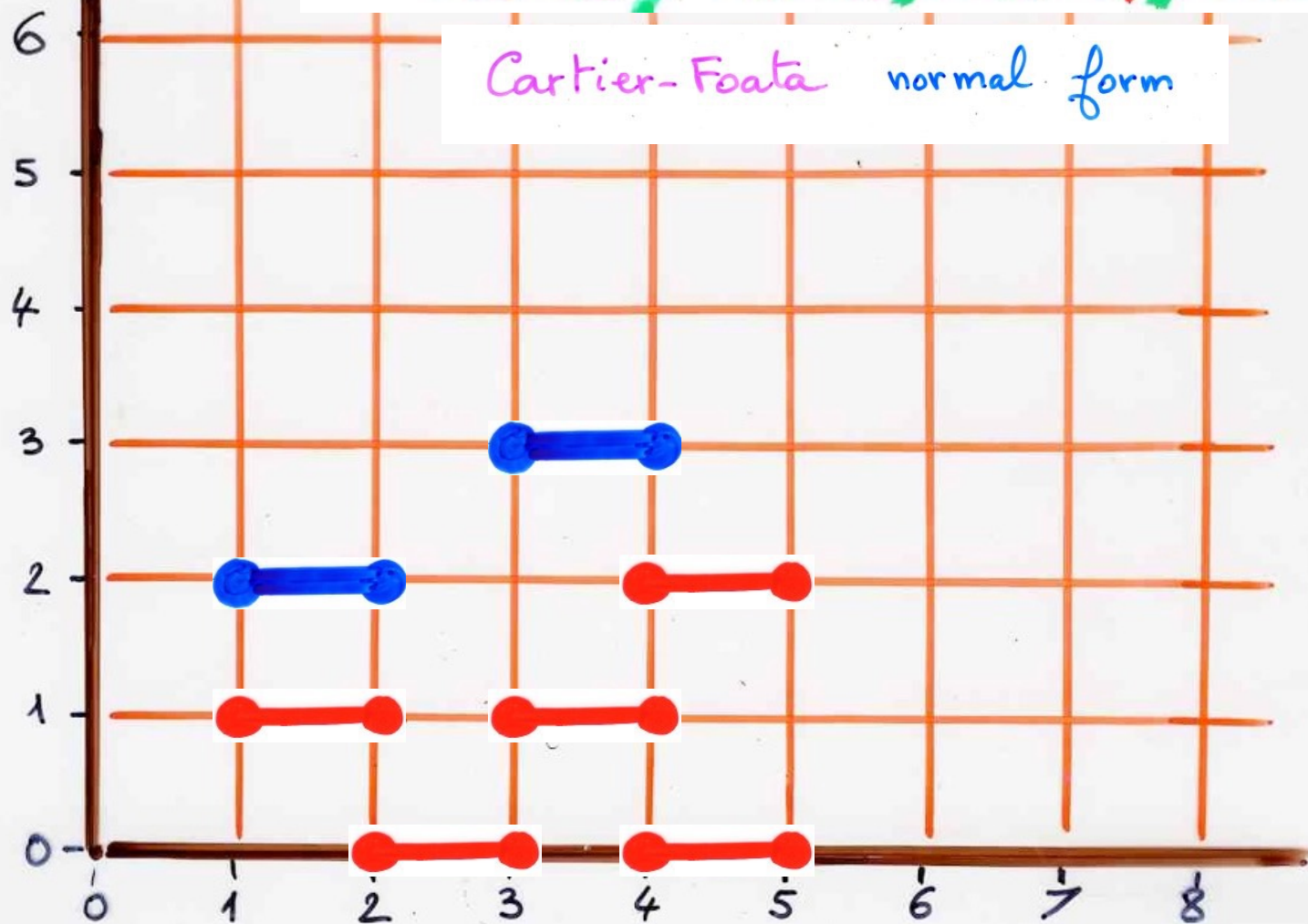
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example

$$w = \sigma_2 \sigma_3 \sigma_5 \sigma_1 \sigma_4 \sigma_1 \sigma_3$$

$$\equiv \sigma_2 \sigma_5 / \sigma_1 \sigma_3 / \sigma_1 \sigma_4 / \sigma_3$$

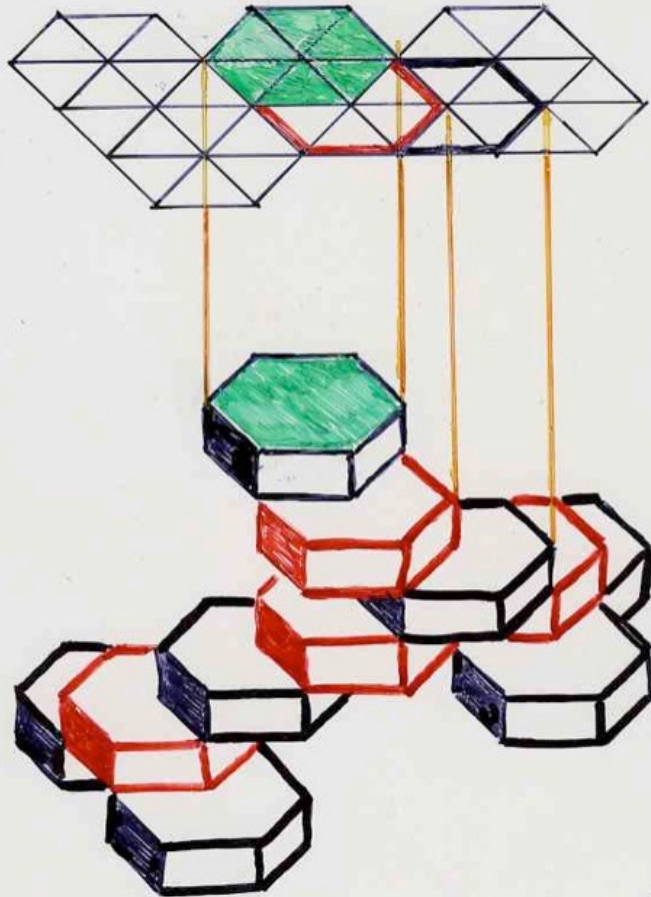


Pyramid

Def- Heap having only one maximal piece



$$-p(-t) = y$$



10.

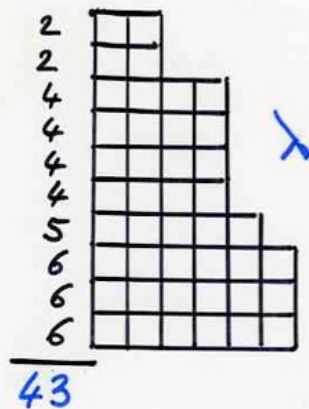
exercise

quasi-partition of integers

partition of an integer n

$$\lambda = (6, 6, 6, 5, 4, 4, 4, 4, 2, 2)$$

$$n = 43 = 6 + 6 + 6 + 5 + 4 + 4 + 4 + 4 + 2 + 2$$



Ferrers
diagram

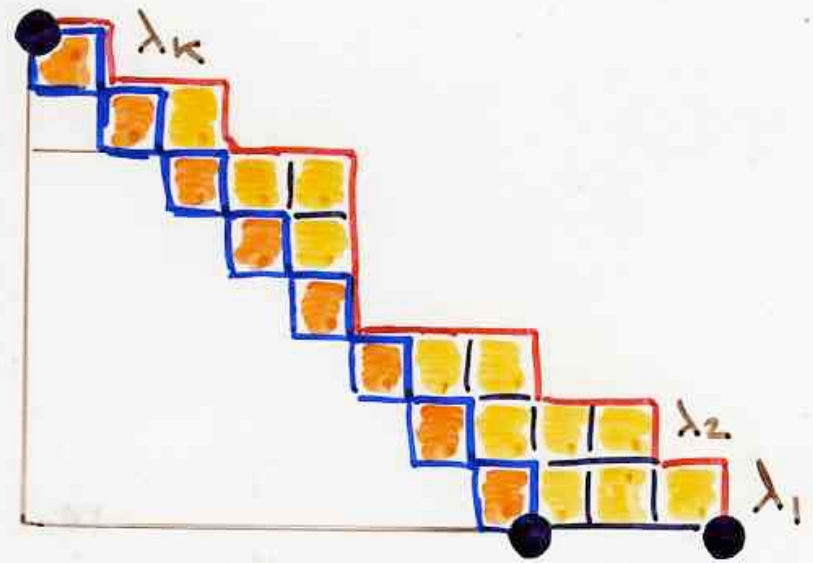
quasi-partition de n

$$n = \lambda_1 + \lambda_2 + \dots + \lambda_k$$

$$\lambda_i \geq \lambda_{i+1} - 1$$

$$i = 1, \dots, k-1$$

$$\lambda = (4, 4, 3, 1, 2, 3, 2, 1)$$



quasi-partitions

Auluck 1951

Andrews 1981

reciprocal of

Rogers-Ramanujan identities

exercise 1 using **lexicographic normal form**

find a **bijection** between **heaps** of

dimers on $\mathbb{N}_+ = \{1, 2, \dots\}$ and

quasi-partitions $(\lambda_1, \dots, \lambda_k)$

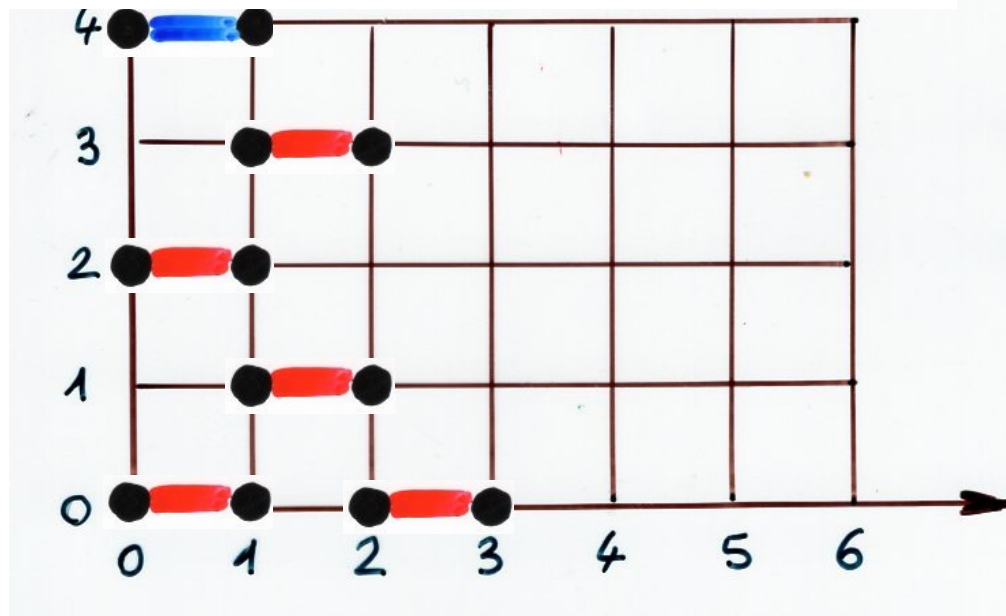
- the number **k** of parts will be the number of **dimers** of the **heap**

- find an interpretation with the corresponding **heap** of **dimers** of $n = \lambda_1 + \dots + \lambda_k$

exercises

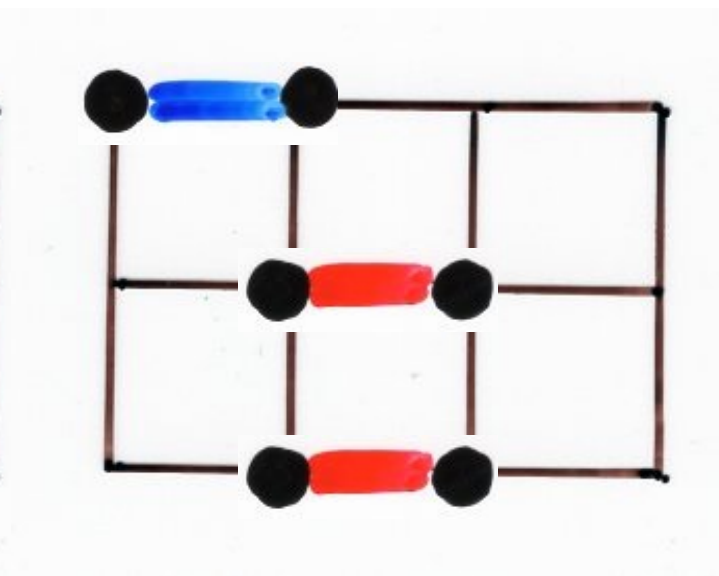
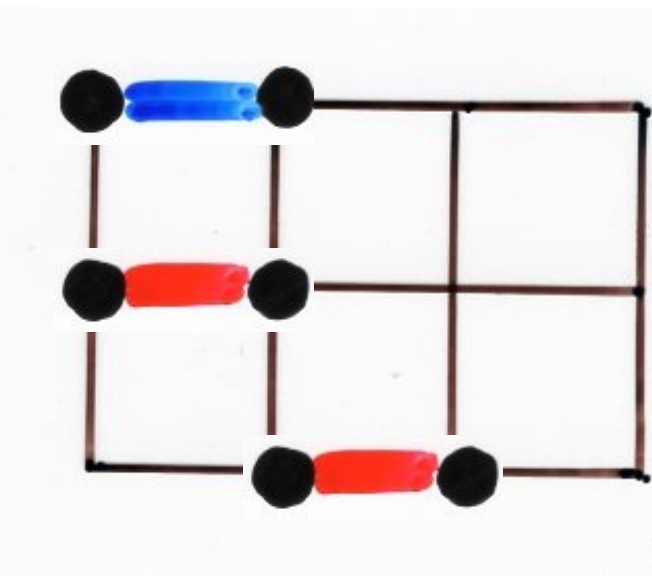
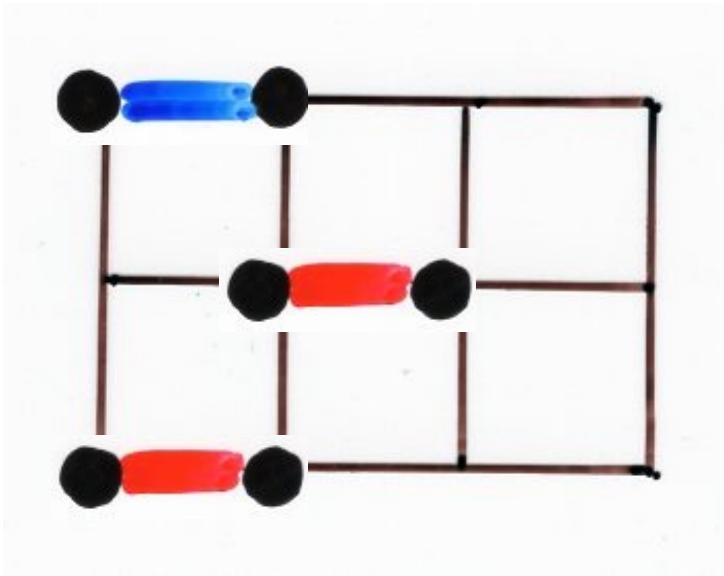
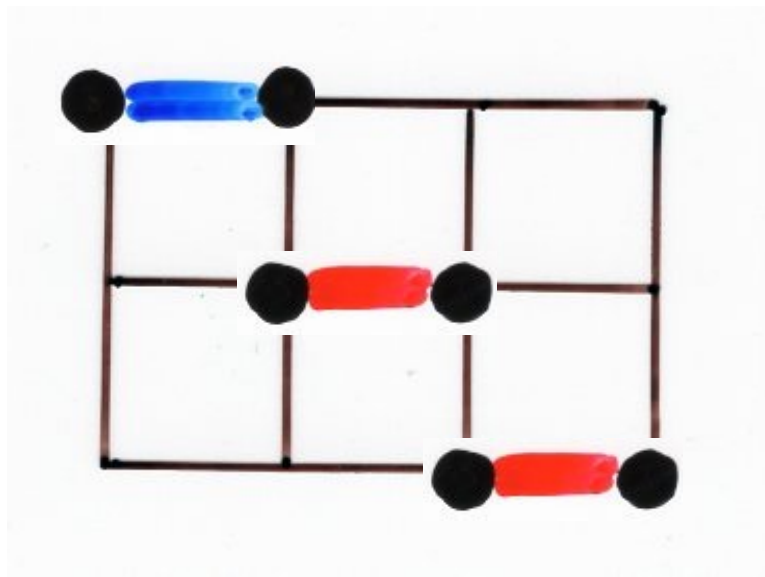
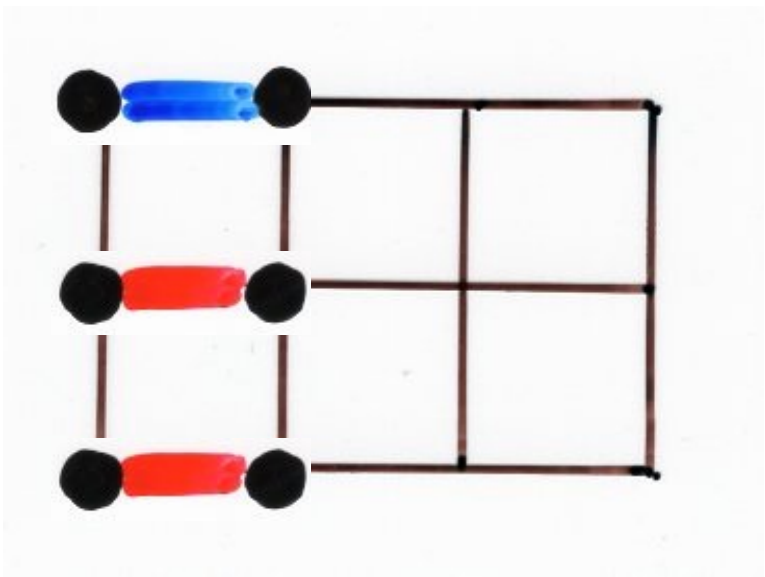
pyramids and semi-pyramids
of dimers

exercise semi-pyramid of dimers
on \mathbb{N}
the unique maximal piece has
projection $[0, 1]$



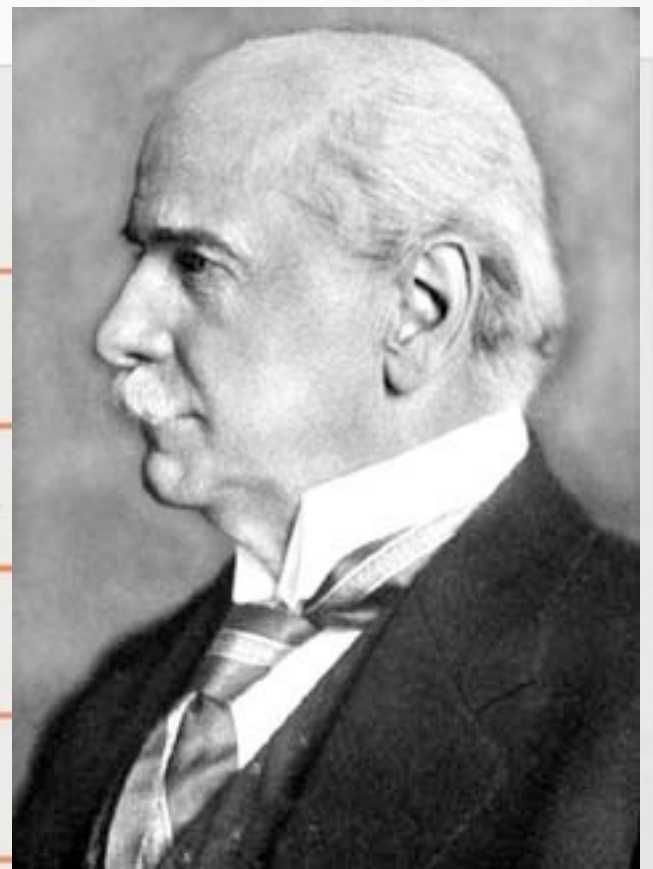
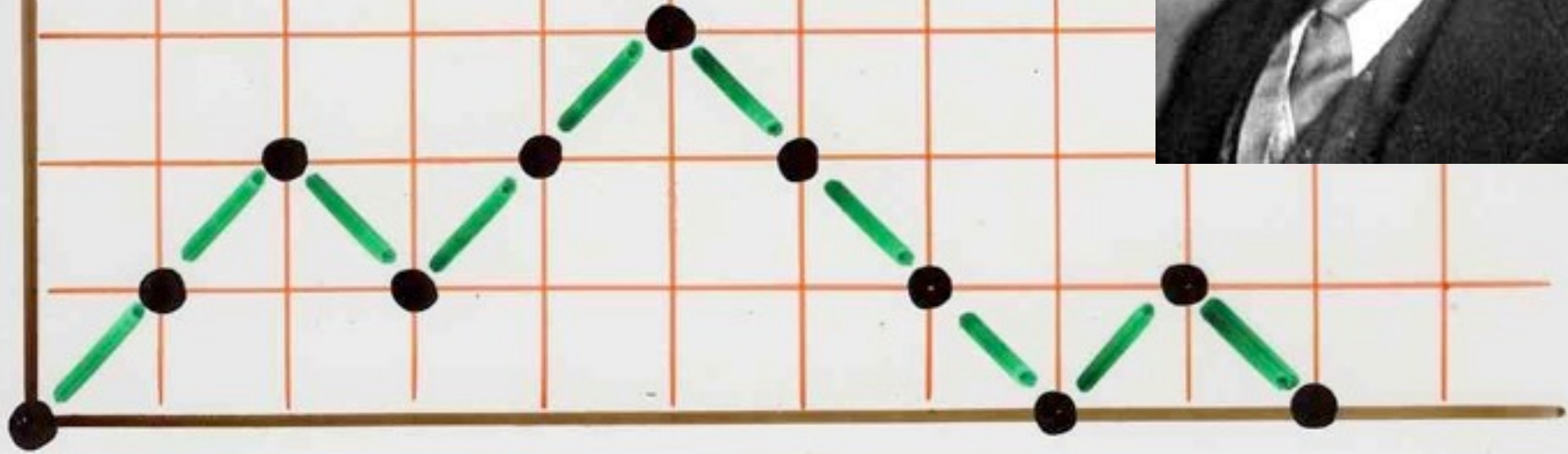
Catalan number

$$C_n = \frac{1}{(n+1)} \binom{2n}{n}$$



$$C_3 = 5$$

Dyck path



Catalan number

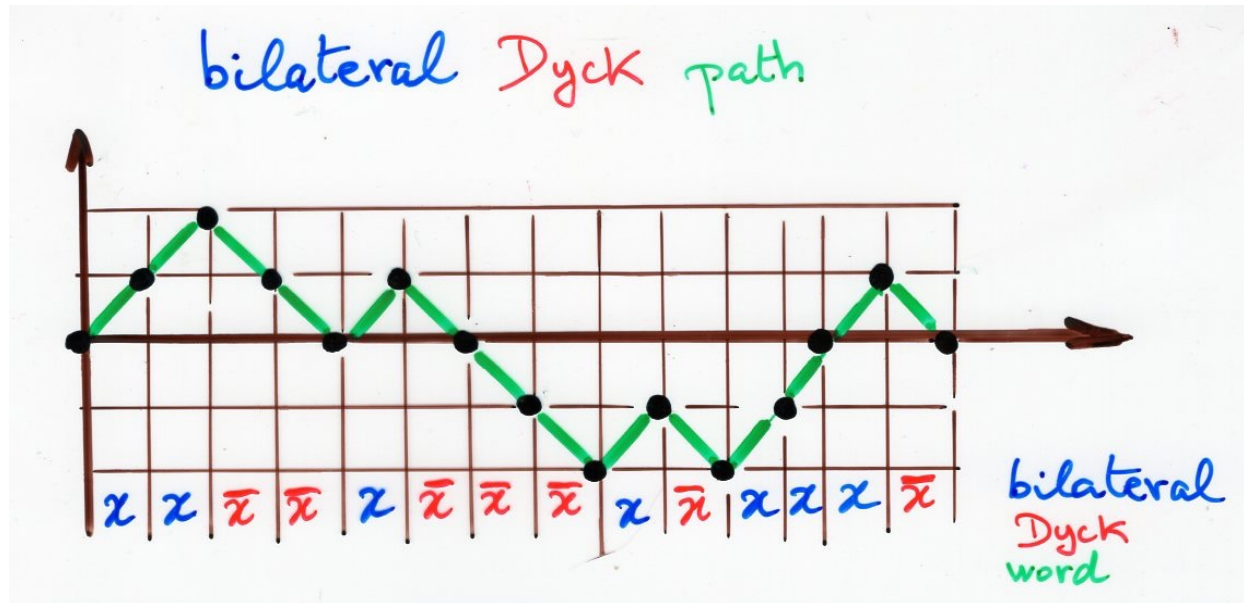
$$C_n = \frac{1}{(n+1)} \binom{2n}{n}$$

exercise 2 Using exercise 1 (about quasi-partitions and lexicographic normal form) find a bijection between semi-pyramids of dimers on \mathbb{N} having n dimers and Dyck paths of length $2n$.

exercise 3

(more difficult)


pyramid of dimers on \mathbb{Z}
up to translation



number of
bilateral Dyck paths
w of length $2n$

$$= \binom{2n}{n}$$

exercise 3 pyramid of dimers on \mathbb{Z}
(more difficult) up to translation

- using exercise 2 find a bijection between pyramids of dimers on \mathbb{Z} such that the projection of the maximal piece is $[0,1]$ and bilateral Dyck paths starting with step 

- thus the number of pyramid of dimers on \mathbb{Z} up to translation is

$$\frac{1}{2} \binom{2n}{n}$$

Posets

Poset (partially ordered set)

(E, \preceq) \preceq order relation

\preceq order relation on E

- reflexive $x \preceq x$ all $x \in E$
- antisymmetric $x \preceq y$ and $y \preceq x \Rightarrow x=y$
- transitive $x \preceq y$ and $y \preceq z \Rightarrow x \preceq z$

for all $x, y, z \in E$

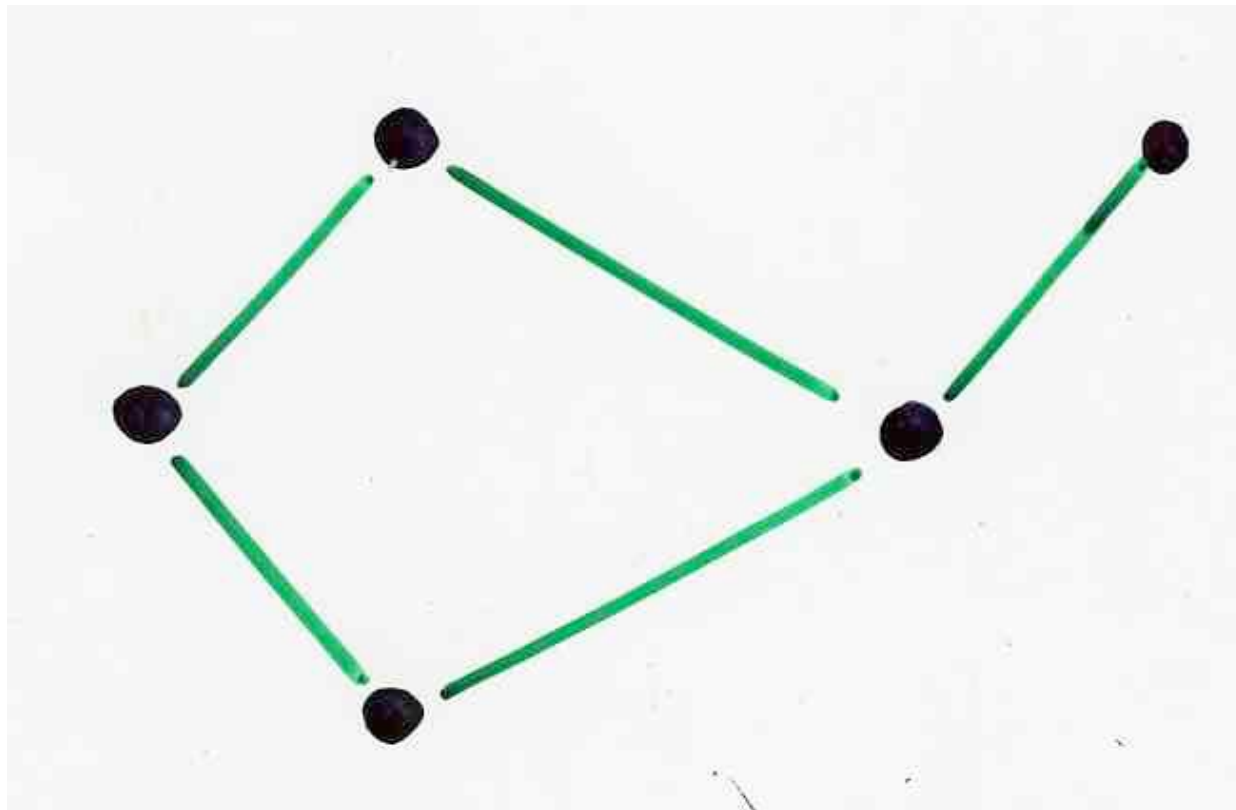
Poset (partially ordered set)
 (E, \preceq) \preceq order relation

covering relation

$x, y \in E$, y covers x
iff $x \prec y$ (strict) and $x \preceq z \preceq y \Rightarrow \begin{cases} z=x \\ \text{or} \\ z=y \end{cases}$

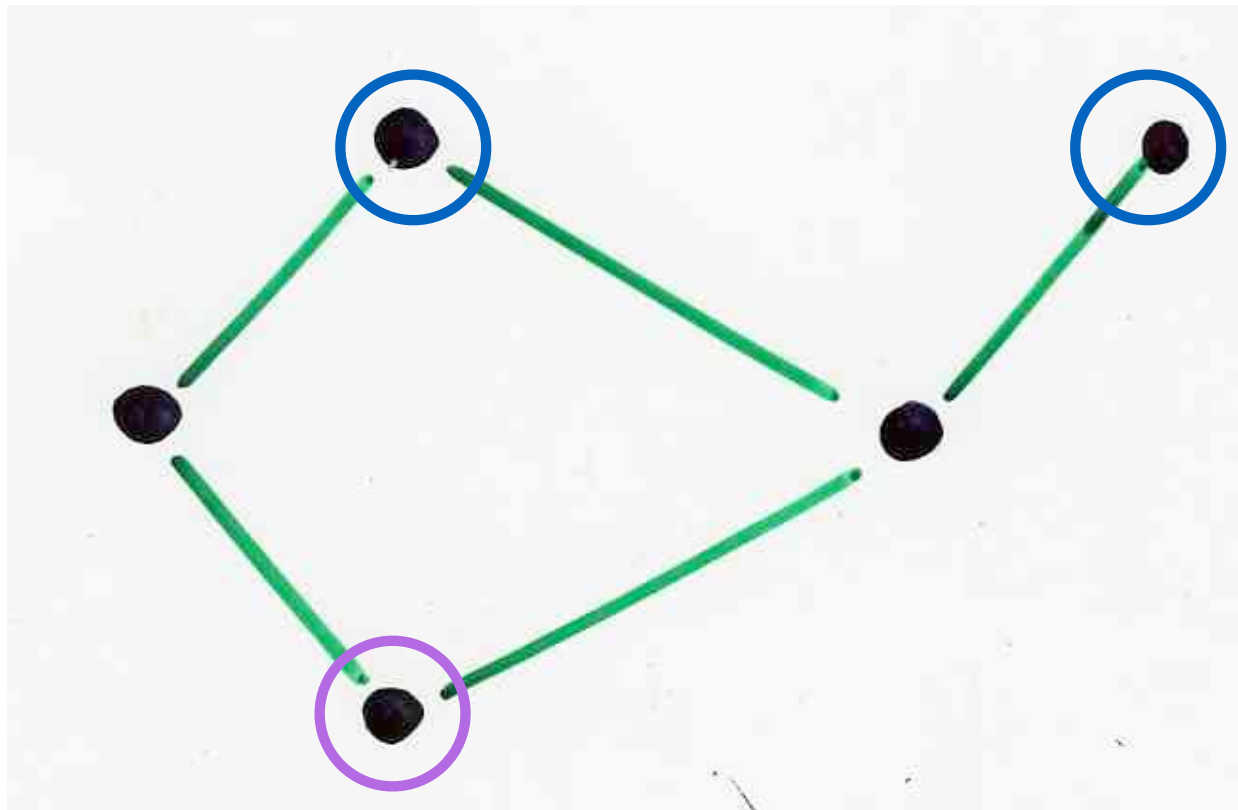
the interval $[x, y]$ is reduced to $\{x, y\}$

Hasse diagram
of a poset



minimal
maximal

element of a poset



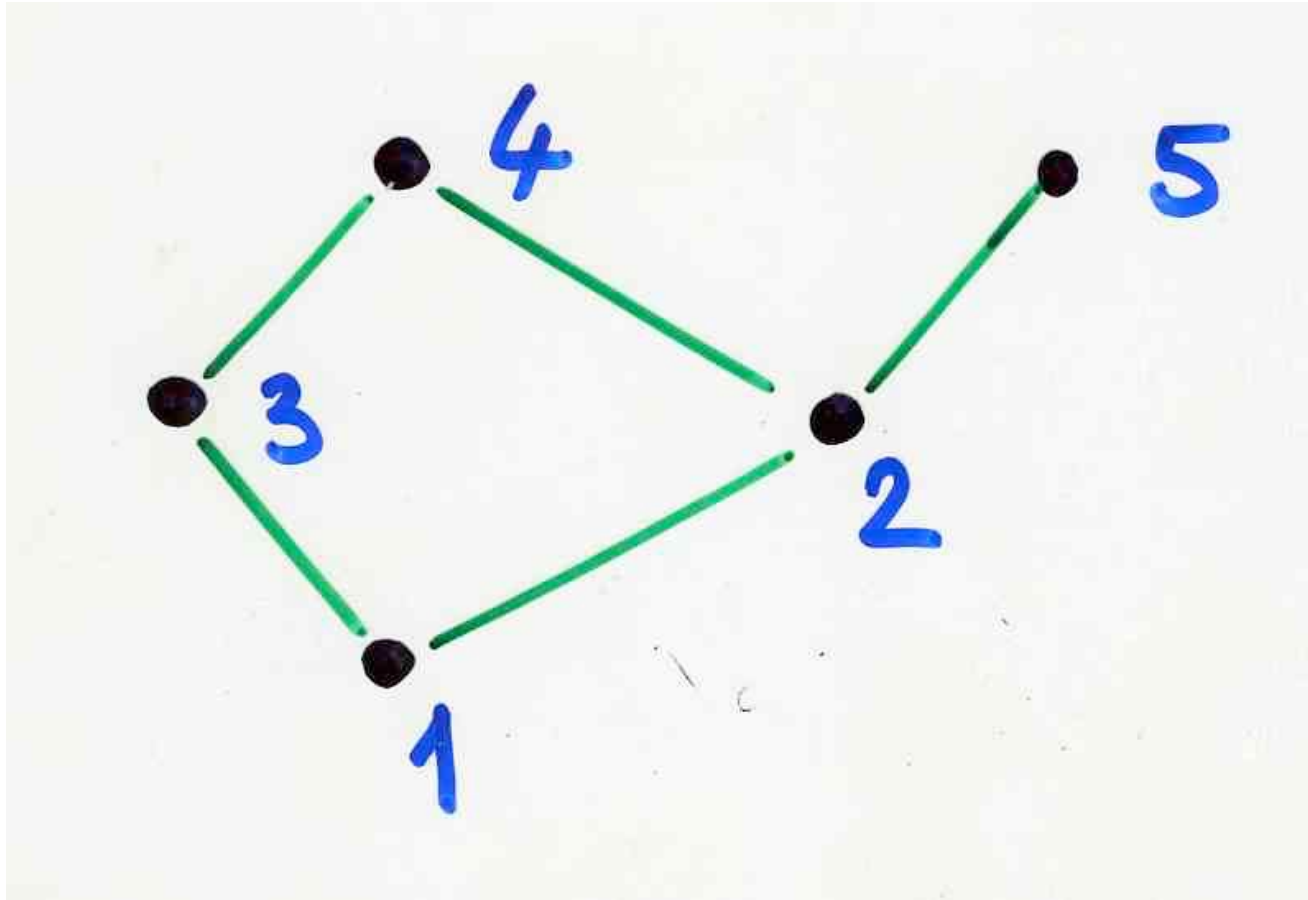
linear
extension

of a

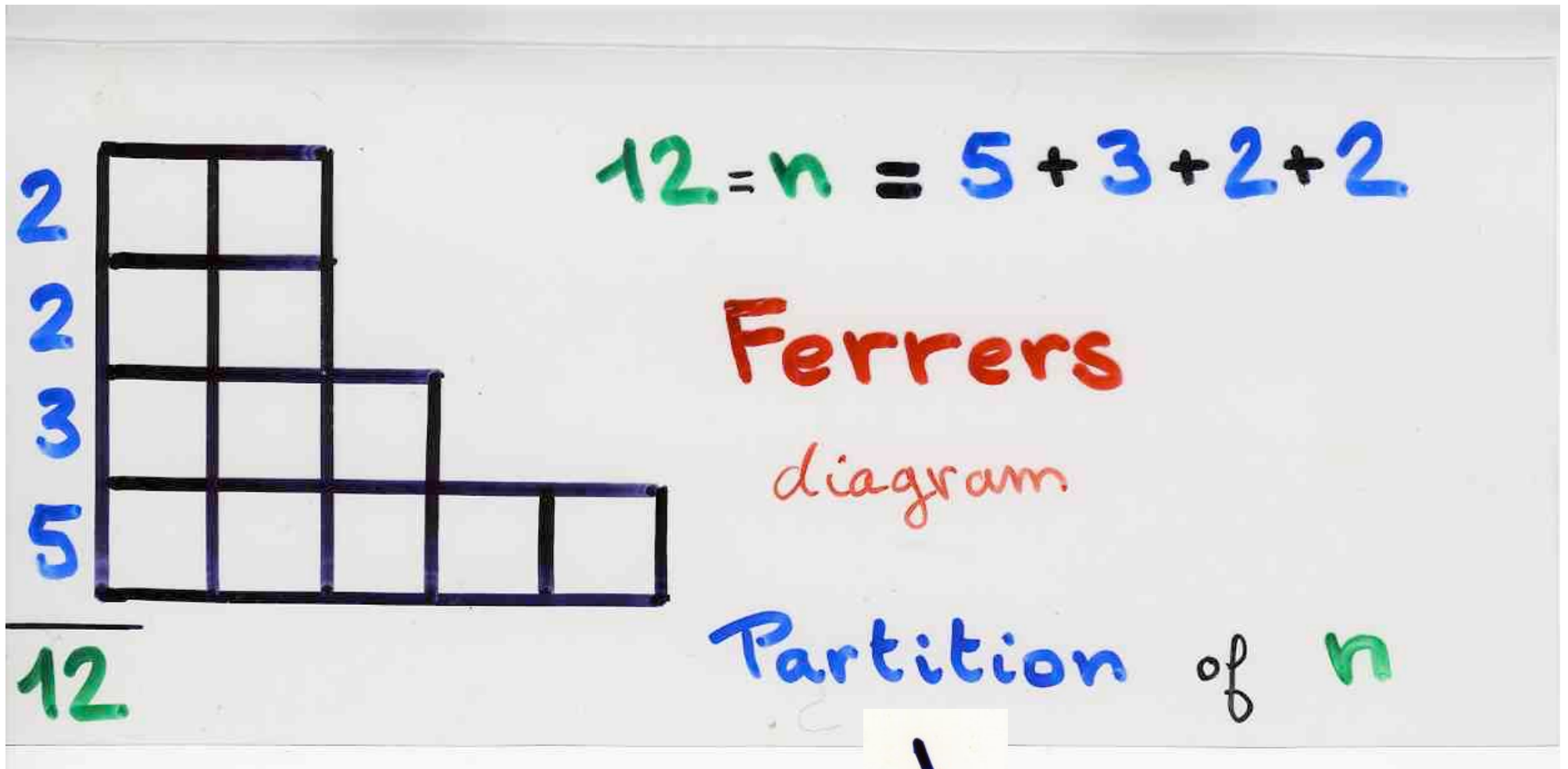
poset

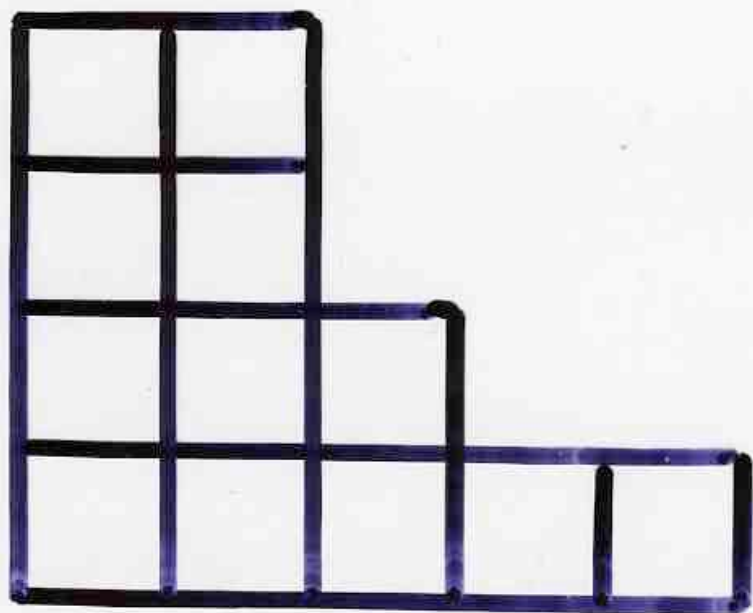
(E, \preceq)

Def $f : E \longrightarrow [1, n]$ is bijection
 $x \preceq y \implies f(x) \leq f(y)$



some examples





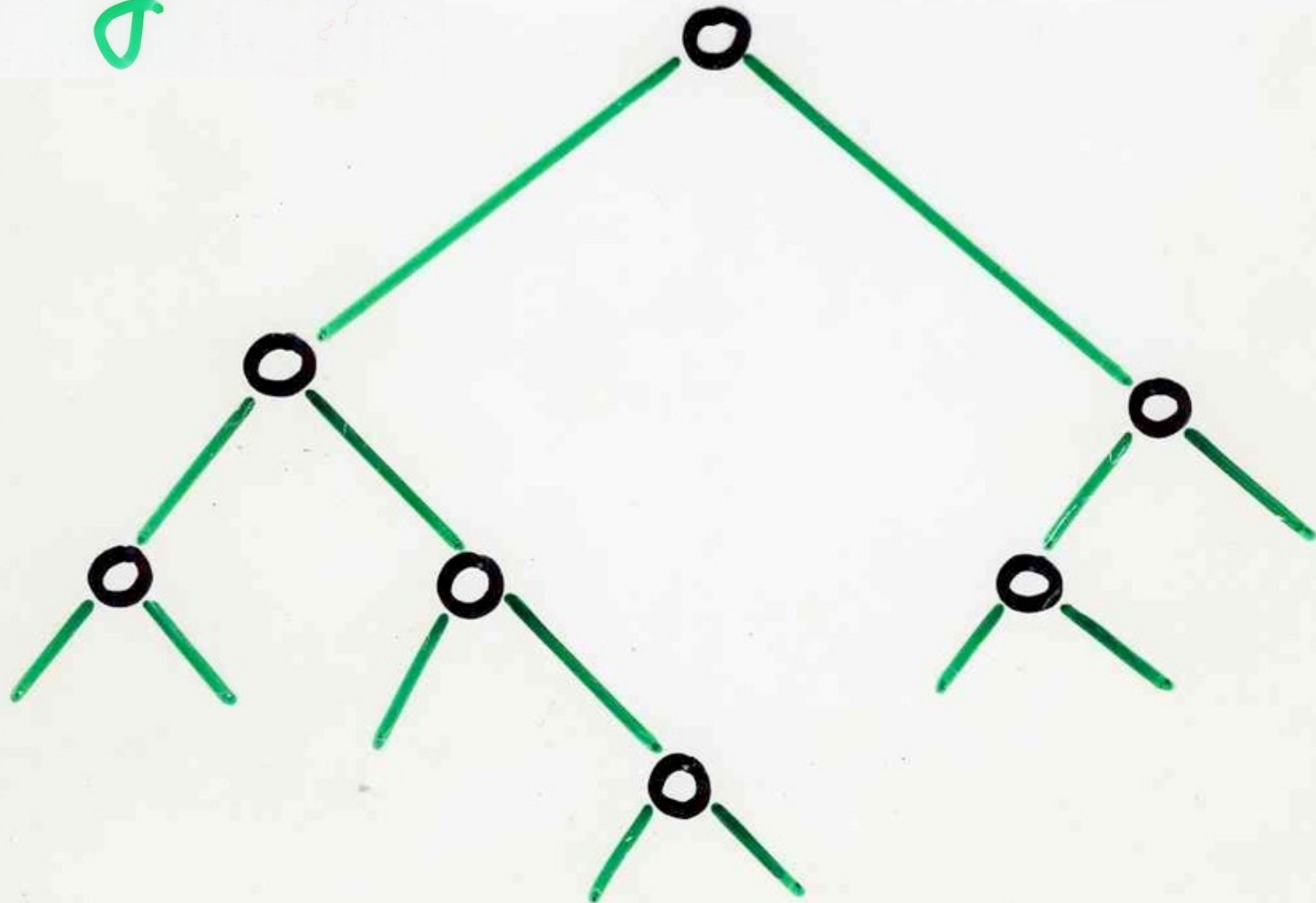
7	12			
6	10			
3	5	9		
1	2	4	8	11

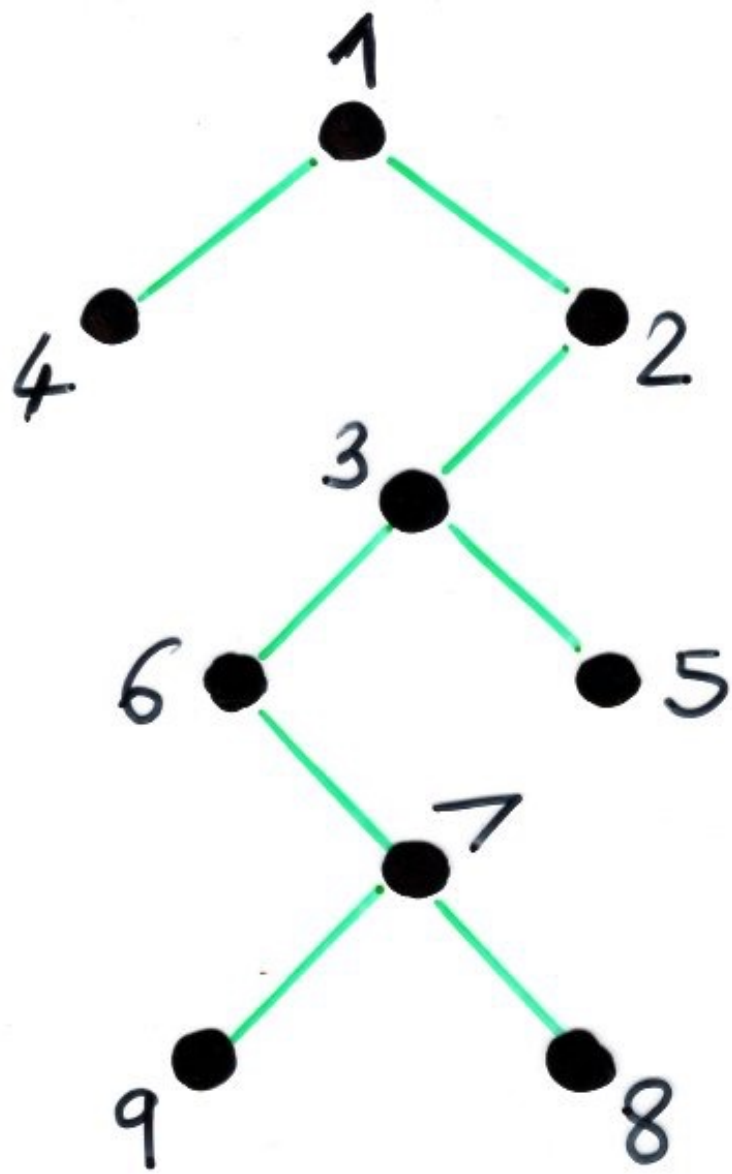
Young
tableau

shape



binary tree

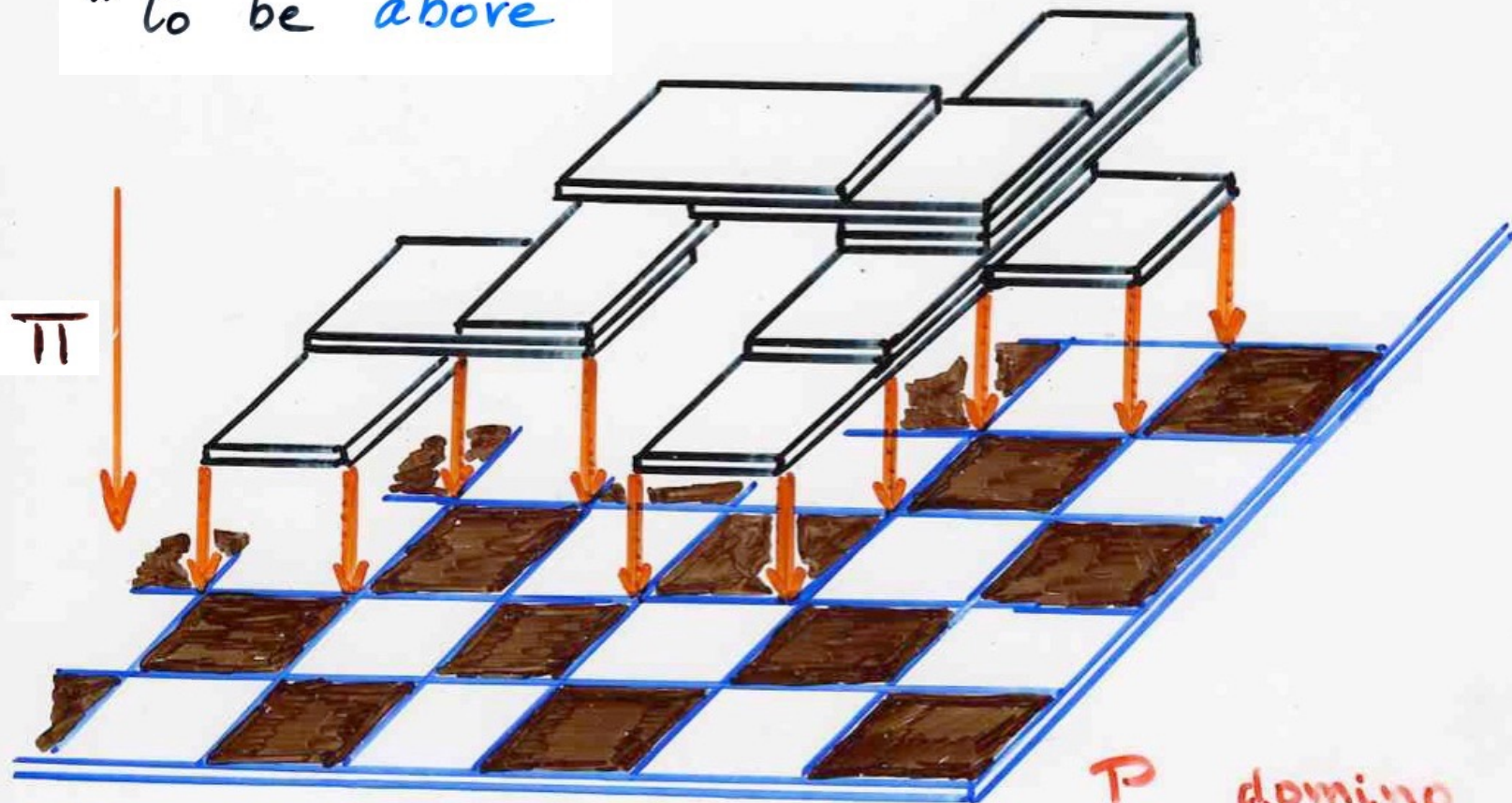




Heaps and posets

poset associated to a heap

"to be above"



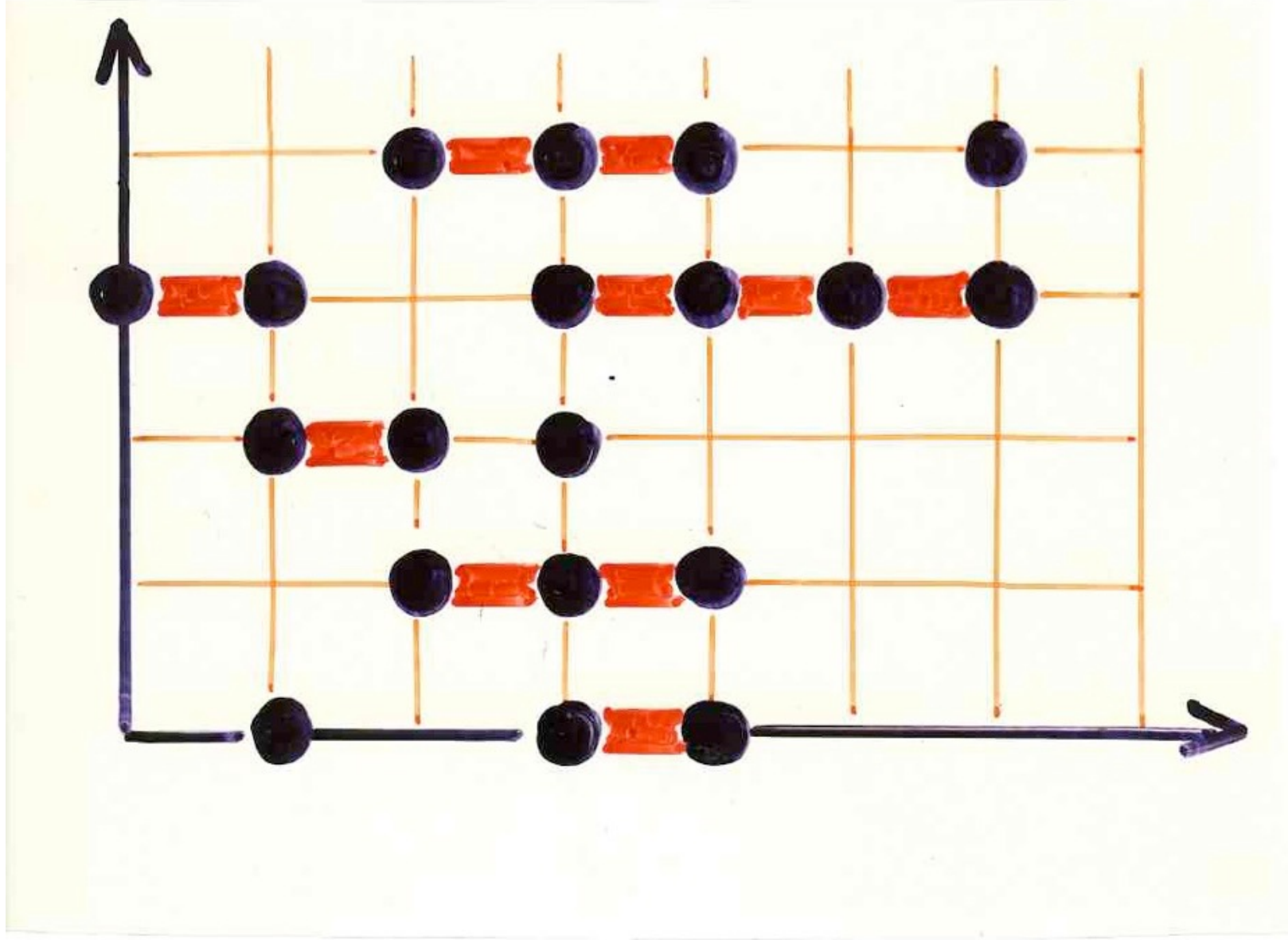
$$B = R \times R$$

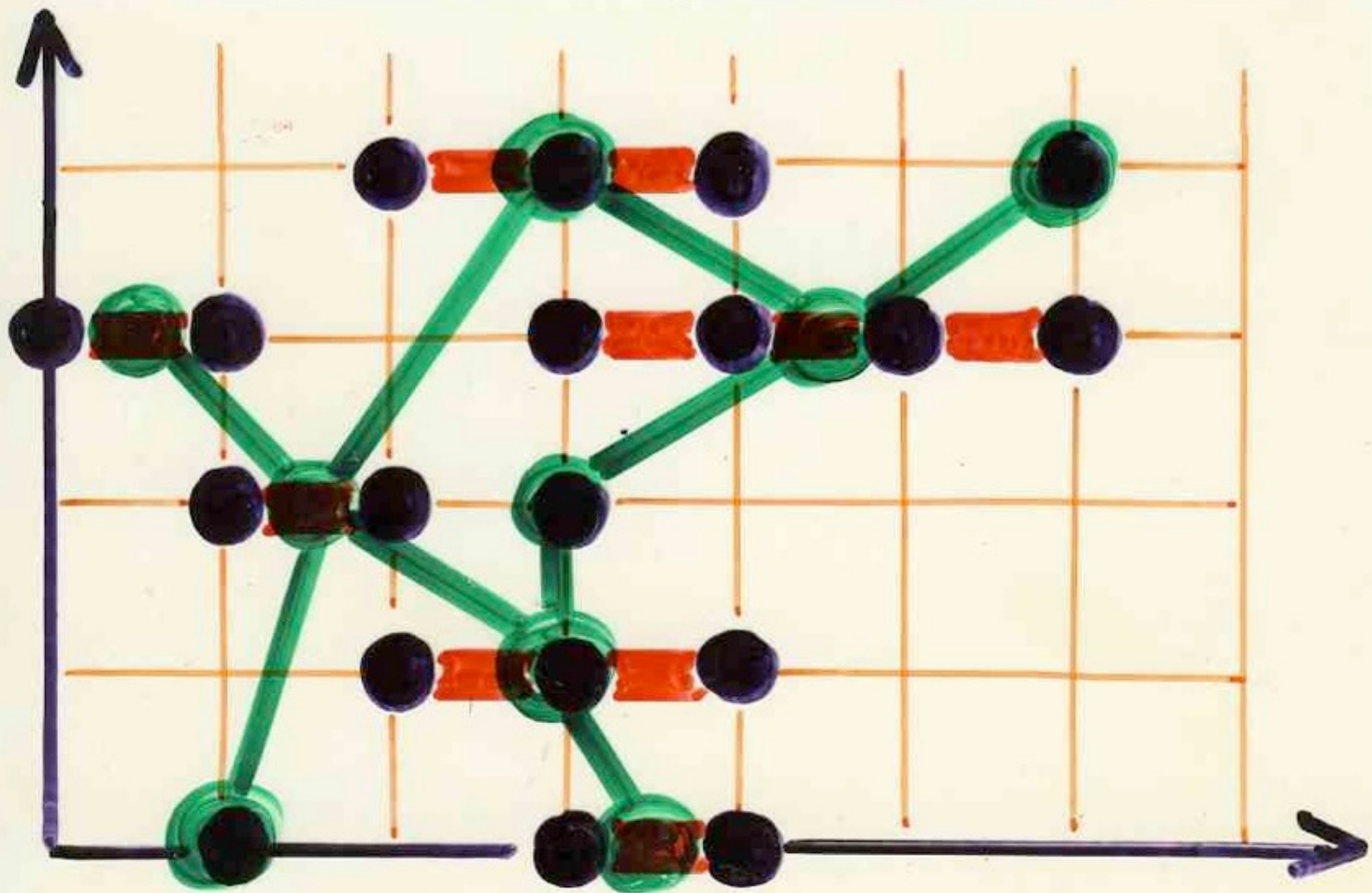
P domino

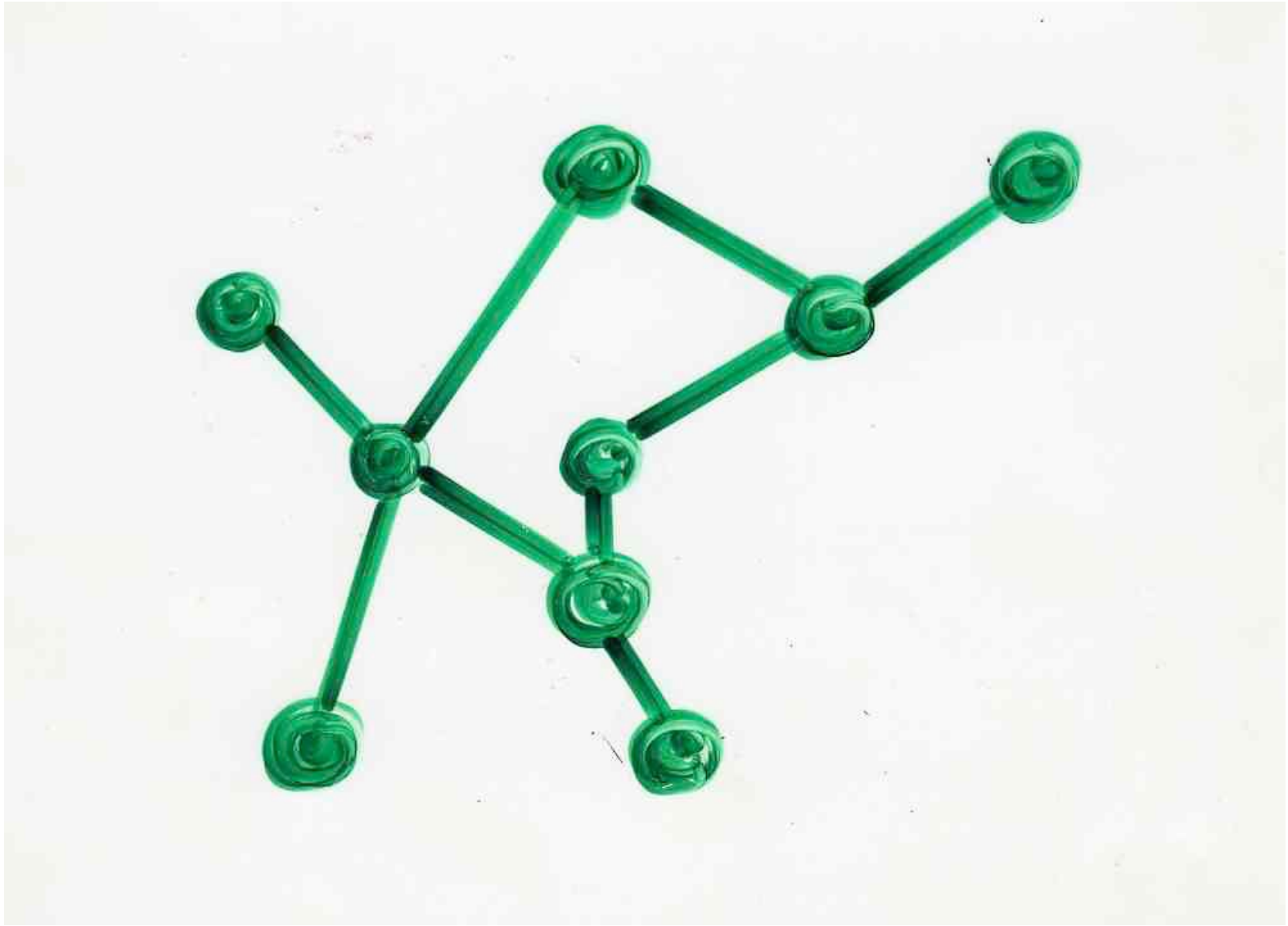
Def. Poset (E, \preceq) associated
to a heap E

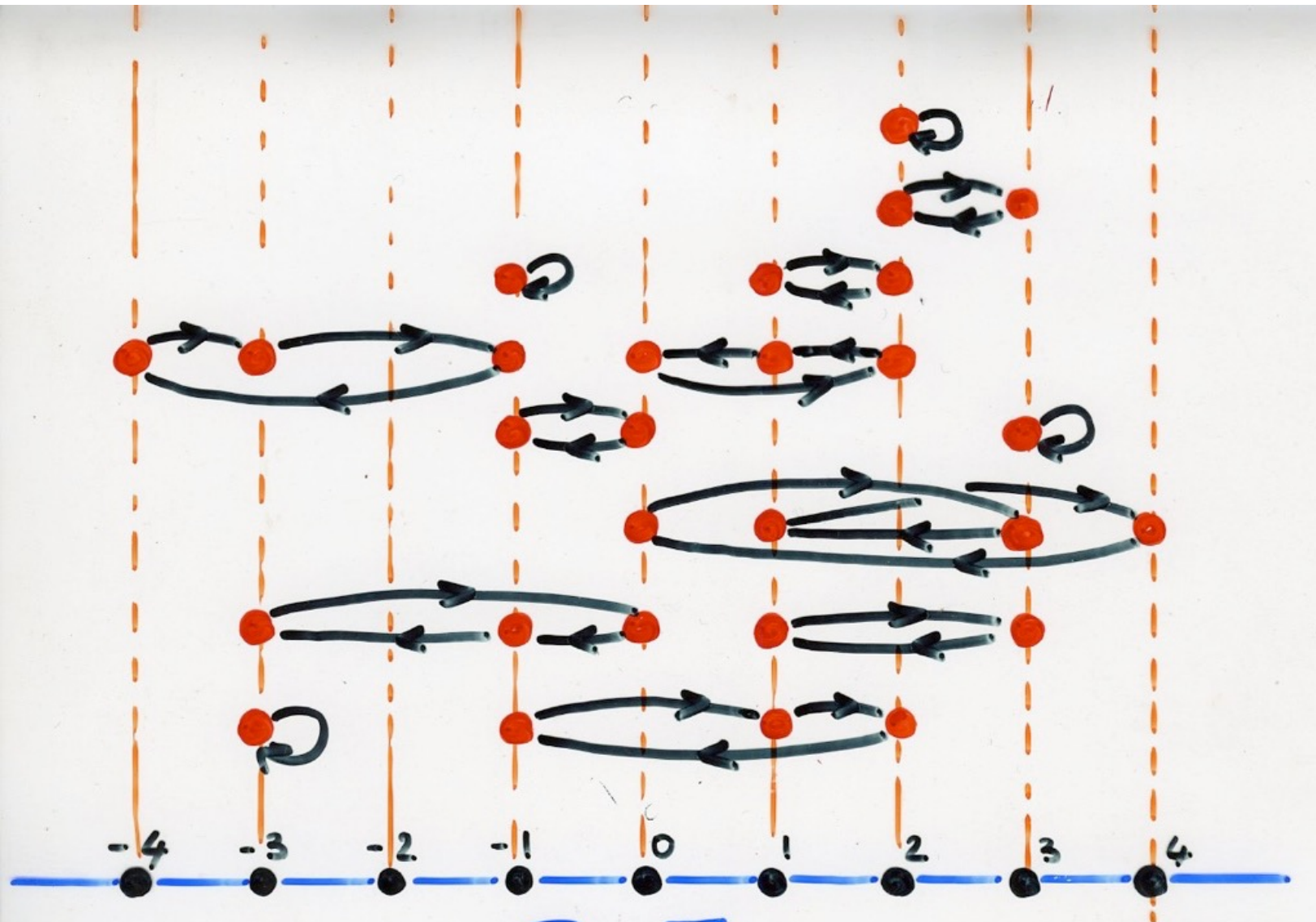
\preceq transitive closure of
the relation $\preceq_{\mathcal{E}}$

$$(\alpha, i) \preceq_{\mathcal{E}} (\beta, j) \iff \alpha \mathcal{E} \beta, i < j$$









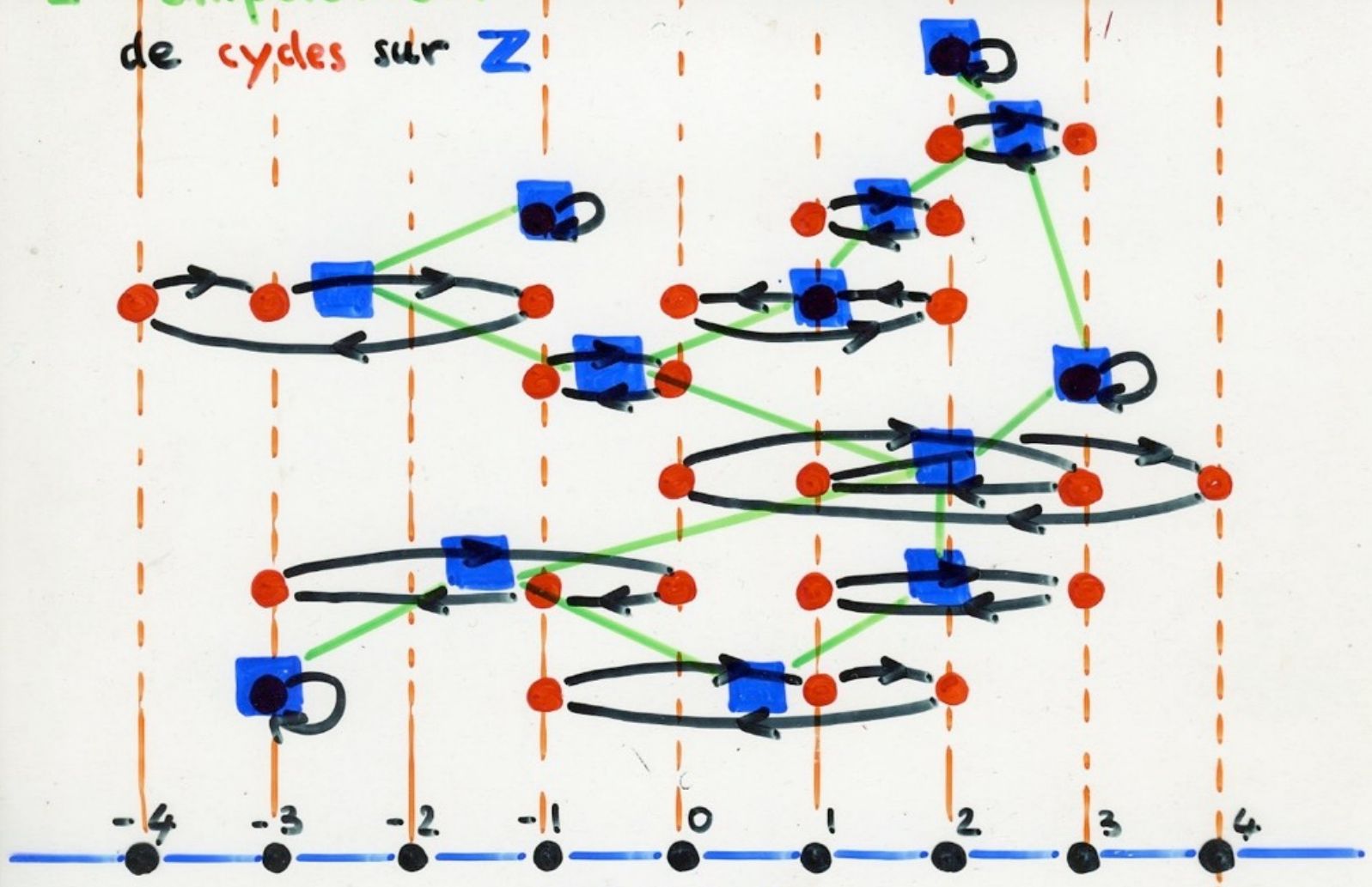
$B = \mathbb{Z}$

P

\mathcal{C}

cycles on \mathbb{Z}
intersection

E empilement
de cycles sur \mathbb{Z}

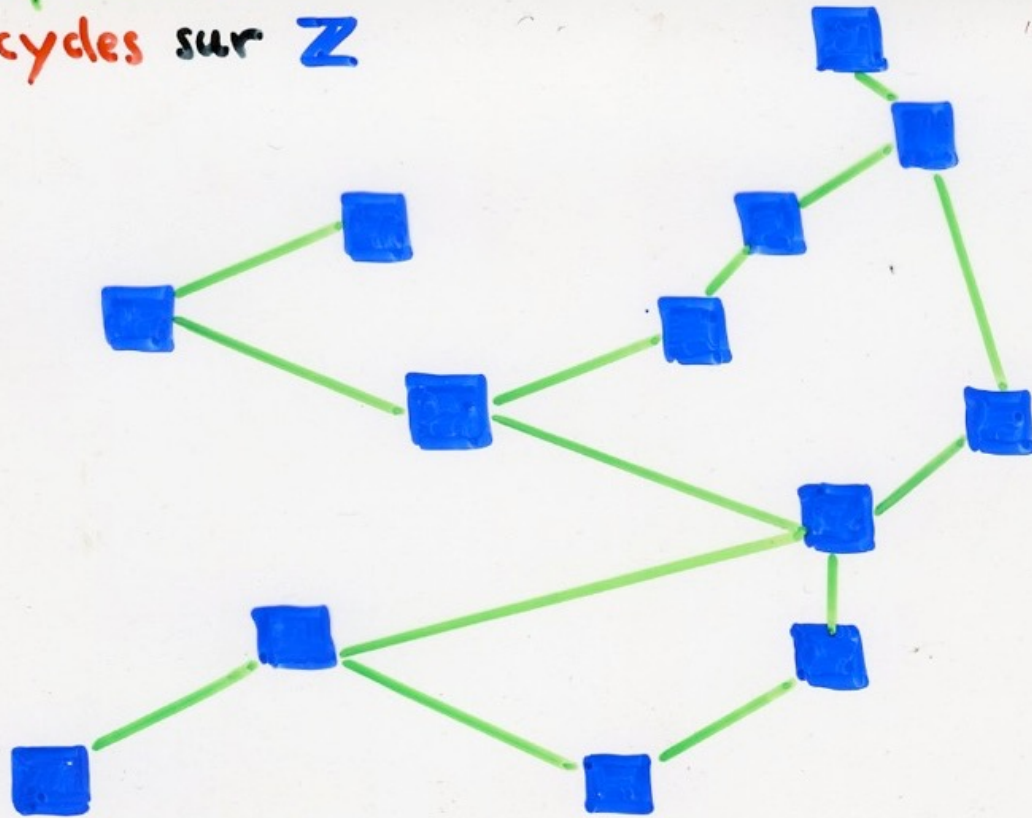


$B = \mathbb{Z}$

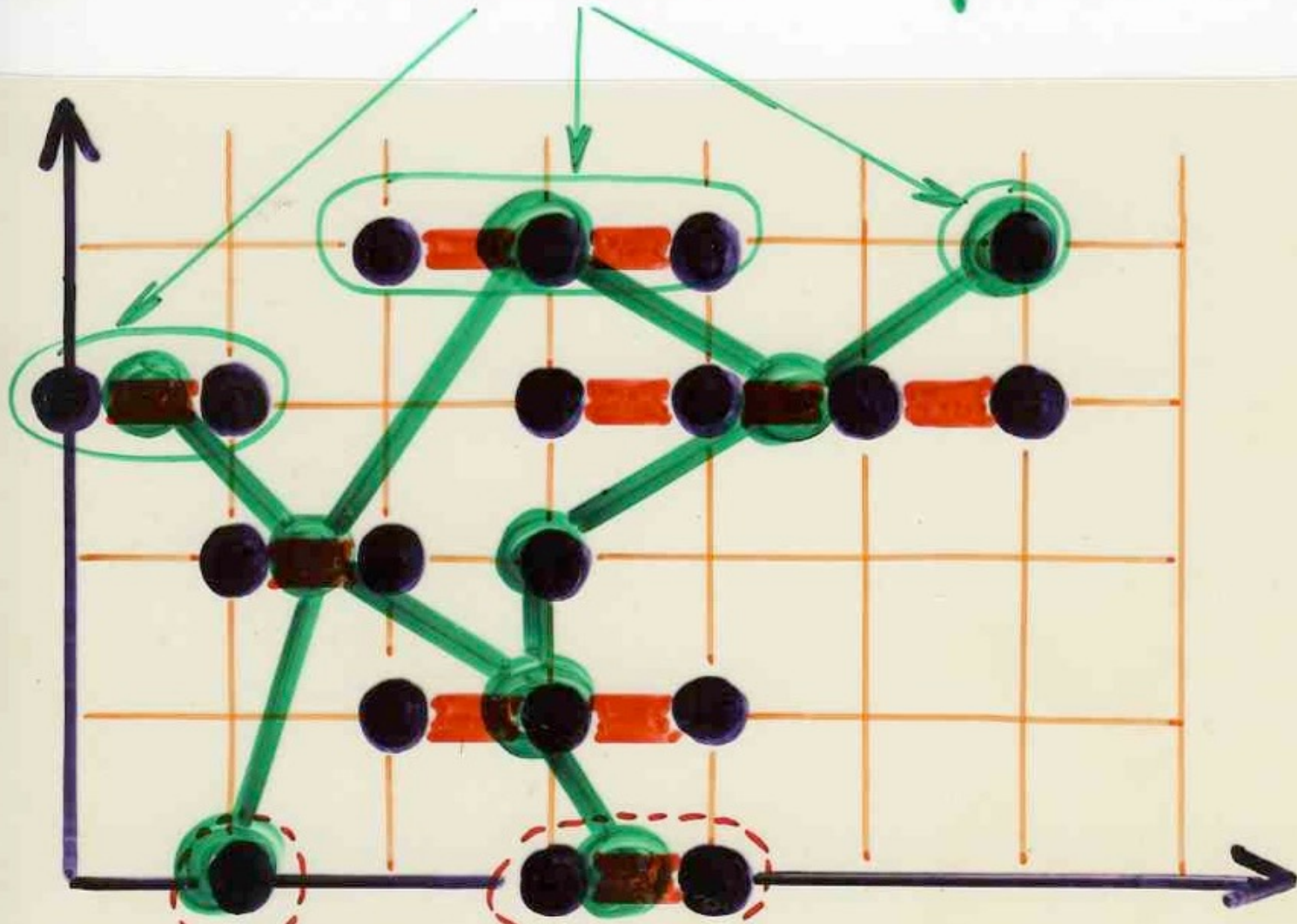
\mathcal{P}
 \mathcal{C}

cycles on \mathbb{Z}
intersection

E empilement
de cycles sur Z



maximal pieces



minimal pieces

minimal

letter of a class $[w]$

$$[w] = [y, y_1]$$

maximal

letter of a class $[w]$

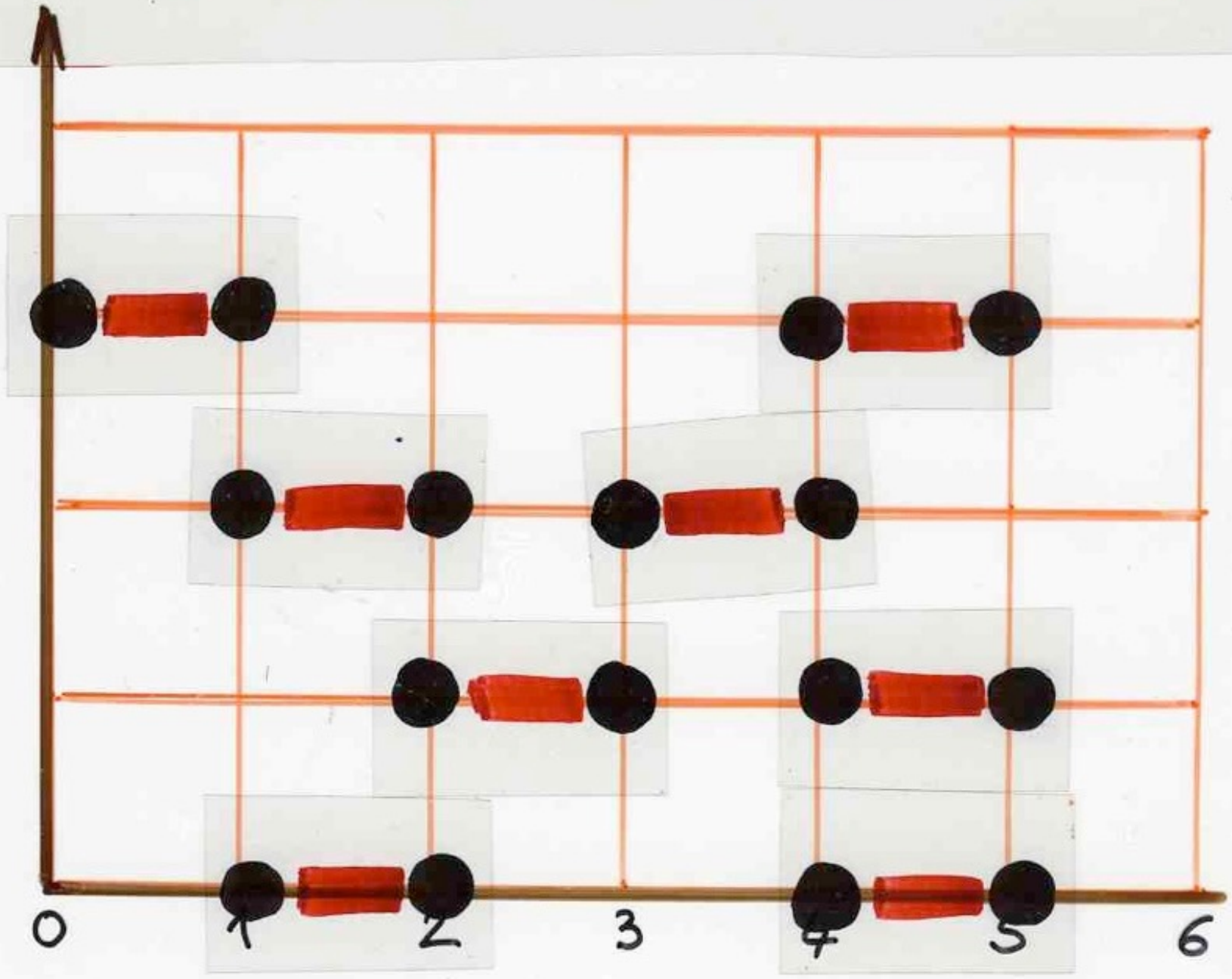
$$[w] = [u_1, z]$$

heaps and linear extensions

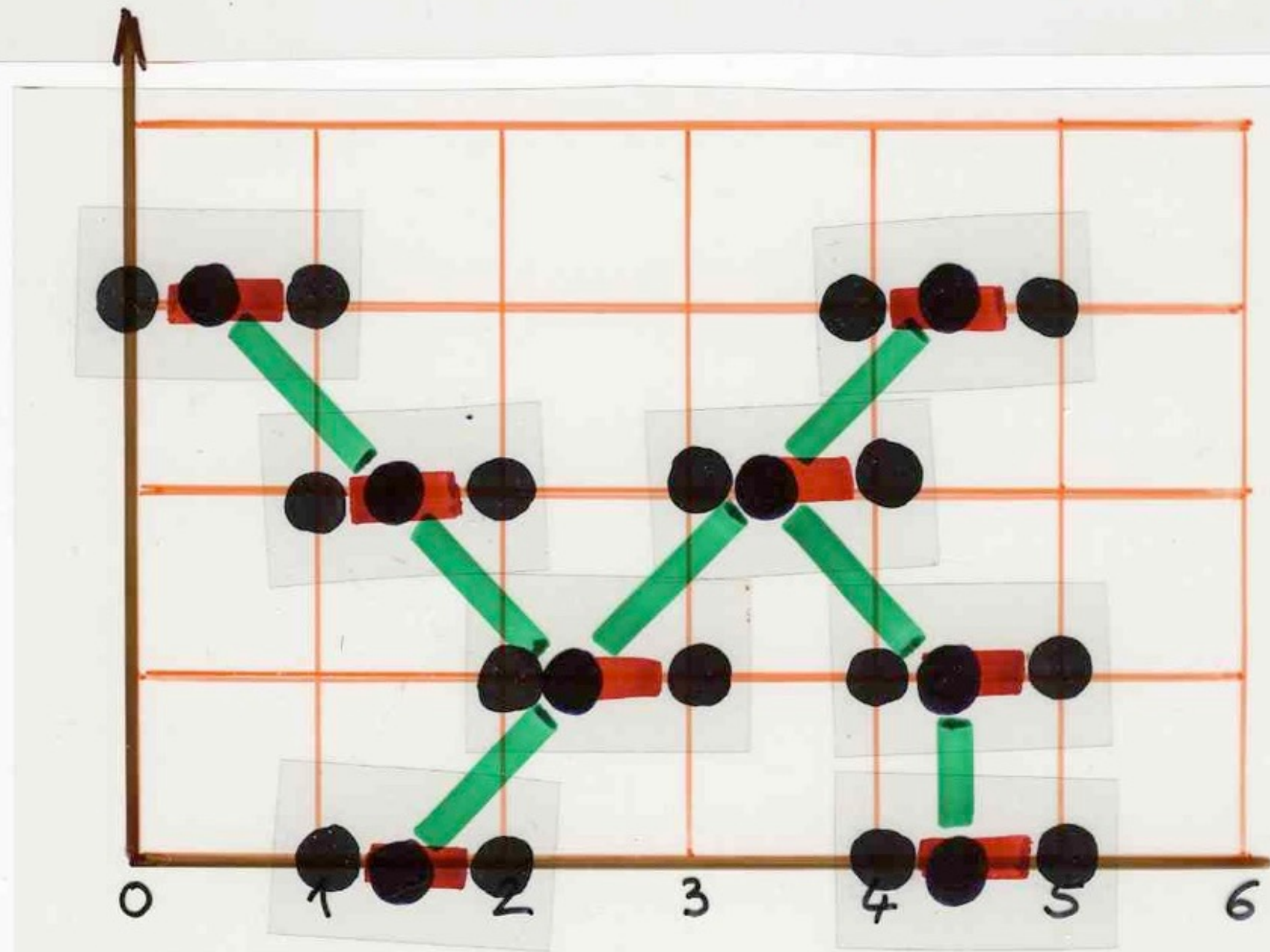
Proposition Let $w \in A^*$ and $[w]$
the equivalence class $[w] \in L(A, C)$.
For $P = A$ and $E = \bar{C}$, let $F \in H(P, E)$
the associated heap $F = \bar{\varphi}(w)$ and
poset (F, \preceq) .

The words of $[w]$ are in bijection
with the linear extensions of (F, \preceq)

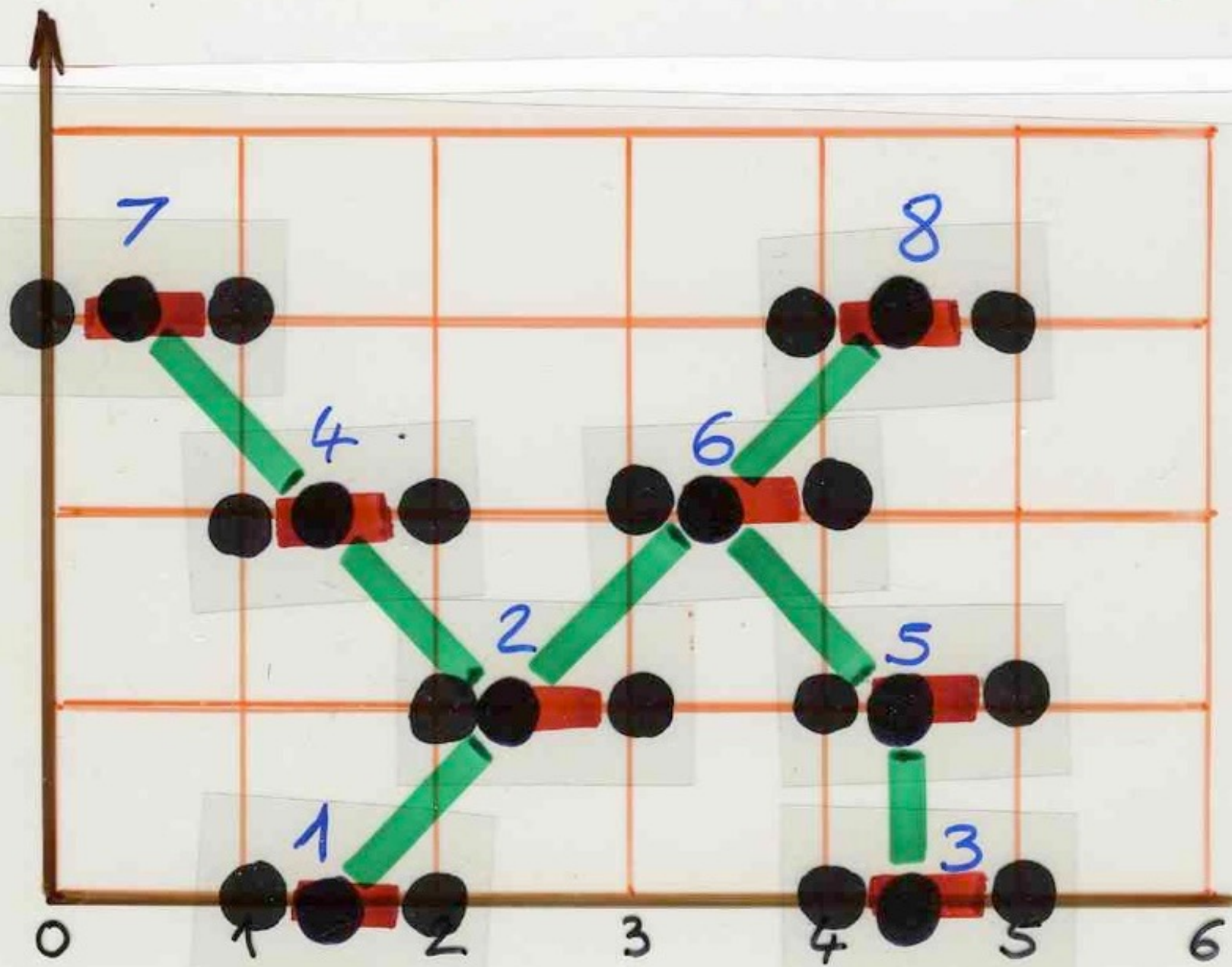
$$W = \sigma_1 \sigma_2 \sigma_4 \sigma_1 \sigma_4 \sigma_3 \sigma_0 \sigma_4$$

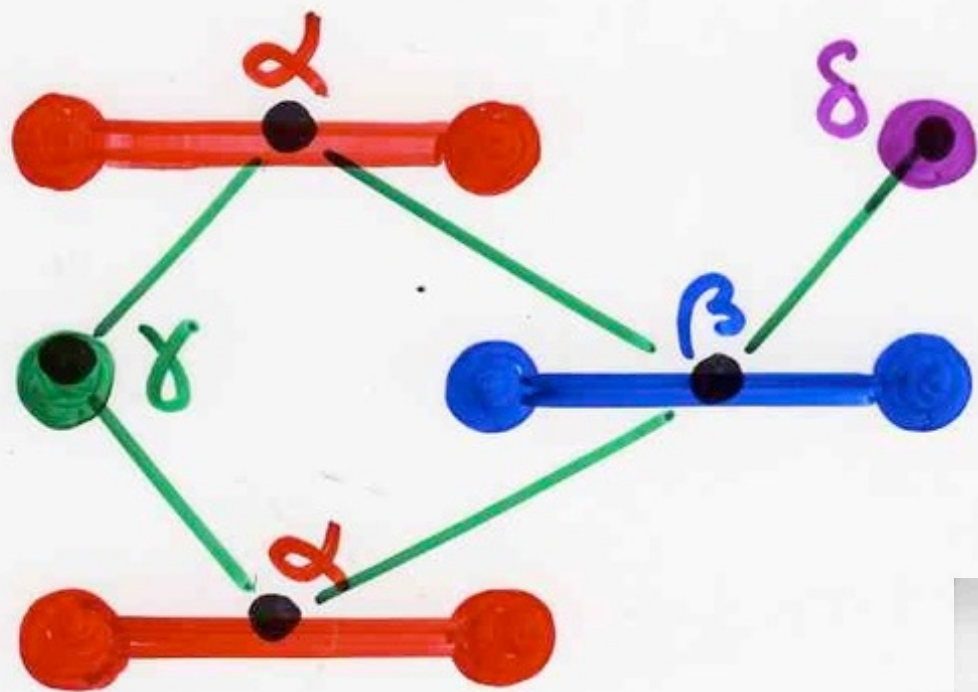


$$W = \sigma_1 \sigma_2 \sigma_4 \sigma_1 \sigma_4 \sigma_3 \sigma_0 \sigma_4$$

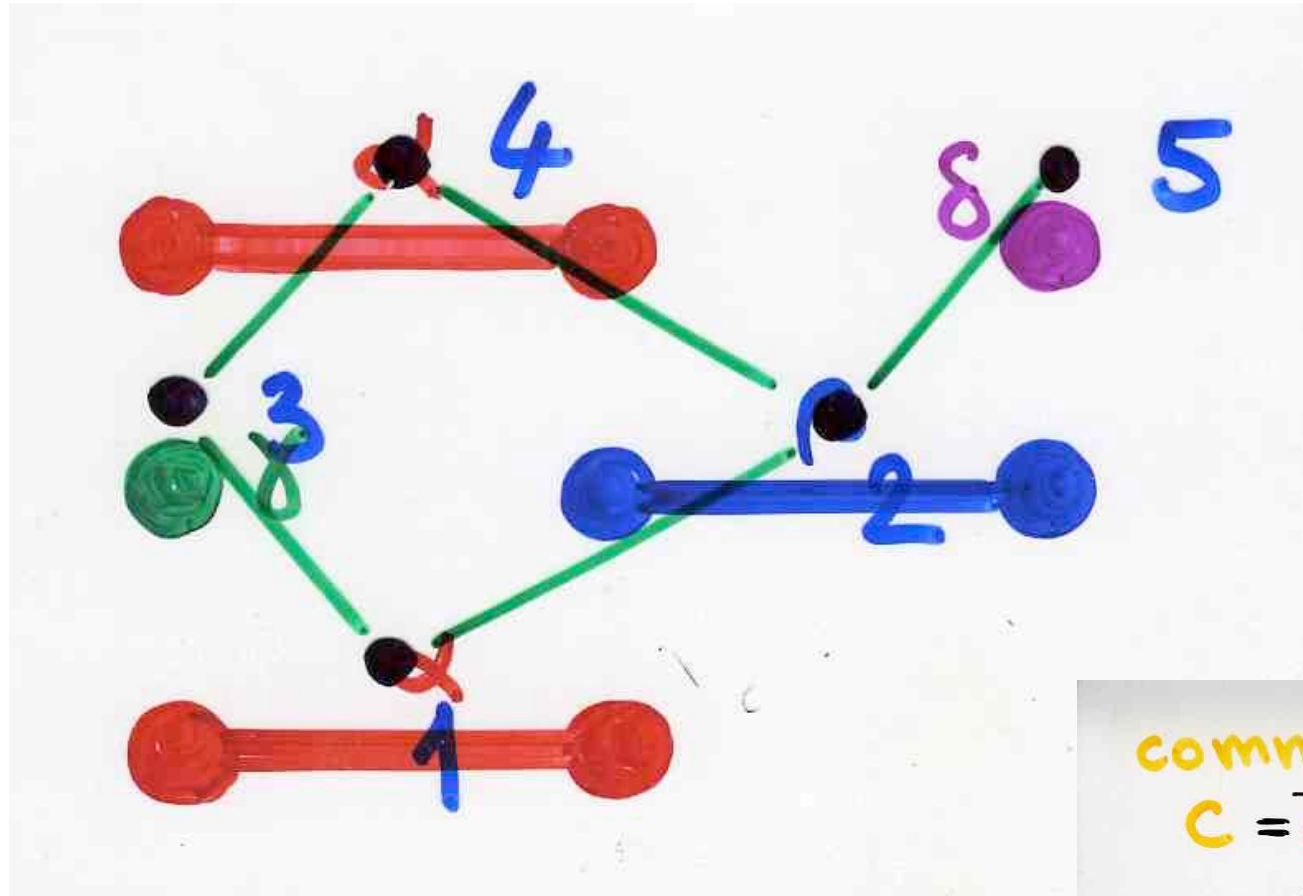


$$W = \rho_1 \rho_2 \rho_4 \rho_1 \rho_4 \rho_3 \rho_0 \rho_4$$



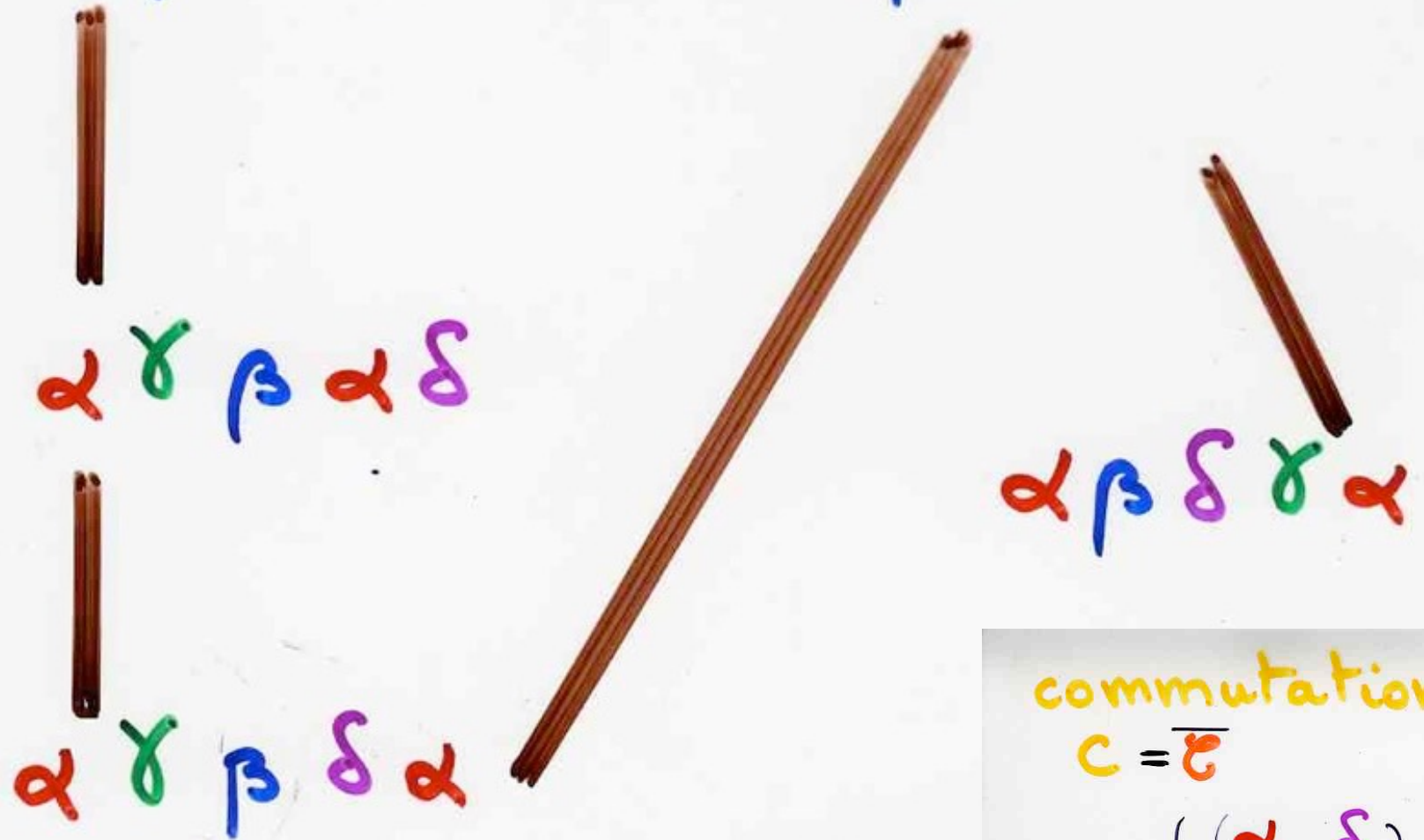


commutations
 $C = \overline{C}$
 $C \left\{ \begin{array}{l} (\alpha, \delta) \\ (\beta, \gamma) \\ (\gamma, \delta) \end{array} \right.$



commutations
 $C = \overline{C}$
 $C \begin{cases} (\alpha, \delta) \\ (\beta, \gamma) \\ (\gamma, \delta) \end{cases}$

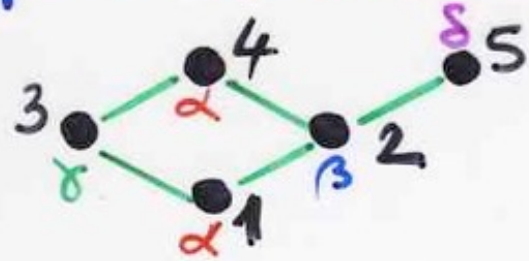
$W = \alpha \beta \gamma \alpha \delta \text{ --- } \alpha \beta \gamma \delta \alpha$



commutations
 $C = \overline{C}$

$C = \left\{ \begin{array}{l} (\alpha, \delta) \\ (\beta, \gamma) \\ (\gamma, \delta) \end{array} \right.$

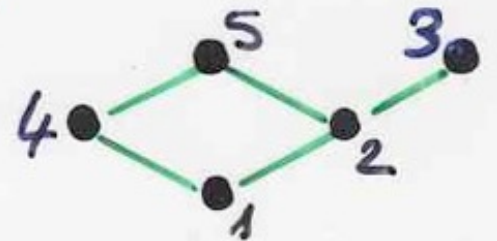
$W = \alpha \beta \gamma \alpha \delta \quad \text{---} \quad \alpha \beta \gamma \delta \alpha$



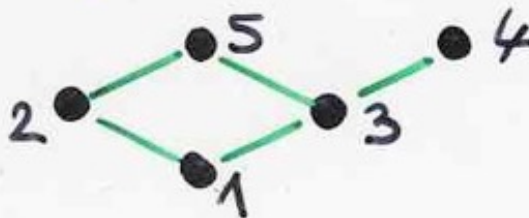
$\alpha \gamma \beta \alpha \delta$

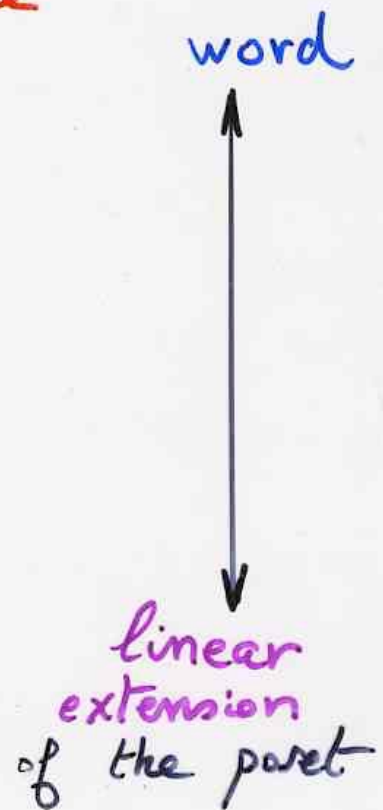
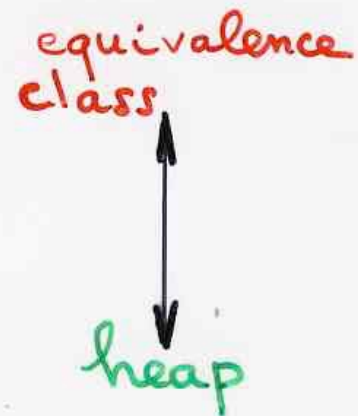
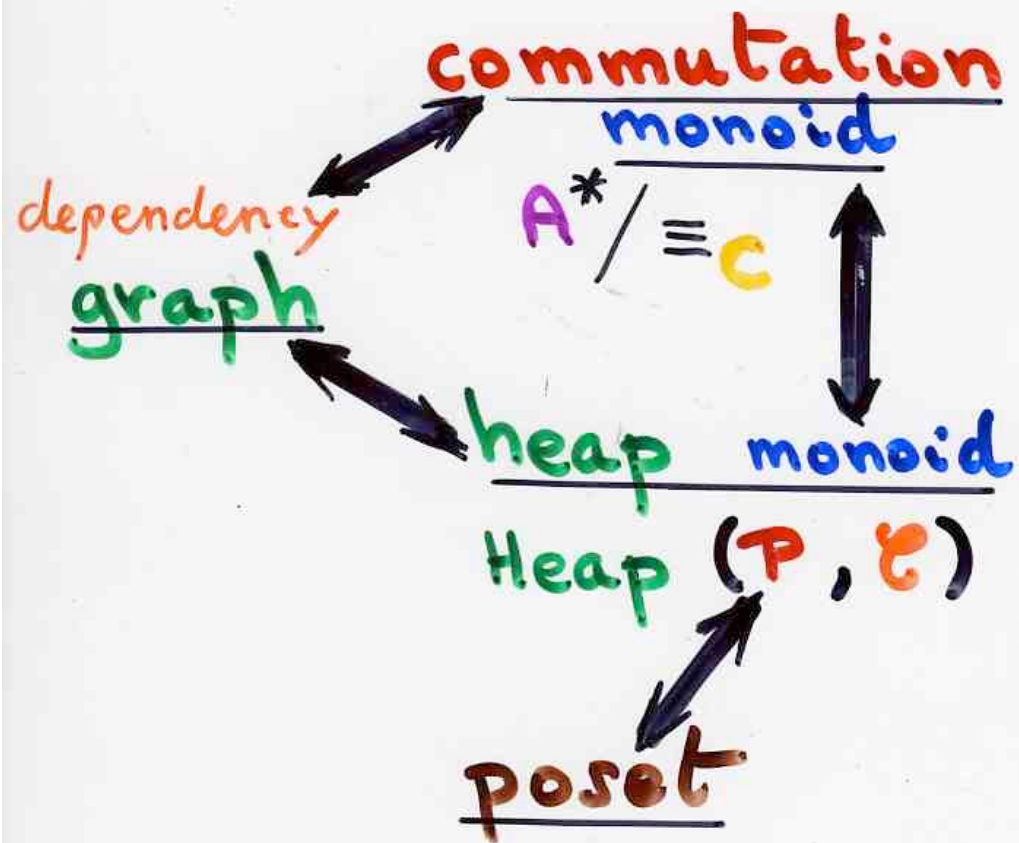


$\alpha \beta \delta \gamma \alpha$



$\alpha \gamma \beta \delta \alpha$





Complements

Heaps, poset and graphs

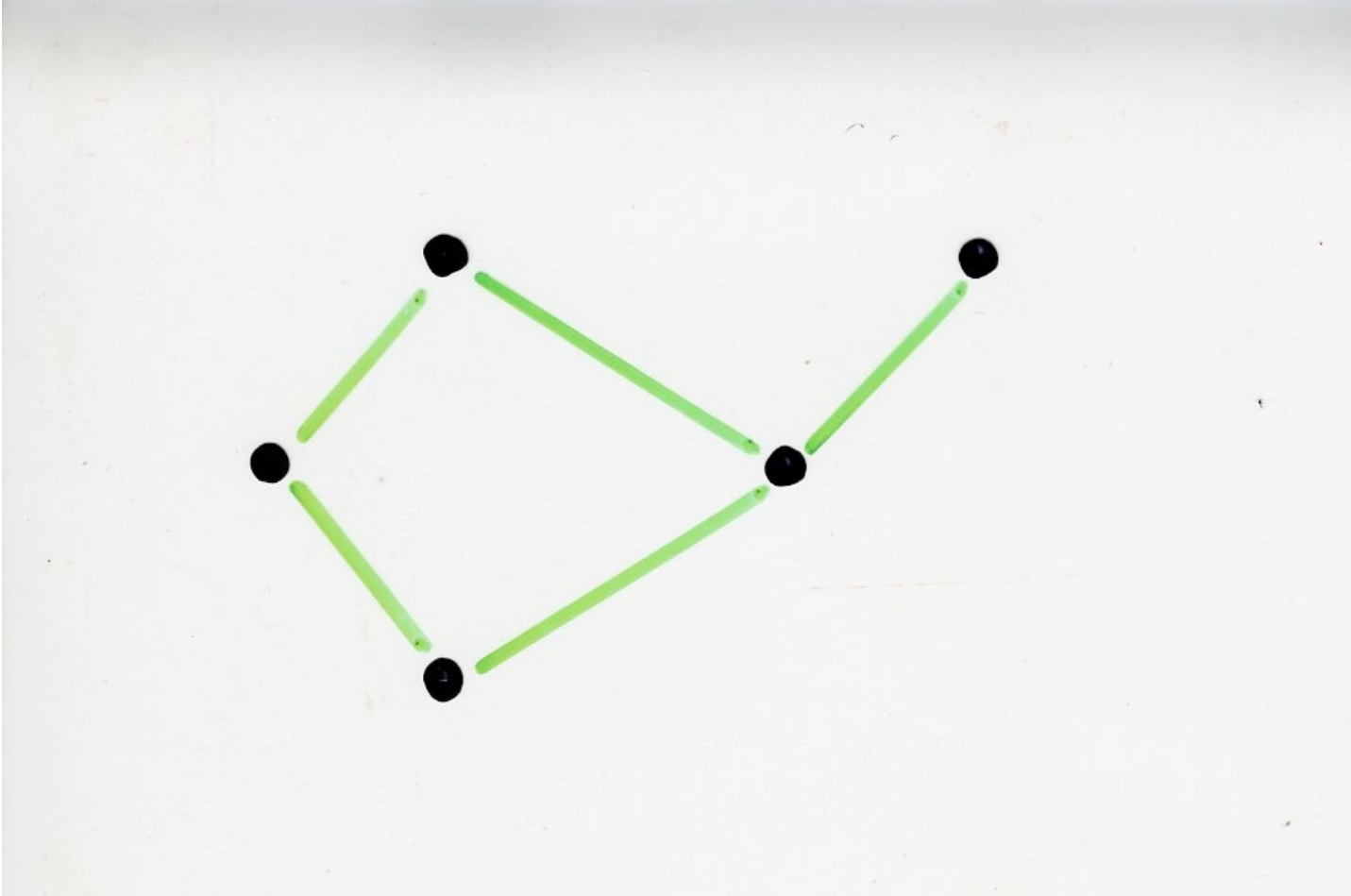
Prop. every poset can be realized as a heap of pieces

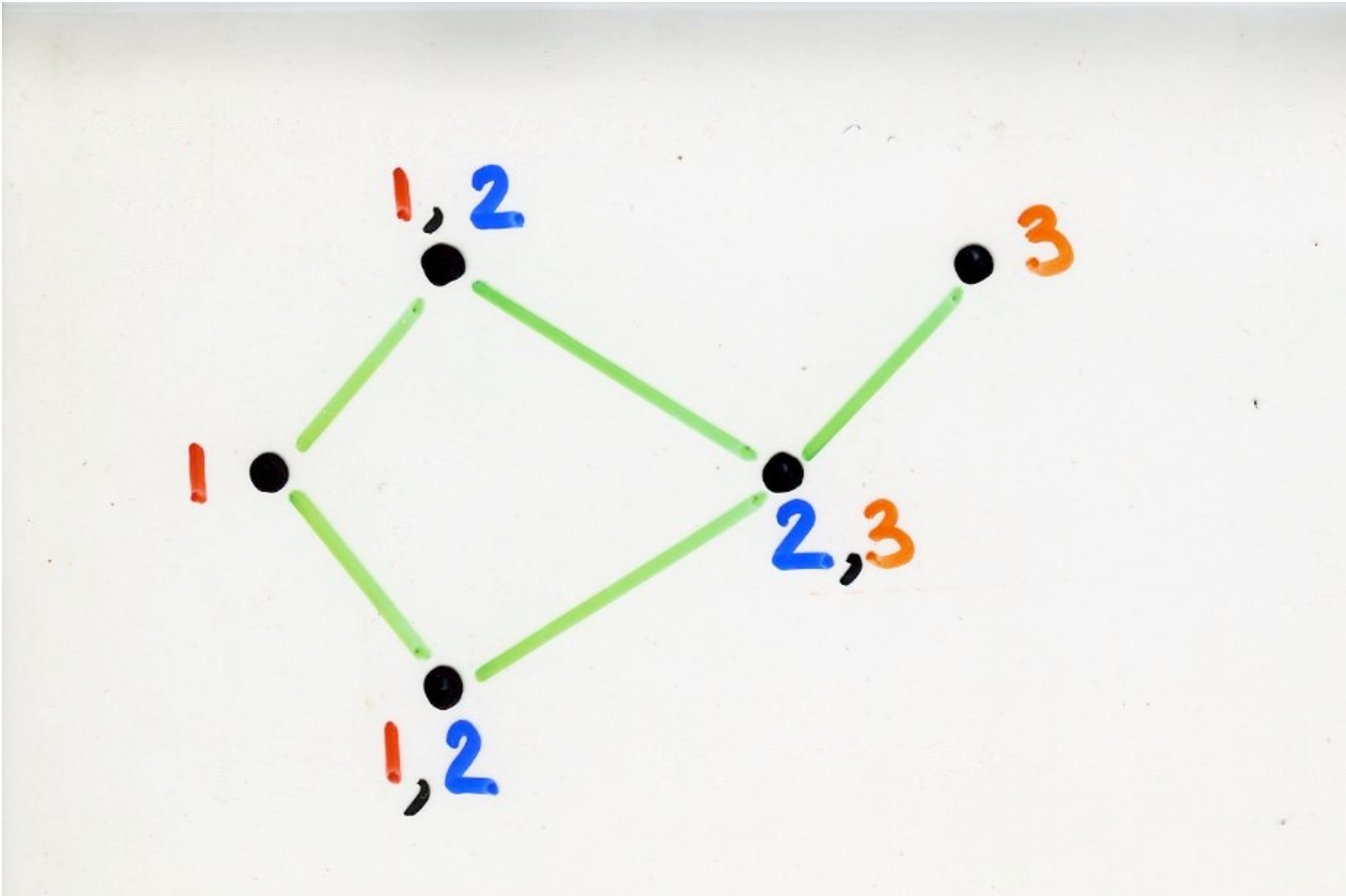
Def E poset

Γ set of chains of E

Γ strongly covers E iff:

$\forall s, t \in E, s \prec t$ and t covers $s, \exists \gamma \in \Gamma$
such that $s, t \in \gamma$





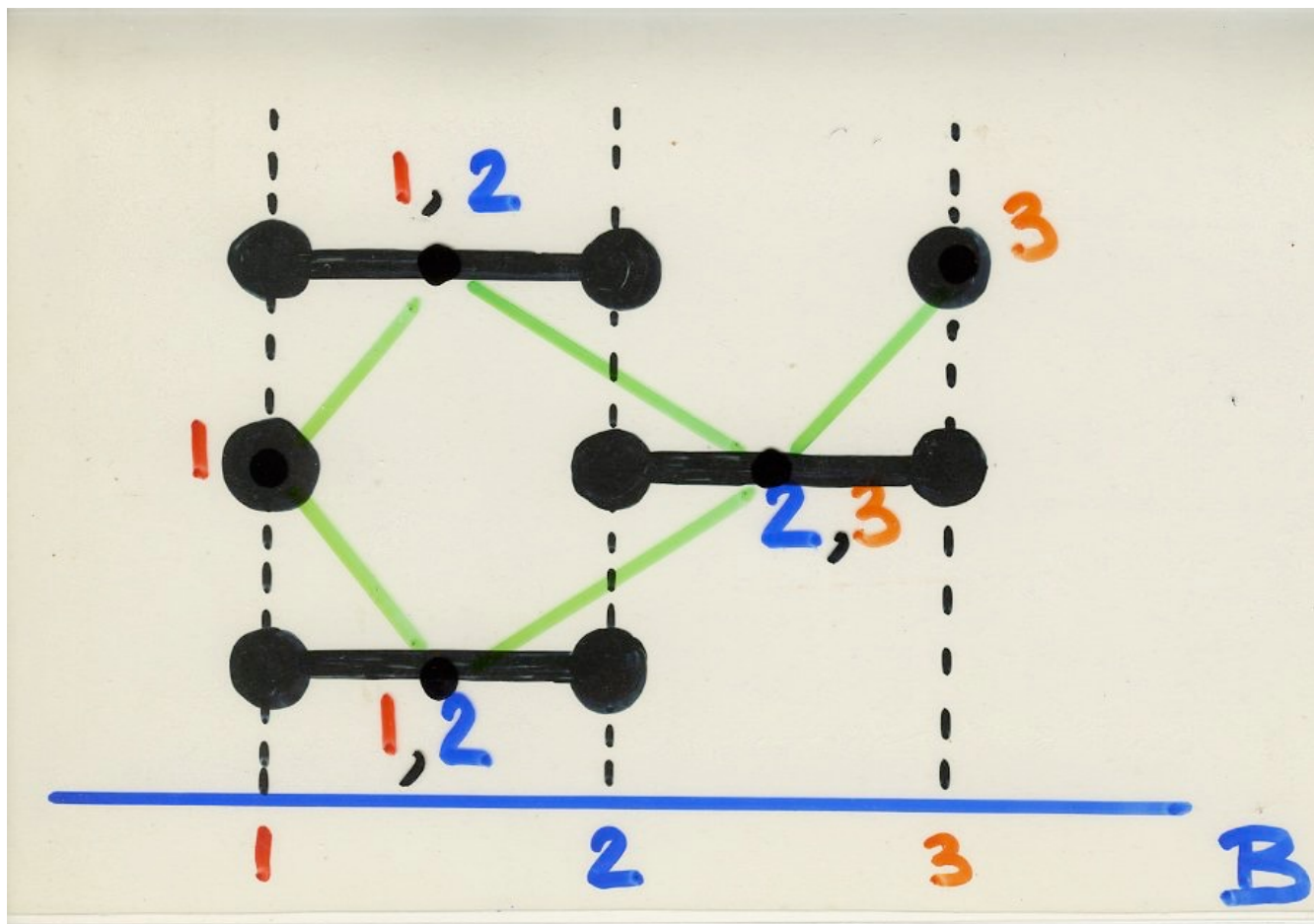
E

base
pieces

$$B = \Gamma$$
$$P = \mathcal{P}(B)$$

$$\pi = Id$$

$$\lambda \in E \rightarrow P_\lambda = \{ \gamma \in \Gamma, \lambda \in \gamma \}$$



commutations

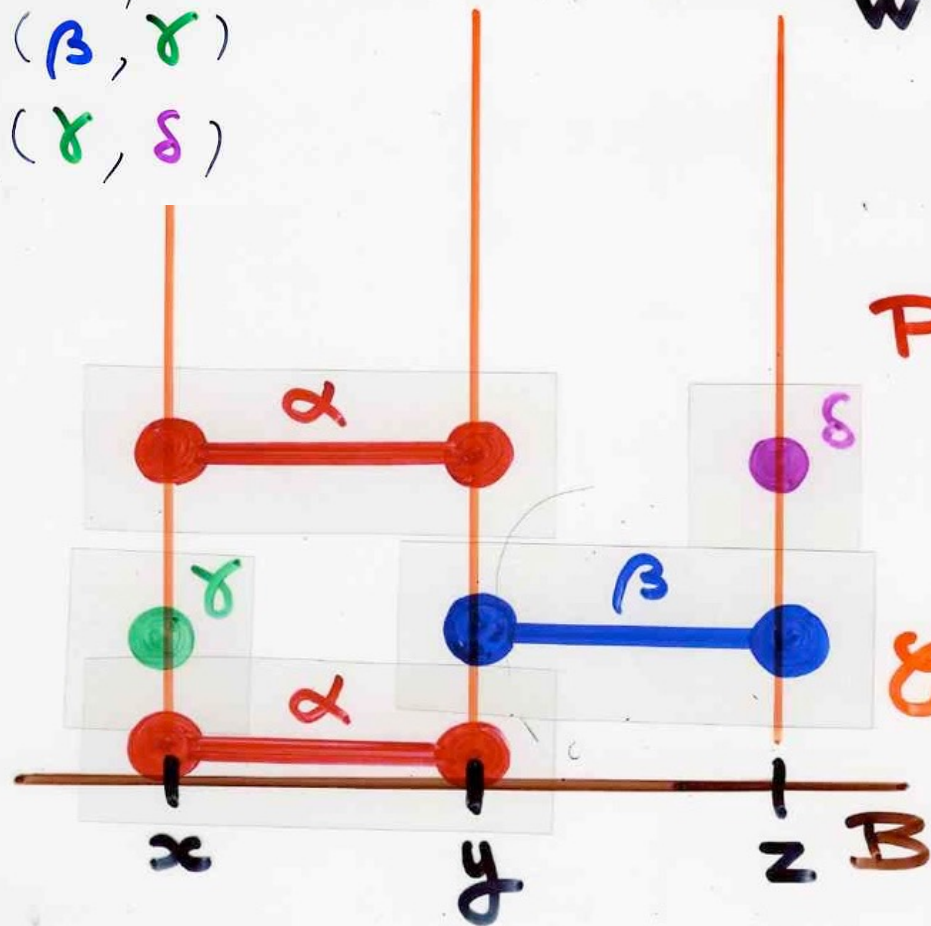
$$C = \overline{C}$$

$$C \begin{cases} (\alpha, \delta) \\ (\beta, \gamma) \\ (\gamma, \delta) \end{cases}$$

$$W = \alpha\beta\gamma\alpha\delta$$

$$P \begin{cases} \alpha = \{x, y\} \\ \beta = \{y, z\} \\ \gamma = \{x\} \\ \delta = \{z\} \end{cases}$$

$$C \begin{cases} (\alpha, \beta) \\ (\alpha, \gamma) \\ (\beta, \delta) \end{cases}$$



Corollary For any poset E , counting the number of linear extensions of E is the same problem as counting the number of words in a commutation class of a commutation monoid.

→ number of Young tableaux
hook-length formula ...

Proposition Every heap monoid is isomorphic to a "heap of subsets of a set X " monoid.

ex: subsets of a set X

• \mathcal{P} set of subsets of X
basic pieces
 $\mathcal{P} \subseteq \mathcal{P}(X)$

• \mathcal{C} dependency relation
 $A, B \in \mathcal{P}, A \mathcal{C} B \Leftrightarrow A \cap B \neq \emptyset$

